# Introduction to Kokkos

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#### What is Kokkos?

- Kokkos is a C++\* performance portability library
  - Write a single source implementation
  - Descriptive programming model
  - Compile for CPUs or GPUs
  - Kokkos is a shared memory programming model (works in conjunction with MPI)
- Major buy-in by DOE and national labs
  - LAAMPS, Trilinos, ORNL Raptor
  - Over 100 projects using kokkos
  - Contributing to the C++ standard
  - Active kokkos developers community via slack (invite only)

Python and Fortran bindings are available

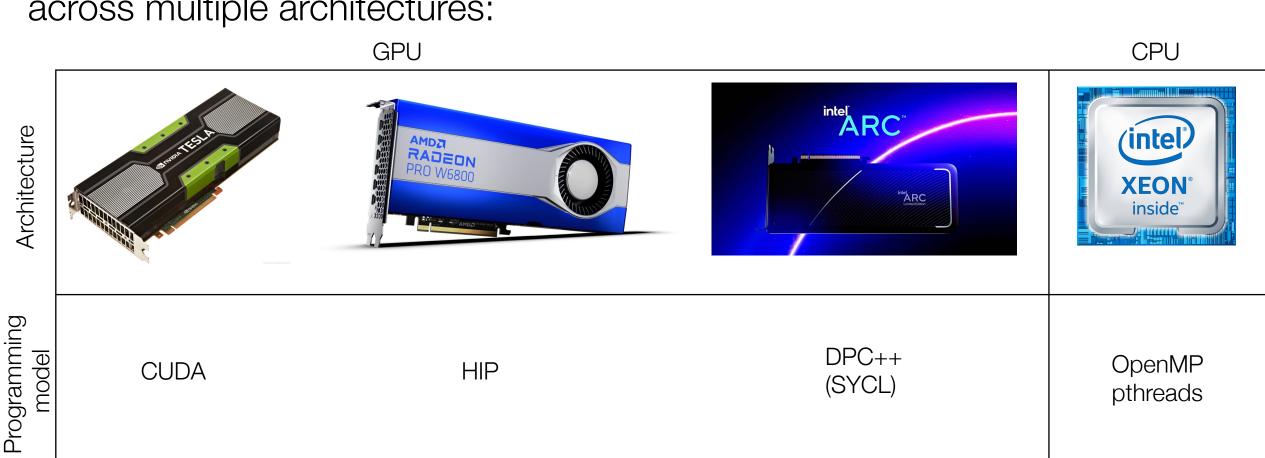
#### Goals

The goal here is to motivate a use case for Kokkos. Given time limit this is not an in-depth tutorial of Kokkos (probably sometime in the future).

- Outline of the talk:
  - Need for performance portability
  - Understanding performance
    - Caching (CPUs) vs Coalescing(GPUs)
  - Kokkos views
  - Live demonstration
  - Why use kokkos?

## Need for performance portability

Modern high-performance computing applications need to run efficiently across multiple architectures:



## Need for performance portability

Modern high-performance computing applications need to run efficiently across multiple architectures:

#### Pre-Exascale



LANL Trinity
Intel Haswell/ Intel KNL
OpenMP 3



ORNL Summit NVIDIA Volta100 OpenMP/CUDA

#### Exascale



ORNL Frontier AMD GPUs OpenMP/HIP



ANL Aurora Intel GPUs DPC++

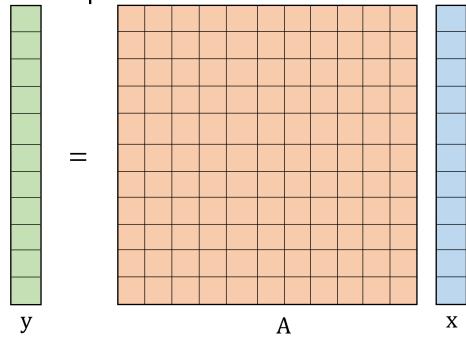
Moving into the exascale era will require adapting your applications to utilize modern architectures

#### Need for performance portability

- Problem: Porting applications for various architectures is time consuming
  - Typical HPC application: 300k 600k lines of code
    - Smaller scale applications: SpecFEM2D + SpecFEM3D ~ 100k lines of code (conservative estimate)
    - Porting requires ~10% rewrite of the application
    - Typical software engineer writes about 20k LOC/year
- Potential portability options: OpenMP 5 or OpenACC

What about performance??

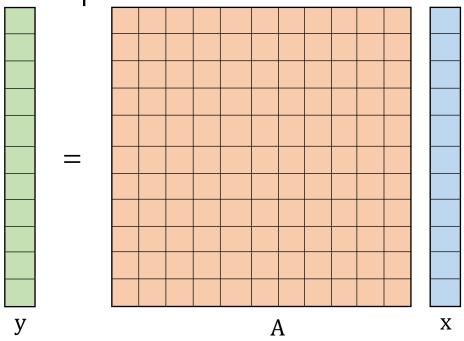
Example: Matrix – vector multiplication



Serial Implementation

```
for (int j = 0; j < N; j++){
    for (int i = 0; i < M; i++){
        y(j) += A(j,i)*x(i);
    }
}</pre>
```

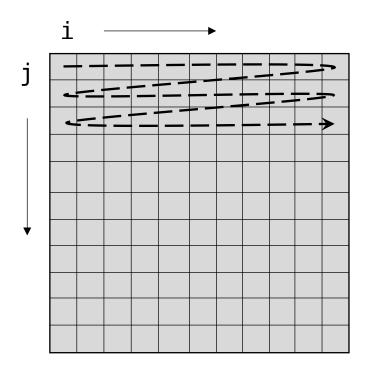
Example: Matrix – vector multiplication



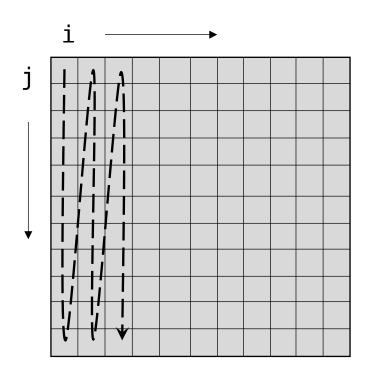
Parallel Implementation

```
parallel_for (int j = 0; j < N; j++){
// distribute different rows to different threads
    for (int i = 0; i < M; i++){
        y(j) += A(j,i)*x(i);
    }
}</pre>
```

• Performance is dependent on matrix layout

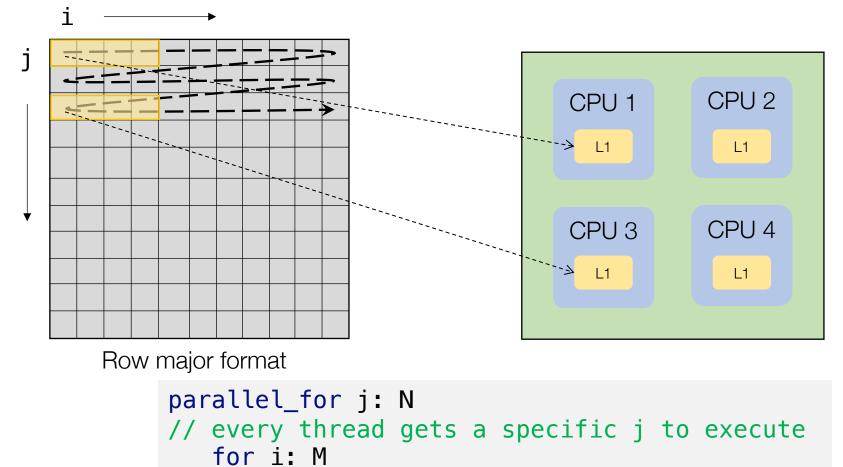


Row major format

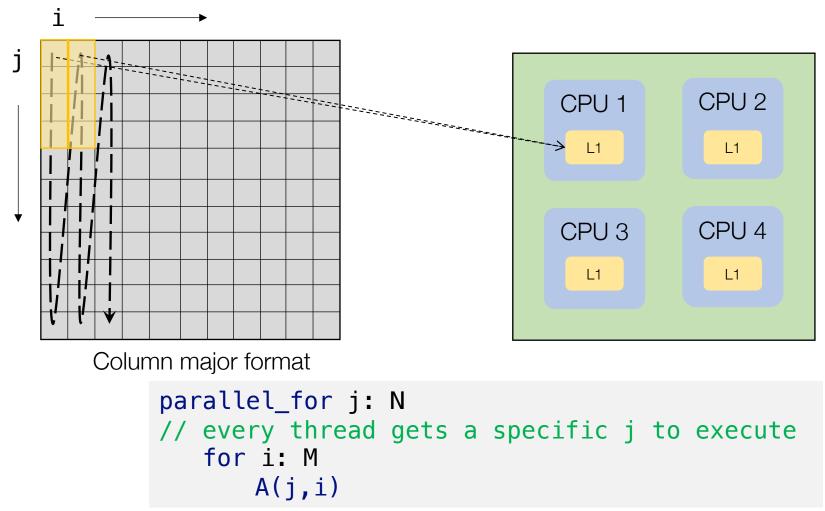


Column major format

Good performance

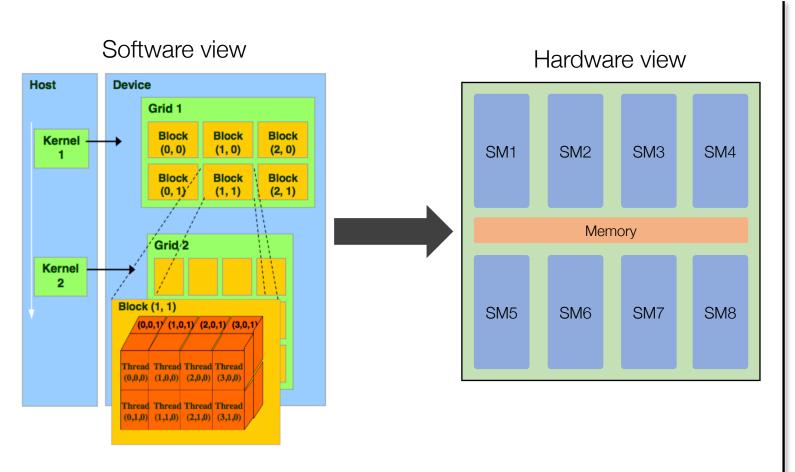


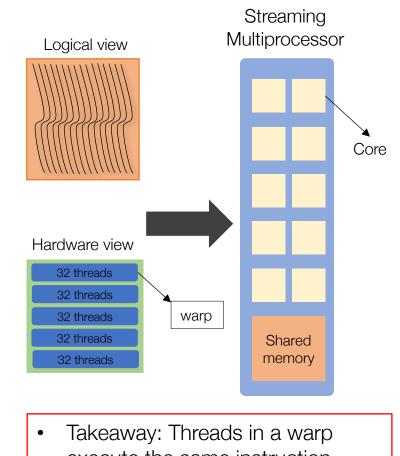
A(j,i)



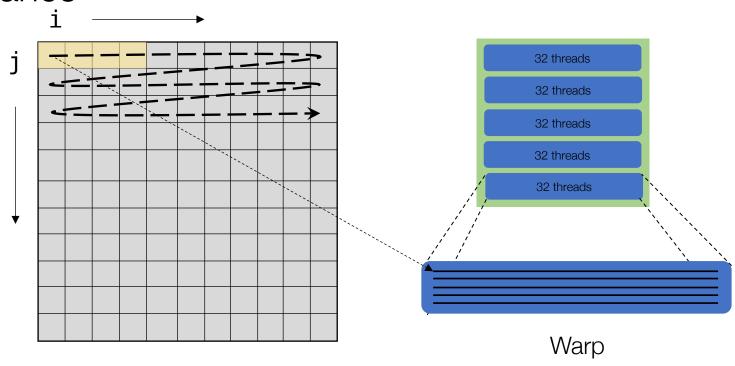
#### GPUs 1 min intro

GPU execution model

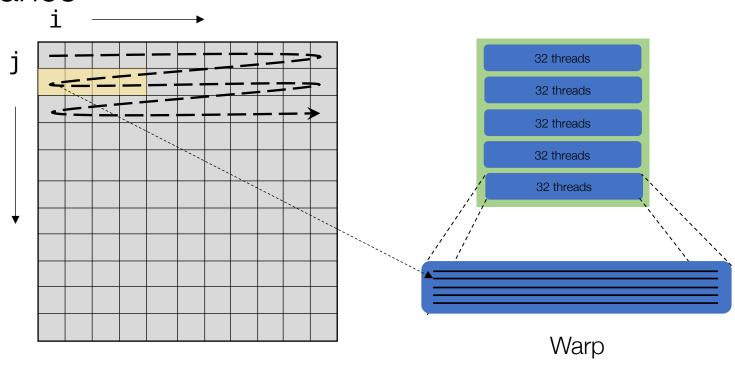




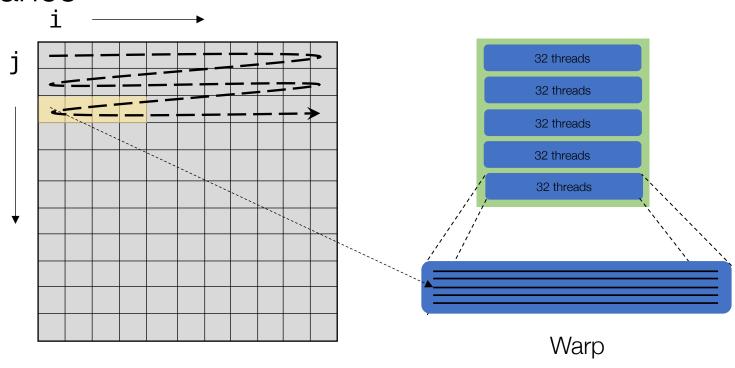
execute the same instruction



```
parallel_for j: N
// every thread gets a specific j to execute
   for i: M
        A(j,i)
```

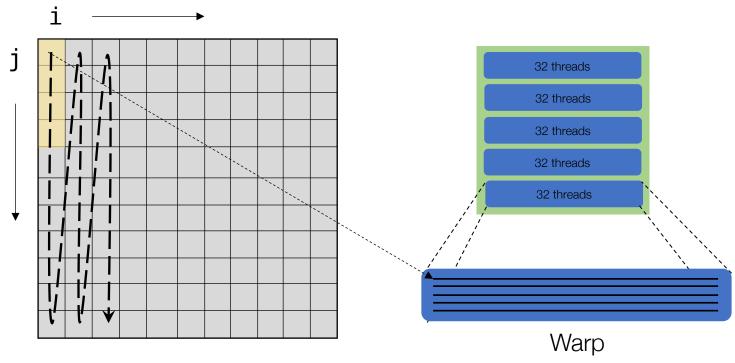


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Good performance



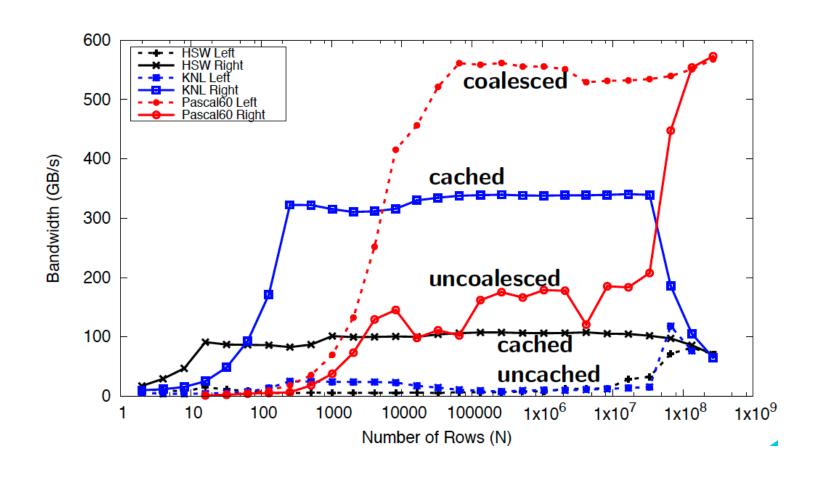
```
parallel_for j: N
// every thread gets a specific j to execute
   for i: M
        A(j,i)
```

- CPUs require row major format
  - $A(i_1, i_2, ..., i_n) \rightarrow i_n$  needs to be continuous in memory (Layout Right)
  - Results in data caching

- GPUs require a column major format
  - A( $i_1$ ,  $i_2$ , ...,  $i_n$ )  $\rightarrow i_1$  needs to be continuous in memory (Layout Left)
  - Results in data coalescing

#### Performance benchmarks

- vector matrix vector multiplication
- Total number of elements in matrix is a constant (M\*N = constant)
- Bandwidth ~ amount of data processed per second



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Data layout needs to be determined at compile time based on target architecture

#### Solution: Kokkos views

- Kokkos views are multidimensional arrays
- Can reside on CPUs or GPUs
  - Use "mirror" of your arrays to sync between CPUs and GPUs
  - deep\_copy can be initialized between original and mirror views
- View layout can be specified at compile time
  - Default behavior is to optimized layout based on execution space
    - GPUs = Layout Left
    - CPUs = Layout Right
- Views are reference counted
  - Easy memory management
- Copy construction leads to both views pointing to the same data
  - deep copy requires use of "deep\_copy" routine rather than a copy construction

#### Kokkos view

```
Kokkos::View<float **, Kokkos::LayoutLeft, Kokkos::CudaSpace>
A("A_matrix", N, M);
```

## Live Example

## Why use Kokkos?

	CUDA/HIP/DPC++	OpenMP 5/ OpenACC	Kokkos
Portability across architectures	No. Need to write separate kernels for every architecture	Yes. Single source code with pragma-based approach	Yes. Single source code implemented using Kokkos functions
Performance	Optimized performance	Tough to optimize for performance	Very good performance
Cost of portability	Very high	hy use Kokkos?	High
Cost of maintenance	Very high. Newer architectures might require tuning of kernels	Low. Assuming compilers do a good job of implementing the standard	Low. Assuming Kokkos backend is always optimized
Compiler dependence	N/A	These are standards, vendors have a flexibility on implementation	N/A
Fortran support	No. Could use bindings	Yes	No. Could use bindings

#### Summary

- Kokkos is a *performance portability* library
- Performance is dependent on data layout
  - Kokkos Views are easy way to manage data layout
- Compile time definition of data using C++ template meta-programming
- Low maintenance overhead for future architectures

#### Useful resources:

- Documentation
- Tutorials
- Slack (invite only): ctrott [a] sandia.gov
- Local help cses@princeton.edu

# Thank you