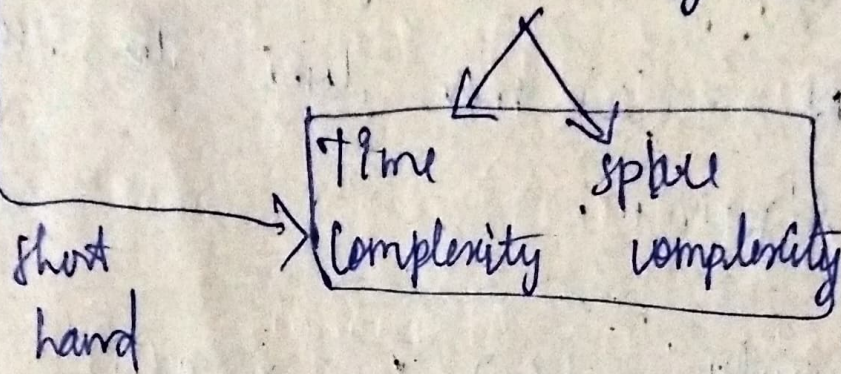


unit -1 :

Asymptotic Notations

An algorithm \rightarrow best an algorithm na.
efficiency check pannam.



Definition:

> It is a mathematical notations used to describe the running time of an Algorithm.

> Various notations. Θ , O , Ω , θ .

1) Big Oh Notation :

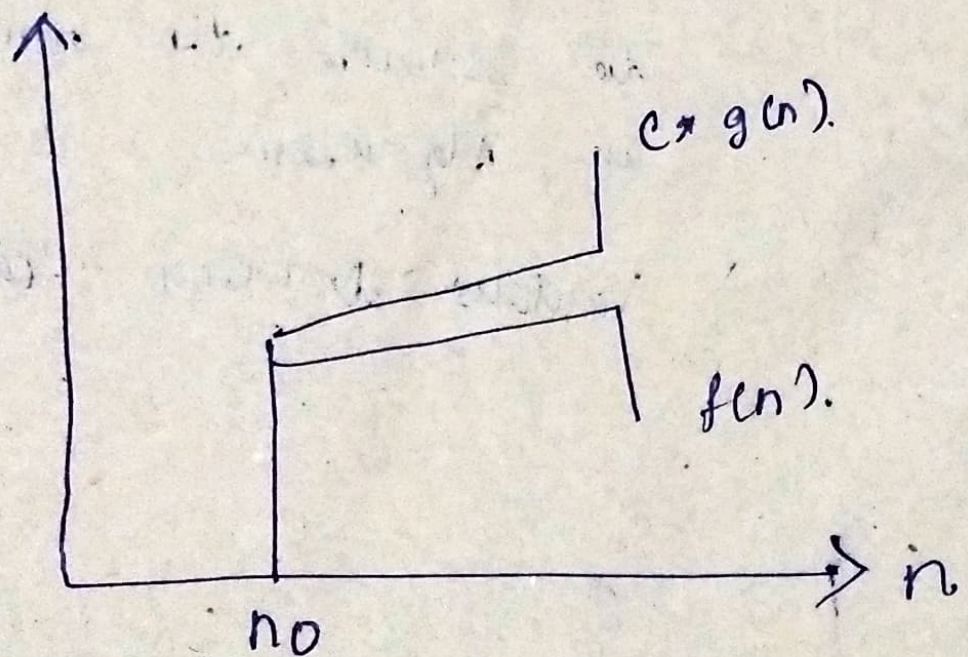
> It is denoted by 'O'

> It represents the upper bound of Algorithm running time.

> It takes more time to run an Algorithm.

Formula : $f(n) \leq c * g(n)$

Diagram :



2). Omega Notation :

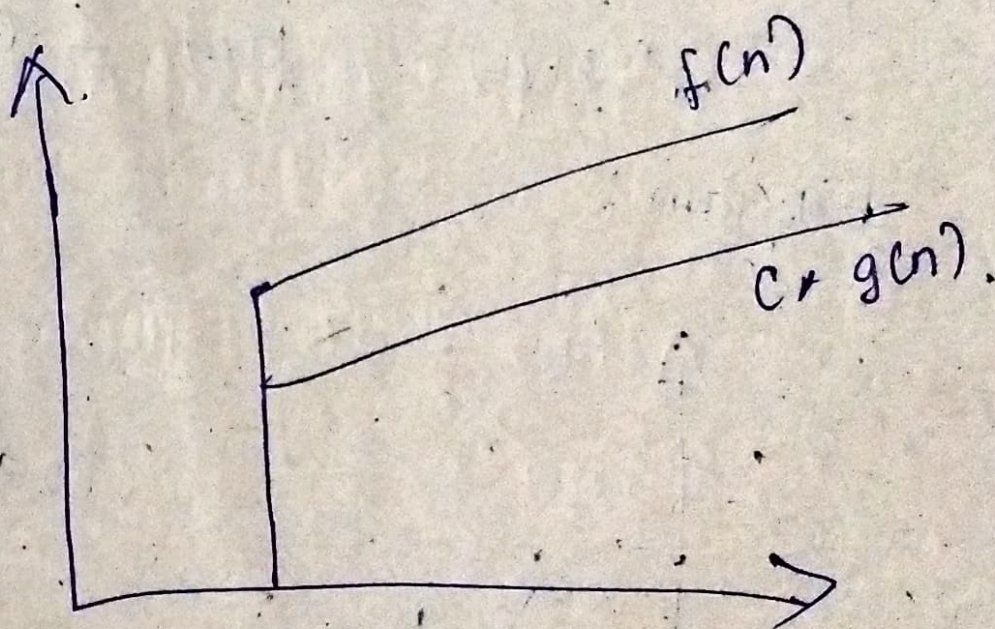
> It is denoted by Ω

> It represents the lower bound of Algorithm running time.

> It takes less time to run an Algorithm.

Formula : $f(n) \geq c * g(n)$

Diagram :



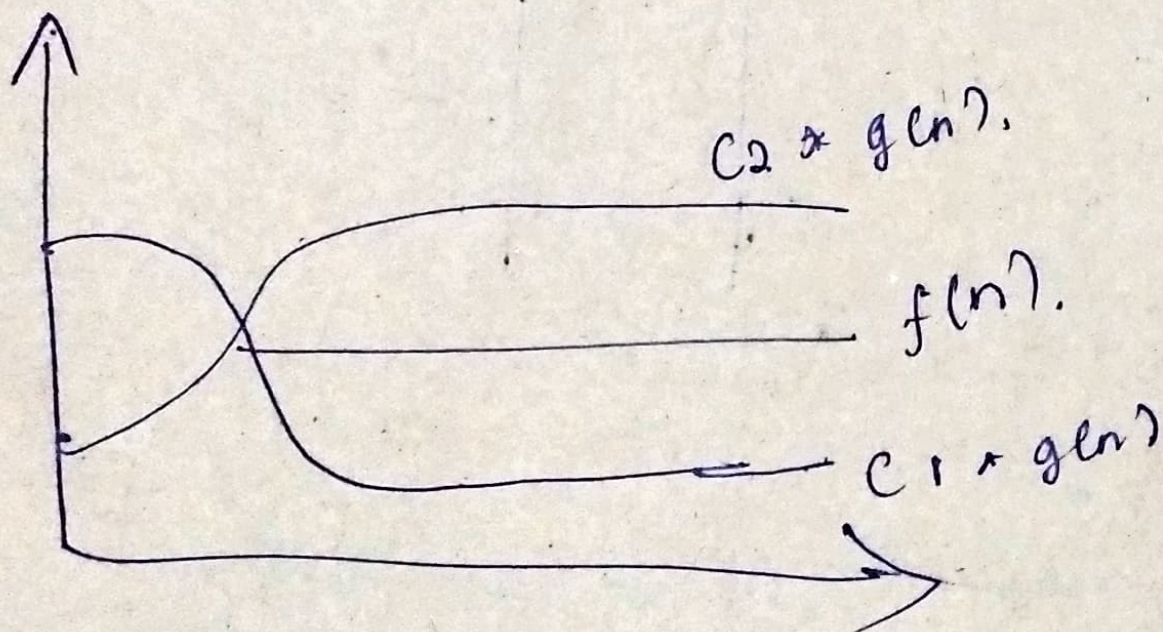
3) Big Theta Notation :

- > It is denoted by ' Θ '.
- > It is the method the running time is between the upper and Lower bound.
- > It takes Average time to run an Algorithm.

Formula :

$$C_1 * g(n) \leq f(n) \leq C_2 * g(n).$$

Diagram :



properties of order of growth

1. $f_1(n)$ is order of $g_1(n)$. &
 $f_2(n)$ is order of $g_2(n)$ then.

$$f_1(n) + f_2(n) \in O(\max(g_1(n), g_2(n))).$$

2. Polynomials of degree $m \in \Theta(n^m)$.

$$3. O(1) < O(\log n) < O(n) < O(n^2) < O(2^n)$$

4. Exponential functions are have different orders of growth for different values of a .

Properties of Asymptotic Notations:

1. General Properties:

i) $f(n) = O(g(n))$ then $a * f(n) = O(g(n))$

ii) $f(n) = \Omega(g(n))$ then $a * f(n) = \Omega(g(n))$

iii) $f(n) = \Theta(g(n))$ then $a * f(n) = \Theta(g(n))$

2. Reflexive:

$$f(n) = O(f(n)).$$

$f(n) = n^2$ then $f(n) = O(n^2)$.

similarly

$$f(n) = \Omega(f(n)).$$

$$f(n) = \Theta(f(n)).$$

3. Symmetry :

$f(n) = \Theta(g(n))$ if and only if $g(n) = \Theta(f(n))$.

Eg :

$f(n) = n^3$ and $g(n) = n^3$ then.

$f(n) = \Theta(n^3)$ and $g(n) = \Theta(n^3)$.

4. Transitivity :

$f(n) = O(g(n))$ & $g(n) = O(h(n))$

then $f(n) = O(h(n))$.

5. Transpose symmetry :

$f(n) = O(g(n))$ if and only if

$g(n) = O(f(n))$.

b). $f(n) = O(g(n))$ and, furthermore,

$f(n) = \Omega(g(n))$ then,

$f(n) = \Theta(g(n))$

$$1) \quad \Theta(f(n) + g(n)) = \max(\Theta(f(n)), \Theta(g(n)))$$