

Binary Search

Algorithm :-

arr = [2, 4, 6, 9, 11, 12, 14, 20, 36, 48]
 ↗ sorted array
 ↑ ascending order

target = 36

1. Find the middle element

2. Check:

if target > middle \Rightarrow search in right

else \Rightarrow search in left

3. If target == middle \Rightarrow we found element

Example :-

In the above array

arr = [2, 4, 6, 9, 11, 12, 14, 20, 36, 48]
 0 1 2 3 4 5 6 7 8 9

target = 36

1st \rightarrow find the middle element

$$\text{mid} = \frac{\text{start} + \text{end}}{2} = \frac{0+9}{2} = 4 \quad [\text{The element at index } 4 \text{ is the middle element}]$$

2nd \rightarrow lets check.

Is target > middle $\Rightarrow 36 > 11 \Rightarrow$ yes! \Rightarrow check right side

Now, arr = [2, 4, 6, 9, 11, 12, 14, 20, 36, 48]

$$\text{mid} = \frac{5+9}{2} = 7 \quad [\text{The element at index } 7 \text{ is the middle element}]$$

Lets check:

Is target > middle $\Rightarrow 36 > 20 \Rightarrow$ yes! \Rightarrow check right side

i.e., arr = [2, 4, 6, 9, 11, 12, 14, 20, 36, 48]

mid = $\frac{8+9}{2} = 8$ [The element at index 8 is the middle element]

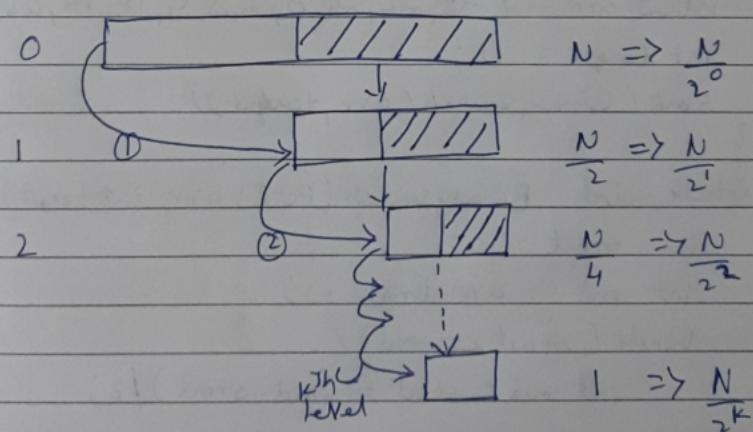
Here, target == middle $\Rightarrow 36 == 36 \Rightarrow$ Element found at index 8

\Rightarrow Time Complexity:

Best case : O(1)

Worst case : O(log n)

Explanation: Find the maximum number of comparisons



$$\frac{N}{2^k} = 1$$

$$N = 2^k$$

$$\log N = \log(2^k)$$

$$\log N = k \log 2$$

$$k = \frac{\log N}{\log 2}$$

$$k = \log_2 N \rightarrow \text{size of array}$$

\hookrightarrow Total no. of comparisons in worst case

=> Order agnostic Binary Search

Let's search say if we don't know that the array is sorted in ascending (or) descending order

arr: [90, 75, 18, 12, 6, 4, 3, 1] target = 75

Here, target > middle => Search left

Here, Start > end → descending order

int low Start < end → Ascending order

Code :- BS code

PSVM () {

 int arr = {-18, -12, -4, 0, 2, 3, 4, 15, 16, 45, 89};

 int target = 22;

 cout (BinarySearch(arr, target))

}

static int BinarySearch (int arr[], int target) {

 int start = 0;

 int end = arr.length - 1;

 while (start <= end) {

 int mid = start + (end - start) / 2;

 if (target < arr[mid]) {

 end = mid - 1;

 } else if (target > arr[mid]) {

 start = start + 1;

 } else {

 return mid;

 }

 return -1;

}

Q8: Ceiling of a number

$\text{arr} = [2, 3, 5, 9, 14, 16, 18]$ target = 14

Find the smallest element in array greater than (or) equal to target

$$\text{Eg: } \text{ceil}(\text{arr}, \text{target} = 14) = 14$$

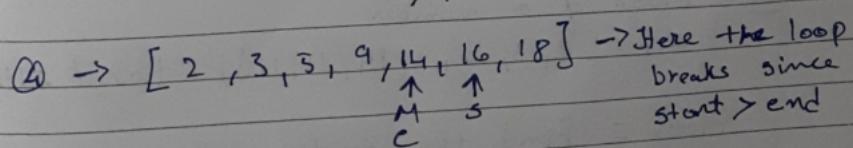
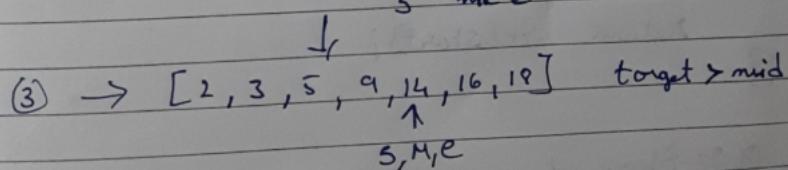
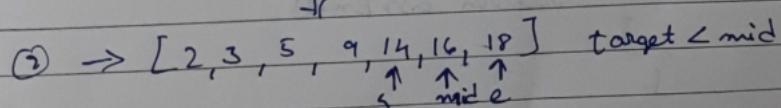
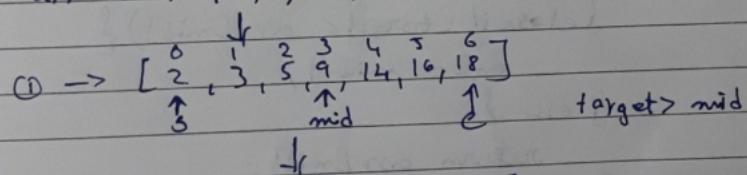
$$\text{ceil}(\text{arr}, \text{target} = 15) = 16$$

$$\text{ceil}(\text{arr}, \text{target} = 4) = 5$$

Hint :- Since, you are searching an element in an array of sorted elements you can apply Binary Search. If the target element is found in the array you can return the element else you can return array element of start index for next greater than the element target.

Explanation :-

Eg:- $\text{arr} = [2, 3, 5, 9, 14, 16, 18]$ target = 15



Here, next greater than the target element will be number in start position

Code :-

PSVM () {

int [] arr = { 2, 4, 6, 9, 11, 12, 14, 20, 36, 48 };

int target = 13;

System.out.println(ceilNumber(arr, target));

}

static int ceilNumber(int [] arr, int target) {

int start = 0;

int end = arr.length - 1;

if (target > arr[end]) {

return arr[end];

}

while (start <= end) {

int mid = start + (end - start) / 2;

if (target > arr[mid]) {

start = mid + 1;

else if (target < arr[mid]) {

end = mid - 1;

} else {

return arr[mid];

}

return arr[start];

}

Q :- Floor of a number

arr = [2, 3, 5, 9, 14, 16, 18]; target = 14

Floor = find the greatest element smaller (or) equal to the target

Eg :- floor(arr, 14) = 14

floor(arr, 10) = 9

floor(arr, 6) = 5

Hint :- Since, you are searching for an element in a sorted array you can apply Binary search. If the target element is found in the array you can return the element else you can return the array element of end index for smaller than the target element.

Explanation :-

$$\text{arr} = [2, 3, 5, 9, 14, 16, 18] \quad \text{target} = 15$$

↑
s ↑
m ↑
e

$$[2, 3, 5, 9, 14, 16, 18] \quad \text{target} > \text{mid}$$

↑
s ↑
m ↑
e

$$[2, 3, 5, 9, 14, 16, 18] \quad \text{target} > \text{mid}$$

↑ ↑
s e

↓,

$$[2, 3, 5, 9, 14, 16, 18] \quad \text{target} < \text{mid}$$

↑ ↑
c sM

Here, the start > end loop breaks and return end index element for floor value

Code :-

PSVM C

```
int arr[] = {2, 4, 6, 9, 11, 12, 14, 20, 36, 48};
```

```
int target = 10;
```

```
cout << floor(arr, target))
```

}

```
Static int floor(int arr[], int target) {
```

```
    int start = 0;
```

```
    int end = arr.length - 1;
```

```
    if (target < arr[start]) {
```

```
        return arr[start];
```

}

```
while (start <= end) {  
    int mid = start + (end - start) / 2;  
  
    if (target > arr[mid]) {  
        start = mid + 1;  
    } else if (target < arr[mid]) {  
        end = mid - 1;  
    } else {  
        return arr[mid];  
    }  
}  
return arr [end];  
}
```

Q :- Find the smallest letter greater than the target

Eg :- letters = ["c", "f", "j"] target = "a"

O/P = "c"

Note :- letters wrap around

Eg :- letters = ['a', 'b'] target = 'z'

O/P :- 'a' → Should start from begin

Hint :- Same as ceiling of a number but return only greater than & remove equal sign

Code :-

```
PSVM () {  
    char [] arr = {'c', 'f', 'j'};  
    char target = 'j';  
    cout (nextChar (arr, target))  
}
```

```
Static char nextCharacter (char[] arr, char target) {
    int start = 0;
    int end = arr.length - 1;
```

```
While (start <= end) {
    int mid = start + (end - start) / 2;
    if (+target > arr[mid]) {
        start = mid + 1;
    } else if (+target < arr[mid]) {
        end = mid - 1;
    } else {
        start = mid + 1;
    }
}
return arr [start % arr.length];
```

\curvearrowleft For Wrap around

Q8- Find first and last position of element in sorted array

Given an array of integers nums Sorted in non-decreasing order, find the starting and ending position of given target value. If target is not found in the array return [-1, -1]

Eg:- nums : [5, 7, 7, 8, 8, 10], target = 8

O/P : [3, 4]

Hint 8- Since, given array is a sorted array and target element to be searched we can use Binary Search. To find the first and last position. Find the index of first and last occurrence and return the indexes.

Code:-

```
PSVm () {
```

```
int [] arr = {5, 7, 7, 8, 8, 10};
```

```
int target = 4;  
int[] ans = {-1, -1};  
ans[0] = SearchIndex(arr, target, isForward);  
ans[1] = SearchIndex(arr, target, !isForward);  
System.out.println(Arrays.toString(ans));  
}  
  
static int SearchIndex(int[] arr, int target, boolean isForward){  
    int start = 0;  
    int end = arr.length - 1;  
    int ansIndex = -1;  
  
    while (start <= end){  
        int mid = start + (end - start) / 2;  
  
        if (target > arr[mid]){  
            start = mid + 1;  
        } else if (target < arr[mid]){  
            end = mid - 1;  
        } else {  
            ansIndex = mid;  
            if (isForward){  
                start = mid + 1;  
            } else {  
                end = mid - 1;  
            }  
        }  
    }  
    return ansIndex;  
}
```

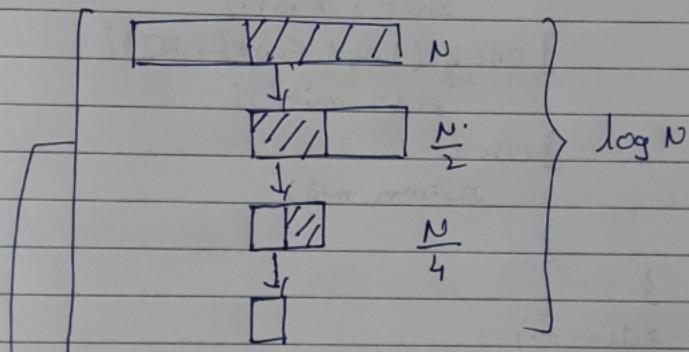
Q :- Find position of an element in a sorted array of infinite numbers

Suppose you have a sorted array of infinite array numbers, how would you search an element in the array

Hint :- Since, Searching is on sorted array we can use Binary Search. Here we don't have starting and ending bound so, we have to search the elements in chunks (like reversing of binary search)

Explanation :-

arr: [2, 3, 5, 6, 7, 8, 10, 11, 12, 15, 20, 23, 30] target = 15



→ Reverse this process to get chunks and increase the size of chunks.

Eg :-

arr: [2, 3, 5, 6, 7, 8, 10, 11, 12, 15, 20, 23, 30] target = 15

size 1st chunk → first chunk with size of chunk is 2, since the target is not in this chunk leave this chunk and create one more chunk with double the size of chunk

Code :-

PSVm () {

int arr = {1, 2, 3, 5, 6, 7, 8, 10, 12, 13, 18, 20, 22};

int target = 18;

int start = 0;

int end = 1;

while (target > arr[end]) {

int temp = end + 1;

```

end = end + (end - start + 1) * 2;
start = temp;
}
cout (binary_search (arr, target, start, end));
    
```

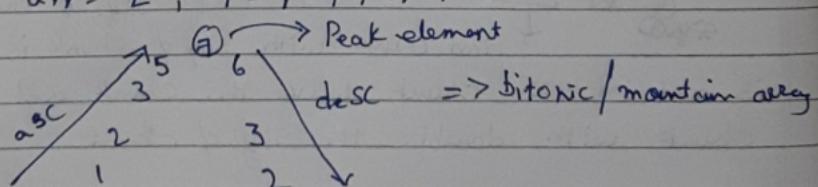
```

Static int binary_search (int arr[], int target, int start, int end)
{
    while (start < end) {
        int mid = start + (end - start) / 2;
        if (target > arr[mid]) {
            start = mid + 1;
        } else if (target < arr[mid]) {
            end = mid - 1;
        } else {
            return mid;
        }
    }
    return -1;
}
    
```

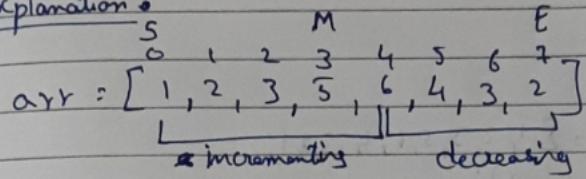
Q :- Find the peak element in mountain array

Mountain array (or) Bitonic array :-

arr = [1, 2, 3, 5, 7, 6, 3, 2]



Hint :- Since the given array is sorted array we can make use of binary search. here the main concept is to use binary search without knowing the target.

Explanation :-Cases :-

(1) If $\text{arr}[\text{mid}] > \text{arr}[\text{mid}+1] \Rightarrow$ you are in the dec part of array

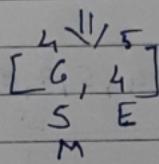
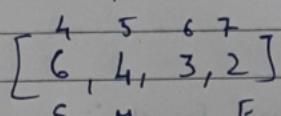
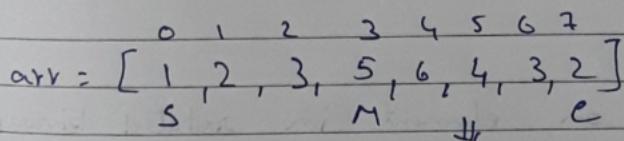
$S \underset{\substack{\text{ans lies} \\ \text{at mid}}} M \underset{\text{E}}{\swarrow} \Rightarrow e = \text{mid}$
 $\text{|| you can stop looking right part}$

(2) $\text{arr}[\text{mid}] < \text{arr}[\text{mid}+1] \Rightarrow$ you are in the asc part of the array

$$\Rightarrow S = \text{mid} + 1$$

$\Rightarrow S \underset{\substack{\text{ans lies in this} \\ \text{part}}}{\underset{\text{M+1}}{\swarrow}} E$

(3) When will loop break?



\Downarrow
 $\begin{bmatrix} 6 \end{bmatrix}$ \Rightarrow In the end S & E both
 S, E will point to largest of
Ans number

Code :-

PSVM() {

int arr = {1, 2, 3, 4, 5, 6, 7};

Sout(peak(arr));

}

Static int peak(int arr) {

int start = 0;

int end = arr.length - 1;

While (start < end) {

int mid = start + (end - start) / 2;

if (arr[mid] > arr[mid + 1]) {

end = mid;

} else {

start = mid + 1;

}

return start;

}

Q :- Search an element in rotated binary search

Rotated binary search :-

Eg :-

arr = [2, 4, 5, 7, 8, 9, 10, 12]
 ↑ 1st

After 1st rotation :- arr = [12, 2, 4, 5, 7, 8, 9, 10]
 ↑ 2nd

After 2nd rotation :- arr = [10, 12, 2, 4, 5, 7, 8, 9]

Hint :-

(1) Find peak element

(2) Binary search in asc array => (0, pivot)

(3) if not found, binary search in [pivot, end]

Explanation :-

Eg :-

$$\text{arr} = [3, 4, 5, 6, \boxed{7}, 0, 1, 2]$$

→ Pivot ⇒ from where next numbers are asc

Case 1 :-

To find pivot

In the rotated sorted array when you find that $\text{mid} > \text{mid} + 1$ i.e pivot

Eg :- [3, 4, 5, 6, 7, 0, 1, 2]

ans = mid → only these 2 will be in dsc order

Case 2 :-

if mid element < mid-1 element

i.e also my ans ⇒ ans = m-1

Eg :- [3, 4, 5, 6, 7, 0, 1, 2]

mid mid-1

[0 < 7] ←

Case 3 :-

Compare the mid with start element to decrease the range

Start element \geq mid element

In this case all the elements from mid will be less than start

Eg :- [4, 5, 6, 3, 2, 1, 0] Smaller than start so we can ignore left side & right side

s

mid

c

left

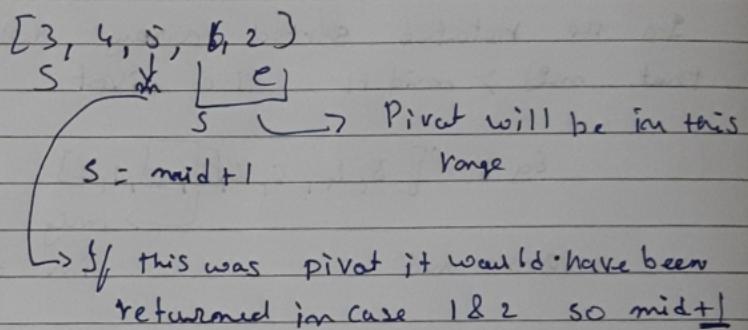
right side

Hence, we can ignore all these elements, since we are looking for peak i.e. largest element

$$\text{end} = \text{mid} - 1$$

Case 4 :- Compare with start element to decrease the range

Start element < mid element



Do binary search to search the target in the range

Case 1 :- Pivot = target // ans

Case 2 :- target \rightarrow start element
Search space (start, Pivot-1)

Case 3 :- target < start element

i.e. we know that all elements from start to pivot are going to be bigger than target

Search space (Pivot+1, till end)

*
Case 4 :- For duplicate elements in array this algo doesn't work in this case if start element & end elements are equal you can skip those elements to find the pivot
Eg :- [3, 9, 2, 2, 1, 2, 1, 2]

Code :-

```
PSVM() {
    int arr = {1, 2, 3, 4, 5, 6, 7, 7};
    cout << find_pivot_with_duplicates(arr);
}
```

```
static int SS_PWD pivot_with_duplicates(int arr) {
```

```
    int start = 0;
```

```
    int End = arr.length - 1;
```

```
    while (start <= end) {
```

```
        int mid = start + (end - start) / 2;
```

// 4 case over here

```
        if (mid < end && arr[mid] > arr[mid + 1]) {
            return mid;
        }
```

```
        if (mid > Start && arr[mid] < arr[mid - 1]) {
            return mid - 1;
        }
```

// If start == mid == end then skip duplicates

```
        if (arr[mid] == arr[start] && arr[mid] == arr[end]) {
```

```
            if (start < end && arr[start] > arr[start + 1]) {
                return start;
            }
```

```
            start++;
        }
```

```
        if (end > start && arr[end] < arr[end - 1]) {
            return end - 1;
        }
```

```
}
```

```
        end++;
    }
```

```
    else if (arr[start] < arr[mid] || arr[start] == arr[mid]
             && arr[mid] > arr[end]) {
```

```
        start = mid + 1;
    }
```

```
    else {
```

```
        end = mid - 1;
    }
```

```
}
```

```
return -1; }
```

Q :- Split array largest sum

Given an array nums which consists of non-negative integers and an integer m , you can split the array into m non-empty continuous subarrays.

Write an algorithm to minimize the largest sum among these m subarrays.

Explanation :-

$$\text{arr} = [7, 2, 5, 10, 8], m=2$$

Sub arrays

$$\begin{array}{c} \overbrace{7, 2, 5, 10}^{\substack{| \\ 24}} \quad | \quad \overbrace{8}^{\substack{| \\ 8}} \\ \text{largest} \end{array} \quad 24$$

$$\begin{array}{c} \overbrace{[7, 2, 5]}^{\substack{| \\ 14}} \quad | \quad \overbrace{[10, 8]}^{\substack{| \\ 18}} \\ \text{Ans min} \\ \text{the largest} \end{array} \quad 18$$

$$\begin{array}{c} \overbrace{[7, 2]}^{\substack{| \\ 9}} \quad | \quad \overbrace{[5, 10, 8]}^{\substack{| \\ 23}} \\ 23 \end{array}$$

$$\begin{array}{c} [7] \quad | \quad \overbrace{[2, 5, 10, 8]}^{\substack{| \\ 25}} \\ 25 \end{array}$$

Cases :-

1) Minimum no. of partitions that we can make = 1

$$\begin{array}{c} \overbrace{[7, 2, 5, 10, 8]}^{\substack{| \\ 32}} \quad \text{if } m=1 \\ 32 \end{array}$$

2) More number of partitions is N because subarrays cannot be empty

[7], [2], [5], [10], [8] if m = N

10

* * *
50, sum of sub arrays lie in the range

$[10, 32] \rightarrow$ The min possible sum of the currency is in this range

So, here we can use Binary Search in this range

Break dream :-

$$\text{Range} = [10, 32]$$

$$\text{start} = 10 \quad ; \quad \text{end} = 32 \quad \Rightarrow \text{mid} = \frac{\text{start} + \text{end}}{2} = \frac{10 + 32}{2} = 21$$

Try to see if we can split the array with 21 as the max sum

[7, 2, 5, 8], 10] Pieces
?

$$\left[\begin{smallmatrix} 7 & 12 & 5 \\ \underbrace{14} & & \end{smallmatrix} \right] \quad \left[\begin{smallmatrix} 8 & 10 \\ & \underbrace{18} \end{smallmatrix} \right]$$

Check 1: if ($\text{pieces} \leq m$) $\Rightarrow \text{end} = \text{mid}$

Now's Start = 10 ; end = 21 ; mid = 15

$$[7, 2, 3], [8], [10] \quad \underline{3}$$

Check 2: If ($\text{pieces} > m$) $\Rightarrow \text{start} = \text{mid} + 1$

Now:- start = 16 ; end = 21 ; mid = 18

$[7, 2, 5], [8, 10]$ pieces = 2

Now:- start = 16 ; end = 18 ; mid = 17

$[7, 2, 5], [8], [17]$ pieces = 3

$$S = m + 1 = 18$$

$$S = 18 ; e = 18$$

$$m = 18 ;$$

Ans

// The ans exists definitely, here by the above 2 checks we will reach

Code :-

```
Public int SplitArray(int[] nums, int m) {
```

```
    int start = 0;
```

```
    int end = 0;
```

```
    for (int i = 0; i < nums.length; i++) {
```

```
        Start = Math.Max(start, nums[i])
```

// in the end of the loop this will contain
the max item of the array

```
        end += nums[i];
```

}

```
    while (start < mid) {
```

```
        int mid = Start + (end - start) / 2;
```

```
        int sum = 0; int pieces = 0;
```

```
        for (int num : nums) {
```

```
            if (sum + num > mid) {
```

// you can't add this in this very m

new one

```
            sum = num;
```

3 pieces++

```

    else {
        sum += num;
    }
    if(pieces > m) {
        start = mid + 1;
    } else {
        end = mid;
    }
}
return end; // here start = end;
}

```

Binary Search in 2D array

Searching in matrices :-

18	9	12	Target = 91
36	-4	91	
44	33	16	

We can do Linear Search for this matrix since, the matrix matrix is not sorted in any manner

Code :-

```

Static int[] search (arr, target) {
    for(r=0; r<n; r++) {
        for(c=0; c<n; c++) {
            if(arr[r][c] == target) {
                return new int[] {r, c};
            }
        }
    }
    return -1;
}

```

$O(N^2)$, $O(1)$

Q :- Matrix is sorted in a row-wise & column-wise manner.

Eg:-

	0	1	2	3	asc
0	10	20	30	40	
1	15	25	35	45	
2	28	29	37	49	
3	33	34	38	50	

Target = 37

Hint :- Since the rows and columns are sorted
We can make use of binary search and
decrease the search space

Explanation :- Start Searching from $\text{row} = 0$; $\text{col} = \text{last col}$

Case 1 :- If $\text{element} == \text{target}$
 \Rightarrow ans found

Case 2 :- If $\text{element} < \text{target}$

\Rightarrow Row ++

Case 3 :- If $\text{element} > \text{target}$
 \Rightarrow Column --

LB

10	20	30	40	Upper bound
15	25	35	45	Compare search
28	29	37	49	element with the
33	34	38	50	Up & l/b and eliminate rows & cols accordingly

Code :-

PSVm () {

 in [] [] arr = {

 { 10, 20, 30, 40 },

 { 15, 25, 35, 45 },

 { 28, 29, 37, 49 },

 { 33, 34, 38, 50 });

 sort(search(arr, target));

}

```

Static int[] Search (int[][] matrix, int target) {
    int r = 0;
    int c = matrix.length - 1;

    while (r < matrix.length && c >= 0) {
        if (matrix[r][c] == target) {
            return new int[] {r, c};
        }
        if (matrix[r][c] < target) {
            r++;
        } else {
            c--;
        }
    }
    return new int[] {-1, -1};
}

```

Q :- Search in a sorted matrix which is strictly sorted

	0	1	2	3
0	1	2	3	4
1	5	6	7	8
2	9	10	11	12
3	13	14	15	16

mid of column mid of row

Target = 2

we can eliminate these rows since mid > target

* Take middle Col & perform BS on it

* If mid of the column == target
 ans found
 if mid of col < target {
 }
 * Start search element from row mid of row &
 mid of col

Cases :-

Case 1 :- if element == target \Rightarrow ans found

Case 2 :- if element > target \Rightarrow // ignore rows after it

Case 3 :- if element < target \Rightarrow // ignore rows above it

In the end 2 rows are remaining :-

①	②	③
① 5	② 6	③ 7 8

- 1) Check whether the mid column you are at that contains the answer [2, 6]
rows found in this case

in case target = 3 ?

- 2) Consider the 4 parts and perform bsr on these parts

Code :-

```
PSym() {
    int I3 I3 > {
        {1, 2, 3},
        {4, 5, 6},
        {7, 8, 9}
    }
    Sout(Search(aarr, a))
}
```

// Search in rows provided between cols provided
 static int C3 binary search(int I3 I3 matrix, int Cstart, int CEnd,
 int target)

```
while (Cstart <= CEnd) {
    int mid = Cstart + (CEnd - Cstart) / 2;
    if (matrix [row] [mid] == target) {
        return new int {row, mid};
    }
    if (matrix [row] [mid] < target) {
        Cstart = mid + 1;
    } else {
        CEnd = mid - 1;
    }
}
```



```
return int[] {-1, -1};  
}  
  
static int[] Search(int[][] matrix, int target) {  
    int rows = matrix.length;  
    int cols = matrix[0].length - 1;  
    if (cols == 0) {  
        return new int[] {-1, -1};  
    }  
    if (rows == 1) {  
        return binarySearch(matrix, 0, 0, cols - 1, target);  
    }  
    int rStart = 0;  
    int rEnd = rows - 1;  
    int cMid = cols / 2;  
  
    // Run the loop till 2 rows are remaining  
    while (rStart < (rEnd - 1)) {  
        int mid = rStart + (rEnd - rStart) / 2;  
        if (matrix[mid][cMid] == target) {  
            return new int[] {mid, cMid};  
        }  
        if (matrix[mid][cMid] < target) {  
            rStart = mid;  
        } else {  
            rEnd = mid;  
        }  
    }  
    // now we have two rows (check in two rows)  
    if (matrix[rStart][cMid] == target) {  
        return new int[] {rStart, cMid};  
    }  
    // Search in 1st half  
    if (target <= matrix[rStart][cMid - 1]) {  
        return binarySearch(matrix, rStart, 0, cMid - 1, target);  
    }
```

// Search in 2nd half

if ($\text{target} \geq \text{matrix}[\text{rstart}][\text{mid}] \& \& \text{target} \leq \text{matrix}[\text{rstart}][\text{cols}-1]$)

{

return binarySearch(matrix, rStart, cMid+1, cols-1, target);

}

// Search in 3rd half

if ($\text{target} \geq \text{matrix}[\text{rstart}+1][\text{cMid}-1]$)

return binarySearch(matrix, rstart+1, 0, cMid-1, target);

else {

return binarySearch(matrix, rstart+1, cMid-1, cols-1, target);

}

{

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