

First Assignment

Samstag, 22. Mai 2021 09:39

1. Assignment 1 in Latex

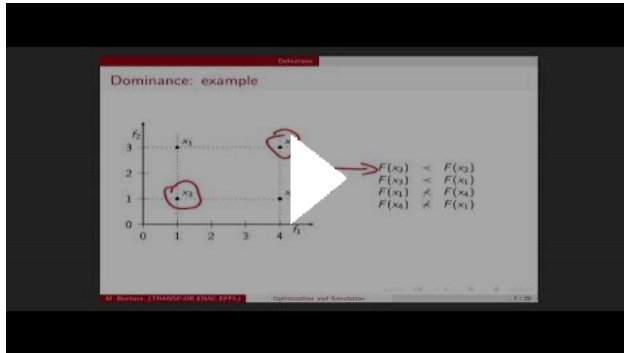
<https://www.overleaf.com/project/609a5574e7374883e5d3f3b5>

2. <https://developers.google.com/optimization/routing/vrp>

3. <https://www.sciencedirect.com/topics/computer-science/weighted-sum-method#:~:text=One%20of%20the%20most%20commonly,generate%20multiple%20Pareto%20optimal%20solutions%2C>

4. [Optimization and simulation. Multi-objective optimization - part 1](#)

5.



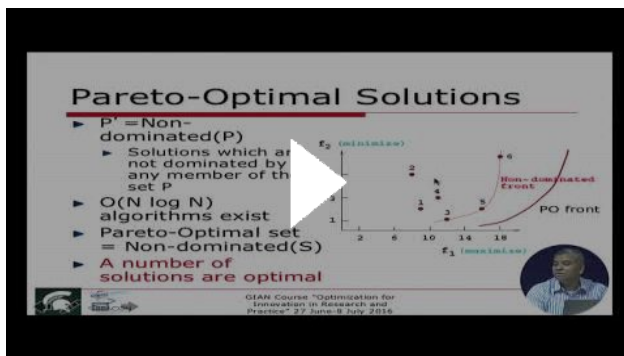
6. <https://de.mathworks.com/help/optim/ug/generate-and-plot-a-pareto-front.html>

7. **Soft constraint:** another approach for solving supplying problem with preference extension recommended from Xie!!! --> try it later

<https://or.stackexchange.com/questions/1050/soft-constraints-and-hard-constraints>

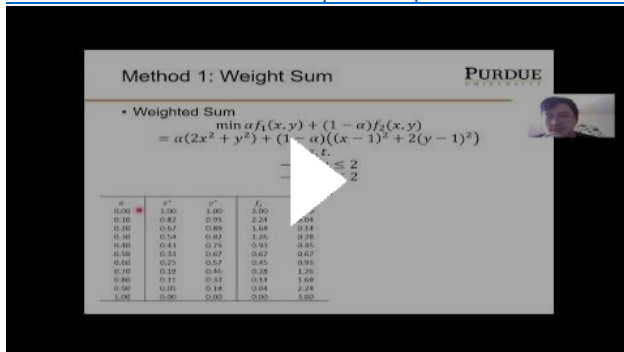
8. Pareto optimal set

[23. Multiobjective Optimization](#)



9. Good for looking for Pareto-optimal set

10. [MET 503 Lecture 18: Multi-Objective Optimization Problem](#)



11. Soft constraint

12. https://www.gams.com/latest/docs/UG_EMP_SoftConstraints.html

13. Non-dominated sorting genetic algorithm 21:08

<https://www.youtube.com/watch?v=Hm2LK4vJzRw>

14.

In order to ensure that there is only one truck travelling on one arc at a time, one more decision variable $s_{ij}^{v'k}$ is defined. $s_{ij}^{v'k} = 1$ if truck v travelling on arc (i, j) before truck v' , and it takes value 0 otherwise. Using big M technique, this constraint is presented as follows:

- Only one truck travelling on one arc at a time

$$w_j^{rk} \leq w_i^{v'k} + M(2 - y_{ij}^{rk} - y_{ij}^{v'k}) + M(1 - s_{ij}^{v'k}), \quad (i, j) \in A, v, v' \in V, k \in K \quad (1.16)$$

1. Because truck v going on arc (i, j) before truck v' , the time windows for two trucks have to be satisfied the following time axis



To ensure that the 2 time intervals (windows) should not be overlapped.

2. About the relation between y and s :

$$\begin{aligned} \Rightarrow \text{If } y_{ij}^{rk} = 1, y_{ij}^{v'k} = 1 &\Rightarrow s_{ij}^{v'k} = \begin{cases} 0 \\ 1 \end{cases} \\ \Rightarrow \text{If } y_{ij}^{rk} + y_{ij}^{v'k} \leq 1 &\Rightarrow s_{ij}^{v'k} = 0 \end{aligned}$$

Using big M technique: we have the constraints for y and s is:

$$y_{ij}^{rk} + y_{ij}^{v'k} \leq 1 + s_{ij}^{v'k} \cdot M$$

M needed to be larger than 1!

$$\begin{aligned} \text{Cobot 0} &: \left\{ [outD_0, 0, 3, 7, \boxed{10}, 12, 17, 22], [outD_0, 1, 3, 4, 11, 13, 16, 23, outD_0] \dots \right\} \\ \text{Cobot 1} &: \left\{ [outD_1, 23, 20, 17, \boxed{12}, 9, 6, 1, outD_1], [outD_1, 22, 10, 9, 0, outD_1] \dots \right\} \end{aligned}$$

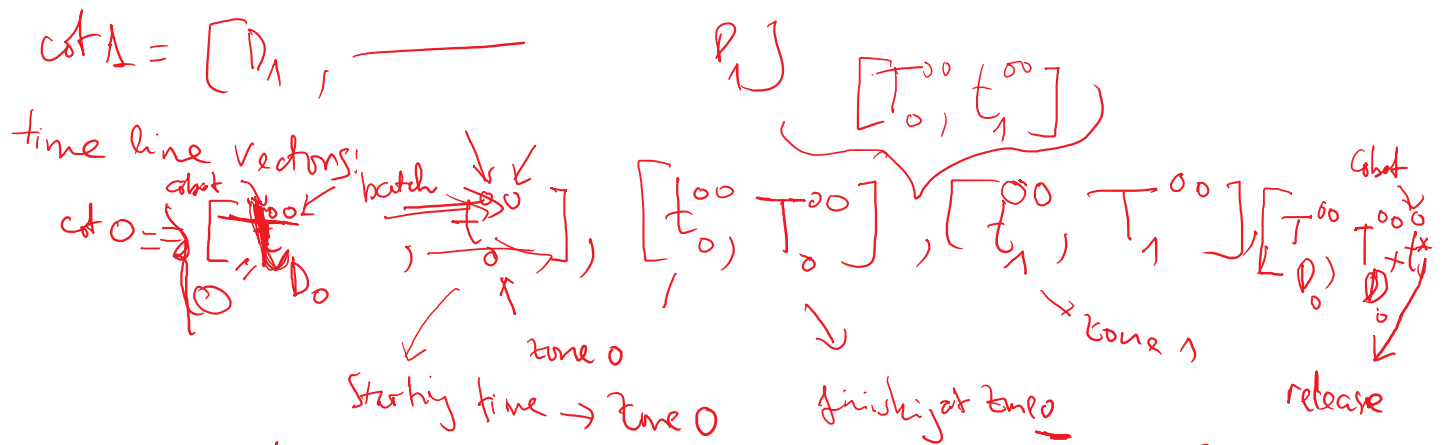
* Schedule policy:



27.7.21

$$\text{Cobot 0} = [D_0, \text{Zone 0}, \text{Zone 1}, D_0]$$

$$\text{Cobot 1} = [D_1, \text{Zone 0}, \text{Zone 1}, D_1]$$



$$[T_0^{00} + t_0^{01}, t_0^{01}], [t_0^{01}, T_0^{01}], \dots$$

([['2', '0', '3'], ['OutD0', '1', '2', '10', '9', '15', 'OutD0'], 1),
 ([['8'], ['OutD0', '15', '16', '11', '2', '7', '1', 'OutD0'], 2),
 ([['6'], ['OutD0', '15', '10', '7', 'OutD0'], 3),
 ([['1'], ['OutD0', '7', '1', '4', '9', '16', 'OutD0'], 4])

From http://localhost:8888/notebooks/Desktop/3.OR%20Analysing%20networks/2.%20Project/3.Code/Task21distance_final.ipynb

Timeline:

$$\rightarrow cot_0 = [T_0^{00}, t_0^{00}, T_0^{00}, t_1^{00}, T_1^{00}, T_0^{00}, t_0^{01}, t_0^{01}, T_0^{01}, t_1^{01}, T_1^{01}, t_1^{01}]$$

$$\rightarrow cot_1 = [t_1^{00}, t_0^{10}, T_0^{10}, t_1^{10}, T_1^{10}, T_0^{10}, \dots]$$

4.2 \rightarrow balancing # batches in each depot
 (Done)

$$\rightarrow cot_0 = [[T_0^{00}, T_0^{00}], [t_1^{00}, T_1^{00}], [t_0^{01}, T_0^{01}], [t_1^{01}, T_1^{01}], \dots]$$

Cobot1 and cobot2
 Could be lists of lists

looping
 to calculate
 waiting time
 (applys queue policy)

$$cot_1 = [T_0^{10}, T_0^{10}], [t_1^{10}, T_1^{10}], [(None, None), (None, None)]$$

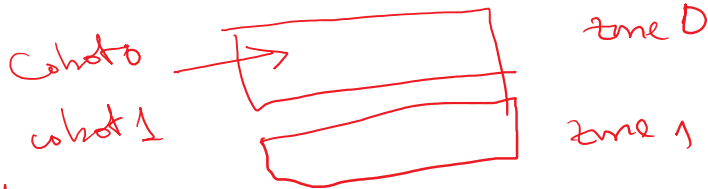
How to calculate the waiting time and the time points in the time line?

$$T_0^{10}, T_0^{10}, T_1^{10}, T_1^{10}, \dots, waiting time$$

How to calculate the waiting time and the time points in the time line?

- ① Calculate those t, T, T_0, T_1, t_0, t_1 without waiting time
 → need to extract shelves in zone 0, 1 for each batch.
2. Query policy:
 → zone 0 → zone 1

Keep the order of the shelves to
Calculate the initial timeline
Without waiting time



3. if else statement for the first pair of intervals of cobot 0 and cobot 1:

if $(t_0^{00} \leq t_0^{10})$:
 $\text{cobot}_1[1:] = T_0^{10} + = T_0^{00} - t_0^{00}$ (servicing time in zone 0 for cobot 0)
 $\text{cobot}_1[2:] += T_0^{00} - t_0^{00}$ (from the time point T_0^{10} on, timeline of cobot 1 is updated by the servicing time in zone 0 for cobot 0)

else:
 $\text{cobot}_0[1:] += T_0^{10} - t_0^{00}$

4. Working with the loops and the two timeline vectors of the cobots

4. a) make the two cobot vectors same length
 len

Reminder

an

If we have a tour like this

$(['1'], ['OutD0', '7', '1', '4', '9', 'OutD0'], 4)$

zone 0

→ Then the interval for zone 1 of this tour would be (None, None)

Reminder

a2: $(['2', '0', '3'], ['OutD0', '1', '2', '10', '9', '15', 'OutD0'], 1)$

← subsection 4.1

(which items → subsection 4.2)
 How many of ~~these~~ items picked from this shelf

→ Remember to save the quantities of the items in the step of eliminating duplication in the Function greedyTour

b) Loop over the batches (or tours)

For batch_0 in cobot0:

3. if else statement for
 $\text{if } (t_0^{00} \leq t_0^{10}) :$

for batch_1 in cobot1:

if batch_0[0][0] <= batch_1[0][0] (see 3. in if statement):
(update)

$\text{cobot}_1[1] = T_0^{10} += T_0^{00} - t_0^{00}$ (servicing in zone 0 for cobot 0)
 $\text{cobot}_1[2:] += T_0^{00} - t_0^{00}$ (from the time point T_0^{10} on,
timeline of cobot 1 is updated by
the servicing time in zone 0 for cobot 0)

5. Structure of the project

- Task 2.1: generate the initial solution s . Please accomplish the last function with for loop for the depots.
- Task 2.2: generate a neighborhood of s (subsection 4.2) + local search: compare $f(s)$ and $f(s')$, where s' is from the neighborhood of s (subsection 4.3 in paper, set some maximal number of iterations): we need to calculate the makespan(f).
- Task 2.3: would be difficult: mixed shelf policy.