Tutorial 2

Scheduling

Real Time Scheduling: Rate Monotonic

Q. Verify the schedulability and construct the schedule according to the RM policy for the following set of periodic tasks. Here C_i and T_i are the execution time and periods respectively.

| | C_i | T_i |
|---------|-------|-------|
| $	au_1$ | 2 | 6 |
| $	au_2$ | 2 | 8 |
| $	au_3$ | 2 | 12 |

Ans.

Utilization (U) = 2/6 + 2/8 + 2/12 = 0.75Processor utilization upper bound $U_{max} = n(2^{1/n} - 1) = 0.78$ U < U_{max} The task set is RM schedulable

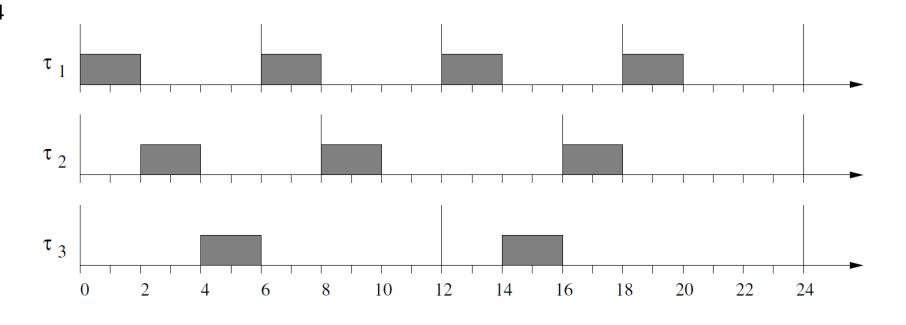
Real Time Scheduling: Rate Monotonic

Q. Write down the RM schedule for the given task set.

Hyper-period = lcm (6,8,12) = 24

| | C_i | T_i |
|---------|-------|-------|
| $	au_1$ | 2 | 6 |
| $	au_2$ | 2 | 8 |
| $	au_3$ | 2 | 12 |





RM Schedule

RM Schedule: $T_1^{\ 1} \ T_2^{\ 1} \ T_3^{\ 1} \ T_1^{\ 2} \ T_2^{\ 2} \ T_1^{\ 3} \ T_3^{\ 2} \ T_2^{\ 3} \ T_1^{\ 4}$

Real Time Scheduling: Rate Monotonic

Consider the following set of tasks

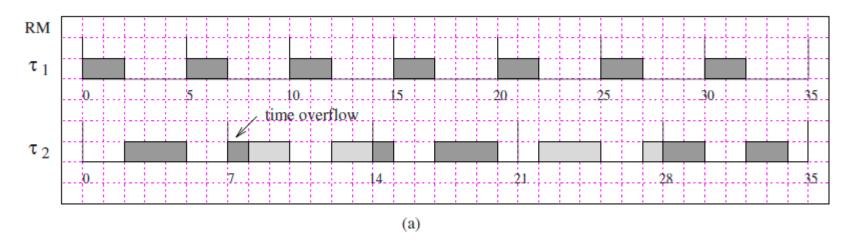
| Tasks | C _i | T _i |
|----------|----------------|----------------|
| τ_1 | 2 | 5 |
| τ_2 | 4 | 7 |

Utilization (U) = 2/5 + 4/7 = 0.97Processor utilization upper bound $U_{max} = n(2^{1/n} - 1) = 0.83$ U > U_{max} The task set may not be RM schedulable

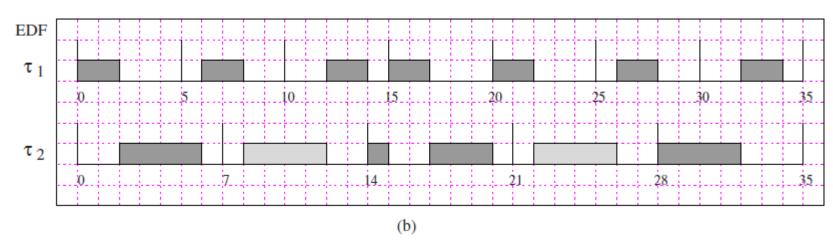
Lets check!!

Real Time Scheduling: Earliest Deadline First (EDF)

Hyper-period = lcm(5, 7) = 35



| Tasks | C _i | T _i |
|----------|----------------|----------------|
| τ_1 | 2 | 5 |
| τ_2 | 4 | 7 |



Exercise - 1

Q. Check if the following task set is RM schedulable? If not, is it EDF schedulable?

| Tasks | Execution Time | Period |
|-------|----------------|--------|
| T1 | 20 | 100 |
| T2 | 30 | 150 |
| Т3 | 90 | 200 |

Real Time Scheduling

Q. Check if the following task set is RM schedulable?

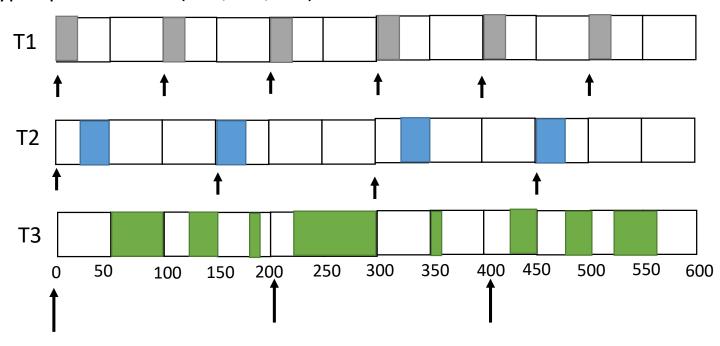
| Tasks | Execution Time | Period |
|-------|----------------|--------|
| T1 | 20 | 100 |
| T2 | 30 | 150 |
| Т3 | 90 | 200 |

Ans.

Utilization (U) = 20/100 + 30/150 + 90/200 = 0.85Processor utilization upper bound $U_{max} = n(2^{1/n} - 1) = 0.78$ U > $U_{max} \rightarrow RM$ schedulable may not be feasible

Real Time Scheduling

Hyper-period = LCM(100,150,200) = 600



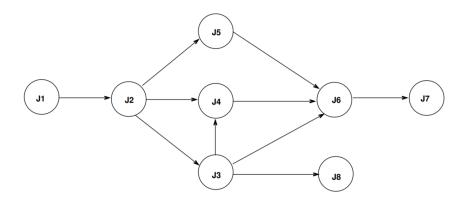
| Tasks | Execution Time | Period |
|-------|-------------------|--------|
| T1 | 20 | 100 |
| T2 | 30 | 150 |
| T3 | 90 | 200 |

RM Schedule: T1¹ T2¹ T3¹ T1² T3¹ T2² T3¹ T1³ T3² T1⁴ T2³ T3² T1⁵ T3³ T2⁴ T3³ T1⁶ T3³

LDF Scheduling

Given the precedence graph in following figure and the following table of task execution times (C_i) and deadlines (D_i) , determine a Latest Deadline First (LDF) schedule.

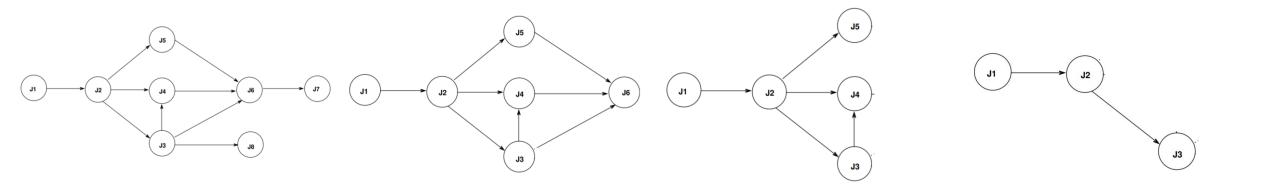
| | J_1 | J_2 | J_3 | J_4 | J_5 | J_6 | J_7 | J_8 |
|-------|-------|-------|-------|-------|-------|-------|-------|-------|
| C_i | 3 | 4 | 2 | 3 | 3 | 2 | 2 | 1 |
| D_i | 5 | 8 | 11 | 15 | 12 | 18 | 19 | 20 |



Schedule: J_1 J_2 J_3 J_5 J_4 J_6 J_7 J_8

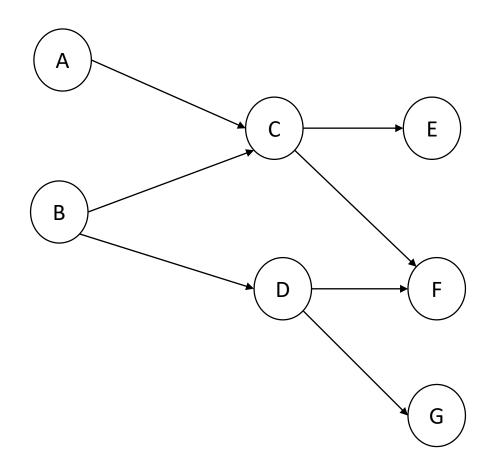
LDF Scheduling

| | J_1 | J_2 | J_3 | J_4 | J_5 | J_6 | J_7 | J_8 |
|-------|-------|-------|-------|-------|-------|-------|-------|-------|
| C_i | 3 | 4 | 2 | 3 | 3 | 2 | 2 | 1 |
| D_i | 5 | 8 | 11 | 15 | 12 | 18 | 19 | 20 |



Schedule: J₁ J₂ J₃ J₅ J₄ J₆ J₇ J₈

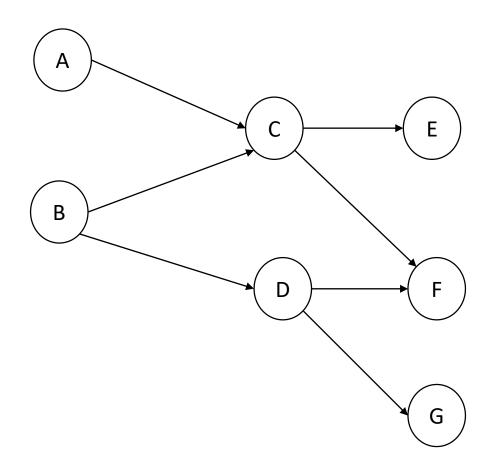
EDF* Scheduling



All tasks arrive at t=0. They all have deadline d=20. Their execution times are given below. Determine the EDF* schedule.

| | Α | В | С | D | Е | F | G |
|-------|---|---|---|---|---|---|---|
| C_i | 3 | 2 | 4 | 3 | 2 | 5 | 1 |

EDF* Scheduling



| | Α | В | С | D | Е | F | G |
|-------|---|---|---|---|---|---|---|
| C_i | 3 | 2 | 4 | 3 | 2 | 5 | 1 |

$$d'_i = \min(d_i, \min_{j \in D(i)} (d'_j - e_j)).$$

$$d'_{E} = 20$$

$$d'_{F} = 20$$

$$d'_{G} = 20$$

$$d'_{C} = min(20, 20-2, 20-5) = 15$$

$$d'_D = min(20, 20-5, 20-1) = 15$$

$$d'_{B} = min(20, 15-4, 15-3) = 11$$

$$d'_A = min(20, 15-4) = 11$$

EDF* schedule: A,B,C,D,E,F,G

Q. Consider two tasks to be scheduled periodically on a single processor using Rate Monotonic (RM) scheduling policy. Task T1 has periodicity p1 = 4 and task T2 has periodicity p2 = 6. If execution time of T1 is e1 = 1

a. what will the execution time of T2 to get a near maximum processor utilization?

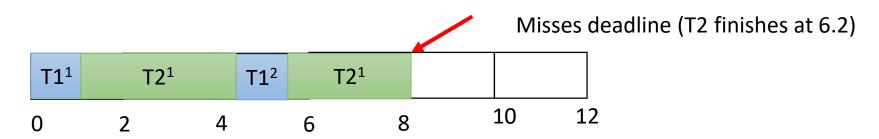
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Sol.
$$\frac{e1}{p1} + \frac{e2}{p2} \le 1 \Rightarrow \frac{1}{4} + \frac{e2}{6} \le 1 \Rightarrow e2 \le 4.5$$

If e2 = 4.5 => Not RM schedulable (Check)

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- Sol. Not RM schedulable



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- c. Had we used Earliest Deadline First (EDF) scheduling policy instead of RM with the same set-up, how would context switch have affected the schedule?

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Sol. EDF schedulable

