

Problem 0:

Implement the fibonacci sequence

$x = \text{fib}(n)$

if $n == 0$
return 0

if $n == 1$
return 1

return $\text{fib}(n-1) + \text{fib}(n-2)$

* To get $\text{fib}(5)$

The first recursive call is $\text{fib}(4)$

The second recursive call is $\text{fib}(3)$

* For $\text{fib}(4)$ which is received from $\text{fib}(5)$

The first rec-call is $\text{fib}(3)$

The second rec-call is $\text{fib}(2)$

* $\text{fib}(3)$,

The first rec-call is $\text{fib}(2)$

The second rec-call is $\text{fib}(1)$

* $\text{fib}(2)$

The first rec-call is $\text{fib}(1)$

The second call is $\text{fib}(0)$

* For $\text{fib}(1)$ - It is taken from $\text{fib}(2)$

$\text{fib}(1)$ returns 1.

* $\text{fib}(0)$ also derived from $\text{fib}(2)$

$\text{fib}(0)$ will return 0.

Hence, $\text{fib}(1)$ returns 1

$\text{fib}(0)$ returns 0.

$\text{fib}(2) = \text{fib}(0) + \text{fib}(1)$

$\text{fib}(2) = 0 + 1 = 1$

$\Rightarrow \text{fib}(3) = \text{fib}(2) + \text{fib}(1)$

$= 1 + 1 = 2$

$\text{fib}(4) = \text{fib}(3) + \text{fib}(2)$

$= 2 + 1 = 3$

$\text{fib}(5) = \text{fib}(4) + \text{fib}(3) = 3 + 2 = 5$

Time Complexity of fibonacci:

The time complexity is $O(2^n)$ and it is exponential.

This is exponential ~~because~~ ^{because} it recursively generates and calls two or more recursive calls. Since it is exponential it is inefficient for larger values of n .

Ways to improve fibonacci:

Memoization:-

A method to optimize the algorithm is by using Memoization, which remembers already calculated values. and this avoids redundant computations and reduces time complexity to $O(n)$.

Iterative Approach:-

In this method fibonacci numbers are computed using a loop, then recursion. This gives $O(n)$ time complexity with $O(1)$ space complexity.

Problem-1

Time Complexity:-

To take elements and store it in a single array, we can assume time complexity of $O(N \times k)$, then have sorted the array based on merge sort. Total over all time complexity is $O(NK) + O(NK \log N)$
 $= O(NK \log(NK))$.

Improved Approach:-

Since each array is sorted could have used a min heap that would return smallest element in constant time.

Since the array is sorted we can take one element from each array and compare the first element of the array. After finding the smallest element we must take next element from the same array & then compare again.

Worst case time complexity is $O(Nk \log k)$.

Problem - 2

Time Complexity:

Array is traversal is done Once so time complexity is $O(n)$.

Improved Approach:

↳ This itself is the efficient approach as both time & space complexity is efficient.