

P1 Print All the Subsequence whos Sum is Equal to the K

→ Here we again used the concept of **Pick and Not Pick**

ans = [1, 2, 2], sum = 2
(Now only task is print all the subsequence and sum is equal to the given sum)

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Thought Process

```
myfunction(qindex, mylist, k) {
```

→ my
base
case

②

2

Revision Tree

$$m = [1, 2, 7]$$


Time Complexity = ?

 $O(2^n)$

→ Low ability to steer the element

Space Complexity = $O(N) + O(N) \approx O(N)$




```

// P3 Total Count the Subsequence whos Sum is Equal to the K
// Time Complexity: O(2^n)
// Space Complexity: O(2^n)

// Approach 1: Brute Force
// Generate all possible subsequences and check if their sum is equal to K.

// Function to generate all possible subsequences
void generateSubsequences(int arr[], int n, int index, vector<int> subseq) {
    if (index == n) {
        // Check if the sum of the subsequence is equal to K
        if (sum(subseq) == K) {
            // Count the subsequence
            count++;
        }
        return;
    }

    // Include the current element
    subseq.push_back(arr[index]);
    generateSubsequences(arr, n, index + 1, subseq);

    // Exclude the current element
    subseq.pop_back();
    generateSubsequences(arr, n, index + 1, subseq);
}

// Function to calculate the sum of a subsequence
int sum(vector<int> subseq) {
    int s = 0;
    for (int i = 0; i < subseq.size(); i++) {
        s += subseq[i];
    }
    return s;
}

// Main function
int main() {
    int arr[] = {1, 2, 1};
    int n = sizeof(arr) / sizeof(arr[0]);
    int K = 3;
    int count = 0;
    generateSubsequences(arr, n, 0, vector<int>());
    cout << count << endl;
    return 0;
}

```

```

// Approach 2: Dynamic Programming
// We will use a 2D array to store the count of subsequences with a given sum.

// Function to calculate the count of subsequences with a given sum
int countSubsequences(int arr[], int n, int sum) {
    // Create a 2D array to store the count of subsequences
    int dp[n][sum + 1];

    // Initialize the array
    for (int i = 0; i < n; i++) {
        for (int j = 0; j <= sum; j++) {
            dp[i][j] = 0;
        }
    }

    // Base case
    dp[0][0] = 1;

    // Fill the array
    for (int i = 1; i <= n; i++) {
        for (int j = 0; j <= sum; j++) {
            // If the current element is greater than the sum, then the count is 0
            if (arr[i - 1] > j) {
                dp[i][j] = dp[i - 1][j];
            } else {
                // If the current element is less than or equal to the sum, then the count is the sum of the count of subsequences with the current element and the count of subsequences without the current element
                dp[i][j] = dp[i - 1][j] + dp[i - 1][j - arr[i - 1]];
            }
        }
    }

    // Return the count of subsequences with the given sum
    return dp[n][sum];
}

// Main function
int main() {
    int arr[] = {1, 2, 1};
    int n = sizeof(arr) / sizeof(arr[0]);
    int K = 3;
    int count = countSubsequences(arr, n, K);
    cout << count << endl;
    return 0;
}

```

P3 Total Count the Subsequence whos Sum is Equal to the K

Now in this problem we only **Count** and all the possible Subsequence and **Sum** is **equal to the K**.

Thought Process

→ Now previously we learned one concept **pick and Not** pick and used **memo** again.

→ previously we maintain the Array for containing the element and sum is equal to the **K** and then array print it.

→ But here in place of **memo** we used **recursion** and **check**.

If Variable == **K** :
Count += 1
→ That is one case approach for this problem.

Example

arr = {1, 2, 1} K = 2

{1, 2} {2} → Count = 2 → **answery**

recursion(arr, count, k, arr)

If index >= len(arr) :
If count == K :
return 1
return 0

Count += arr[index]
pick = recursion(arr, count + 1, k, arr);
Count -= arr[index]
Notpick = recursion(arr, count, k, arr);
return pick + notpick;

Let's Dry Run and understand the Recursion

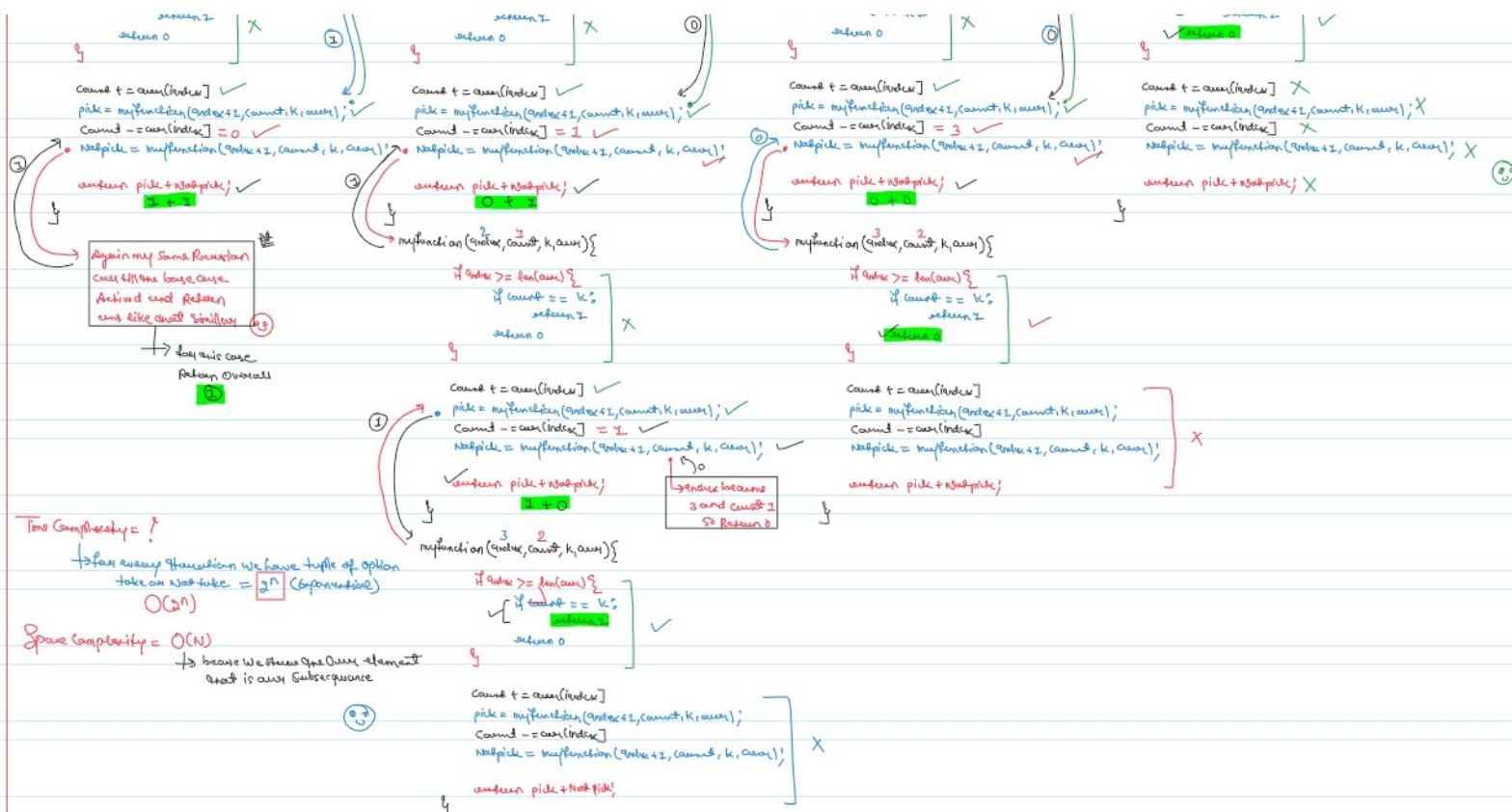
arr = {1, 2, 1}
K = 2
Count = 0
→ **Return 2** **answery**

recursion(arr, count, k, arr)
If index >= len(arr) :
If count == K :
return 1
return 0
Count += arr[index]
pick = recursion(arr, count + 1, k, arr);
Count -= arr[index]
Notpick = recursion(arr, count, k, arr);
return pick + notpick;

recursion(arr, count, k, arr)
If index >= len(arr) :
If count == K :
return 1
return 0
Count += arr[index]
pick = recursion(arr, count + 1, k, arr);
Count -= arr[index]
Notpick = recursion(arr, count, k, arr);
return pick + notpick;

recursion(arr, count, k, arr)
If index >= len(arr) :
If count == K :
return 1
return 0
Count += arr[index]
pick = recursion(arr, count + 1, k, arr);
Count -= arr[index]
Notpick = recursion(arr, count, k, arr);
return pick + notpick;

recursion(arr, count, k, arr)
If index >= len(arr) :
If count == K :
return 1
return 0
Count += arr[index]
pick = recursion(arr, count + 1, k, arr);
Count -= arr[index]
Notpick = recursion(arr, count, k, arr);
return pick + notpick;



```

// Code for the recursive solution
// Function to find the maximum sum of a subsequence with no two adjacent elements
// arr: array of integers, index: current index, count: number of elements picked, k: maximum number of elements to pick
// Returns: maximum sum of a subsequence

int maxFunction(int index, int count, int k, int arr[]) {
    // Base case: if index is out of bounds, return 0
    if (index >= arr.length) {
        return 0;
    }

    // If count is equal to k, return 0 as no more elements can be picked
    if (count == k) {
        return 0;
    }

    // Pick the current element
    int pick = arr[index] + maxFunction(index + 1, count + 1, k, arr);

    // Not pick the current element
    int notpick = maxFunction(index + 1, count, k, arr);

    // Return the maximum of pick and notpick
    return Math.max(pick, notpick);
}

// Main function to test the recursive solution
public static void main(String[] args) {
    int arr[] = {1, 2, 3, 4, 5};
    int k = 3;
    int index = 0;
    int count = 0;
    int maxSum = maxFunction(index, count, k, arr);
    System.out.println("Maximum sum of a subsequence: " + maxSum);
}

```

```

// Code for the iterative solution using dynamic programming
// arr: array of integers, k: maximum number of elements to pick
// Returns: maximum sum of a subsequence

int maxFunction(int arr[], int k) {
    // Base case: if k is 0, return 0
    if (k == 0) {
        return 0;
    }

    // Create a DP array to store the maximum sum of a subsequence for each index and count
    int dp[][] = new int[arr.length + 1][k + 1];

    // Fill the DP array
    for (int index = arr.length - 1; index >= 0; index--) {
        for (int count = 0; count < k; count++) {
            // Pick the current element
            int pick = arr[index] + dp[index + 1][count + 1];

            // Not pick the current element
            int notpick = dp[index + 1][count];

            // Store the maximum of pick and notpick in the DP array
            dp[index][count] = Math.max(pick, notpick);
        }
    }

    // Return the maximum sum of a subsequence
    return dp[0][0];
}

// Main function to test the iterative solution
public static void main(String[] args) {
    int arr[] = {1, 2, 3, 4, 5};
    int k = 3;
    int maxSum = maxFunction(arr, k);
    System.out.println("Maximum sum of a subsequence: " + maxSum);
}

```

