

# DBMS Assignment — Unit 5 (Solved)

Prepared by: ChatGPT (answers based on the uploaded assignment).

## Fill in the blanks

**Q:** 1. A \_\_\_\_\_ dependency is a constraint between two attributes in a database.

**A:** functional dependency

**Q:** 2. A functional dependency of the form  $A \rightarrow A$  is called a \_\_\_\_\_ dependency.

**A:** trivial dependency

**Q:** 3. Armstrong's axioms include Reflexivity, Augmentation, and \_\_\_\_\_.

**A:** Transitivity

**Q:** 4. The closure of a set of functional dependencies is the set of all functional dependencies that can be \_\_\_\_\_ from the given set.

**A:** inferred

## Short Questions

### 1. What is a trivial functional dependency?

A functional dependency  $X \rightarrow Y$  is trivial if  $Y$  is a subset of  $X$ . For example,  $\text{RollNo}, \text{Name} \rightarrow \text{RollNo}$  is a trivial dependency because  $\text{RollNo} \subseteq \{\text{RollNo}, \text{Name}\}$ .

### 2. What is a candidate key?

A candidate key is a minimal set of attributes that uniquely identifies a tuple in a relation. There may be multiple candidate keys, one of which becomes the primary key.

### 3. What is a lossless decomposition?

A decomposition of a relation  $R$  into  $R_1$  and  $R_2$  is lossless if no information is lost during decomposition and join. It ensures that  $R$  can be perfectly reconstructed from  $R_1$  and  $R_2$  using a natural join.

### 4. How does BCNF differ from 3NF?

3NF allows some redundancy if the non-prime attribute is transitively dependent on a key through a superkey. BCNF is stricter — every determinant must be a candidate key, ensuring minimal redundancy.

### 1. What is functional dependency? Explain all its types with example.

A **functional dependency (FD)** is a constraint between two sets of attributes in a relation. It states that if two tuples have the same values for attribute(s)  $X$ , then they must also have the same values for attribute(s)  $Y$ , written as  $X \rightarrow Y$ . **Types of Functional Dependencies:**

1. **Trivial FD:**  $Y \subseteq X$  (e.g.,  $\text{RollNo}, \text{Name} \rightarrow \text{RollNo}$ ).

2. **Non-trivial FD:**  $Y \not\subseteq X$  (e.g.,  $\text{RollNo} \rightarrow \text{Name}$ ).

3. **Transitive FD:** If  $X \rightarrow Y$  and  $Y \rightarrow Z$ , then  $X \rightarrow Z$ .

4. **Multivalued Dependency:** A dependency where one attribute determines multiple independent attributes.

Example: In  $\text{Student}(\text{RollNo}, \text{Name}, \text{Dept})$ ,  $\text{RollNo} \rightarrow \text{Name}$  means  $\text{Name}$  depends on  $\text{RollNo}$ .

## 2. What is normalization in DBMS? Also discuss the need of normalization. Explain 2NF, 3NF, BCNF with suitable example.

**Normalization** is the process of organizing data in a database to minimize redundancy and improve data integrity. It divides large tables into smaller, related tables and defines relationships among them. **Need for Normalization:**

- To eliminate redundancy
- To prevent anomalies (update, insertion, deletion)
- To improve consistency and efficiency. **2NF (Second Normal Form):** A relation is in 2NF if it is in 1NF and no non-prime attribute is partially dependent on a candidate key.

Example: Student(RollNo, Subject, Marks). Here, RollNo → Name and (RollNo, Subject) → Marks. To remove partial dependency, split into Student(RollNo, Name) and Marks(RollNo, Subject, Marks). **3NF (Third Normal Form):** A relation is in 3NF if it is in 2NF and there is no transitive dependency of non-prime attributes on a key.

Example: Emp(EmpID, DeptID, DeptName). DeptName depends on DeptID, not directly on EmpID, so it's transitive. Decompose into Emp(EmpID, DeptID) and Dept(DeptID, DeptName).

**BCNF (Boyce-Codd Normal Form):** A relation is in BCNF if every determinant is a candidate key. It removes all redundancy allowed in 3NF.

Example: In R(A,B,C) with FDs A → B and B → A, only BCNF decomposition ensures both A and B are keys.

## 3. What is the need for Normalization? Explain 1NF, 2NF and BCNF.

Normalization is needed to remove redundancy, prevent anomalies, and maintain data integrity. **1NF:** A table is in 1NF if it contains only atomic (indivisible) values and each column contains values of a single type. Example: Repeating groups are removed. **2NF:** A table is in 2NF if it is in 1NF and all non-key attributes are fully functionally dependent on the primary key.

**BCNF:** A table is in BCNF if for every functional dependency X → Y, X is a superkey. It removes all anomalies possible in 3NF.

## 4. What is closure of a set of attributes? Explain how to find closure of a set of attributes.

The **closure** of an attribute set X, denoted  $X^+$ , is the set of all attributes that can be functionally determined from X using a given set of functional dependencies. **Steps to find closure:**

1. Start with  $X^+ = X$ .
2. For each functional dependency  $Y \rightarrow Z$ , if  $Z \subseteq X^+$ , then add Z to  $X^+$ .
3. Repeat until no new attributes can be added.

**Example:** R(A, B, C, D), FDs = {A → B, B → C}. Then  $A^+ = \{A, B, C\}$ .

## 5. What is meant by normalization? Write its need. List and discuss database anomaly during database design.

Normalization is the process of structuring a database to reduce redundancy and improve integrity. **Need:** To remove duplication and anomalies. **Database anomalies:**

- **Insertion anomaly:** Cannot insert data due to missing related data.
- **Deletion anomaly:** Deleting a record unintentionally removes useful data.
- **Update anomaly:** Inconsistency due to data redundancy when updating records.

## 6. List and explain Armstrong's axioms.

Armstrong's axioms are a set of inference rules used to derive all possible functional dependencies from a given set:

1. **Reflexivity:** If  $Y \subseteq X$ , then  $X \rightarrow Y$ .
2. **Augmentation:** If  $X \rightarrow Y$ , then  $XZ \rightarrow YZ$ .
3. **Transitivity:** If  $X \rightarrow Y$  and  $Y \rightarrow Z$ , then  $X \rightarrow Z$ .

**Additional derived rules:**

- **Union:** If  $X \rightarrow Y$  and  $X \rightarrow Z$ , then  $X \rightarrow YZ$ .
- **Decomposition:** If  $X \rightarrow YZ$ , then  $X \rightarrow Y$  and  $X \rightarrow Z$ .
- **Pseudotransitivity:** If  $X \rightarrow Y$  and  $YZ \rightarrow W$ , then  $XZ \rightarrow W$ .

## 7. What is an anomaly in database design? How it can be solved.

An **anomaly** is an inconsistency or problem in database operations caused by data redundancy.

- **Update anomaly:** Multiple updates required to maintain consistency.
- **Insertion anomaly:** Inability to add new data due to dependency on other data.
- **Deletion anomaly:** Loss of data when deleting unrelated records.

**Solution:** Anomalies can be resolved through normalization (1NF, 2NF, 3NF, BCNF), which ensures that data is stored in well-structured, non-redundant tables.

References:

Source: Assignment Unit 5 uploaded by user.