

# Unit 2

Difference between Array and linked list :-

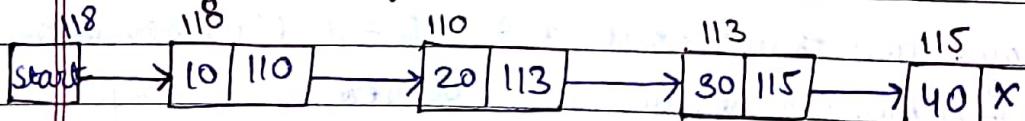
Array	Linked List
1. Array is a collection of elements having same data type with common name.	1. Linked list is an ordered collection of elements which are connected by links / pointers.
2. In array, elements can be accessed using index / sub-script value, i.e. elements can be randomly accessed like arr[0], arr[1] etc. So array provides fast and random access	2. In linked list, element can't be accessed randomly but can be accessed only sequentially and accessing element takes $O(n)$ time.
3. In array, elements are stored in consecutive manner in memory.	3. In linked list, elements can be stored at any available place as address of node is stored in previous node.
4. Insertion & deletion takes more time in array as elements are stored in consecutive memory locations.	4. Insertion & deletion are fast and easy in linked list as only value of pointer is needed to change.

## Array

## Linked List

- |  |   |
|--|---|
| 5. Array can be single dimensional, two dimensional or multidimensional.                           | 5. Linked list can be singly doubly or circular linked list.  |
| 6. In array, each element is dependent, no connection with previous element or with its location.  | 6. In linked list, location or address of element is stored in the link part of previous element's node.  |
| 7. In array, no pointer are used like linked list so no need of extra space in memory for pointer. | 7. In linked list, adjacency between the elements are maintained using pointers or links, so pointers are used and for that extra memory space is needed. |

## Structure of linked list :-



struct node

{

int info;

struct node \*link;

};

Struct node \* start;

**Linked list:-** Linked list is a collection of nodes. There are two parts in single node. First part contain information and second part

contain link of the next node. In last node link part ~~contain~~ ~~will~~ contain NULL value and first node contain START.

`struct node *P = (struct node *) malloc ( size of (struct node));`

## Operations

### 1. Traversing in a linked list :-

Algo

1. Set `Ptr := start;`
2. Repeat step 3 & 4 while `Ptr != NULL` do
  3. Apply PROCESS to `info [Ptr]`
  4. Set `Ptr := link [Ptr];`
5. Exit

### 2. Counting in a linked list :-

Algo.

1. Set `Ptr := start and count := 0.`
2. Repeat step 3 to 4 while `Ptr != NULL` do
  3. `count := count + 1`
  4. Set `Ptr := link [Ptr];`
5. Print count
6. Exit.

### 3. Searching in a linked list :-

Algo

1. Set `Ptr := start`
2. Repeat 3 while `Ptr != NULL` do and Algo
  3. If `info [Ptr] = ELE`  
set `Loc := Ptr` and go to step 5.

- else  
set  $\text{ptr} := \text{link}[\text{ptr}]$
4. Set  $\text{LOC} := \text{NULL}$
  5. Print  $\text{LOC}$
  6. Exit

#### 4. Searching in sorted linked list

Algo

1. Set  $\text{ptr} := \text{start}$
2. Repeat step 3 while  $\text{ptr} \neq \text{NULL}$ 
  3. If  $\text{ele} > \text{info}[\text{ptr}]$  then  
set  $\text{ptr} := \text{link}[\text{ptr}]$
  - else  
if  $\text{ele} = \text{info}[\text{ptr}]$  then  
set  $\text{LOC} := \text{ptr}$  and exit
  - else  
 $\text{LOC} := \text{NULL}$  and exit
4. Set  $\text{LOC} := \text{NULL}$
5. Exit.

#### 5. Insertion in linked list :-

- (i) at the beginning of linked list:-

Algo.

1. If  $\text{Avail} = \text{NULL}$  then  
write: Memory overflow and Exit
2. Set  $\text{new} := \text{Avail}$  and  $\text{Avail} := \text{link}[\text{Avail}]$
3. Set  $\text{info}[\text{new}] := \text{ele}$
4.  $\text{link}[\text{new}] := \text{start}$
5.  $\text{start} := \text{new}$
6. Exit.

(ii) at the end of the linked list:-

Algo

1. If  $\text{Avail} = \text{NULL}$  then  
    write: memory overflow and exit
2. Set  $\text{new} := \text{Avail}$  and  $\text{Avail} := \text{link}[\text{Avail}]$
3. Set  $\text{info}[\text{New}] := \text{Ele}$  and  $\text{link}[\text{new}] := \text{NULL}$
4. Set  $P := \text{start}$
5. Repeat while  $\text{link}[P] \neq \text{NULL}$   
         $P := \text{link}[P]$
6. Set  $\text{link}[P] := \text{new}$
7. Exit.

(iii) after a given node in linked list:-

Algo

1. If  $\text{Avail} = \text{NULL}$  then  
    write: memory overflow and exit
2. Set  $\text{new} := \text{Avail}$  and  $\text{Avail} := \text{link}[\text{Avail}]$
3. Set  $\text{info}[\text{New}] := \text{Ele}$  and  $\text{link}[\text{new}] := \text{NULL}$
4. Set  $P = \text{start}$
5. Repeat while ( $P \neq \text{NULL}$  and  $\text{info}[P] \neq x$ )  
         $P := \text{link}[P]$
6. If  $P = \text{NULL}$  then  
        write: element not found
7. else  
        (a)  $\text{link}[\text{new}] := \text{link}[P]$   
        (b)  $\text{link}[P] := \text{new}$
7. Exit.

(iv) Before a given node in linked list :-

Algo

1. If  $\text{Avail} := \text{NULL}$  then  
    write: memory overflow and exit
  2. Set  $\text{New} := \text{Avail}$  and  $\text{Avail} := \text{link}[\text{Avail}]$
  3. Set  $\text{Info}[\text{New}] := \text{ele}$  and  $\text{link}[\text{new}] := \text{NULL}$
  4. Set  $\text{Ptr} := \text{start}$  and  $\text{temp} := \text{NULL}$
  5. Repeat while ( $\text{Ptr}! = \text{NULL}$  and  $\text{Info}[\text{Ptr}]! = \text{X}$ )

Set temp: = ptr

Set  $\text{ptr} := \text{link}[\text{Ptr}]$

3

6. If  $\text{ptr} = \text{NULL}$  then  
    write element not found  
else  
(a)  $\text{link}[\text{new}] := \text{ptr}$

7. Exit.

Algo

1. If  $\text{Avail} := \text{NULL}$  then  
    write: memory overflow and exit
  2. Set  $\text{New} := \text{Avail}$  and  $\text{Avail} := \text{link}[\text{Avail}]$
  3. Set  $\text{Info}[\text{New}] := \text{ele}$  and  $\text{link}[\text{New}] = \text{NULL}$
  4. Set  $\text{temp} := \text{NULL}$  and  $\text{Ptr} := \text{Start}$
  5. Repeat while  $(\text{Ptr}! = \text{NULL})$  and  $\text{Info}[\text{Ptr}] < \text{ele}$ 
    - {

temp := ptr

ptr := link [ptr]

3

6. If  $\text{ptr} = \text{NULL}$

$\text{link}[\text{temp}] := \text{new}$

else if  $\text{temp} == \text{NULL}$

$\text{start} := \text{new}$

else

$\text{link}[\text{new}] := \text{ptr}$

$\text{link}[\text{temp}] := \text{new}$

7. Exit.

6. Deletion in linked list:-

(i) at the beginning of linked list:-

Algo

1. If  $\text{start} = \text{NULL}$  then

    write: underflow and exit

2. set  $\text{temp} := \text{start}$

3.  $\text{start} := \text{link}[\text{start}]$

4.  $\text{link}[\text{temp}] := \text{Avail}$  and  $\text{Avail} := \text{temp}$

5. Exit.

(ii) at the ending of linked list :-

Algo

1. If  $\text{start} = \text{NULL}$  then

    write: underflow and exit

2. Set  $\text{temp} := \text{start}$  and  $\text{Pre} := \text{NULL}$

3. Repeat while  $\text{link}[\text{temp}] \neq \text{NULL}$

{

$\text{pre} := \text{temp}$

$\text{temp} := \text{link}[\text{temp}]$

}

4. If ( $\text{pre} = \text{NULL}$ )

$\text{start} := \text{NULL}$

Else

$\text{link}[\text{Pre}] := \text{NULL}$

5.  $\text{link}[\text{temp}] := \text{Avail}$  and  $\text{Avail} := \text{temp}$

6. Exit.

(iii) A given node in linked list :-

Algo

1. If  $\text{start} = \text{NULL}$  then

    write: Underflow and exit

2. Set  $\text{temp} := \text{start}$  and  $\text{pre} := \text{NULL}$

3. Repeat while  $\text{temp} \neq \text{NULL}$  and  $\text{temp}[\text{temp}] \neq \text{Elle}$

{

$\text{pre} := \text{temp}$

$\text{temp} := \text{link}[\text{temp}]$

}

4. If  $\text{temp} = \text{NULL}$  then

    write: "Element not found" and exit.

else

    If  $\text{pre} = \text{NULL}$  then

$\text{start} := \text{link}[\text{start}]$

    else

$\text{link}[\text{Pre}] := \text{link}[\text{temp}]$

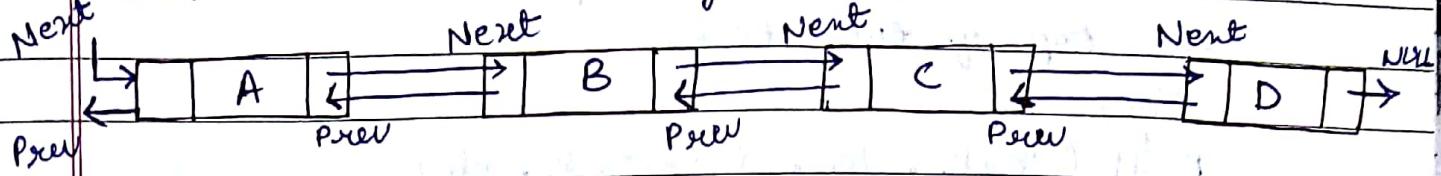
5.  $\text{free}(\text{temp})$

6. Exit.

\* Doubly Link list :-

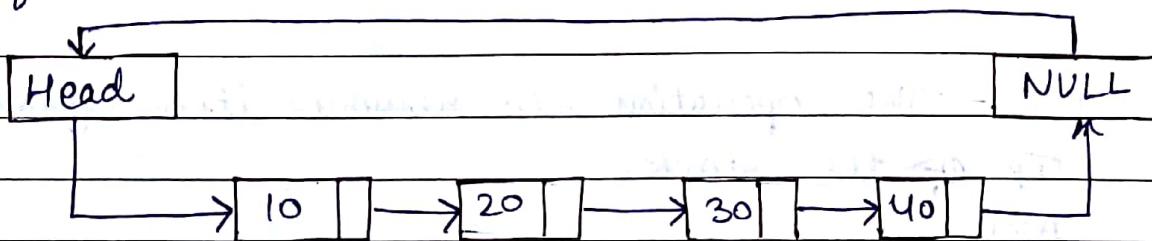
It contains an extra pointer, typically called previous pointer together with the next pointer and data which are there in singly linked list.

It is also k/a two way list.



Circular Linked List :-

Circular list is a list in which the link field of the last node is made to point to the start / first node of the list.



Doubly circular Linked list :-

- In doubly linked list, the next pointer of the last node points to the first node and the previous pointer of the first node points of the last node making the circular in both direction.
- Insertion and deletion at specified position is same as doubly link list.

-X-

STACK :-

- Stack is a collection of elements, where element last to be inserted is first to be taken out.
- Stack implement principle of FILO or LIFO.
- Only access to the stack is the top element.

Operations on Stack :-

1. PUSH
2. POP.

1. PUSH :- The operation to place a new item at the top of the stack.

Algo.

PUSH (stack, Top, MaxStack, Ele)

1. If TOP = MaxStack then  
write: stack overflow and exit.
2. Set TOP := TOP + 1
3. Stack [TOP] := Ele.
4. Return.

2. POP :- The operation to remove item from the top of the stack.

Algo.

POP (stack, top, ele)

1. If TOP = 0 then  
write: stack underflow and exit
2. Set Ele := stack [TOP]
3. Set TOP = TOP - 1
4. Return.

### Application of Stack:-

1. Function calling / Recursion
2. Matching Parenthesis
3. Compilation of Arithmetic Expressions
4. Convert no. bases.
5. Depth first search.

X

## Arithmetic expressions:-

1. Infix expression ( notation )
2. Prefix expression ( Polish notation )
3. Postfix expression ( Reverse Polish notation ).

## Evaluation of postfix expression :-

Suppose  $P$  is an arithmetic expression written in postfix notation, the following algorithm which uses a STACK to hold operands evaluates  $P$ .

**Algorithm :-** This algorithm finds the value of an arithmetic expression and written in postfix notation.

1. Add a right parenthesis ")" at the end of  $P$ .
2. Scan  $P$  from left to right and repeat steps 3 and 4 if for each element of  $P$  unit ")" is encountered.
3. If an operand is encountered then put it on stack.
4. If an operator is encountered then
  - (a) Remove the two top element of stack where A is the top and B is the next of top element.
  - (b) evaluate  $B \otimes A$
  - (c) Place the result B back on the STACK.

5. Set value := top element of STACK

6. Exit.

Conversion of Infix to Postfix :-  $(Q_1, P_1)$

Algorithm:-

1. PUSH "C" onto stack and add ")" to the end of  $Q_1$ .

2. Scan  $Q_1$  from left to right and repeat step 3 to 6 for each element of  $Q_1$ , until stack is empty.

3. If an operand is encountered add it to  $P_1$ .

4. If a "(" parenthesis is encountered then PUSH it in the STACK.

5. If an operator  $(X)$  is encountered then:-

(a) Repeatedly pop from stack and add to  $P_1$  each operator (on the top of stack) which has the same precedence as or higher precedence than  $(X)$ .

(b) Add  $(X)$  to STACK.

6. If a ")" parenthesis is encountered then:-

(a) Repeatedly pop from stack and add a  $P_1$  each operator (to the top of stack) until a left parenthesis is encountered.

(b) Remove the left parenthesis.

7. Exit.

Recursion :-

Recursion have 2 properties:-

1. These must be certain criteria called base criteria for which recursion doesn't call itself.
2. Each time the procedures does call itself it must be closer to base criteria.

factorial function :-

If  $n=1$ , then  $n!=1$ .

If  $n > 1$ , then  $n! = n \times (n-1)!$

FACTORIAL (Fact, n)

1. If  $N=1$  then

Set fact := 1 and return fact

2. Call FACTORIAL (fact, n-1)

3. Set Fact := n \* FACT

4. Return.