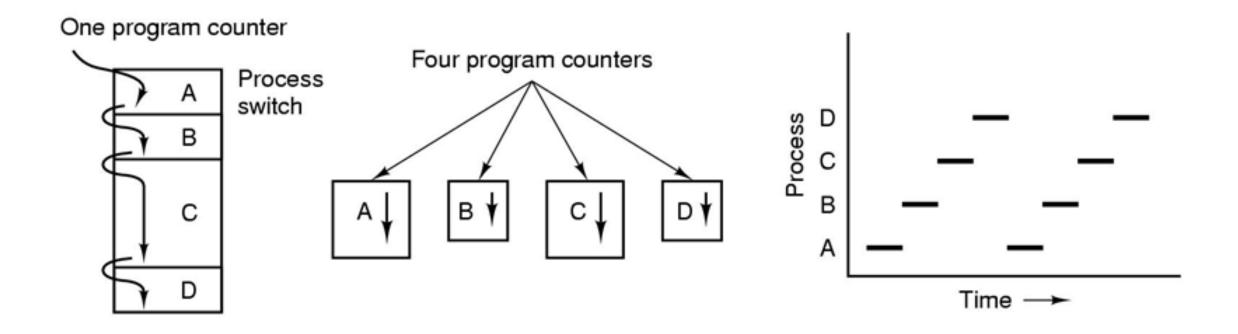


Session Outline

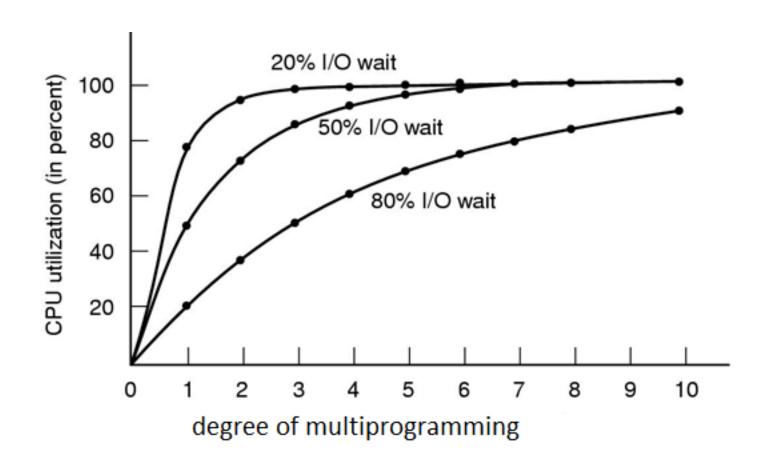
- Concept of Multiprogramming
- CPU vs I/O bound Processes
- CPU vs Job Schedulers
- CPU Scheduling
- Dispatcher
- Preemptive vs Non-preemptive Scheduling
- Scheduling Criteria

Multiprogramming

- In a uni-processor system, CPU is running only one process.
- In a multiprogramming system, CPU switches from processes quickly.

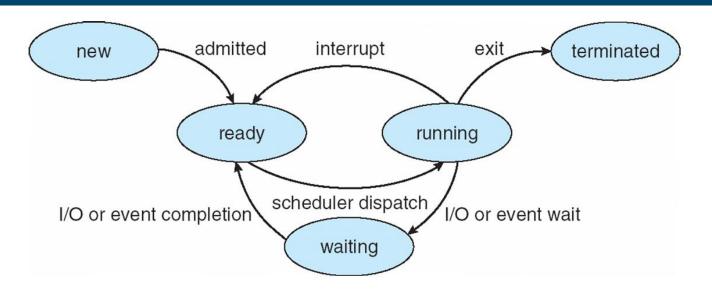


Modeling Multiprogramming



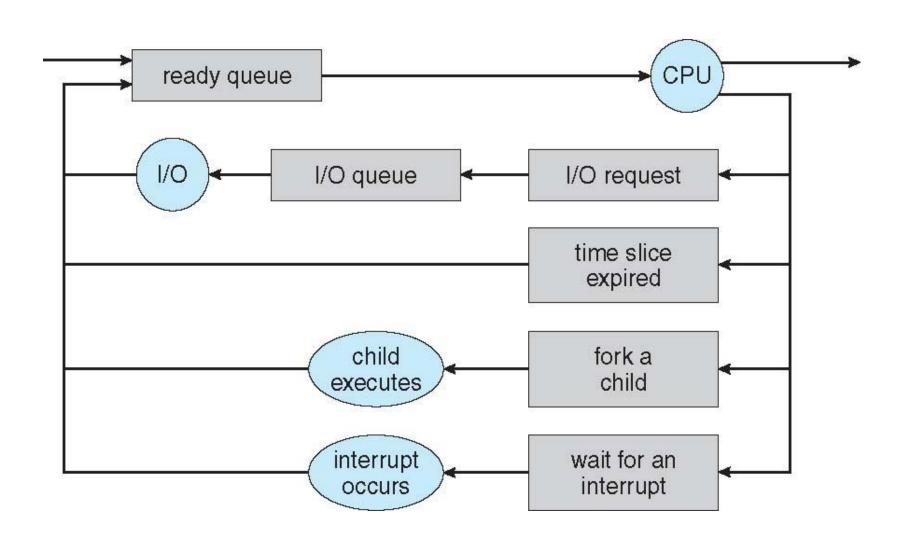
CPU utilization as a function of the number of processes in memory.

Process State Diagram

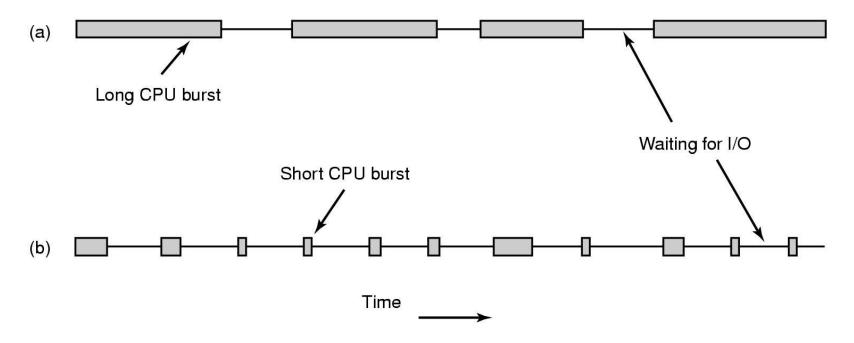


- new: The process is being created
- running: Instructions are being executed
- waiting: The process is waiting for some event to occur
- ready: The process is waiting to processor assignment
- terminated: The process has finished execution

Representation of Process Scheduling



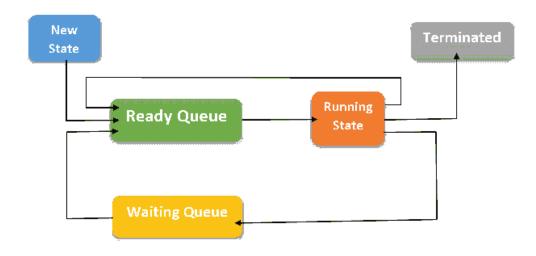
CPU vs I/O Bound Processes



Bursts of CPU usage alternate with periods of waiting for I/O. (a) A CPU-bound process. (b) An I/O-bound process.

CPU vs Job Schedulers

- Short-term scheduler (CPU scheduler) selects from among the processes that are ready to execute and allocates the CPU to one of them.
 - Selection from Ready state to Running state
 - Short-term scheduler is invoked frequently



- Long-term scheduler (Job scheduler) selects which processes should be brought into the ready queue from the job queue
 - Long-term scheduler is invoked less frequently
 - It controls the degree of multiprogramming

CPU Scheduler

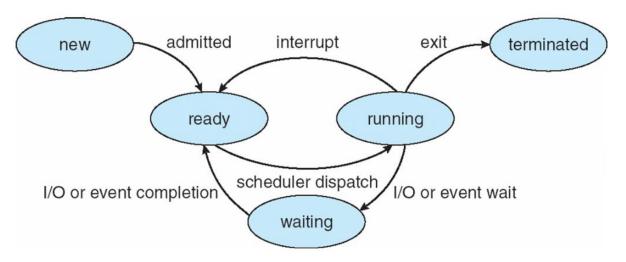
- Whenever the CPU becomes idle, the OS must select one of the processes in the ready queue to be executed.
- The selection process is carried out by the CPU scheduler of the OS.
- The scheduler selects a process that is in ready state and allocates the CPU to that process.
- Ready queue can be implemented as
 - FIFO queue
 - Priority queue
 - Tree
 - Unordered linked list

Dispatcher

- The dispatcher is the module that gives control of the CPU to the process selected by the CPU scheduler.
- Function of dispatcher involves the following:
 - Switching context
 - Switching to user mode
 - Initiate restarting/ resumption of selected user program
- Dispatcher should be as fast as possible, since it is invoked during every process switch.
- ❖ The time it takes for the dispatcher to stop one process and start another running is known as the dispatch latency.

Preemptive vs Non-preemptive Scheduling

CPU-scheduling decisions happen during the four circumstances:



- When a process switches from the running state to the waiting state
- 2. When a process switches from the running state to the ready state
- 3. When a process switches from the waiting state to the ready state
- 4. When a process terminates

Preemptive vs Non-preemptive Scheduling

- CPU-scheduling decisions happen during the four circumstances:
 - When a process switches from the running state to the waiting state
 - 2. When a process switches from the running state to the ready state
 - 3. When a process switches from the waiting state to the ready state
 - 4. When a process terminates
- 4 1 & 4, the scheduling is non-preemptive (cooperative)
- ❖ A running process keeps the CPU until it releases the CPU (terminating or by switching to the waiting state).
- 2 & 3, the scheduling is pre-emptive
- The process is removed from running state forcefully.

Scheduling Criteria

- Different CPU-scheduling algorithms have different properties.
- Certain characteristics/criteria are used for comparing various CPU scheduling algorithms.
 - CPU Utilization
 - Throughput
 - Turnaround time
 - Waiting Time
 - Response Time

Scheduling Criteria

- ❖ CPU Utilization Percentage time CPU is busy executing process.
 ↑
- Throughput Number of processes that are completed per time unit.
- Turnaround time The interval from the time of submission of a process to the time of completion.
- Waiting Time Amount of time that a process spends waiting in the ready queue.

- Response Time Time from the submission of a request until the first response is produced.



Thank You