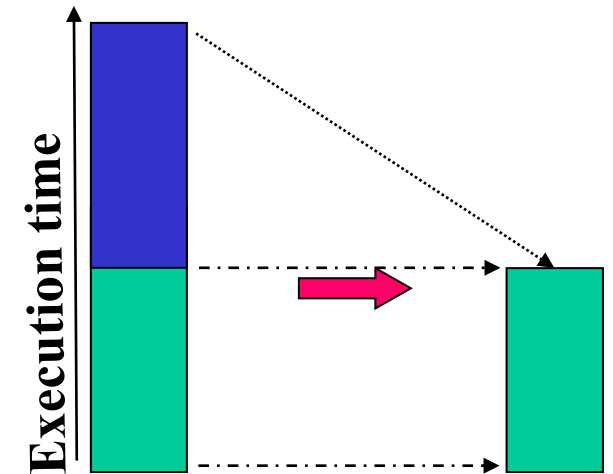


Lecture 4

14-01-2021

Amdahl's Law: Exercise 2

- A software engineer decides to rewrite a portion of a sequential program that accounts for 60% of execution time of the program:



- So that this portion can be run on multiple processors in parallel.
- What is the maximum speedup that the software engineer can hope to achieve?

- $100/40=2.5$

Performance Measurements

- Performance measurement is important:
 - Helps us to determine if one processor (or computer) works faster than another.
 - **Intuitively, a computer has higher performance if it executes programs faster...**

How to Select A Computer System?

- Suppose you are working in a company.
- You have been asked by your manager to select a computer for some specific application:
 - Lets say: to run a web service for your organization...
- How will you proceed?

Purchasing Decision by a Novice...

- Computer A has a 2.5Ghz processor
- Computer B has a 3.2GHz processor
- Which one is faster?



B is naturally faster...

Obviously WRONG!

Where is the fallacy?

Comparing Computer Performance

- We say computer X is n times faster than Y , iff

$$\frac{\text{Execution time}_Y}{\text{Execution time}_X} = n$$

Relative Performance

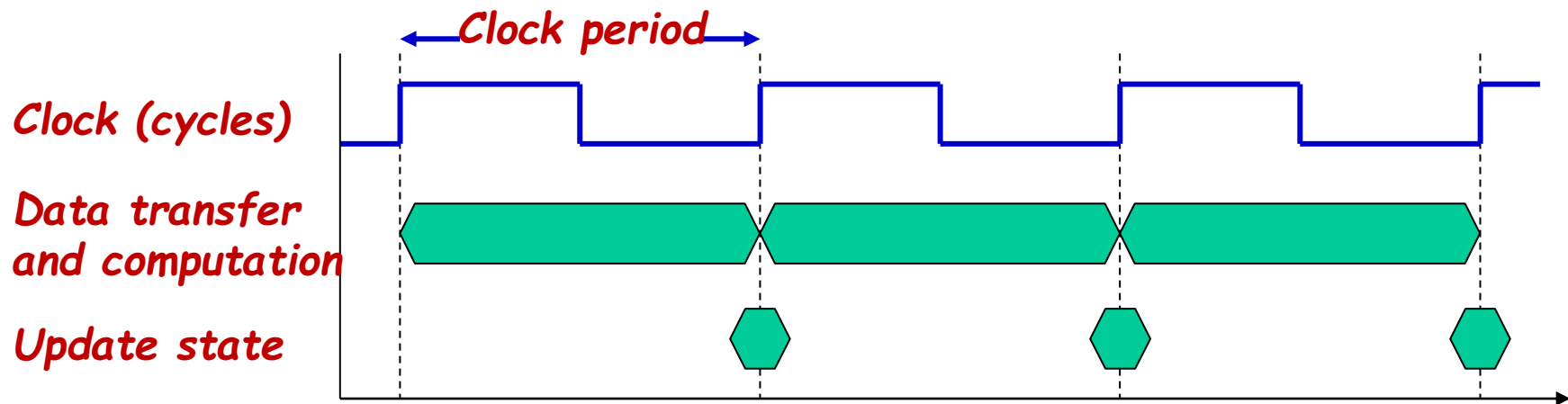
- Define Performance = $1/\text{Execution Time}$
- "X is n time faster than Y"

$$\begin{aligned} & \text{Performance}_X / \text{Performance}_Y \\ &= \text{Execution time}_Y / \text{Execution time}_X = n \end{aligned}$$

- *Example: time taken to run a program*
 - 10s on A, 15s on B
 - $\text{Execution Time}_B / \text{Execution Time}_A$
 $= 15s / 10s = 1.5$
 - So A is 1.5 times faster than B

CPU Clocking

- All operation of computer are governed by a constant-rate clock...



- *Clock period: duration of a clock cycle*
 - *e.g., $250\text{ps} = 0.25\text{ns} = 250 \times 10^{-12}\text{s}$*
- *Clock frequency (rate): cycles per second*
 - *e.g., $4.0\text{GHz} = 4000\text{MHz} = 4.0 \times 10^9\text{Hz}$*

CPU Time

$$\begin{aligned}\text{CPU Time} &= \text{CPU Clock Cycles} \times \text{Clock Cycle Time} \\ &= \frac{\text{CPU Clock Cycles}}{\text{Clock Rate}}\end{aligned}$$

- Performance can be improved by:
 - Reducing number of clock cycles
 - Increasing clock rate
 - Hardware designer must often trade off clock rate against cycle count

CPU Time Exercise

- Computer A: 2 GHz clock, takes 10s CPU time
- Computer B:
 - Completes the job in 6s CPU time
 - Has faster clock, but uses $1.2 \times$ clock cycles
 - What is Computer B's clock rate?

$$\text{Clock Rate}_B = \frac{\text{Clock Cycles}_B}{\text{CPU Time}_B} = \frac{1.2 \times \text{Clock Cycles}_A}{6\text{s}}$$

$$\begin{aligned}\text{Clock Cycles}_A &= \text{CPU Time}_A \times \text{Clock Rate}_A \\ &= 10\text{s} \times 2\text{GHz} = 20 \times 10^9\end{aligned}$$

$$\text{Clock Rate}_B = \frac{1.2 \times 20 \times 10^9}{6\text{s}} = \frac{24 \times 10^9}{6\text{s}} = 4\text{GHz}$$

Instruction Count and CPI

Clock Cycles = Instruction Count \times Cycles per Instruction

CPU Time = Instruction Count \times CPI \times Clock Cycle Time

$$= \frac{\text{Instruction Count} \times \text{CPI}}{\text{Clock Rate}}$$

- Instruction Count (IC) for a program:
 - Determined by program, ISA and compiler
- Average cycles per instruction (CPI):
 - Determined by CPU hardware
 - If different instructions have different CPI
 - Average CPI affected by instruction mix

Exercise

- Computer A: Cycle Time = 250ps, CPI = 2.0
- Computer B: Cycle Time = 500ps, CPI = 1.2
- Same ISA
- Which is faster, and by how much?

$$\begin{aligned}\text{CPU Time}_A &= \text{Instruction Count} \times \text{CPI}_A \times \text{Cycle Time}_A \\ &= 1 \times 2.0 \times 250\text{ps} = 1 \times 500\text{ps}\end{aligned}$$

$$\begin{aligned}\text{CPU Time}_B &= \text{Instruction Count} \times \text{CPI}_B \times \text{Cycle Time}_B \\ &= 1 \times 1.2 \times 500\text{ps} = 1 \times 600\text{ps}\end{aligned}$$

$$\frac{\text{CPU Time}_B}{\text{CPU Time}_A} = \frac{1 \times 600\text{ps}}{1 \times 500\text{ps}} = 1.2$$

*A is
faster...*

*...by this
much*

Exercise

- A given application written in Java runs in 15 seconds on a desktop processor.
 - A new Java compiler is released that requires only 0.6 as many instructions as the old compiler.
 - Unfortunately, it increases the CPI by 1.1.
- How fast can we expect the application to run using this new compiler?

Ans: $15 \text{ sec} \times 0.66 = 9.9 \text{ sec}$

Clock-Rate Based Performance Measurement

- Comparing performance based on clock rates alone, is obviously meaningless:

$$\text{Execution time} = \text{CPI} \times \text{Clock cycle time} \times \text{Number of instructions}$$

– Caveat:

- A processor with a higher clock rate may actually execute programs much slower!

Exercise: Which Machine is Faster and what is the speedup?

- Machine A for prog X

- Clock cycle time = 10ns/cycle
- CPI = 2.0

- Machine B for prog X

- Clock cycle time = 30ns/cycle
- CPI = 0.5

Assume same ISA.

Let I = number of instructions in the program.

$$\begin{aligned}\text{CPU clock cycles (A)} &= I * 2.0 \\ \text{CPU time (A)} &= \text{CPU clock cycles} * \\ &\quad \text{clock cycle time} \\ &= I * 2.0 * 10 \\ &= I * 20 \text{ ns}\end{aligned}$$

$$\begin{aligned}\text{CPU clock cycles (B)} &= I * 0.5 \\ \text{CPU time (B)} &= \text{CPU clock cycles} * \\ &\quad \text{clock cycle time} \\ &= I * 0.5 * 30 \\ &= I * 15 \text{ ns}\end{aligned}$$

$$\frac{\text{Performance (A)}}{\text{Performance (B)}} = \frac{\text{Execution (B)}}{\text{Execution (A)}} = 0.75$$

Exercise: Calculate Overall CPI (Cycles per Instruction)

Operation	Freq	CPI(i)
ALU	40%	1
Load	27%	2
Store	13%	2
Branch	20%	5

Typical Instruction Mix

$$\begin{aligned}\text{Overall CPI} &= 1*0.4 + 2*0.27 + 2*0.13 + 5*0.2 \\ &= 2.2\end{aligned}$$

MIPS and MFLOPS

- Used extensively 30 years back.
- **MIPS**: millions of instructions processed per second.
- **MFLOPS**: Millions of FLoating point OPerations completed per Second

$$\text{MIPS} = \frac{\text{Instruction Count}}{\text{Exec. Time} \times 10^6} = \frac{\text{Clock Rate}}{\text{CPI} \times 10^6}$$

Problems with MIPS

- MIPS suffers from several severe shortcomings.
- So severe, made some one coin:
 - “Meaningless Information about Processing Speed”
- **A Major Problem:**
 - MIPS depends on the program being executed to measure MIPS.

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TESTS, WE'LL SAY THEY
USED OLD DRIVERS.



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WHENEVER I TALK TO
YOU, I FEEL LIKE I
SHOULD BE WEARING
A WIRE.



SINCE WHEN
IS MARKETING
A CRIME?



Exercise 4

- A computer A-100 executes a certain benchmark program with an average CPI of 2.0.
- Another computer A-101 has the same instruction set but has an enhanced compiler.
- The compiler for A-101 generated 20% less instructions for the same benchmark program.
- The compiled code when run on A-101 achieved a CPI of 0.5 at 800MHz.
- In order for the two machines to have the same MIPS rating, what should be the clock rate of the computer A-100?

Exercise 4: Answer

- Let I be the number of instructions in the benchmark program
- A-101 executes $0.8 \cdot I$ instructions in $0.8 \cdot I \cdot \text{CPI}(A-101) \cdot C(A-101)$ seconds
- $= 0.8 \cdot I \cdot 0.5 / (800 \cdot 10^6)$ seconds
- $\text{MIPS} = 800 / 0.5 = 1600$
- $\text{MIPS of A-100} = I \cdot 10^{-6} / (I \cdot 2 \cdot C) = 10^{-6} / (2 \cdot C)$
- Frequency of A-100 = 3200 MHz = 3.2GHz

Toy Benchmarks

- The performance of different computers can be compared by running some standard programs:
 - Quick sort, Merge sort, etc.
- But, the basic problem remains:
 - Even if a computer performs well for a toy benchmark, it may not perform well for the required applications.
 - What can be a solution then?

Synthetic Benchmarks

- **Basic Principle:** Analyze the distribution of instructions over a large number of practical programs.
- **Synthesize a program that has the same instruction distribution as a typical program:**
 - Need not compute something meaningful.
- Dhrystone, Khornerstone, Linpack are some of the older synthetic benchmarks:
 - More recent is **SPEC...**

Dhrystone

- Short synthetic benchmark program
 - Largely uses integer programming.
 - Based on published statistics on use of programming language features;
 - original publication in *CACM* 27,10 (Oct. 1984),
- **Problems:** Due to its small size (100 HLL statements, about 1.5 KB code), **the memory system outside the cache is not tested;**
 - **Compilers can easily optimize for Dhrystone.**

Linpack

- Developed from the "LINPACK" package of linear algebra routines.
 - Originally written in FORTRAN; a C version also exists.
- Almost all of the benchmark's time is spent in a subroutine executing the inner loop for frequent matrix operations:
 - Such as $y(i) = y(i) + a * x(i)$
- The standard version operates on 100x100 matrices;
 - There are also versions for sizes 300x300 and 1000x1000, with different optimization rules.
- **Problems:**
 - Code mostly concerns for matrix computation.
 - LINPACK is easily vectorizable on most systems.

Problems with Benchmarks

- Linpack consists of different types of operations on 300×300 matrices.
 - Optimization of this inner-loop resulted in performance improvement by a factor of 9.
- Optimizing compilers can discard 25% Dhrystone code
- Solution: **Benchmark suite**

SPEC Benchmarks

- **SPEC:** Standard Performance Evaluation Corporation:
 - A non-profit organization (www.spec.org)
 - **Uses a suite of programs**
- CPU-intensive benchmark for evaluating processor performance of workstation:
 - Generations: SPEC89, SPEC92, SPEC95, and SPEC2000 ...
 - Emphasizes memory system performance in SPEC2000.

Other SPEC Benchmarks

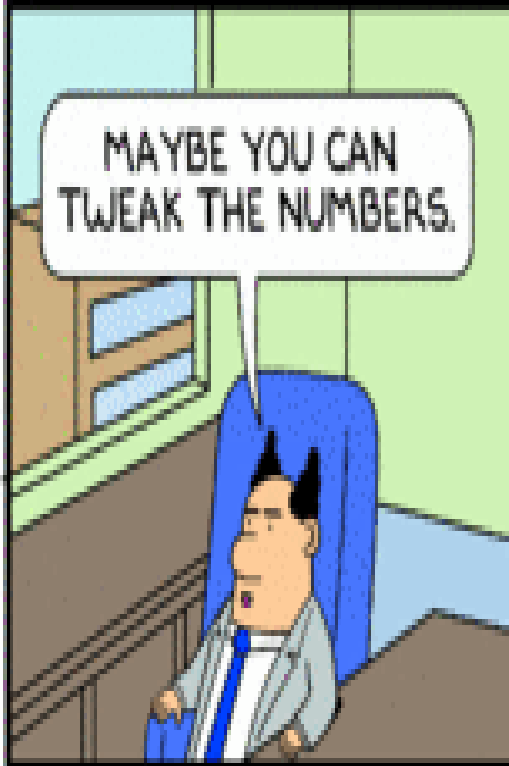
- **SPECviewperf**: 3D graphics performance
 - For applications such as CAD/CAM, visualization, content creations, etc.
- **SPEC JVM98**: performance of client-side Java virtual machine.
- **SPEC JBB2000**: Server-side Java application
- **SPEC WEB2005**: evaluating WWW servers
 - Contains multiple workloads utilizing both http and https, dynamic content implemented in PHP and JSP.

Summary of SPEC Benchmarks

- CPU: CPU2006
- Graphics: SPECviewperf
- Java Client/Server: jAppServer2004
- Mail Servers: MAIL2001
- Network File System: SDS97_R1
- Power (under development)
- Web Servers: WEB2005



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