CS217 - Algorithm Design and Analysis
Homework Assignment 1

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1 Recursion and Dynamic Programming

Exercise 1.1. This complexity of the recursion process mainly depends on the additions. Let's define $F_{i,j}$ as when such recursion process start at calculating $\binom{i}{j}$, the number of additions it will do.

obviously, $F_{i,j}$ consists of three parts below:

- $F_{i-1,j-1}$: add the number of addition operations in $\binom{i-1}{j-1}$ to the current answer.
- $F_{i-1,j}$: add the $\binom{i-1}{j-1}$'s answer.
- 1: the number of additions which used to add $\binom{i-1}{j-1}$ and $\binom{i-1}{j}$ up.

Thus we have

$$F_{i,j} = F_{i-1,j-1} + F_{i-1,j} + 1 \tag{1}$$

This is a quite strange and Pascal-like equation. Consider

$$(F_{i,j} + 1) = (F_{i-1,j-1} + 1) + (F_{i-1,j} + 1)$$

We can easily find out that

$$F_{i,j} + 1 = \binom{i}{j} + \binom{i}{j-1} + \binom{i}{j+1}$$

Therefore

$$F_{i,j} = \binom{i}{j} + \binom{i}{j-1} + \binom{i}{j+1} - 1 \tag{2}$$

Hence, the total number of addition operation is

$$\binom{n}{k} + \binom{n}{k-1} + \binom{n}{k+1} - 1$$

The lower bound of the complexity of adding two numbers together being $O(1)$ and the upper bound being $O(n)$, we can easily figure out that the lower bound of the complexity of this algorithm is

$$O\left(\binom{n}{k}+\binom{n}{k-1}+\binom{n}{k+1}\right)=O\left(\binom{n}{k}\right)$$

While the upper bound is

$$O\left(n\binom{n}{k}\right)$$

When $n = 50, k = 25$, this algorithm needs to do about 1.3×10^{16} additions, which is too slow. Hence this algorithm is inefficient.

Exercise 1.2. Since we use dynamic programing, the runing time depends on how many binomial coefficients we calculate. Consider the formula

$$\binom{n}{k}=\binom{n-1}{k-1}+\binom{n-1}{k}$$

It's easy to see that the area we calculate is an parallelogram in Pascal's triangle. More specificly, in order to get the answer of $\binom{n}{k}$, there are exactly $n-k$ numbers we calculate in the columns from 1 to k, so we need to do $O(k \cdot (n-k))$ additions.

The time we need for a single addtion depends on the length of the result, namely $\log\binom{n}{k}$. Since

$$\sum\binom{n}{k}=2^n$$

We can easily see that $\binom{n}{k}$ has a length up to n. So the upper bound of the running time is $O(nk \cdot (n-k))$, which is not very efficient.

Exercise 1.3. The additions we need is the same as Ex2, $O(k \cdot (n-k))$. But this time the addition can be done in $O(1)$ time because we only need to keep the result modulo 2. So the running time is $O(k \cdot (n-k))$, which is quite efficient.

Exercise 1.4.

1. Proof by introduction

   Proof. We can easily prove the lemma using induction:

   (a) Basic step
   When $n = 2^1$, the lemma obviously holds because

   $$\binom{2}{1}=0$$

   (b) Induction step
   Assume that the lemma holds where n equals to 2^d, namely

   $$\binom{n}{k}=0$$

   when $1 \leq k \leq n-1$.
   Then we will show that the lemma holds when n equals to 2^{d+1}.

According to the hints, on the one hand, the number of paths to the n-th line and the k-th grid equals to

$$\binom{n}{k}$$

On the other hand, we can also calculate the number by Multiplication Principle and Addition Principle, namely

$$\binom{n}{k} = \sum_{i=0}^{n/2} \binom{n/2}{i} * \binom{n/2}{k-i} \tag{3}$$

Then our proof divides into three cases:

- When $1 \leq k \leq n/2 - 1$, it is obviously that

  $$\binom{n/2}{k} = 0$$

  and

  $$\binom{n/2}{0} = \binom{n/2}{n/2} = 1$$

  According to the equation (3), we have

  $$\binom{n}{k} = \binom{n/2}{k-0} + \binom{n/2}{k-n/2}$$

  By assumption, when $1 \leq k < n/2$, we have

  $$\binom{n/2}{k} = 0$$

  Notice that $k - n/2 < n/2$, so we have

  $$\binom{n/2}{k-n/2} = 0$$

  Therefore

  $$\binom{n}{k} = 0$$

- When $k = n/2$, it is obviously that

  $$\binom{n/2}{k} = \binom{n/2}{k-n/2} = 1$$

  Therefore

  $$\binom{n}{k} = (1+1) = 2 \equiv 0 \pmod 2$$

- When $n/2 \leq k < n$, it is easy to see that

  $$\binom{n/2}{k} = 0$$

And by assumption

$$\binom{n/2}{k-n/2} = 0$$

So that we have

$$\binom{n}{k} = 0$$

(c) Conclusion
By induction, we have

$$\binom{n}{k} \equiv 0 \pmod 2$$

where $1 \leq k \leq n-1$ and $n = 2^d$ $(d \geq 1)$.

$\square$

2. Direct Proof

Proof. We can denote the binomial coefficient as $2^p * q$, namely

$$\binom{n}{k} = \frac{n!}{k!(n-k)!} = 2^p * q \tag{4}$$

where $2 \nmid q$, $n = 2^d (d \geq 1)$ and $1 \leq k \leq (n-1)$.
We can easily calculate the exponent p with the formular below

$$p = \sum_{i=1}^{d} \left\lfloor \frac{n}{2^i} \right\rfloor - \sum_{i=1}^{d} \left\lfloor \frac{k}{2^i} \right\rfloor - \sum_{i=1}^{d} \left\lfloor \frac{n-k}{2^i} \right\rfloor$$

Specially when $i = d$, we have

$$\left\lfloor \frac{n}{2^i} \right\rfloor = 1 \text{ and } \left\lfloor \frac{k}{2^i} \right\rfloor = \left\lfloor \frac{n-k}{2^i} \right\rfloor = 0$$

According to the equation (4), we have

$$\begin{aligned} p &= 1 - 0 - 0 + \sum_{i=1}^{d-1} \left\lfloor \frac{n}{2^i} \right\rfloor - \left\lfloor \frac{k}{2^i} \right\rfloor - \left\lfloor \frac{n-k}{2^i} \right\rfloor \\ &\geq 1 + \sum_{i=1}^{d-1} \frac{n}{2^i} - \frac{k}{2^i} - \frac{n-k}{2^i} = 1 \end{aligned}$$

Hence we have

$$2 \mid \binom{n}{k}$$

That is

$$\binom{n}{k} \equiv 0 \pmod 2$$

$\square$

Exercise 1.5. We can easily come out an algorithm using the formular which is proved in Ex5. And here comes out of a theorem proved below.

Theorem 1. *Let p represents the maximum 2^d which is less than n, then either $m > n - p$ or $m - p \leq 0$ holds.*

Proof. It is obvirusly that $p \geq \frac{n}{2}$ according to the pseudocode and the definition. Then the proof divided into two cases:

1. If $m \leq \frac{n}{2}$, then $m - p \leq 0$;
2. If $m > \frac{n}{2}$, then $m > \frac{n}{2} \geq n - p$.

□

And with the theorem proved above, we can come up with a algorithm:

Algorithm 1 An efficient method calculating the binomial coefficient

```
function BINOM(n, m, p)
    if m > n or m < 0 then
        return 0
    end if
    if n < 2 then :
        return 1
    end if
    if p = 0 then :
        p = 1
        while p ∗ 2 < n do
            p ← p ∗ 2
        end while
    else
        while p ≥ n do
            p ← p/2
        end while
    end if
    return (BINOM(n − p, m − p, p) + BINOM(n − p, m, p)) mod2
end function
```

The implementation in python is displayed below

```python
def binom(n, m, p = 0):
    if(m > n or m < 0):
       return 0;
    if(n < 2):
        return 1;
    if(p == 0):
        p = 1;
        while(p * 2 < n):
            p = p * 2;
    else:
```

```
        while(p >= n):
            p = p / 2;
    return binom(n - p, m - p, p) ^ binom(n - p, m, p);
```

Exercise 1.6. Assume it will form a rho-like process. Considering i is the entrance of the "circle" of the process, we have

$$F'_i = F'_j \text{ and } F'_{i+1} = F'_{j+1}$$

But

$$F'_{i-1} \neq F'_{j-1}$$

If not, you can easily see that $i - 1$ is the entrace of the "circle".
But according to the formula

$$F'_{i+1} = F'_i + F'_{i-1}$$

We have

$$F'_i = F'_j \text{ and } F'_{i+1} = F'_{j+1} \Rightarrow F'_{i-1} = F'_{j-1}$$

which is a contradictory. Thus the assumption does not hold.

Consider every F'_i, F'_{i+1} can form a pair $\{F'_i, F'_{i+1}\}$. Because $F'_i \in [0, k)$, the number of such pair will not exceed k^2. According to the Pigeon's Theorm, there must be a cycle-like process whose length doesn't exceed k^2, so we can find j which satisfiied $F'_0 = F'_j$, $F'_1 = F'_{j+1}$ and $F'_n = F'_{n\,mod\,j}$ within the complexity $O(k^2)$.

Exercise 1.7. The theorem in the homework is the direct inferrence of Lucas's Theorem. However Lucas's Theorem holds only when the module number is a prime. Our question is that what if the module number is a composite number?