

Kruskal's Algo

→ Sorting edges in order of increasing cost
(rename edges $1, 2, 3, \dots, n$ so that $c_1 < c_2 < \dots < c_n$)

Disjoint set

two operations

- Union → joining
- find

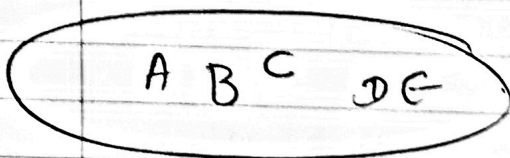
two set into one set

find(B) = S1

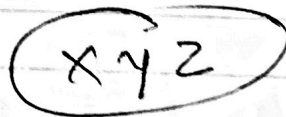
find(D) = S1

find(x) = S3

Exo



S1



S3

Cycle detection

Key Point :- each edge
make subset
using the ^{both} vertices of the edge
if (both vertices are in same subset)
{
 a cycle is found
}

Pseudo code:

for each unvisited edge (u, v) in set E

 if $[find(u) == find(v)]$

 // cycle is detected

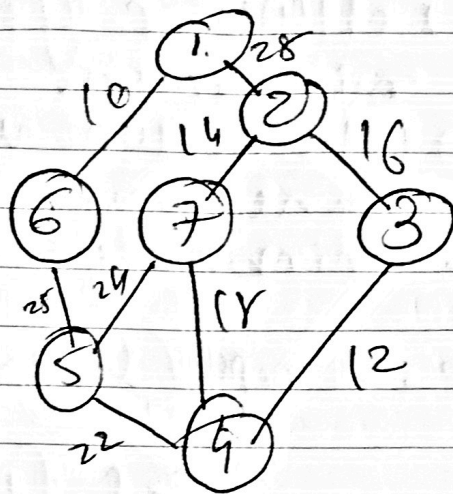
 else

 union (x, y) ;

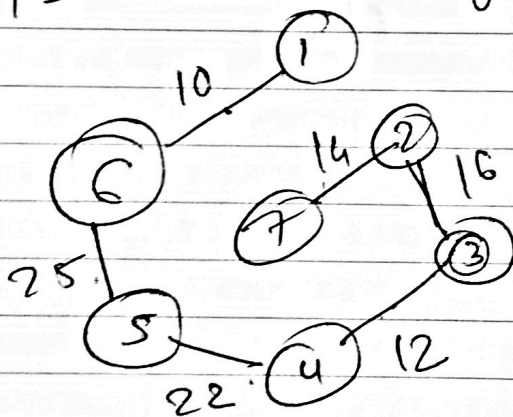
 }

in graph files \rightarrow is cycle function

Kruskal's



Kruskal's smallest edge



→ always choose the minimum weighted edge but if it will make a cycle don't choose that edge.

→ so (18) will make a cycle will not choose that
also don't include (24)

number of edges $|E|$

for spanning tree

it will select same prim's algo

~~it will~~ $|V| - 1$

$$O(|V||E|)$$

$$O(n \cdot E) = O(n^2)$$

min-heap \rightarrow always gives minimum node

\rightarrow when we use remove function

when we delete minnode from min-heap

time comp $O(\log n)$

cost of finding a minimum edge

so $E = \text{vertices} - 1$

so, $O(n \cdot e)$

$O(n \log n)$

$O(E \log E)$