

Experiment 5 :

Ex: $E \rightarrow E + T \mid T$
 $T \rightarrow T * F \mid F$
 $F \rightarrow (E) \mid id$

This Grammar contains left recursion because
 $E \rightarrow E$ & $T \rightarrow T$ is there.

But it doesn't contain factoring because
 it doesn't contain $A \rightarrow \alpha_1 A$
 $A \rightarrow \alpha_2 A$
 $B \rightarrow \alpha_1 B$
 $B \rightarrow \alpha_2 B$.

Remove the left recursion by adding E'
 and T' . then the Grammar will become

$E \rightarrow TE'$
 $E' \rightarrow + TE' \mid \epsilon$
 $T \rightarrow FT'$
 $T' \rightarrow * FT' \mid \epsilon$
 $F \rightarrow (E) \mid id$.

first

$first(E) = \{ (, id \}$
 $first(E') = \{ +, \epsilon \}$
 $first(T) = \{ (, id \}$
 $first(T') = \{ *, \epsilon \}$
 $first(F) = \{ (, id \}$

follow

$follow(E) = \{), id, \$ \}$
 $follow(E') = \{), \$ \}$
 $follow(T) = \{ +,), \$ \}$
 $follow(T') = \{ +,), \$ \}$

$follow(F) = \{ *, +,), \$ \}$

parse table

	id	+	*	()	
E	$E \rightarrow TE'$			$E \rightarrow TE'$		
E'		$E' \rightarrow +TE'$			$E' \rightarrow E$	$E' \rightarrow E$
T	$T \rightarrow FT'$			$T \rightarrow FT'$		
T'		$T' \rightarrow E$	$T' \rightarrow *FT'$		$T' \rightarrow E$	$T' \rightarrow E$
F	$F \rightarrow id$			$F \rightarrow (E)$		