



Data Communications and Networking

Fourth Edition

Chapter 19

Network Layer: Logical Addressing

19-1 IPv4 ADDRESSES

An **IPv4 address** is a **32-bit** address that uniquely and universally defines the connection of a device (for example, a computer or a router) to the Internet.

Topics discussed in this section:

Address Space

Notations

Classful Addressing

Classless Addressing

Network Address Translation (NAT)

IPv4 Address

- An IPv4 address is 32 bits long.
- The IPv4 addresses are unique and universal.
- The address space of IPv4 is 2^{32} or 4,294,967,296

IPV4 Address

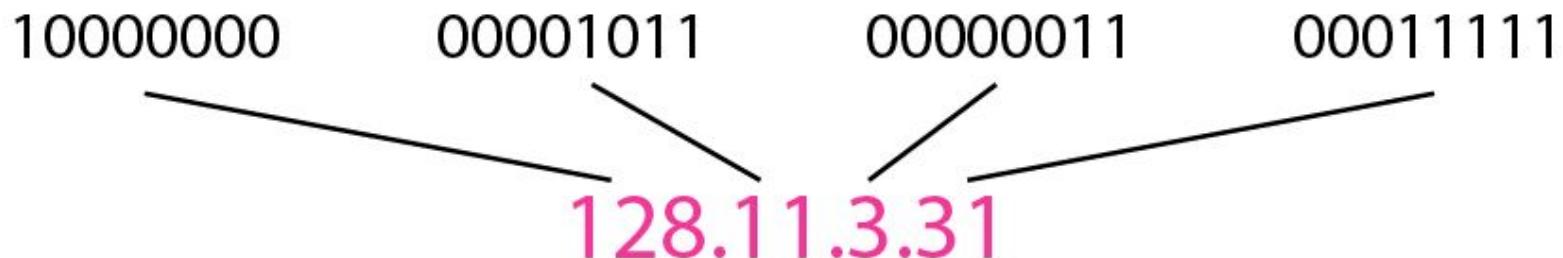
□ Notation

◆ Binary Notation :

- The IPV4 address is displayed as 32 bits.
- Each octet is often referred to as a byte.

◆ Dotted-Decimal Notation

- To make the IPV4 address more compact and easier to read, Internet addresses are written in decimal form with a decimal point (dot) separating the byte.



Notation (cont'd)

□ Hexadecimal Notation

0111 0101 1001 0101 0001 1101 1110 1010

75

95

1D

EA

0x75951DEA

- 8 *hexadecimal digits*
- *Used in network programming*

Classful Addressing

- In classful addressing, the address space is divided into five classes: A, B, C, D, and E.
 - ◆ If the address is given in binary notation, the first few bits can tell us the class of the address.
 - ◆ If the address is given in decimal-dotted notation, the first byte defines the class.

	First byte	Second byte	Third byte	Fourth byte
Class A	0			
Class B	10			
Class C	110			
Class D	1110			
Class E	1111			

a. Binary notation

	First byte	Second byte	Third byte	Fourth byte
Class A	0–127			
Class B	128–191			
Class C	192–223			
Class D	224–239			
Class E	240–255			

b. Dotted-decimal notation

Classful Addressing

□ Classful Addresses

- ◆ **Unicast Communication - A, B, C Class**
(~must be delivered to specific computer)
- ◆ **Multicast Communication – D Class**
(~must be delivered to each member of the group)
- ◆ **For reserve – E Class**

Classful Addressing

□ Classes and blocks

- ◆ One problem with classful addressing is that each class is divided into a fixed number of blocks with each block having a fixed size.

Table 19.1 *Number of blocks and block size in classful IPv4 addressing*

Class	Number of Blocks	Block Size	Application
A	128	16,777,216	Unicast
B	16,384	65,536	Unicast
C	2,097,152	256	Unicast
D	1	268,435,456	Multicast
E	1	268,435,456	Reserved

Classful Addressing

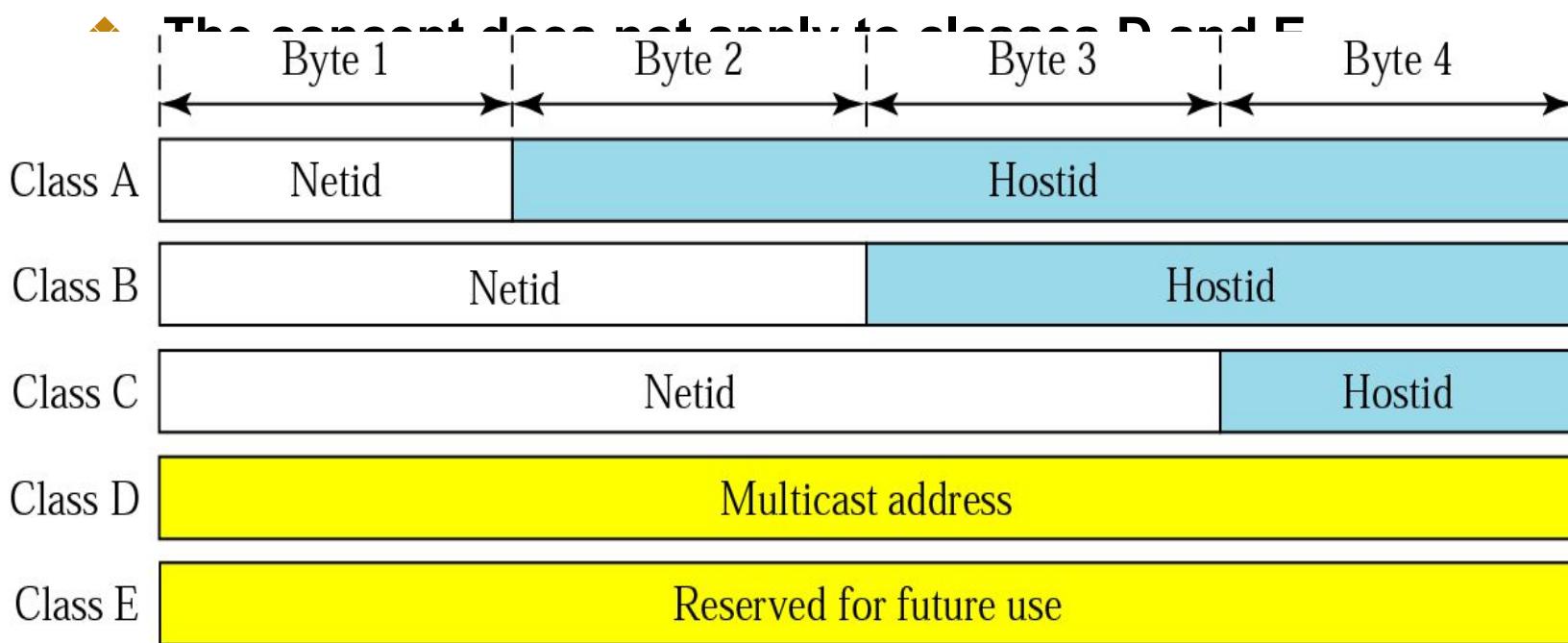
- ❑ Class A addresses were designed for large organizations
 - ◆ The most of the addresses were wasted and were not used.
- ❑ Class B addresses were designed for midsize organizations
 - ◆ Class B is also too large for many organizations.
- ❑ Class C addresses were designed for small organizations
 - ◆ Class C is too small for many organizations.
- ❑ Class D addresses were designed for multicasting
 - ◆ Each addresses in this class is used to define one group of hosts on the Internet.
- ❑ Class E addresses were reserved for future use.
 - ◆ Only a few used, resulting in another waste of addresses.

**In classful addressing, a large part of
the available addresses were wasted.**

Classful Addressing

Netid and Hostid

- IP address in class A,B, or C is divided into netID and hostID.
 - netID defines a Network, and hostID defines a host in the networks.



MASK

□ Mask

- ◆ When a router receives a packet with a destination address, it needs to route the packet.
- ◆ The routing is based on the network address and subnetwork address.
 - The router outside the organization has a routing table with one column based on the network addresses;
 - The router inside the organization has a routing table based on the subnetwork addresses.
- ◆ The mask is a 32-bit binary number, and the mask can help to find the network and subnetwork address.
 - The routers outside the organization use a Default Mask to find the network address and,
 - The routers inside the organization use a Subnet Mask to find the subnetwork address..

Default Mask

□ Default Mask

- ◆ A default mask is a 32-bit binary number, and the default mask for each class are as follows; 255.0.0.0, 255.255.0.0, 255.255.255.0.
- ◆ Default mask gives the network address when ANDed with an address in the block.
 - If the bit in the mask is 1, the corresponding bit in the address is retained in the output (no change)
 - If the mask is 0, a 0 bit in the output is the result.

Table 19.1 Default masks

Class	Binary	Dotted-Decimal	CIDR
A	11111111 00000000 00000000 00000000	255.0.0.0	/8
B	11111111 11111111 00000000 00000000	255.255.0.0	/16
C	11111111 11111111 11111111 00000000	255.255.255.0	/24

Subnet Mask

❑ Subnetting

- ❖ A network is divided into several smaller networks with each subnetwork (or subnet) having its subnetwork address

❑ Subnet Mask :

- ❖ We change some of the leftmost 0s in the default mask to make a subnet mask.
- ❖ The number of subnets is determined by the number of extra 1s.
 - If the number of extra 1s in n, the number of subnets is 2^n .
 - If the number of subnets is N, the number of extra 1s is $\log_2 N$.

255.255.0.0				
Default Mask	11111111	11111111	00000000	00000000
255.255.224.0	11111111	11111111	111	000000000
Subnet Mask	11111111	11111111	111	000000000
	3	13		

$$2^n = 2^3 = 8 \text{ subnets}$$

Supernetting and Address depletion

□ Supernetting

- ◆ Combining several class C addresses to create a larger range of addresses

□ Address Depletion

- ◆ The fast growth of the Internet led to the near depletion of the available addresses.
- ◆ Classful addressing, which is almost obsolete, is replaced with classless addressing.

Network Address

□ Network Addresses

- ◆ The first address in a block is normally not assigned to any device;
- ◆ It is used as the network address that represents the organization to the rest of the world.

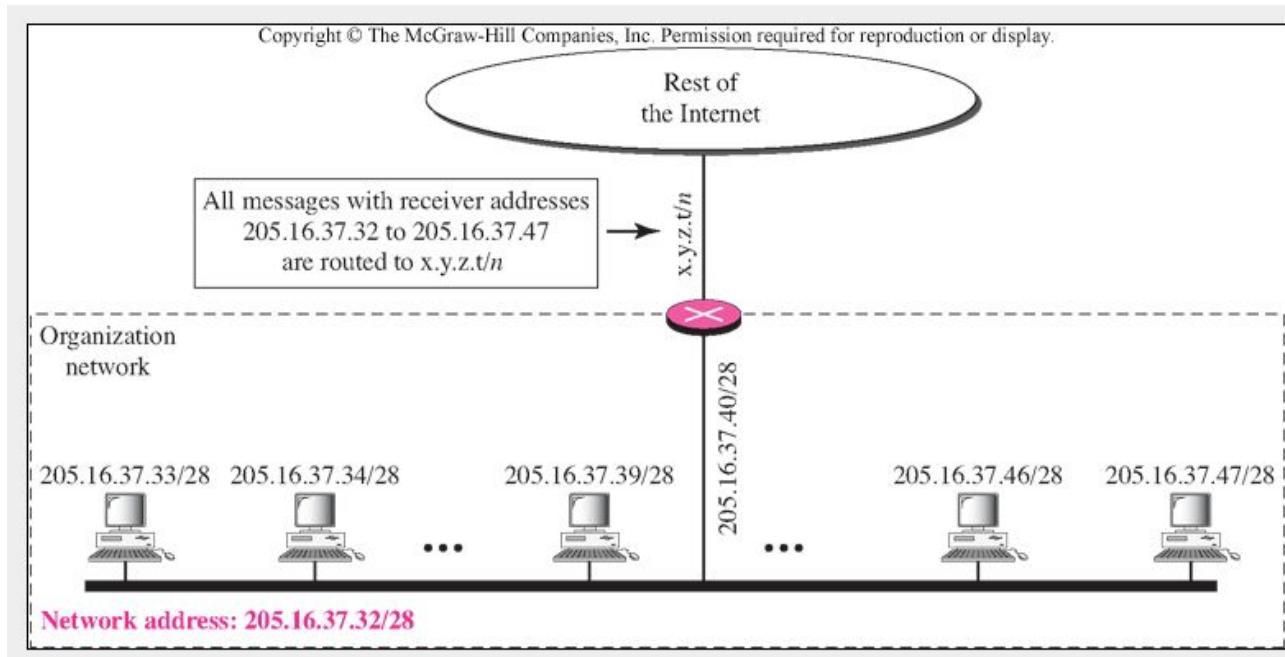


Figure 19.4 A network configuration for the block 205.16.37.32/28

Network Address

□ Hierarchy

- ◆ IP addresses have levels of hierarchy.
- ◆ For example, a telephone network has three levels of hierarchy.
 - The leftmost 3 digits define the area code, the next 3 digits define the exchange, the last 4 digits define the connection of the local loop to the central office.

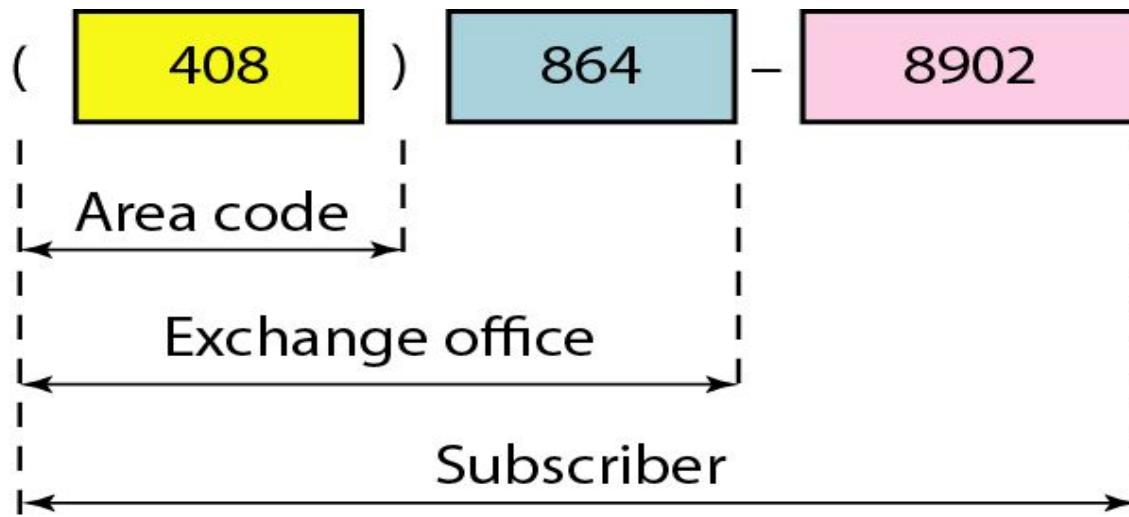


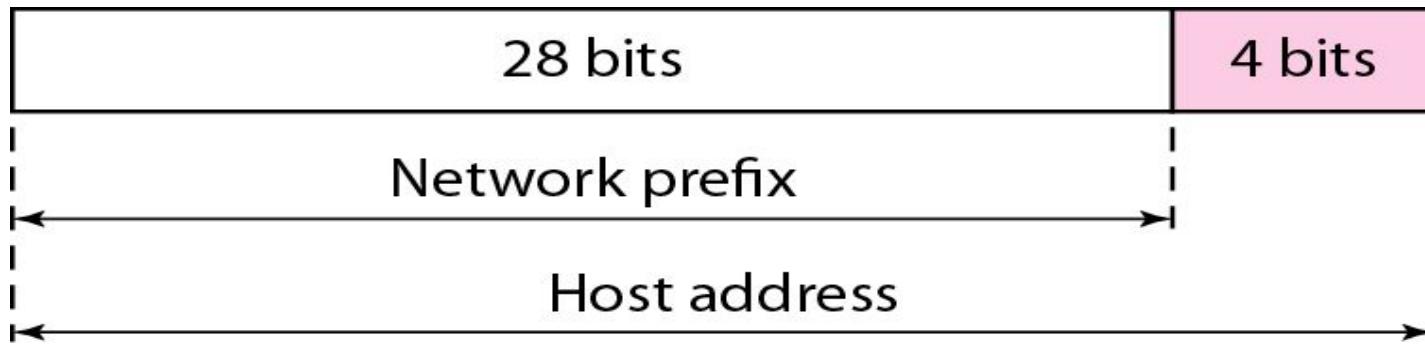
Figure 19.5 *Two levels of hierarchy in an IPv4 address*

Subnetting and Supernetting

□ Two-level Hierarchy : No Subnetting

- ◆ Each IP address in the block can define only two-level of hierarchy when not subnetted.
 - the leftmost n bits (prefix) define the network;
 - the rightmost $32 - n$ bits define the host.
- ◆ The part of the address that defines the network is called the **Prefix**;
- ◆ The part that defines the host is called the **Suffix**.
- ◆ The prefix is **common to all addresses in the network**; the suffix changes from one device to another.

Figure 19.6 A frame in a character-oriented protocol



Subnetting and Supernetting

□ Three-Levels of Hierarchy : Subnetting

- ◆ Creating clusters of networks (called subnets)

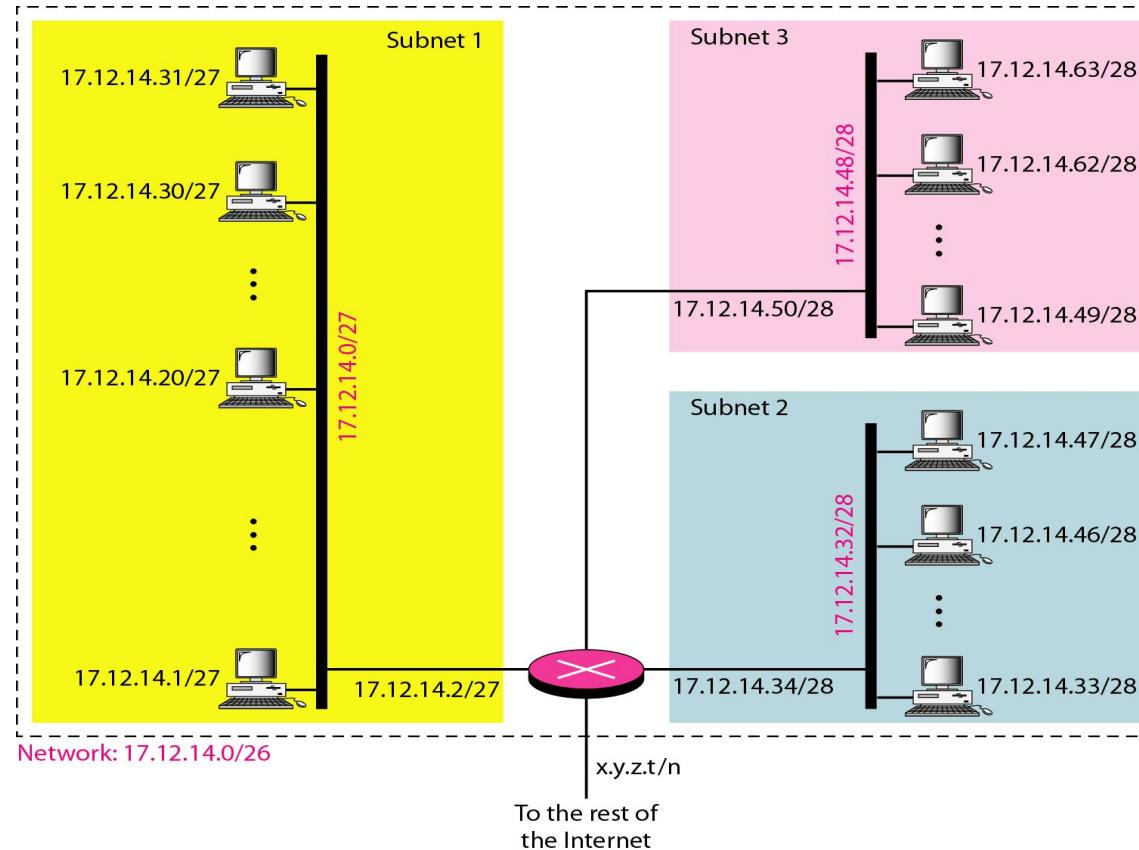


Figure 19.7 Configuration and addresses in a subnetted network

Subnetting and Supernetting

- We have three levels of hierarchy through subnetting.
 - ◆ The subnet prefix length can differ for the subnets.

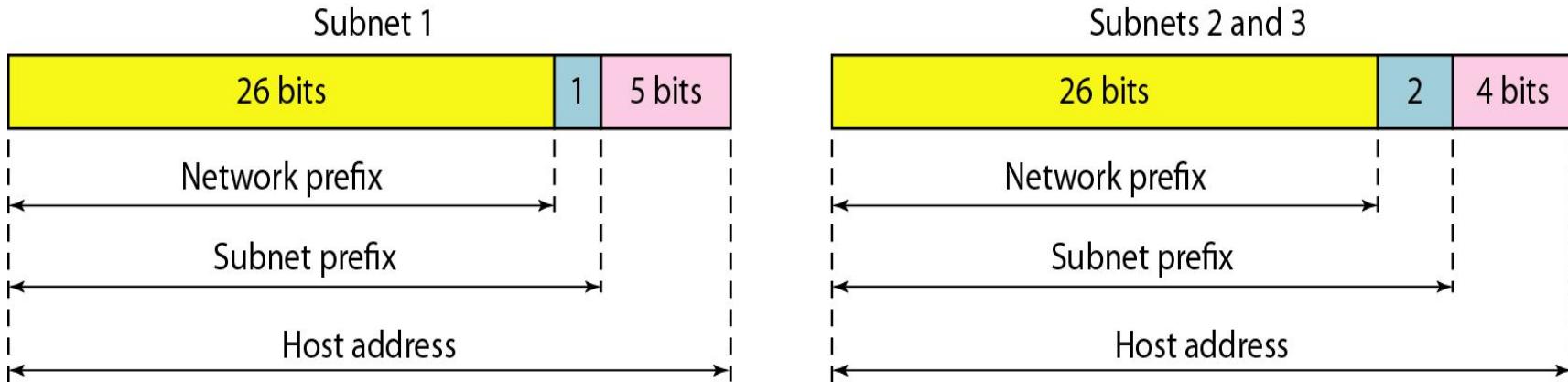
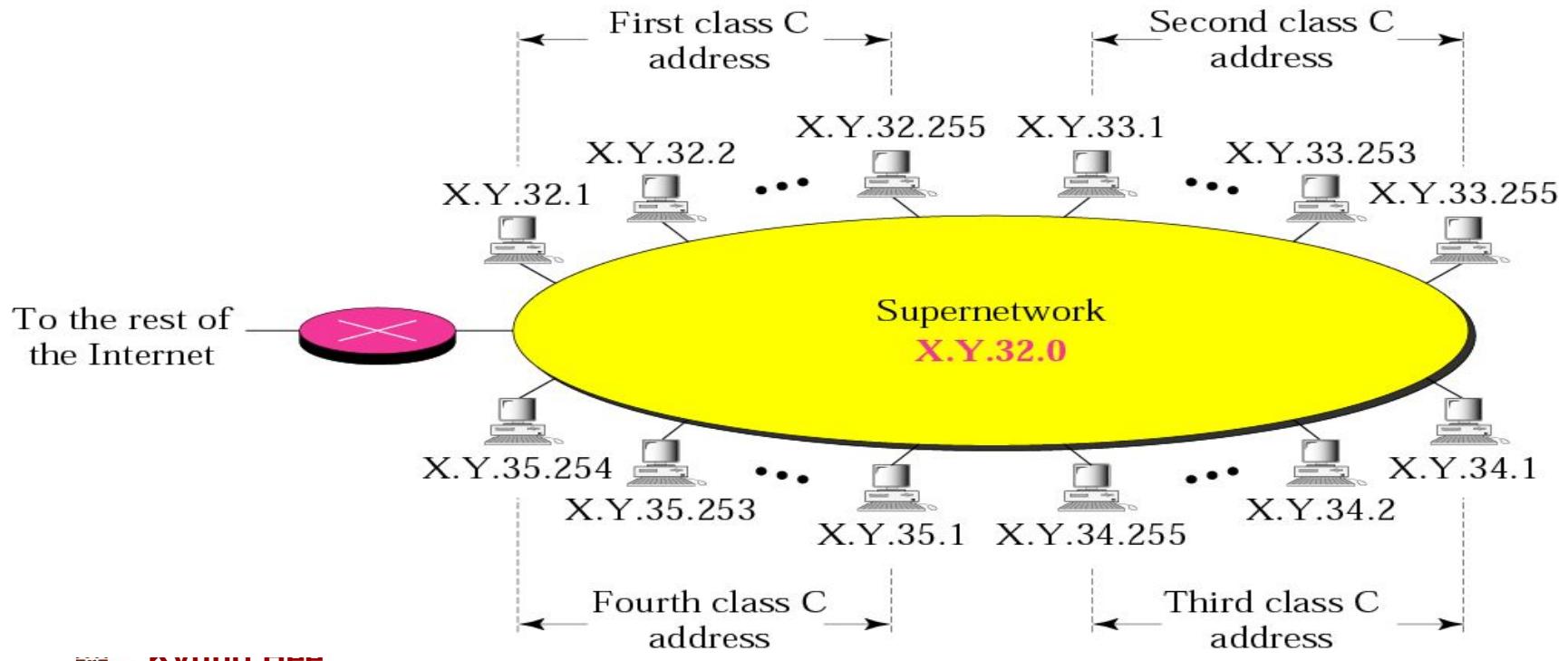


Figure 19.8 *Three-level hierarchy in an IPv4 address*

Supernetting

□ Supernetting

- ◆ A maximum number of Class C is 256 addresses,
- ◆ If organization needed more addresses, The Supernetting can combine several class C blocks to create a larger range of addresses.
(The mask changes from /24 to /22)



Network Address Translation (NAT)

□ Network Address Translation (NAT)

- ◆ NAT enables a user to have a large set of address internally and one address, or a small set of addresses, externally.
- ◆ The Internet authorities have reserved 3 sets of addresses as private addresses.
 - Any organization can use an address out of this set without permission from the Internet authorities.
 - They are unique inside the organization, but they are not unique globally.
 - No router will forward this packet as the destination address.

Table 19.3 *Addresses for private networks*

Range		Total
10.0.0.0	to	2^{24}
172.16.0.0	to	2^{20}
192.168.0.0	to	2^{16}

Network Address Translation (NAT)

❑ NAT Implementation

- ❖ The router that connects the network to the global address uses one private address and one global address.
- ❖ The private network is transparent to the rest of the Internet; the rest of the Internet sees only the NAT router with the address 200.24.5.8.

Site using private addresses

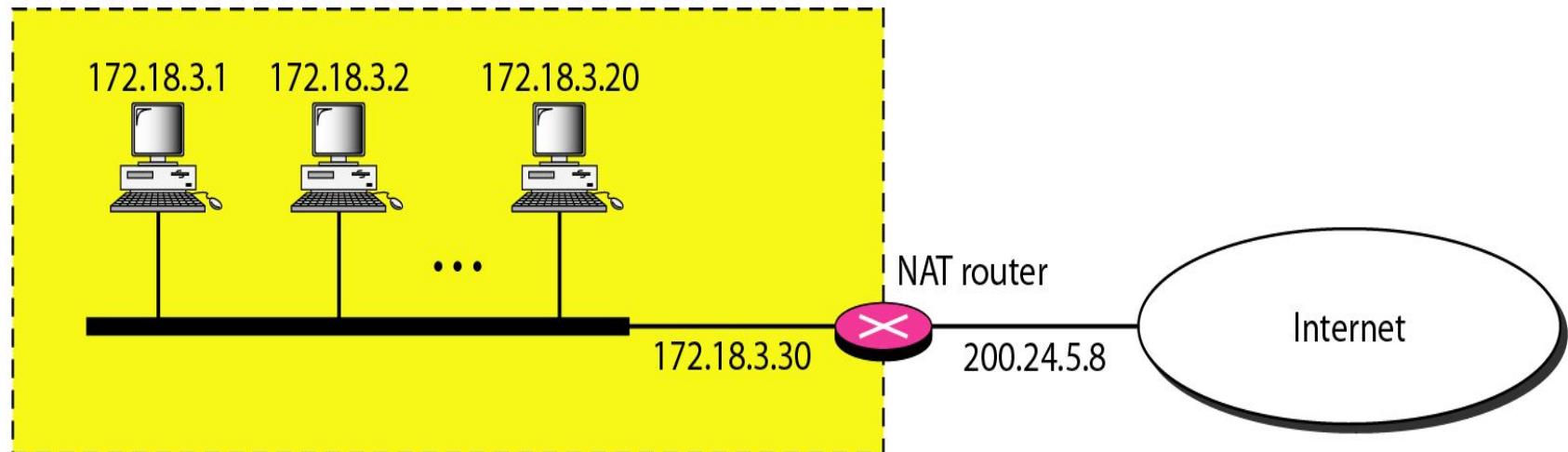


Figure 19.10 *A NAT implementation*

Network Address Translation (NAT)

Address translation

- ◆ All the outgoing packets go through the NAT router, which replaces the source address in the packet with the global NAT address.
- ◆ All incoming packets also pass through the NAT router, which replaces the destination addresses in the packet with the appropriate private address.

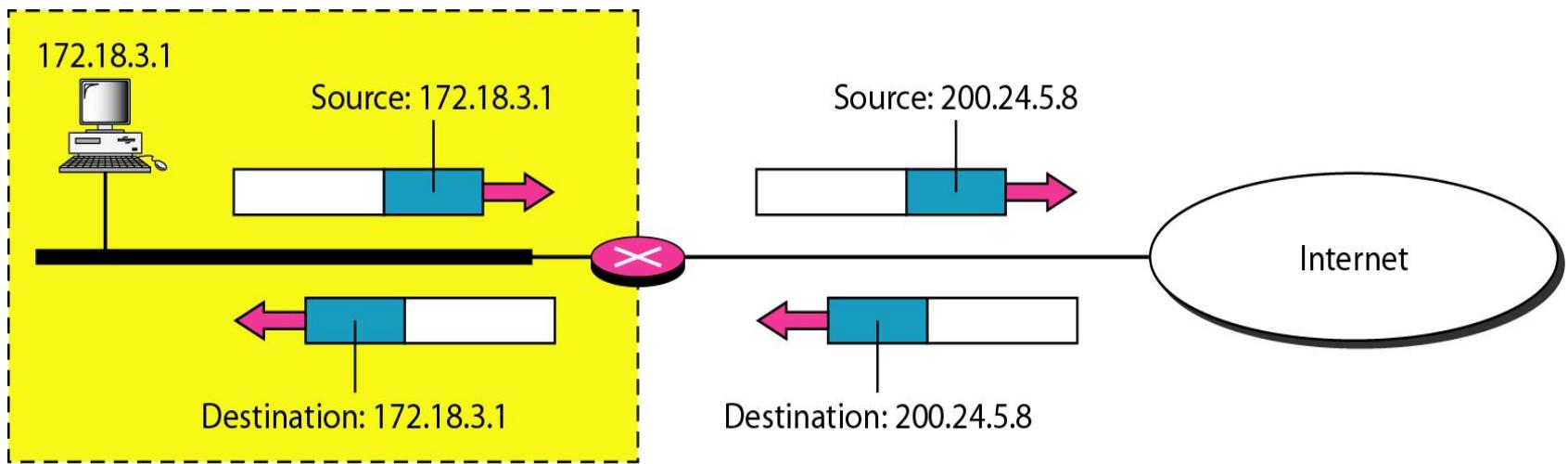


Figure 19.11 *Addresses in a NAT*

Network Address Translation (NAT)

Translation Table

- When the router translates the source address of the outgoing packet, it also makes note of the destination address – where the packet is going.
- When the response comes back from the destination, the router uses the source address of the packet to find the private address of the packet.

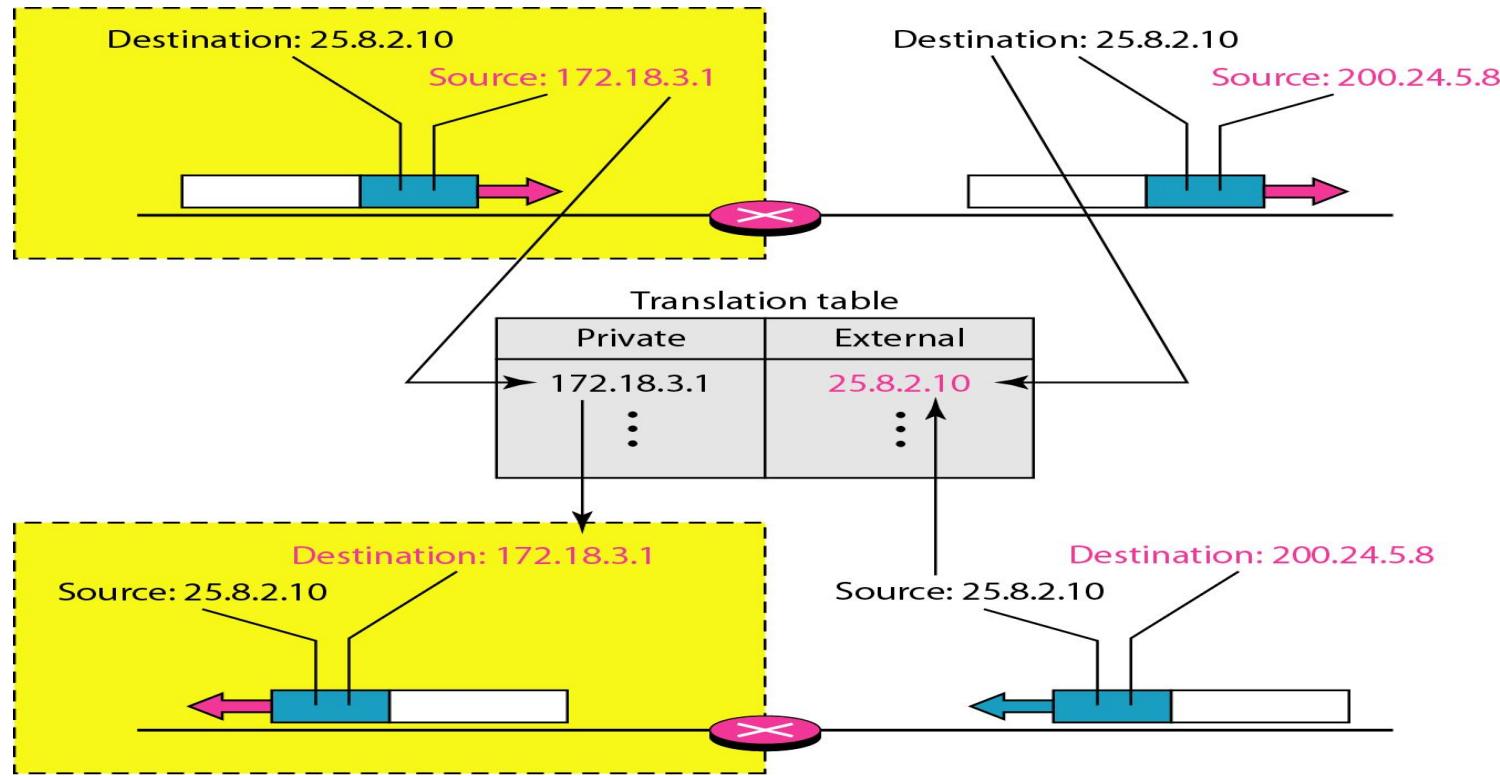


Figure 19.12 *NAT address translation*

Network Address Translation (NAT)

□ Using both IP addresses and port numbers

Table 19.4 *Five-column translation table*

<i>Private Address</i>	<i>Private Port</i>	<i>External Address</i>	<i>External Port</i>	<i>Transport Protocol</i>
172.18.3.1	1400	25.8.3.2	80	TCP
172.18.3.2	1401	25.8.3.2	80	TCP
...

19-2 IPv6 ADDRESSES

Despite all short-term solutions, address depletion is still a long-term problem for the Internet. This and other problems in the IP protocol itself have been the motivation for IPv6.

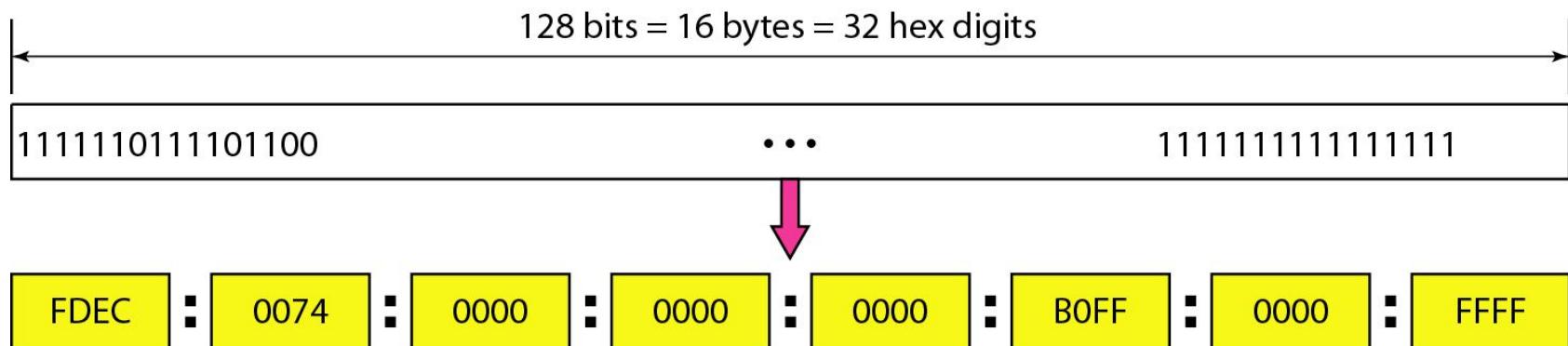
Topics discussed in this section:

Structure

Address Space

Structure - IPv6

- An IPv6 address consists of 16 bytes (Octets); it is 128 bits long.
- Hexadecimal Colon Notation
 - ◆ In this notation, 128 bits is divided into eight sections, each 2 bytes in length.
 - ◆ Therefore, the address consists of 32 hexadecimal digits, with every four digits separated by a colon.



Abbreviation

- ❑ Although the IP address, even in hexadecimal format, is very long, many of the digits are zeros.
- ❑ The leading zeros of a section (four digits between two colons) can be omitted.
 - ◆ Only the leading zeros can be dropped, not the trailing zeros.

Original

```
FDEC :: 0074 :: 0000 :: 0000 :: 0000 :: BOFF :: 0000 :: FFF0
```



Abbreviated

```
FDEC :: 74 :: 0 :: 0 :: 0 :: BOFF :: 0 :: FFF0
```



More abbreviated

```
FDEC :: 74 :: BOFF :: 0 :: FFF0
```



Address Space

- ❑ IPv6 has a much larger address space; 2^{128} addresses are available.

Summary (1)

- ❑ At the Network layer, a global identification system that uniquely identifies every host and router is necessary for delivery of packet from host to host.
- ❑ An IPv4 address is 32 bits long and uniquely and universally defines a host or router on the Internet.
- ❑ In classful addressing, the portion of the IP address that identifies the network is called the netid.
- ❑ In classful addressing, the portion of the IP address that identifies the host or router on the network is called the hosted.
- ❑ An IP address defines a device's connection to a network.
- ❑ There are five classes in IPv4 addresses. Classes A, B, and C differ in the number of hosts allowed per network. Class D is for multicasting and Class E is reserved.

Summary(2)

- The class of an address is easily determined by examination of the first byte.
- Addresses in classes A, B, or C are mostly used for unicast communication.
- Address in class D are used for multicast communication.
- Subnetting devides on large network into several smaller ones, adding an intermediate level of hierarchy in IP addressing.
- Supernetting combines several networks into one large one.
- In classless addressing, we can divided the address space into variable-length blocks.