<u> </u>	TOPIC
]	DATE
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-Part A	Operating System Arignment 01
fuesi	Despite systems ?
Ans. 1	Moderer systems still rely heavily on.
	operating systems because :-
	Resource management > The OS efficiently - manages hardware resources (CPU, memory,
Contract of the Contract of th	storage, T/O devices) and allocates them to
	different programs.
- 1	Iser & Application Interface > Mrs as by sinder -
	a roundine interface blue hardurare and
	making prilability
	rultitasking and security.
fusz. 7	Ou _ why?
Ans 2. 7	Real - Time Operating System (RTOS)
K	TOS ensures timely, predictable, and reliable
\!	espouse to inputs like heart rate signals processes
- da	ata with low latency, provides quint resource
	male , tow - power hardware.
	itual for health monitoring devices.
us3. lui	New why?
n83. Avo	rd a Monolithic Kirkel, while it give fast sestem
	is, they lack moderably and are narder to main-
flau	a /debug. A bug in one service can vaile
- the	a /debug. A bug in one service can varb
leut	ical système.

- ques Refute the claim, because os structure Ans4 directly impacts performance, rediability

scalability and security

for en:- microkernel isolates services for fault tolerance, while a layred structure improves maintainability. Just "running processes" interest enough if the system is low, unsecure or unstable i) Explain states. The PCB stores CPV registers, program country, state, and nemory infor By enamining it, we can detect misinitialized registers, wrong state glags, coros incorrect program country. Values that cause faulty switching involve ? to waiting, content switching saves the current process state (registers, program counter, PCB updates) and loads the state of the next process. It ensures encution resumes correctly later. iii) The and why? Use an asynchronous, non-blocking system call because this allows the process to continue enecution while the I/O is the CPV from idling.

11									
0 10	TOPIC								
Parts	DATE								
-qus:6	Total content switching time,								
a)	Total content it								
	save state = 2ms								
	hoad state = 3ms								
-	Scheduler overhead = 1 ms								
	- LWS								
	Total time = 2+3+1 = 6 ms								
	D I UWW								
	Explain bertognioner								
	content switching is have overhead (no useful								
	work is done during this time)								
-	Explain performance. content switching is pure overhead (no useful work is clone during this time) Higher switching time reduces CPV efficiency, as more time is spent switching than enecuting brocuses.								
	more time is spent switching than executing								
	riocitis.								
	In multitasking, prequent content switches with righ overhead can slow down thoroughput and inverse response time.								
	righ overhead can slow down thoroughput and								
	universe response time.								
Jus 7.									
- given:	- enecution time (single-threaded) = 40 sec								
	- enecution time (single-threaded) = 40 sec Whithreading is used with n threads per process								
÷ .									
	recution time estimate:								
a	r ideal conditions (perfect parallelism, no overhead):								
	To LA								
	Tmulti = Tsingle = 40 seconds								
100	Y								
Example	20 - 2 - 2 00 100								
of of	$\frac{1}{2} = \frac{1}{2} = \frac{1}$								
- In	$n = 8 \rightarrow 5$ Sec								
- It	n - D / D MC								

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	TOPIC
	How multithreading improves performance. It improves performance by running tasks in parallel, reducing encertions time. It keeps the CPU being even during I/O waits avoiding idle time. Threads share resources, making encertion fastir and more efficient. Process: P1 P2 P3 P4 Burst time: 5 3 8 6
(,0,)	FCFS Burst Completion Waiting TAT Frocess Arrival Burst Completion Waiting TAT time (AT) June (BT) June (CT) June (WT) 5 5-5=0 5-0=5
<u></u>	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$
-	WT = Turnavourd - Burst (TAT-BT) TAT = completion - Arrival (CT-AT)
	Grant Chart P1 P2 P3 P4 P4 P3 P4 P4 P4 P4
) (T(P1) CT(P3) CT(P4)

Arr 120 to 10 = (0+5+8+16)/4 = 7.25 mb									
	Avg waiting lime = (0+5+8+16)/4 = 7.25 ms Avg turnaround hime = (5+8+16+22)/4 = 12.75 ms								
	They turnavoura now								
()	Non-	Driem ble	ve SJ	E					
		7			*				
	Process	AT	BT	CT	WT	TAT			
			,	>1					
	P ₁		5	8	8-5 = 3	8=8-0			
	P2	0	_3	3	3-3=0	3 = 3-6			
)	P3	0	8	22	1	22 = 22-0			
	Py	0	6	14	14-6 = 8	14 = 14-0			
4		•							
	TAT = CT - AT								
	WT = TAT - BT								
Grant chart,									
		P1 Pi		3					
	0 3 8 14 22								
•	CT(B) CT(P4) CT(P3)								
	Ava waiting time = (3+0+14+8) /4 = 6.25 mz								
Avg waiting time = $(3+0+14+8)/4 = 6.25 \text{ ms}$ Avg turnaround time = $(8+3+22+14)/4 = 11-75 \text{ ms}$									
T)		× ×							

V							
	TOPIC			+	· 4ms)		
(2)	Round	Robin		T CT	[WI	TAT	; ;
6	Process	# T	BT		16-5 = 11	16-0=16	
	P ₁ P ₂ P ₃ P ₄	0 0	5 3 8 6	16 7 20 22	7-3=4 $20-8=12$ $22-6=16$	7-0 = 7 $20-0 = 20$ $22-0 = 22$	M.
<u></u>	WT = 7 TAT =	AT - BT CT - AT					- - - -
	lyrantt P1 P2 0 4	Chart P3 P 7 11 L, CT(P2)	7 4 P ₁ 15	P3 16 2 16 2 CT(P1)	P4 2 2 CT (P3)	→ CT(P4)	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
<u> </u>	Avg waiti	ng time	= (11+4 = (16+	+12 +16.) /4 = 16 +22) /4 = 1	0.75 ms 6.25 ms	
*	Non-p gives (6.25	reempti the la ms) a thr	re SJ nvest nd ti oughpe	F is averag irnar it re	bist bles e. waiter ound tin mains t	e live Le CII. 75m Le Same balances better	ردن



