



Data Communications and Networking

Fourth Edition

Chapter 19

Network Layer: Logical Addressing



19-1 IPv4 ADDRESSES

An **IPv4 address** is a **32-bit** address that uniquely and universally defines the connection of a device (for example, a computer or a router) to the Internet.

Topics discussed in this section:

Address Space

Notations

Classful Addressing

Classless Addressing

Network Address Translation (NAT)



IPV4 Address

- ❑ An IPv4 address is 32 bits long.
- ❑ The IPv4 addresses are unique and universal.
- ❑ The address space of IPv4 is 2^{32} or 4,294,967,296

IPv4 Address

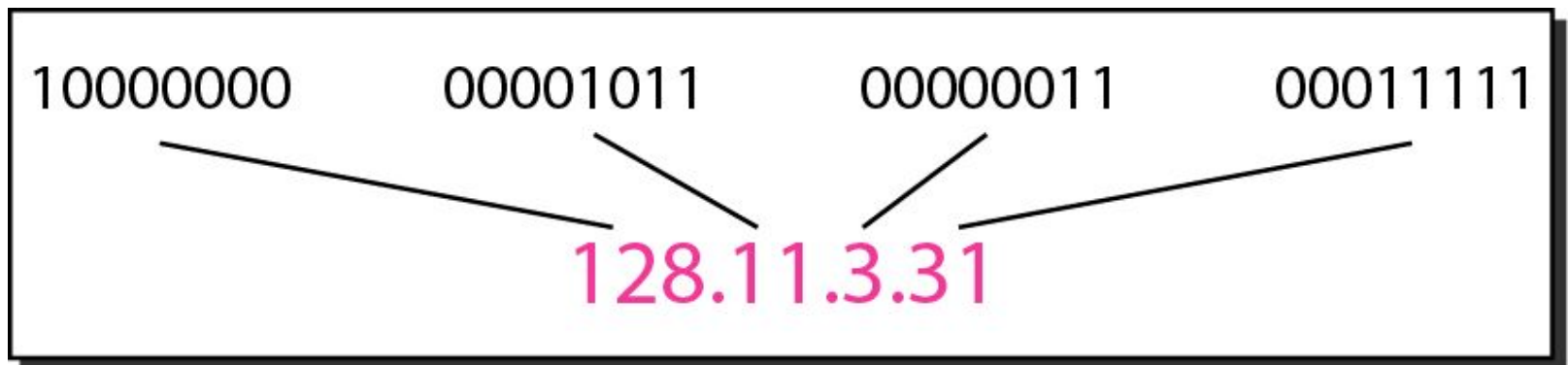
□ Notation

◆ Binary Notation :

- The IPv4 address is displayed as 32 bits.
- Each octet is often referred to as a byte.

◆ Dotted-Decimal Notation

- To make the IPv4 address more compact and easier to read, Internet addresses are written in decimal form with a decimal point (dot) separating the byte.



Notation (cont'd)

□ Hexadecimal Notation

0111 0101 1001 0101 0001 1101 1110 1010

75

95

1D

EA

0x75951DEA

- *8 hexadecimal digits*
- *Used in network programming*

Classful Addressing

- ❑ In classful addressing, the address space is divided into five classes: A, B, C, D, and E.
- ◆ If the address is given in binary notation, the first few bits can tell us the class of the address.
- ◆ If the address is given in decimal-dotted notation, the first byte defines the class.

	First byte	Second byte	Third byte	Fourth byte
Class A	0			
Class B	10			
Class C	110			
Class D	1110			
Class E	1111			

a. Binary notation

	First byte	Second byte	Third byte	Fourth byte
Class A	0–127			
Class B	128–191			
Class C	192–223			
Class D	224–239			
Class E	240–255			

b. Dotted-decimal notation

Classful Addressing

❏ Classful Addresses

- ❖ Unicast Communication - A, B, C Class
(~must be delivered to specific computer)
- ❖ Multicast Communication – D Class
(~must be delivered to each member of the group)
- ❖ For reserve – E Class

Classful Addressing

❏ Classes and blocks

- ◆ One problem with classful addressing is that each class is divided into a fixed number of blocks with each block having a fixed size.

Table 19.1 *Number of blocks and block size in classful IPv4 addressing*

<i>Class</i>	<i>Number of Blocks</i>	<i>Block Size</i>	<i>Application</i>
A	128	16,777,216	Unicast
B	16,384	65,536	Unicast
C	2,097,152	256	Unicast
D	1	268,435,456	Multicast
E	1	268,435,456	Reserved



Classful Addressing

- ❑ Class A addresses were designed for large organizations
 - ◆ The most of the addresses were wasted and were not used.
- ❑ Class B addresses were designed for midsize organizations
 - ◆ Class B is also too large for many organizations.
- ❑ Class C addresses were designed for small organizations
 - ◆ Class C is too small for many organizations.
- ❑ Class D addresses were designed for multicasting
 - ◆ Each addresses in this class is used to define one group of hosts on the Internet.
- ❑ Class E addresses were reserved for future use.
 - ◆ Only a few used, resulting in another waste of addresses.

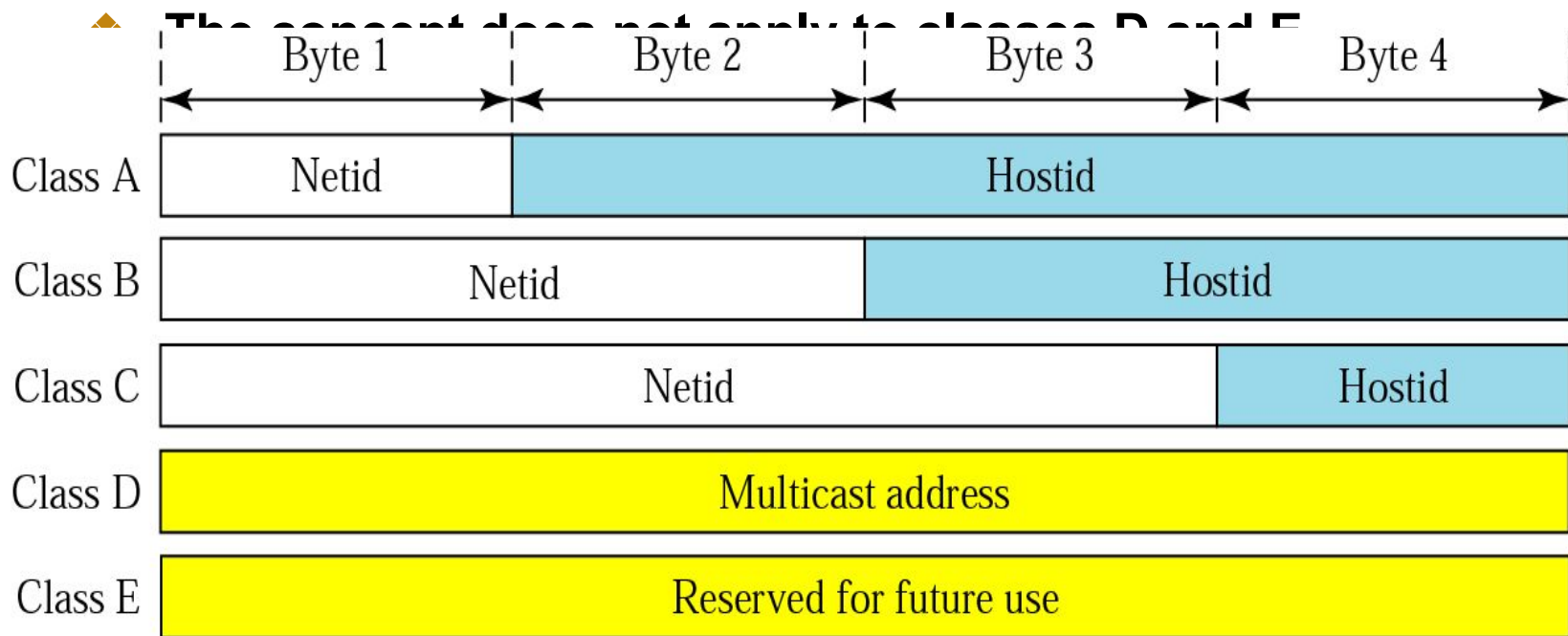
In classful addressing, a large part of the available addresses were wasted.



Classful Addressing

□ Netid and Hostid

- ◆ IP address in class A,B, or C is divided into netID and hostID.
 - netID defines a Network, and hostID defines a host in the networks.



MASK

□ Mask

- ◆ When a router receives a packet with a destination address, it needs to route the packet.
- ◆ The routing is based on the network address and subnetwork address.
 - The router outside the organization has a **routing table with one column based on the network addresses**;
 - The router inside the organization has a **routing table based on the subnetwork addresses**.
- ◆ The mask is a 32-bit binary number, and the mask can help to find the network and subnetwork address.
 - The routers outside the organization use a **Default Mask** to find the network address and,
 - The routers inside the organization use a **Subnet Mask** to find the subnetwork address..



Default Mask

□ Default Mask

- ◆ A default mask is a 32-bit binary number, and the default mask for each class are as follows; 255.0.0.0, 255.255.0.0, 255.255.255.0.
- ◆ Default mask gives the network address when ANDed with an address in the block.
 - If the bit in the mask is 1, the corresponding bit in the address is retained in the output (no change)
 - If the mask is 0, a 0 bit in the output is the result.

Table 19.1 Default masks

<i>Class</i>	<i>Binary</i>	<i>Dotted-Decimal</i>	<i>CIDR</i>
A	11111111 00000000 00000000 00000000	255.0.0.0	/8
B	11111111 11111111 00000000 00000000	255.255.0.0	/16
C	11111111 11111111 11111111 00000000	255.255.255.0	/24



Subnet Mask

□ Subnetting

- ◆ A network is divided into several smaller networks with each subnetwork (or subnet) having its subnetwork address

□ Subnet Mask :

- ◆ We change some of the leftmost 0s in the default mask to make a subnet mask.
- ◆ The number of subnets is **determined by the number of extra 1s.**
 - If the number of extra 1s in n, the number of subnets is 2^n .
 - If the number of subnets is N, the number of extra 1s is $\log_2 N$.

	255.255.0.0			
Default Mask	11111111	11111111	00000000	00000000
			16	
	255.255.224.0			
Subnet Mask	11111111	11111111	111	00000000
			3	13

$$2^n = 2^3 = 8 \text{ subnets}$$



Supernetting and Address depletion

❑ Supernetting

- ◆ Combining several class C addresses to create a larger range of addresses

❑ Address Depletion

- ◆ The fast growth of the Internet led to the near depletion of the available addresses.
- ◆ Classful addressing, which is almost obsolete, is replaced with classless addressing.



Network Address

❏ Network Addresses

- ❖ The first address in a block is normally not assigned to any device;
- ❖ It is used as the network address that represents the organization to the rest of the world.

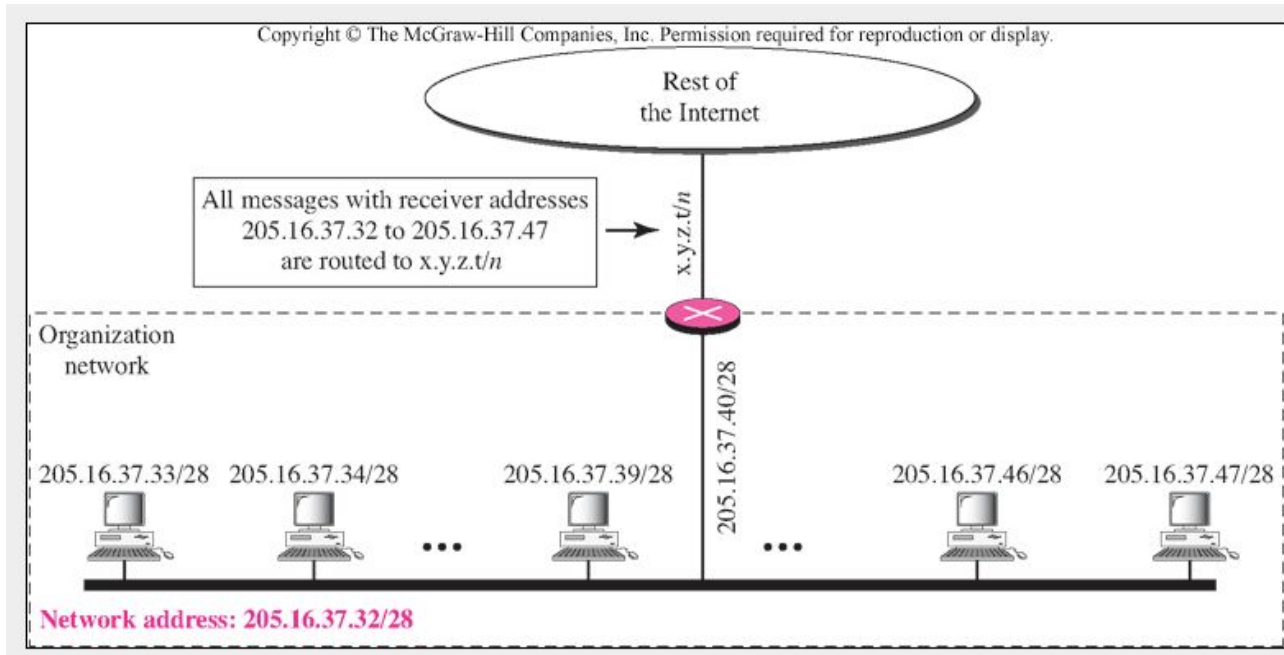


Figure 19.4 *A network configuration for the block 205.16.37.32/28*

Network Address

□ Hierarchy

- ◆ IP addresses have levels of hierarchy.
- ◆ For example, a telephone network has three levels of hierarchy.
 - The leftmost 3 digits define the area code, the next 3 digits define the exchange, the last 4 digits define the connection of the local loop to the central office.

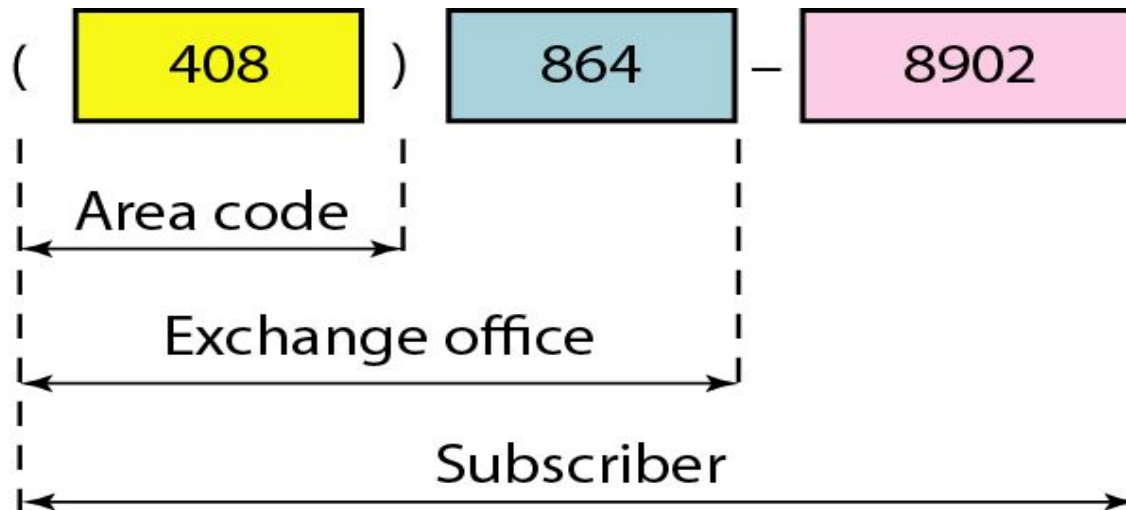


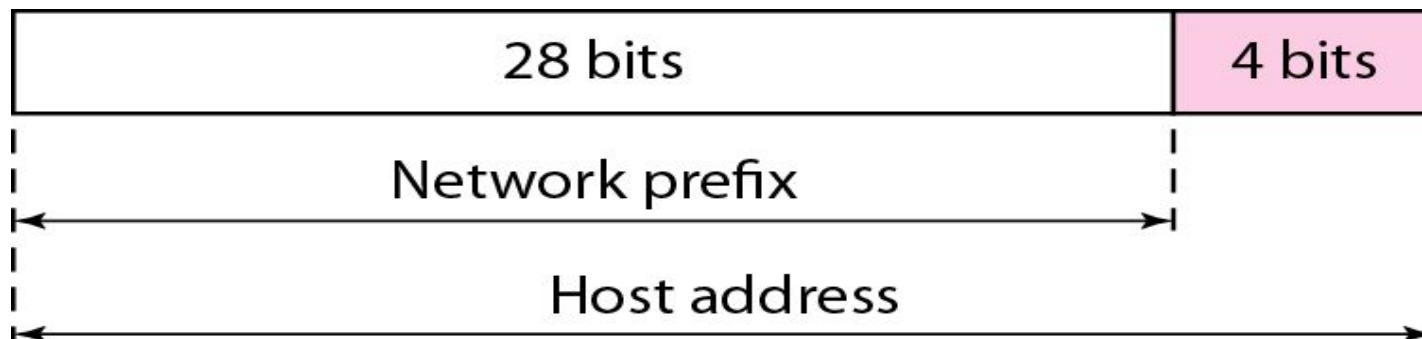
Figure 19.5 *Two levels of hierarchy in an IPv4 address*

Subnetting and Supernetting

❑ Two-level Hierarchy : No Subnetting

- ◆ Each IP address in the block can define only two-level of hierarchy when not subnetted.
 - the leftmost n bits (prefix) define the network;
 - the rightmost $32 - n$ bits define the host.
- ◆ The part of the address that defines the network is called the **Prefix**;
- ◆ The part that defines the host is called the Suffix.
- ◆ The prefix is **common to all addresses in the network**; the suffix changes from one device to another.

Figure 19.6 *A frame in a character-oriented protocol*



Subnetting and Supernetting

Three-Levels of Hierarchy : Subnetting

- Creating clusters of networks (called subnets)

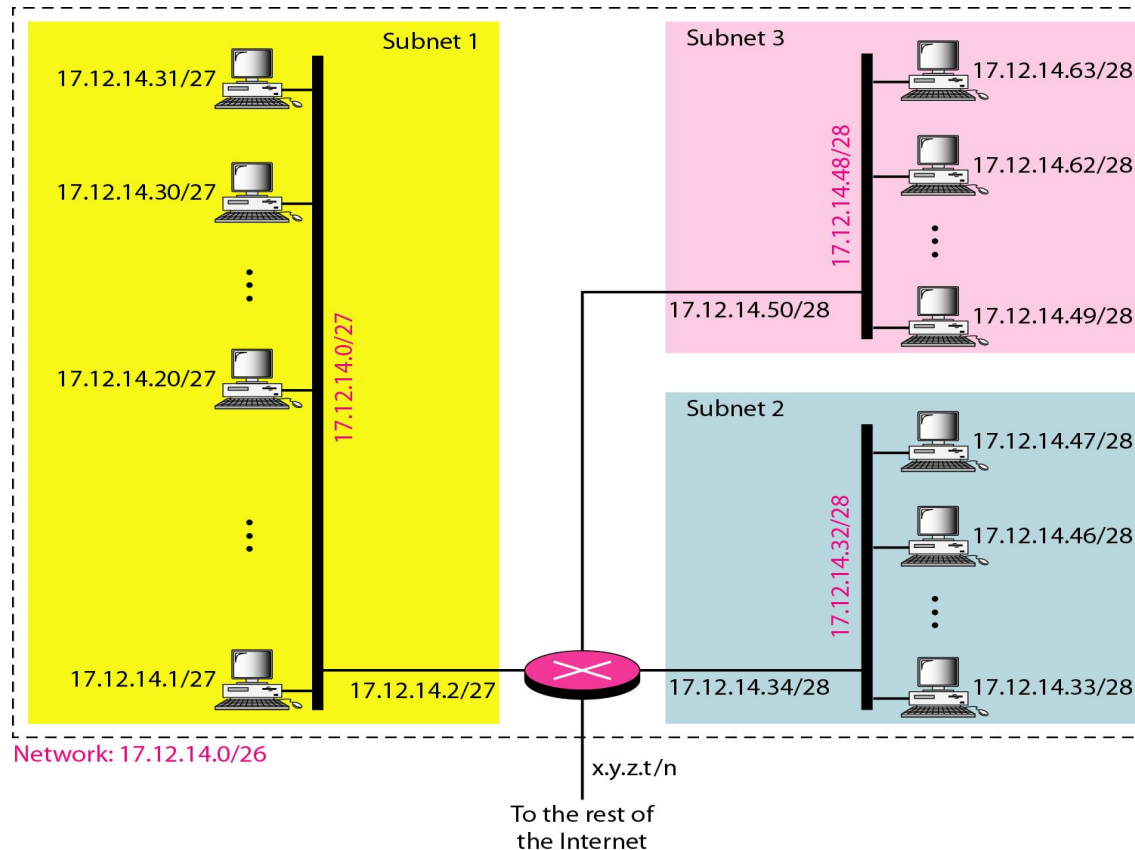


Figure 19.7 Configuration and addresses in a subnetted network

Subnetting and Supernetting

- ❑ We have three levels of hierarchy through subnetting.
 - ❖ The subnet prefix length can differ for the subnets.

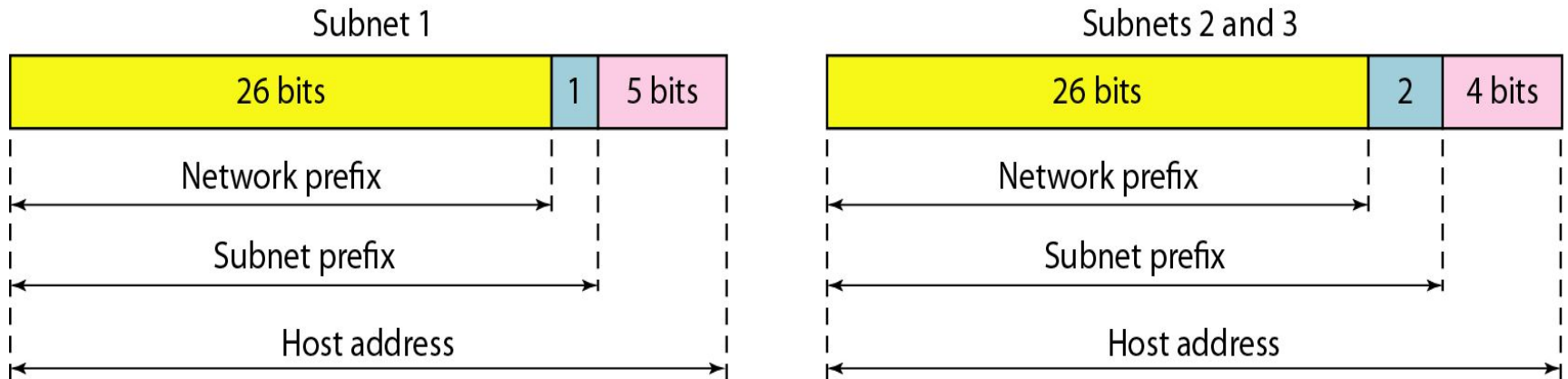
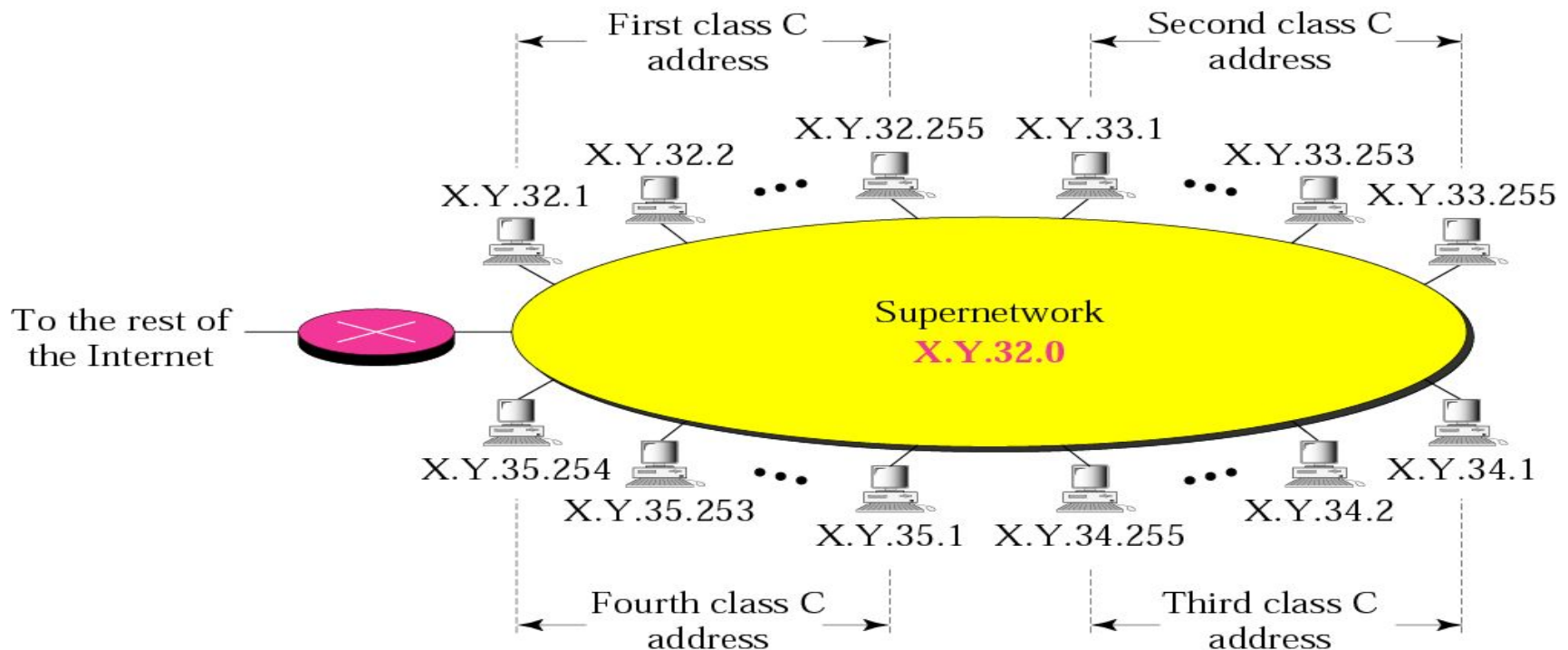


Figure 19.8 *Three-level hierarchy in an IPv4 address*

Supernetting

❏ Supernetting

- ◆ A maximum number of Class C is 256 addresses,
- ◆ If organization needed more addresses, The Supernetting can combine several class C blocks to create a larger range of addresses.
(The mask changes from /24 to /22)



Network Address Translation (NAT)

❑ Network Address Translation (NAT)

- ❖ NAT enables a user to have a large set of address internally and one address, or a small set of addresses, externally.
- ❖ The Internet authorities have reserved 3 sets of addresses as private addresses.
 - Any organization can use an address out of this set without permission from the Internet authorities.
 - They are unique inside the organization, but they are not unique globally.
 - No router will forward this packet as the destination address.

Table 19.3 *Addresses for private networks*

Range			Total
10.0.0.0	to	10.255.255.255	2^{24}
172.16.0.0	to	172.31.255.255	2^{20}
192.168.0.0	to	192.168.255.255	2^{16}



Network Address Translation (NAT)

❑ NAT Implementation

- ❖ The router that connects the network to the global address uses **one private address and one global address.**
- ❖ The private network is transparent to the rest of the Internet; the rest of the Internet **sees only the NAT router with the address 200.24.5.8.**

Site using private addresses

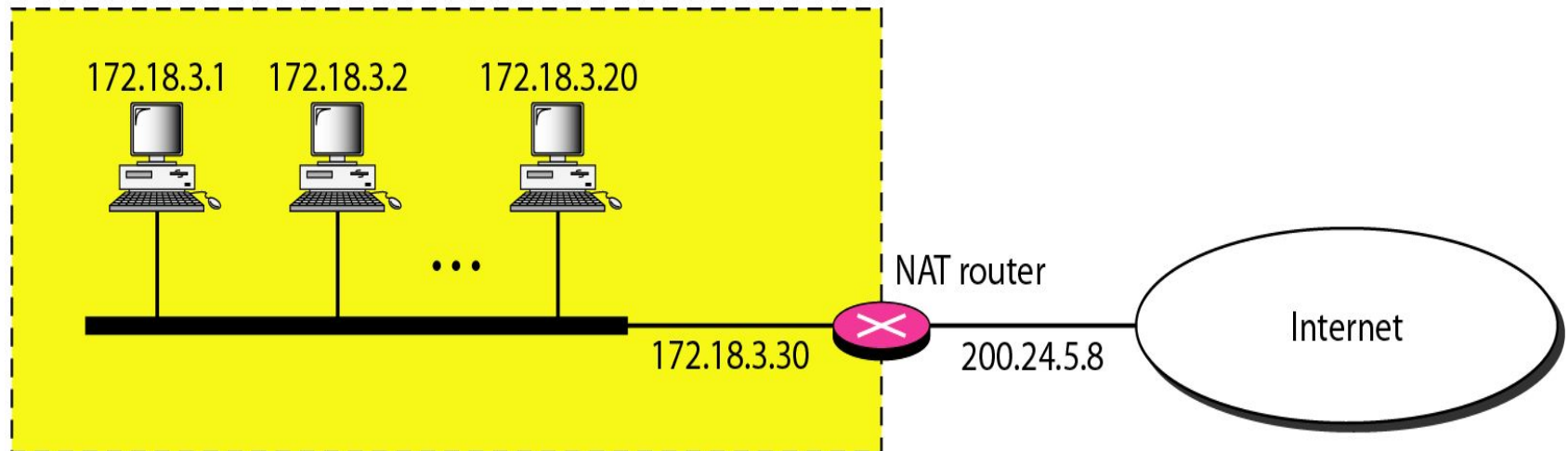


Figure 19.10 *A NAT implementation*

Network Address Translation (NAT)

□ Address translation

- ❖ All the outgoing packets go through the NAT router, which replaces the source address in the packet with the global NAT address.
- ❖ All incoming packets also pass through the NAT router, which replaces the destination addresses in the packet with the appropriate private address.

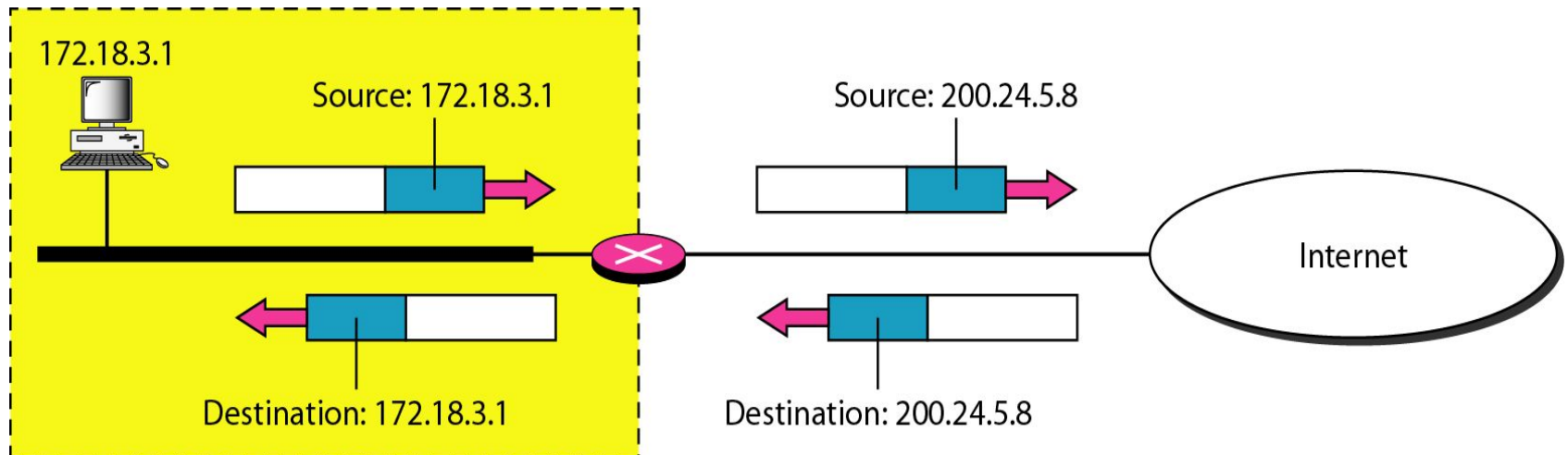


Figure 19.11 *Addresses in a NAT*

Network Address Translation (NAT)

Translation Table

- ◆ When the router translates the source address of the outgoing packet, it also makes note of the destination address – where the packet is going.
- ◆ When the response comes back from the destination, the router uses the source address of the packet to find the private address of the packet.

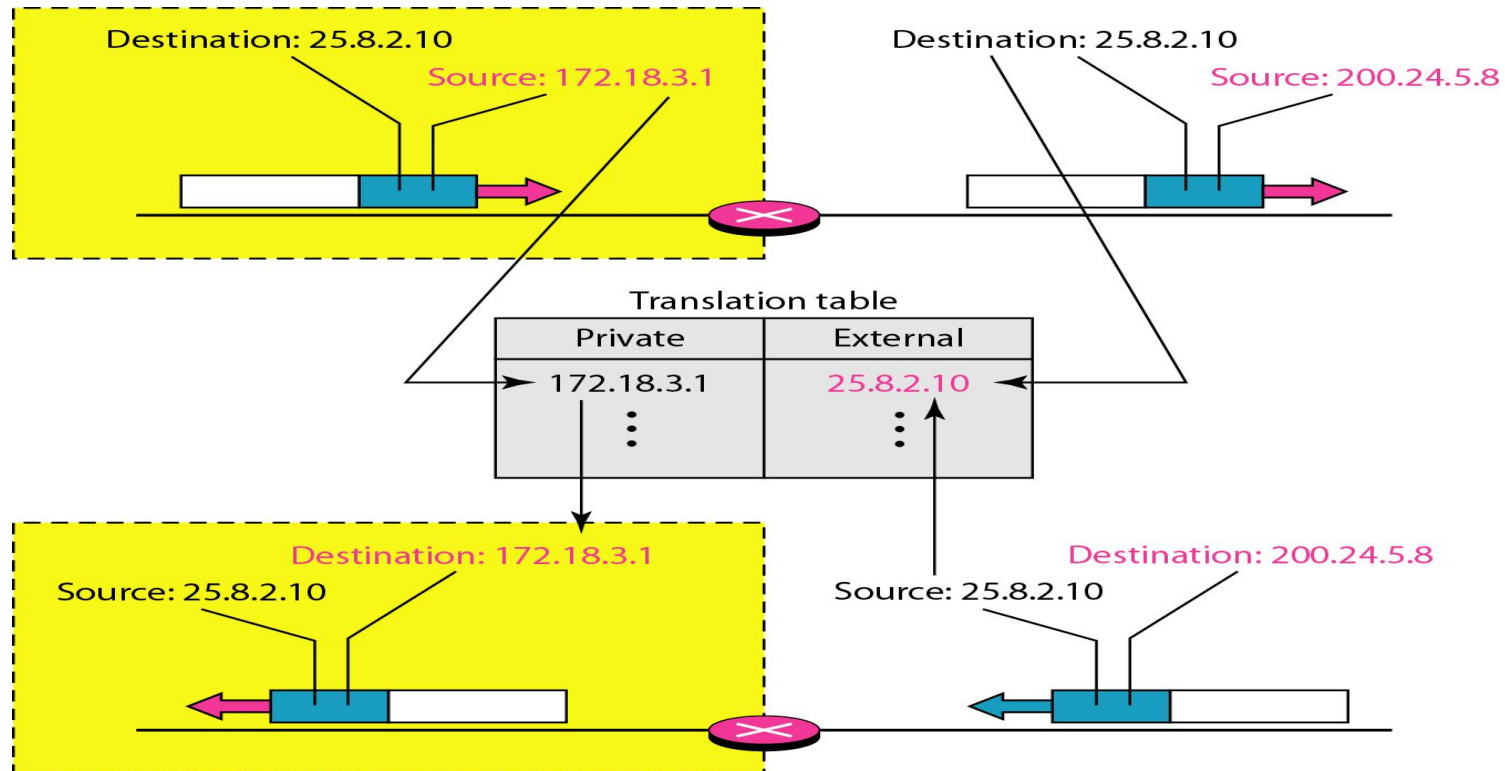


Figure 19.12 NAT address translation

Network Address Translation (NAT)

□ Using both IP addresses and port numbers

Table 19.4 *Five-column translation table*

<i>Private Address</i>	<i>Private Port</i>	<i>External Address</i>	<i>External Port</i>	<i>Transport Protocol</i>
172.18.3.1	1400	25.8.3.2	80	TCP
172.18.3.2	1401	25.8.3.2	80	TCP
...

19-2 IPv6 ADDRESSES

Despite all short-term solutions, address depletion is still a long-term problem for the Internet. This and other problems in the IP protocol itself have been the motivation for IPv6.

Topics discussed in this section:

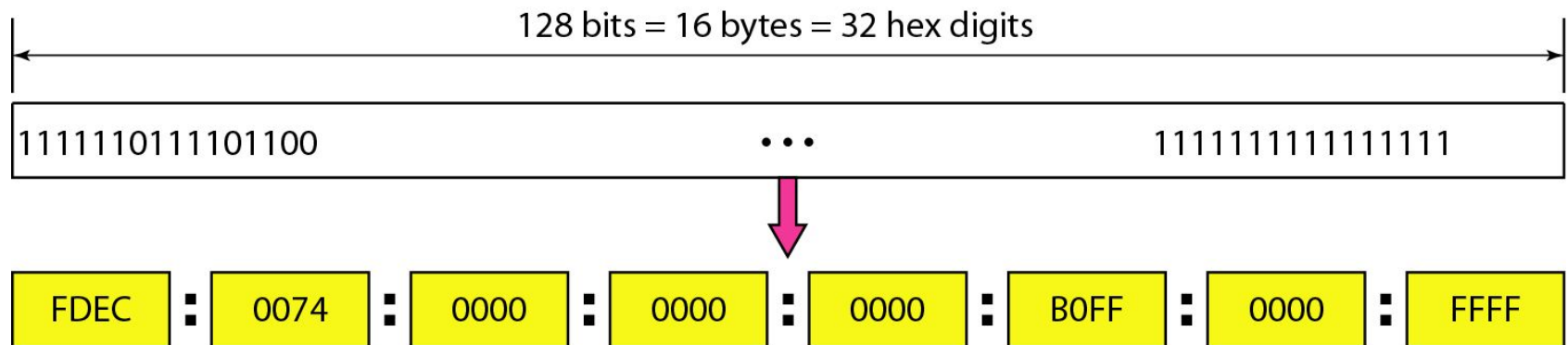
Structure

Address Space



Structure - IPv6

- ❑ An IPv6 address consists of 16 bytes (Octets); it is 128 bits long.
- ❑ Hexadeximal Colon Notation
 - ❖ In this notation, 128 bits is divided into eight sections, each 2 bytes in length.
 - ❖ Therefore, the address consists of 32 hexadecimal digits, with every four digits separated by a colon.



Abbreviation

- ❑ Although the IP address, even in hexadecimal format, is very long, many of the digits are zeros.
- ❑ The leading zeros of a section (four digits between two colons) can be omitted.
 - ◆ Only the leading zeros can be dropped, not the trailing zeros.

Original

FDEC : 0074 : 0000 : 0000 : 0000 : B0FF : 0000 : FFF0

Abbreviated

FDEC : 74 : 0 : 0 : 0 : B0FF : 0 : FFF0

More abbreviated

FDEC : 74 : : B0FF : 0 : FFF0

Gap



Address Space

- ❑ IPv6 has a much larger address space; 2^{128} addresses are available.



Summary (1)

- ❑ At the Network layer, a global identification system that uniquely identifies every host and router is necessary for delivery of packet from host to host.
- ❑ An IPv4 address is 32 bits long and uniquely and universally defines a host or router on the Internet.
- ❑ In classful addressing, the portion of the IP address that identifies the network is called the netid.
- ❑ In classful addressing, the portion of the IP address that identifies the host or router on the network is called the hosted.
- ❑ An IP address defines a device's connection to a network.
- ❑ There are five classes in IPv4 addresses. Classes A, B, and C differ in the number of hosts allowed per network. Class D is for multicasting and Class E is reserved.



Summary(2)

- ❑ The class of an address is easily determined by examination of the first byte.
- ❑ Addresses in classes A, B, or C are mostly used for unicast communication.
- ❑ Address in class D are used for multicast communication.
- ❑ Subnetting divides on large network into several smaller ones, adding an intermediate level of hierarchy in IP addressing.
- ❑ Supernetting combines several networks into one large one.
- ❑ In classless addressing, we can divided the address space into variable-length blocks.