Estimating Population Range by Recurring Online Chunk Bootstrap with Non-cumulative Data in Streaming Data Environment

Abstract—

Index Terms—Article submission, IEEE, IEEEtran, journal, LATEX, paper, template, typesetting.

I. Introduction

The bootstrap method is a powerful and widely used statistical tool for estimating the uncertainty from finite samples of any parameter of interest, such as standard error, confidence intervals, and accuracy, etc. Efron [?] introduced bootstrap methods as a generalization of the jackknife method, where the jackknife was mathematically expressed as a linear approximation of the bootstrap. Empirical results demonstrated that these methods are capable of estimating the standard error of complex estimators effectively. Efron and Tibshirani [?] further explored the basic concepts and applications of bootstrap methods for estimating standard errors, confidence intervals, and other measures of statistical accuracy. For standard error estimation, data from a population is resampled with replacement to create multiple bootstrap samples, allowing for the calculation of statistics and their standard deviations for each sample. The results showed that bootstrap estimations provided reasonably accurate and efficient alignments with theoretical density curves. In terms of confidence intervals, several methods were empirically tested, showing that bootstrap confidence intervals can be more accurate than standard methods, particularly when the distribution of the statistic is nonnormal. Achieving accurate results requires a sufficient number of bootstrap replications. Carpenter and Bithell [?] presented a practical guide for applying bootstrap confidence intervals in healthcare data, addressing when, which, and how to implement these methods. They assessed various bootstrap methods for confidence intervals across three families: pivotal, non-pivotal, and test-inversion. The experiments concluded that bootstrap confidence intervals are a suitable alternative when the

assumptions of the underlying distribution do not hold, such as asymptotic normality, especially with small sample sizes or complex data structures.

Range approximation in one-dimensional data plays an important role in statistics, data analysis, and computational sciences, aiming to estimate the interval between a dataset's minimum and maximum values.

II. STUDIED PROBLEM AND OBJECTIVES

Let a chunk of integer data, a set of bins, and a set of integer data be defined as follows:

- 1) A set of integer data, $\mathbf{D} = \{a \leq d_j \leq b \mid 1 \leq j \leq a\}$ p}, for integer constants p, a, and b.
- 2) An integer data chunk C divided into n smaller chunks, $\mathbf{C} = \{\mathbf{c}_1, \dots, \mathbf{c}_i, \dots, \mathbf{c}_n \mid 1 \le n < \infty\}.$ Each chunk has integers randomly taken from D.
- 3) A bin B for containing incoming all integers in each chunk. There are two attributes, $v_i^{(min)}$ and $v_i^{(max)}$, defining the interval of integer values such that any integer $v_i^{(min)} \leq d_j \leq v_i^{(max)}$ can be assigned to this bin.

Each chunk \mathbf{c}_i sequentially flows into the bin. If some integers whose values are less than $v_1^{(min)}$ or larger than $v_m^{(max)}$, then the width of B must be expanded.

In the beginning, the size of bin B is made large enough to contain all integers in the first incoming chunk c₁. Then size of B is occasionally expanded so that all integers in the other next incoming chunks can be assigned to the intervals of B. Bin B is expanded if the values of some integers in some incoming chunks are either less than $v^{(min)}$ or larger than $v^{(max)}$. Therefore, the studied problem is defined as follows.

Let $min(\mathbf{D})$ and $max(\mathbf{D})$ be the minimum and maximum values of D. After capturing of integers in the first chunk, how to achieve the minimum number of expansions of B so that

- 1) All incoming integers in the next other chunks can be assigned to B.
- 2) $(min(\mathbf{D}) v_1^{(min)}) \ge 0$ is minimum. 0000-0000/00\$00.00 \mathfrak{G} $\mathfrak{Dot}v_1^{(mEE)} max(\mathbf{D})) \ge 0$ is minimum.

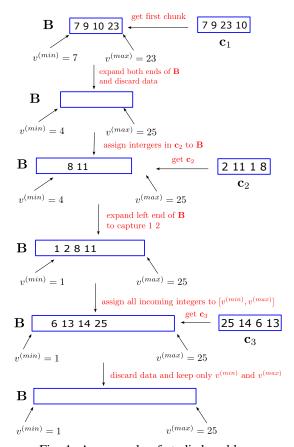


Fig. 1: An example of studied problem.

Figure 1 illustrates an example scenario of capturing chunks. There are 3 sequentially incoming chunks containing these integers: $\mathbf{c}_1 = \{7, 9, 23, 10\}, \mathbf{c}_2 = \{2, 11, 1, 8\},\$ $\mathbf{c}_3 = \{25, 14, 6, 13\}$. After capturing \mathbf{c}_1 , the values of left and right ends of bin **B** are set to $v^{(min)} = 7$ and $v^{(max)} = 23$ and all data are discarded. Then, both ends are expanded in advance to $v^{(min)} = 4$ and $v^{(max)} = 25$. preparing for c_2 . When c_2 enters, all integers except 2 can be captured because the left end value $v^{(min)} = 4$ is larger than 2. Thus, the left end $v^{(min)}$ is expanded to $v^{(min)}=2$. No need to expand the right end. All data in c_2 are discarded. Then, get c_3 . The interval of B is large enough to capture all integers of c_3 . After capturing c_3 , all integers are discarded. In this example, the number of expansions is 2, after chunks 1 and 2. Generally, how to achieve the minimum number of expansions of B in advance so that both expanded values are also minimum.

A. Constraints

The amount of integers in each \mathbf{c}_i is denoted by $|\mathbf{c}_i|$. Bin B can be considered as an interval having left end denoted by L and right end denoted by R. The following constraints are imposed.

- 1) The probability of distribution of each integer d_j in \mathbf{c}_i is unknown.
- 2) $|\mathbf{c}_i| \neq |\mathbf{c}_{i+1}| \text{ or } |\mathbf{c}_i| = |\mathbf{c}_{i+1}|.$
- 3) After assigning all integers of c_i inside the bin, chunk c_i is completely discarded and never reentered the bin assignment.

III. CONCEPT OF RECURRING ONLINE BOOTSTRAP IN STREAMING ENVIRONMENT WITH UNRECORDED DATA

A. Algorithm

explain standard histogram first The algorithm has two phases. The first phase

Algorithm 1: Capturing data chunks and expanding B when it is necessary.

Input: Stream of data chunks, c_i for $1 \le i < \infty$. Output:

Phase 1: Capturing c_1 and get initial expansion of B.

- 1. Get \mathbf{c}_1 and set $total_data = |\mathbf{c}_1|$.
- 2. Let $v^{(min)} = \min(\mathbf{c}_1)$ and $v^{(max)} = \max(\mathbf{c}_1)$.
- 3. Divide **B** into 8 equal sub-intervals \mathbf{b}_i for $1 \le i \le 8$, each of size $(v^{(max)} v^{(min)})/8$.
- 4. Put the integers in c_1 whose values are within sub-intervals b_1 and b_8 into these two sub-intervals.
- 5. Count the number of integers in \mathbf{b}_1 and \mathbf{b}_8 and let $|\mathbf{b}_1|$ and $|\mathbf{b}_8|$ denote these numbers.
- 6. Let $avg = (v^{(max)} + v^{(min)})/2$ be the middle value of B.
- 7. Find the types of probability distribution in list *P* best fitting the data in b₁, and b₈ by using Algorithm 2.1 with *total_data*, b₁, and b₈.
- 8. Compute the standard number of integers in b_1 denoted by lstd, and in b_8 denoted by rstd. from the best fitted probability distribution by using Algorithm 2.2.
- 9. While $|\mathbf{b}_1| > lstd$ or $|\mathbf{b}_8| > rstd$ do
- 10. If $|\mathbf{b}_1| > lstd$ then
- 11. Expand $v^{(min)}$ by using Algorithm 3 with all integers in \mathbf{b}_1 .
- 12. EndIf
- 13. If $|\mathbf{b}_8| > rstd$ then
- 14. Expand $v^{(max)}$ by using Algorithm 4 with

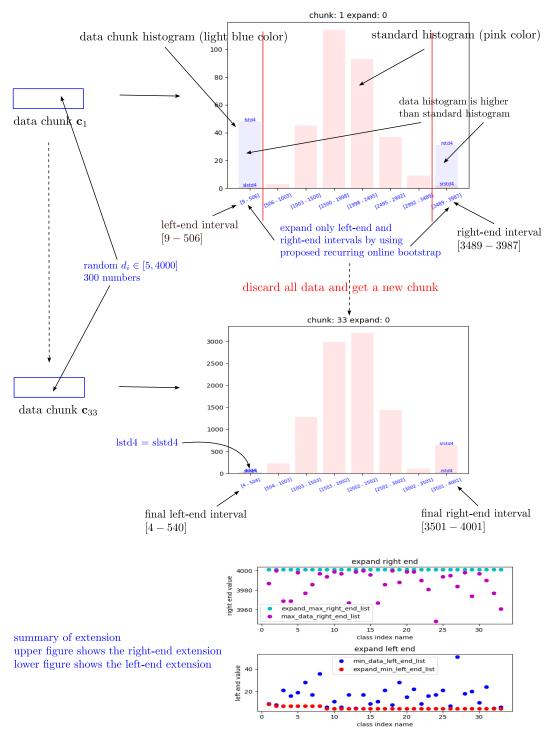


Fig. 2: Framework.

all integers in b₈.

15. EndIf

16. Adjust the width of \mathbf{b}_1 and \mathbf{b}_8 by dividing \mathbf{B} into 8 equal sub-intervals \mathbf{b}_i for $1 \le i \le 8$,

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each of size (v^{(max)} - v^{(min)})/8.
                                                                      20.
                                                                                              EndIf
17. EndWhile
                                                                                              If v^{min}
                                                                                                          and v^{max} do not change.
                                                                      х.
18. Discard c_1 and all integers in b_1 and b_8.
                                                                                                  Go to Line XX.
                                                                      х.
                                                                                              EndIf
                                                                      e.
                                                                                              Adjust the width of b_1 and b_8 by dividing
                                                                      21.
Phase 2: Capturing other c_i and determining the necessity
                                                                                              B into 8 equal sub-intervals \mathbf{b}_i for
                                                                                              1 \le i \le 8, each of size (v^{(max)} - v^{(min)})/8.
of expanding B.
                                                                      22.
                                                                                    EndWhile.
     while there exists a new incoming chunk c_i do
                                                                      23.
                                                                                    Discard c_i and all integers in b_1 and b_8.
2.
             total data = total data + |\mathbf{c}_i|.
                                                                      24. EndWhile.
             If |\mathbf{b}_1| \geq min_B then
х.
                  \mathbf{B}_1 = \{\min(\mathbf{c}_i)\} \cup \mathbf{B}_1.
х.
                 Apply Alg. 3 with \mathbf{B}_1 to get v_B^{(min)}.
                                                                      Algorithm 2.1: Finding the types of probability distribu-
х.
                 If v^{(min)} > v_{B}^{(min)} then
                                                                      tion in list P best fitting the data in b_1, and b_8.
х.
                                                                      Input: (1) A list of standard probability distribution P.
                      v^{(min)} = v_B^{(min)}.
х.
                                                                      (2) total\_data. (3) \mathbf{b}_1. (4) \mathbf{b}_8.
                  EndIf
х.
                                                                      Output: lname and rname.
             Else
х.
                  v^{(min)} = \{\min(\mathbf{c}_i)\}.
х.
             EndIf
х.
                                                                      1.
                                                                            For each type probability distribution p \in P do
             If |\mathbf{b}_8| \geq min_B then
XX.
                                                                      2.
                                                                                 Divide the area under standard probability
                  \mathbf{B}_8 = \{ \max(\mathbf{c}_i) \} \cup \mathbf{B}_8.
х.
                                                                                 distribution p into 8 stripes of equal width.
                  Apply Alg. 4 with \mathbf{B}_8 to get v_B^{(max)}.
                                                                                 Let l_4^{(p)} be the percentage of data in the 4^{th} stripe
Х.
                                                                      3.
                 If v^{(max)} < v_B^{(max)} then v^{(max)} = v_B^{(max)}.
                                                                                 to the left of mean.
х.
                                                                                 Let r_4^{(p)} be the percentage of data in the 4^{th} stripe
                                                                      4.
Х.
                  EndIf
                                                                                 to the right of mean.
х.
             Else
Χ.
                                                                      5.
                                                                                 Compute the difference between standard number
                  v^{(max)} = {\max(\mathbf{c}_i)}.
                                                                                 of integers in the 4^{th} left stripe and |\mathbf{b}_1|:
х.
                                                                                 ld^{(p)} = abs(l_4^{(p)} * total\_data - |\mathbf{b}_1|).
             EndIf
х.
             Divide B into 8 equal sub-intervals b_i for
9.
                                                                                 Compute the difference between standard number
                                                                      6.
              1 \le i \le 8, each of size (v^{(max)} - v^{(min)})/8.
                                                                                 of integers in the 4^{th} right stripe and |\mathbf{b}_8|:
10.
             Put the integers in c_i, whose values are within
                                                                                 rd^{(p)} = abs(r_4^{(p)} * total\_data - |\mathbf{b}_8|).
             sub-intervals b_1 and b_8, into these two
                                                                           EndFor
             sub-intervals.
                                                                           Find lname = \arg\min_{p \in P} (ld^{(p)}).
11.
             Count the number of elements in b_1 and b_8
                                                                           Find rname = \arg\max_{x}(rd^{(p)})
             and let |\mathbf{b}_1| and |\mathbf{b}_8| denote these numbers.
             Find the types of probability distribution in list \mathbb{R}. Return lname and rname.
12.
             best fitting the data in b_1, and b_8 by using
              Algorithm 2.1 with total\_data, \mathbf{b}_1, and \mathbf{b}_8.
             Compute the standard number of elements in balgorithm 2.2: Computing lstd and rstd.
13.
             denoted by lstd, and in b_8 denoted by rstd
             from the best fitted probability distribution by Input: (1) A list of standard probability distribution P.
                                                                      (2) total\_data. (3) lname. (4) rname. (5) l_A^{(p)} and r_A^{(p)};
             using Algorithm 2.2.
14.
              While |\mathbf{b}_1| > lstd or |\mathbf{b}_8| > rstd do
                                                                      \forall p \in P \text{ from Algorithm 2.1.}
                       v_{old}^{(min)}=v^{(min)} and v_{old}^{(max)}=v^{(max)} Output: lstd and rstd.
х.
                       If |\mathbf{b}_1| > lstd then
15.
                                                                      1. If lname is the same as rname then
                           Expand v^{(min)} by using Algorithm \frac{3}{2}.
16.
                                                                             \begin{aligned} & \text{Set } lstd = l_4^{(lname)} * total\_data. \\ & \text{Set } rstd = r_4^{(rname)} * total\_data. \end{aligned}
                           with all integers in b_1.
                                                                      3.
17.
                       EndIf
                                                                      4. EndIf
18.
                       If |\mathbf{b}_8| > rstd then
                           Expand v^{(max)} by using Algorithm ^5\!\!4
                                                                         If lname is different from rname then
19.
                                                                              Set lstd = \max_{p \in P} (l_4^{(p)}) * total\_data.
                           with all integers in b_8.
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Set rstd = \max_{p \in P}(r_4^{(p)}) * total\_data.
7.
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8. EndIf

9. Return lstd and rstd.

Algorithm 3: Recurring online chunk bootstrap for B_1 .

Input: (1) Present set of incoming integers in B_1 ; (2) Number of bootstrap iterations N; (3) $mean(\mathbf{a})$ is a function computing the mean of set a; (4) $std(\mathbf{a})$ is a function computing the standard deviation of set a; (5) abs(x) is the absolute value of constant x.

Output: $v^{(min)}$.

```
1. Let S = \emptyset be a set of bootstrapped samples.
```

- 2. Let $\mathbf{M} = \emptyset$ be a set of mean of each bootstrapped
- Let $\mathbf{Max} = \emptyset$ be a set of maximum values of each bootstrapped samples.
- Let $Min = \emptyset$ be a set of minimum values of each bootstrapped samples.
- 3. Let $P = \emptyset$ be a set of standard deviation of each bootstrapped sample.

```
4.
   prev\_mean = 0.
```

5. For $1 \le i \le N$ do:

Let s_i be a set of randomly sampled integers of 6. size $|\mathbf{B}_1|$ from \mathbf{B}_1 with replacement.

```
7.
            S = S \cup \{s_i\}.
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 $present_mean = (mean(\mathbf{s}_i) + prev_mean)/2.$ 8.

9. $prev_mean = present_mean.$

10. $\mathbf{M} = \mathbf{M} \cup \{present_mean\}.$

 $Min = Min \cup \{min(s_i)\}.$ х.

11. EndFor.

12. $\mu^{(boot)} = mean(\mathbf{M}).$

13. For each $s_i \in S$ do

 $\mathbf{P} = \mathbf{P} \cup \{std(\mathbf{s}_i)\}.$

15. EndFor

16. $\sigma^{(boot)} = mean(\mathbf{P}).$

17. $\mu^{(diff)} = abs(mean(\mathbf{B}_1) - \mu^{(boot)}).$

18. $\sigma^{(diff)} = abs(std(\mathbf{B}_1) - \sigma^{(boot)}).$

If MinmaxBoost

 $min_{left} = meanProbBased(\mathbf{Min}).$ х.

Else X.

 $min_{left} = min(\mathbf{B}_1).$ х.

19. If $mu^{(boot)} < mean(\mathbf{B}_1)$ do

 $v^{(min)} = min_{left} - \mu^{(diff)}.$

21. If $mean(\mathbf{B}_1) < mu^{(boot)}$ do

 $v^{(min)} = min_{left} - \sigma^{(diff)}$.

Algorithm 4: Recurring online chunk bootstrap for B_8 .

Input: (1) Present set of incoming integers in B_1 ; (2) Number of bootstrap iterations N; (3) $mean(\mathbf{a})$ is a function computing the mean of set a; (4) $std(\mathbf{a})$ is a function computing the standard deviation of set a; (5) abs(x) is the absolute value of constant x. Output: $v^{(max)}$.

```
1. Let S = \emptyset be a set of bootstrapped samples.
```

- Let $M = \emptyset$ be a set of mean of each bootstrapped sample.
- Let $P = \emptyset$ be a set of standard deviation of each bootstrapped sample.

```
4.
   prev\_mean = 0.
```

For $1 \le i \le N$ do: 5.

Let s_i be a set of randomly sampled integers of 6. size $|\mathbf{B}_8|$ from \mathbf{B}_8 with replacement.

7. $S = S \cup \{s_i\}.$

8. $present_mean = (mean(\mathbf{s}_i) + prev_mean)/2.$

9. $prev_mean = present_mean.$

10. $\mathbf{M} = \mathbf{M} \cup \{present_mean\}.$

 $\mathbf{Max} = \mathbf{Max} \cup \{ max(\mathbf{s}_i) \}.$ х.

11. EndFor.

12. $\mu^{(boot)} = mean(\mathbf{M}).$

13. For each $s_i \in S$ do

 $\mathbf{P} = \mathbf{P} \cup \{std(\mathbf{s}_i)\}.$ 14.

15. EndFor

16. $\sigma^{(boot)} = mean(\mathbf{P}).$

17. $\mu^{(diff)} = abs(mean(\mathbf{B}_8) - \mu^{(boot)}).$

18. $\sigma^{(diff)} = abs(std(\mathbf{B}_8) - \sigma^{(boot)}).$

If MinmaxBoostх.

 $max_{right} = meanProbBased(\mathbf{Max}).$ х.

Else х.

х. $max_{right} = max(\mathbf{B}_8).$

19. If $mu^{(boot)} < mean(\mathbf{B}_8)$ do

 $v^{(max)} = max_{right} + \sigma^{(diff)}$. х.

21. If $mean(\mathbf{B}_1) < mu^{(boot)}$ do

 $v^{(max)} = max_{right} + \mu^{(diff)}$. х.

IV. EXPERIMENTS

In this study, we evaluate the performance of the range approximation using two scenarios of one-dimensional data populations: simulated and real-world data. We used eight data sets from four established statistical distributions for the simulation, each with different parameter settings as detailed in Table I. The simulation framework encompasses four fundamental distributions: F-distribution, Normal distribution, Wald distribution, and Weibull distribution. We implemented two configurations for the F-distribution to examine the effect of varying denominator degrees of freedom while maintaining a

constant numerator. The first (no. 1) employs 5 degrees of freedom for the numerator and 10 for the denominator (F(5,10)), while the second (no. 2) remains the numerator at 5 but increases the denominator to 20 (F(5,10)). In our examination of the normal distribution, we maintained a consistent mean of zero while varying the standard deviation. Configuration no. 3 uses a standard deviation of 4, (N(0,4)), while configuration No. 4 employs a standard deviation of 1, (N(0,1)) allowing us to investigate the impact of different spread patterns under identical central tendencies. We explored location and scale effects for the Wald distribution through two distinct parameterizations. Configuration no. 5 combines a mean of 1 with a standard deviation of 0.5, while configuration no. 6 pairs with a mean of 0 with a standard deviation of 2. These contrasting specifications comprehensively examine the behavioral characteristics of the distribution under different conditions. For the Weibull distribution, we focused on the influence of the shape parameter on distributional characteristics. Two configurations were implemented: no. 5 with a shape parameter of 1, and no. 6 with a shape parameter of 5. This parameterization allows us to examine how varying shape values affect the distribution's tail behaviour and overall form. In addition, we set the population size to 10,000 for all datasets.

TABLE I: Description of statistical distribution's parameters, minimum and maximum values of the simulation scenario.

No.	Distribution	Minimum	Maximum
1	$F(\nu_1 = 5, \nu_2 = 10)$	0.0083	52.4746
2	$F(\nu_1 = 5, \nu_2 = 15)$	0.0199	83.5992
3	$\chi^2(\nu_1=2)$	0.0000	21.9290
4	$\chi^2(\nu_1 = 10)$	0.7468	33.6357
5	$N(\mu = 0, \sigma^2 = 16)$	-3.6943	4.2593
6	$N(\mu=0,\sigma^2=1)$	-14.2604	18.0044
7	$Wald(\mu = 1, \lambda = 0.5)$	0.0323	15.1948
8	$Wald(\mu=1,\lambda=2)$	0.1090	12.8305
9	Weibull(shape = 1)	0.1090	12.2323
10	Weibull(shape = 5)	0.1209	1.5980

In our experiments on real-world scenarios, we examined four distinct datasets representing various economic sectors, as detailed in Table II. These data sets were selected to evaluate our methodology in different real-world applications and value distributions. The first data set comprises laptop price data, with values ranging from 174.00 to 6,099.00. This substantial price range encompasses the diverse laptop market, from entry-

level consumer devices to high-end professional systems, providing a comprehensive representation of price distribution in the technology retail sector. The second data set focuses on electronic sales transactions from 20.75 to 11,396.80. This extensive range captures the full spectrum of electronic retail activity, from minor accessory purchases to significant investments in premium electronic equipment. Our third data set examines ecommerce sales data, with transactions ranging from 100.30 to 9,995.62. This range reflects the diverse nature of online retail transactions, demonstrating the breadth of consumer purchasing behaviour in the digital marketplace. The dataset's distribution provides valuable insights into e-commerce transaction patterns and consumer spending behaviours. The fourth dataset investigates world tourism economy indicators, ranging from 0.1578 to 28.1923. This dataset's unique scale and distribution characteristics offer an important perspective on our methodology's applicability to macroeconomic indicators.

TABLE II: Description of minimum and maximum values of the real-world scenario.

No.	Distribution	Minimum	Maximum
1	Laptop prices	174.00	6,099.00
2	Electronic sales	20.75	11,396.80
3	E-commerce sales	100.30	9,995.62
4	world tourism economy	0.1578	28.1923

A. Experimental results

For performance evaluation on range approximation in streaming data chunks, given a population \mathbb{P} with size N, let $\mathbf{X} = \{x_1, x_2, \dots, x_n\}$ be the set of samples, randomly selected from the population \mathbb{P} . We created the streaming data chunks with size L.

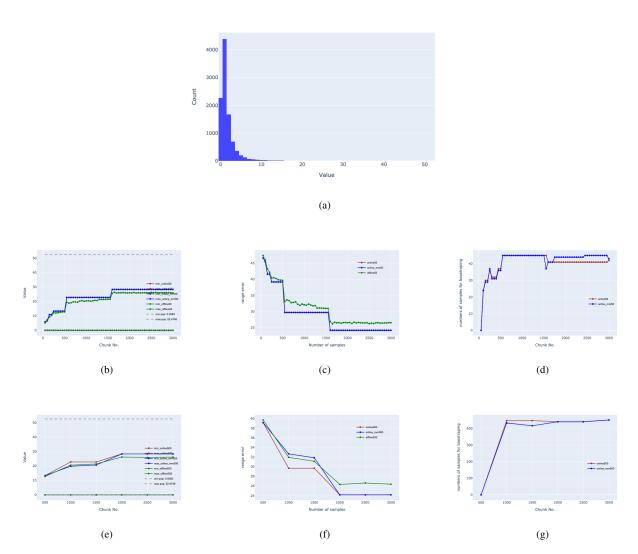


Fig. 3: Range approximation results of $F(\nu_1=5,\nu_2=10)$ including i) histogram of population data as shown in (a), ii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 50 samples/chunk, as shown in (b), (c), and (d), respectively, and iii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 500 samples/chunk, as shown in (e), (f), and (g), respectively.

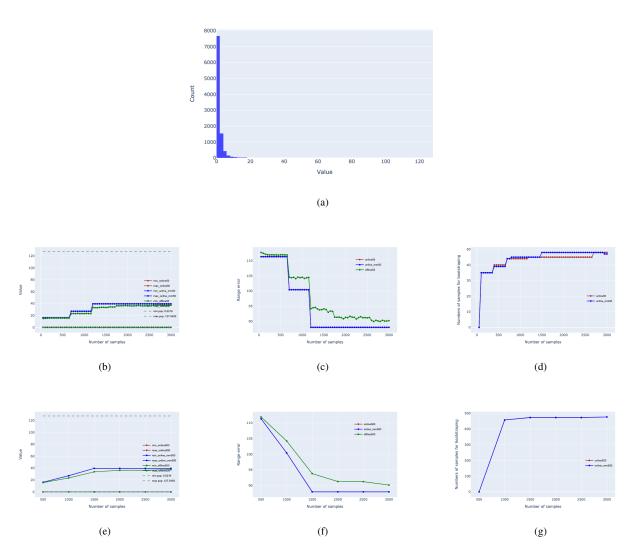


Fig. 4: Range approximation results of $F(\nu_1=5,\nu_2=20)$ including i) histogram of population data as shown in (a), ii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 50 samples/chunk, as shown in (b), (c), and (d), respectively, and iii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 500 samples/chunk, as shown in (e), (f), and (g), respectively.

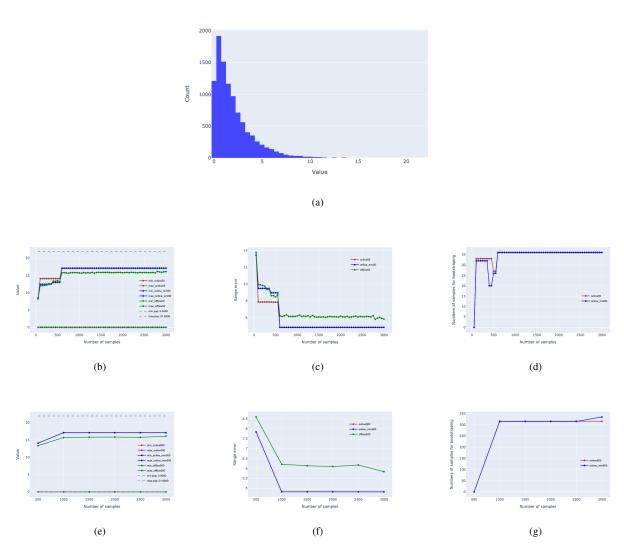


Fig. 5: Range approximation results of $\chi^2(\nu=2)$ including i) histogram of population data as shown in (a), ii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 50 samples/chunk, as shown in (b), (c), and (d), respectively, and iii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 500 samples/chunk, as shown in (e), (f), and (g), respectively.

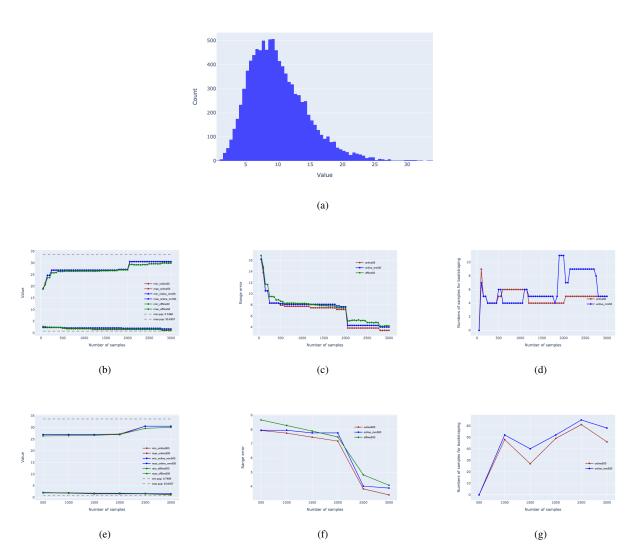


Fig. 6: Range approximation results of $\chi^2(\nu=10)$ including i) histogram of population data as shown in (a), ii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 50 samples/chunk, as shown in (b), (c), and (d), respectively, and iii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 500 samples/chunk, as shown in (e), (f), and (g), respectively.

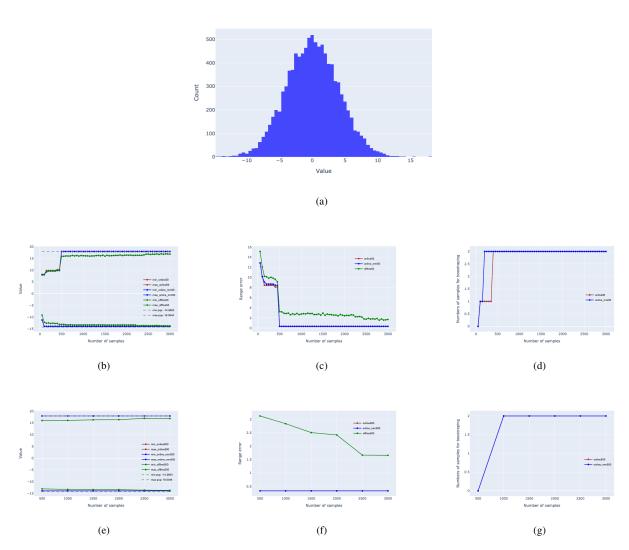


Fig. 7: Range approximation results of $N(\mu=0,\sigma^2=16)$ including i) histogram of population data as shown in (a), ii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 50 samples/chunk, as shown in (b), (c), and (d), respectively, and iii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 500 samples/chunk, as shown in (e), (f), and (g), respectively.

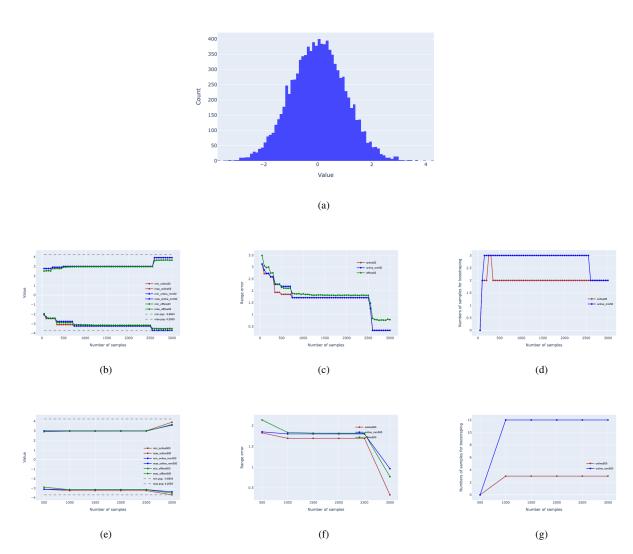


Fig. 8: Range approximation results of $N(\mu=0,\sigma^2=1)$ including i) histogram of population data as shown in (a), ii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 50 samples/chunk, as shown in (b), (c), and (d), respectively, and iii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 500 samples/chunk, as shown in (e), (f), and (g), respectively.

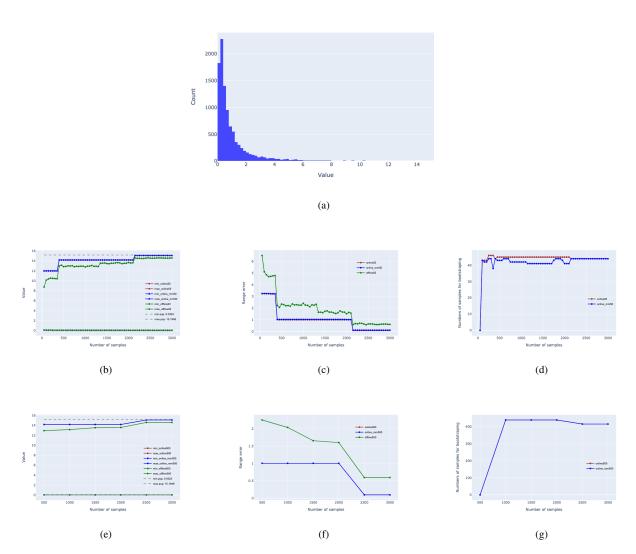


Fig. 9: Range approximation results of $Wald(\mu=1,\lambda=0.5)$ including i) histogram of population data as shown in (a), ii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 50 samples/chunk, as shown in (b), (c), and (d), respectively, and iii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 500 samples/chunk, as shown in (e), (f), and (g), respectively.

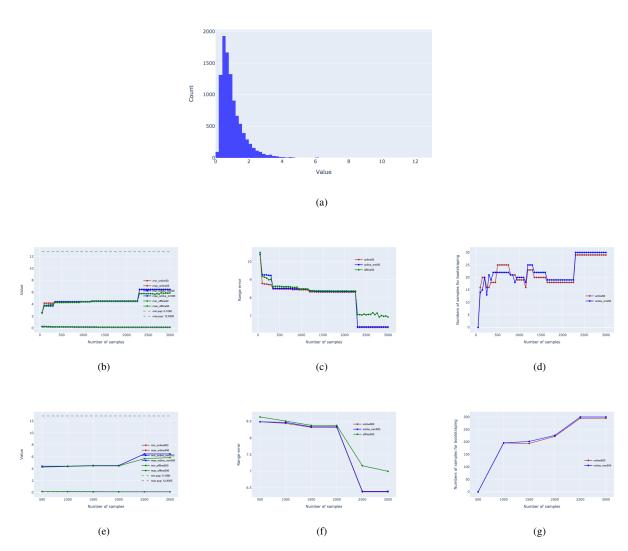


Fig. 10: Range approximation results of $Wald(\mu=1,\lambda=2.0)$ including i) histogram of population data as shown in (a), ii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 50 samples/chunk, as shown in (b), (c), and (d), respectively, and iii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 500 samples/chunk, as shown in (e), (f), and (g), respectively.

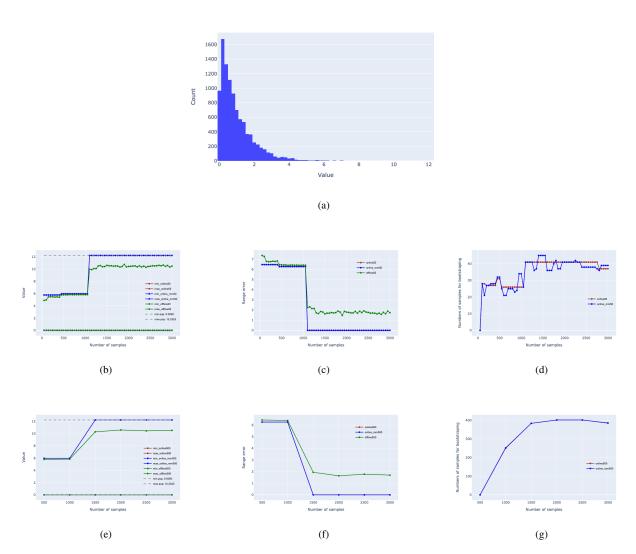


Fig. 11: Range approximation results of Weibull(shape=1) including i) histogram of population data as shown in (a), ii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 50 samples/chunk, as shown in (b), (c), and (d), respectively, and iii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 500 samples/chunk, as shown in (e), (f), and (g), respectively.

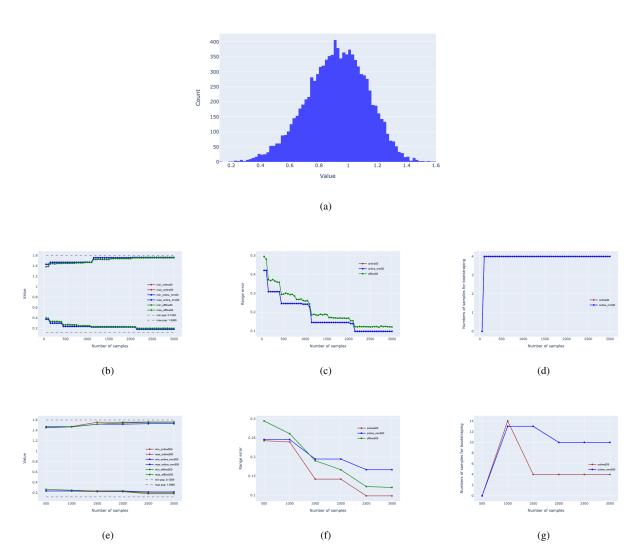


Fig. 12: Range approximation results of Weibull(shape=5) including i) histogram of population data as shown in (a), ii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 50 samples/chunk, as shown in (b), (c), and (d), respectively, and iii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 500 samples/chunk, as shown in (e), (f), and (g), respectively.

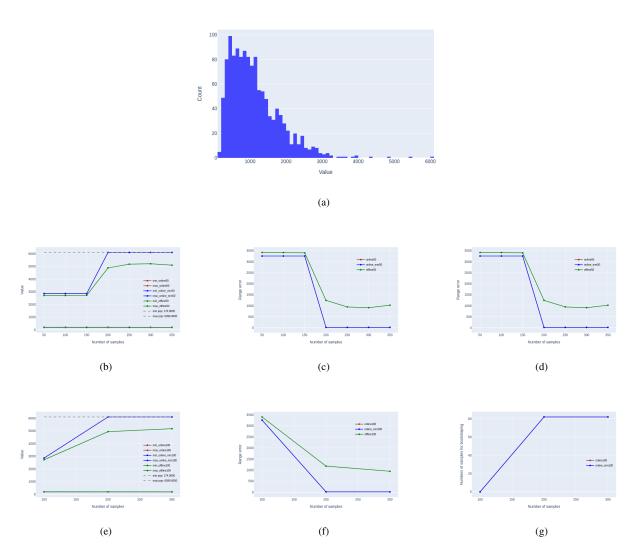


Fig. 13: Range approximation results of laptop prices including i) histogram of population data as shown in (a), ii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 50 samples/chunk, as shown in (b), (c), and (d), respectively, and iii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 500 samples/chunk, as shown in (e), (f), and (g), respectively.

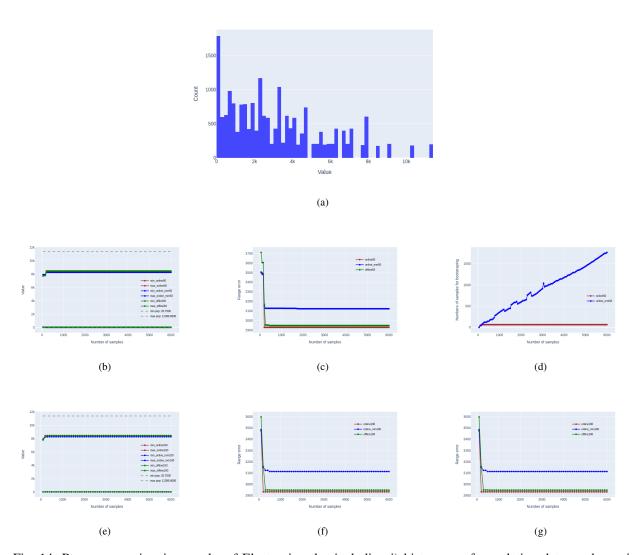


Fig. 14: Range approximation results of Electronic sales including i) histogram of population data as shown in (a), ii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 50 samples/chunk, as shown in (b), (c), and (d), respectively, and iii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 500 samples/chunk, as shown in (e), (f), and (g), respectively.

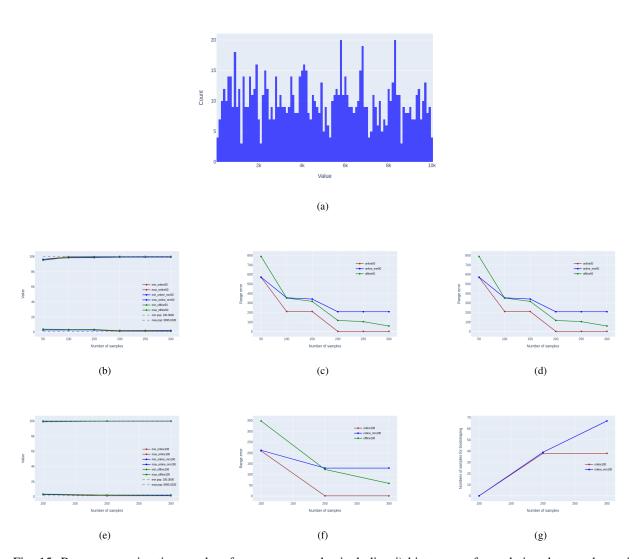


Fig. 15: Range approximation results of e-commerce sales including i) histogram of population data as shown in (a), ii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 50 samples/chunk, as shown in (b), (c), and (d), respectively, and iii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 500 samples/chunk, as shown in (e), (f), and (g), respectively.

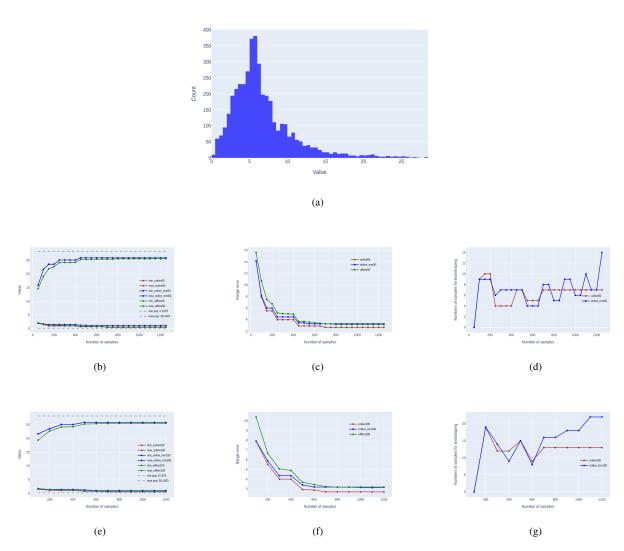


Fig. 16: Range approximation results of world tourism economy including i) histogram of population data as shown in (a), ii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 50 samples/chunk, as shown in (b), (c), and (d), respectively, and iii) min-max approximation, range error and number of leftmost and rightmost bins for bootstrapping, for 500 samples/chunk, as shown in (e), (f), and (g), respectively.