L7 All Kind of Patterns in Recursion

Sunday, December 18, 2022 4:19 PM

In this lecture we learn all about the different - different pattern based on Recursion In the previously lectures we learn in depth all about the Subsequences in depth

P1 Print All the Subsequence whos Sum is Equal to the K

In this problem our task is print/find all the subsequence that sum is equal to the K

→ Here we again used the concept of Pick and Not Pick

Example arr = [1, 2, 1], Sum = 2

(Now our task is print all the subsequence that sum is equal to the given sum)

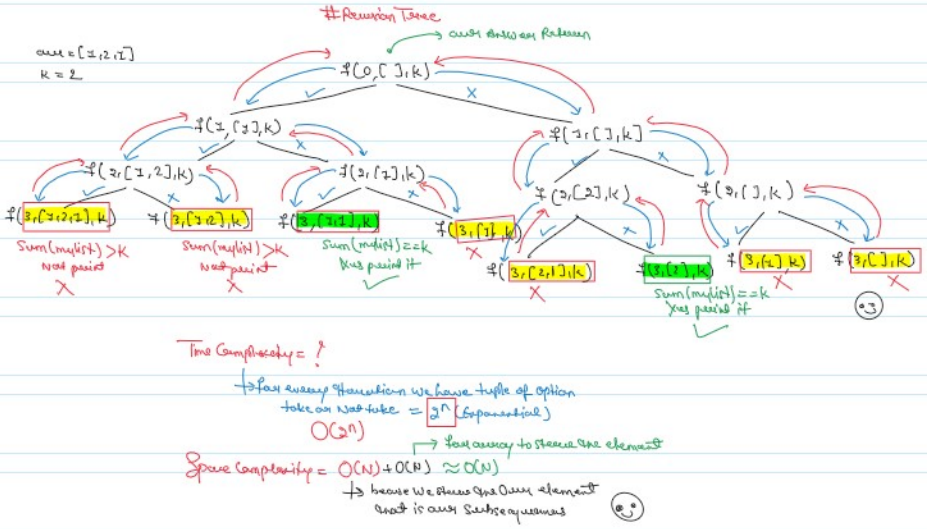
[1,1]
[2] } → that is all the possible Subsequence that sum is equal to the given sum

Thought Process

→ Again we used the concept of pick and Not pick which we learn previously

```
myfunction(index, mylist, k) {
    if(index >= n){
        if(sum(mylist) == k){
            print(mylist)        → my base case
            return;
        else{
            return;
        }
    }
    mylist.add(arr[index])       → Add the element/ Subsequence into mylist
    myfunction(index+1, mylist, k)   → pick case
    mylist.removed(arr[index])
    myfunction(index+1, mylist, k)   → Not pick case
}
```

→ Here we not discuss the Recursion code because that is the same Pseudo code that we learn previously that is easy



Time Complexity = ?

→ For every function we have tuple of option take or not take = 2^n (Exponential)

$O(2^n)$

Space Complexity = $O(N) + O(N) \approx O(N)$

→ for auxiliary to store the element

→ because We store the every element that is every subsequences

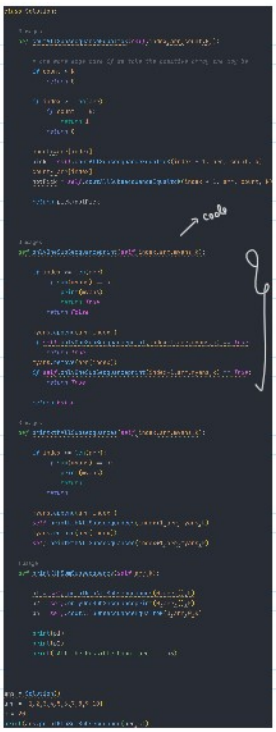
P2 Print Only One Subsequence whos Sum is Equal to the K

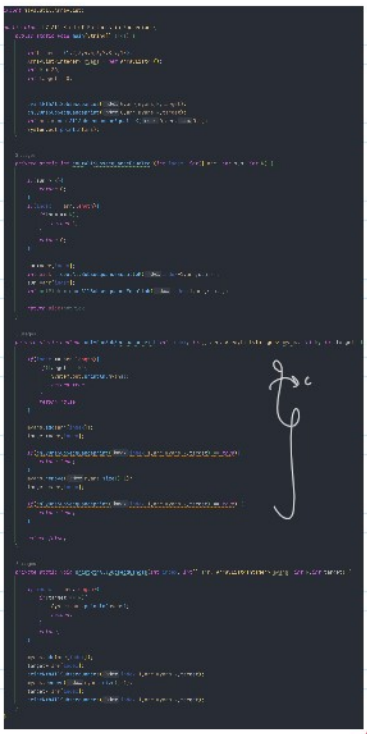
→ this is the modified problem of previous one here our task is print only one Subsequence

```
myfunction(index, mylist, K){

    if(index>=n){
        if(sum(mylist)==k){
            print(mylist)
            return True
        else{
            return False;
        }
                                  → my base case
    mylist.add(arr[index])        → Add the element/Subsequence into the mylist

Pick condition [
    if myfunction(index+1, mylist, K)==True{
        return True;
    }

    mylist.remove(arr[index])

Not pick condition [
    if myfunction(index+1, mylist, K)==True{
        return True;
    }
    return false;
}
```

→ if my Recursion become True than No Need for further Recursion Just Return

→ Recursion Tree, Sudo code and Recursion function are Same that is way Here We Not again discuss





P3 Total Count the Subsequence whos Sum is Equal to the K

Now In this problem we only count the all the possible Subsequence whose Sum is equal to the K

Thought Process

→ Now previously we learned one concept pick and Not pick that used here again

→ previously we maintain the Array for containing the element whose Sum is equal to the (K) and then we only print it

→ But here in place of Array we used one Variable and Check

if Variable == (K): Count += 1 → that is the easy approach for this problem

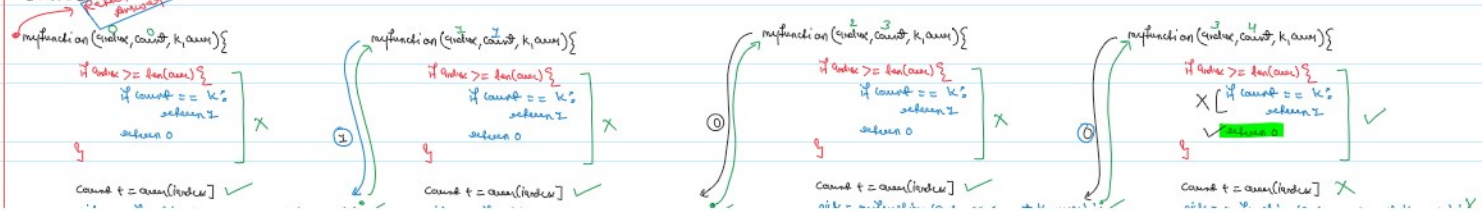
example arr = [1,2,1] K = 2

[1,1] [2] → Count = 2 → answer

```
myfunction(index, count, k, arr){
    if index >= len(arr){
        if count == k:
            return 1
        return 0
    }
    count += arr[index]
    pick = myfunction(index+1, count+k, arr);
    count -= arr[index]
    Notpick = myfunction(index+1, count, k, arr);

    return pick + Notpick;
}
```

Let's Dry Run and understand the Recursion

arr = [1,2,1]
K = 2
Count = 0 → return 2 answer



```
        return 1
    return 0
}

count += arr[index]
pick = myfunction(index+1, count, k, arr);
count -= arr[index]   = 0
Notpick = myfunction(index+1, count, k, arr);

return pick + Notpick;
        1 + 1
}
```

Again my Same Recursion call till the base case Activated and Return
ans like that Similar

→ for this case Return Overall **2**

```
        return 1
    return 0
}

count += arr[index]
pick = myfunction(index+1, count, k, arr);
count -= arr[index]   = 1
Notpick = myfunction(index+1, count, k, arr);

return pick + Notpick;
        0 + 1
}

myfunction(index=2, count=1, k, arr){
    if index >= len(arr) {
        if count == k:
            return 1
        return 0
    }

    count += arr[index]
    pick = myfunction(index+1, count, k, arr);
    count -= arr[index]   = 1
    Notpick = myfunction(index+1, count, k, arr);

    return pick + Notpick;
            1 + 0
}

myfunction(index=3, count=2, k, arr){
    if index >= len(arr) {
        if count == k:
            return 1
        return 0
    }

    count += arr[index]
    pick = myfunction(index+1, count, k, arr);
    count -= arr[index]
    Notpick = myfunction(index+1, count, k, arr);

    return pick + Notpick;
}
```

→ index became 3 and count 1 So Return 0

```
        return 0
}

count += arr[index]
pick = myfunction(index+1, count, k, arr);
count -= arr[index]   = 3
Notpick = myfunction(index+1, count, k, arr);

return pick + Notpick;
        0 + 0
}

myfunction(index=3, count=2, k, arr){
    if index >= len(arr) {
        if count == k:
            return 1
        return 0
    }

    count += arr[index]
    pick = myfunction(index+1, count, k, arr);
    count -= arr[index]
    Notpick = myfunction(index+1, count, k, arr);

    return pick + Notpick;
}
```

```
        return 0
}

count += arr[index]
pick = myfunction(index+1, count, k, arr);
count -= arr[index]
Notpick = myfunction(index+1, count, k, arr);

return pick + Notpick;
}
```

Time Complexity = ?

→ for every iteration we have tuple of option take or Not take = 2^n (Exponential)

$O(2^n)$

Space Complexity = $O(N)$

→ because we store the array element that is array Subsequence