Pritam Mogal TE(B)-17

Assignment No 1

PASS-I ASSEMBLER.

Aim: To implement Pass-1 assembler:

Problem Statement: Design suitable data structure and implement Pass-1 of a two pass assembler for pseudo-machine in Java using object Orient feature. Implementation should consist of a few instructions from each category and few assembler directives.

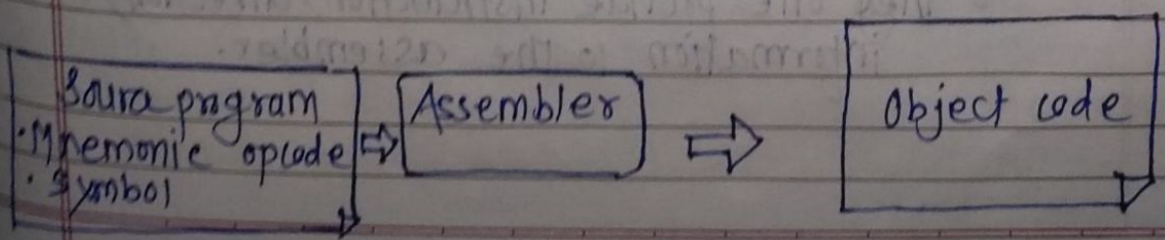
Theory:

Assembler language:

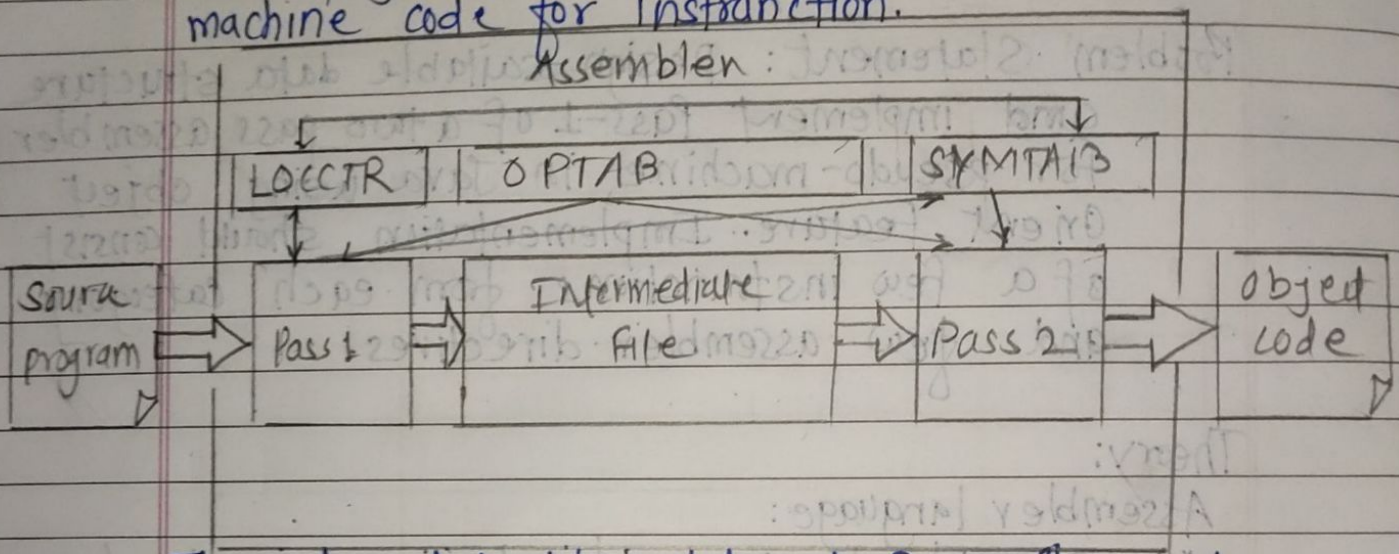
An assembly language is low level programming language for computer, or the programmable device, in which there is very strong (generally one-to-one) correspondance between the language and the architectures machine code instructions. Each assembly language is specific to a Particular computer architecture, in contrast to most high-level Programming language.

Assembler:

Assembly language is converted into executable machine code by utility program reffered to an an assembler The conversion process is reffered to as assembly, or assembling the code.



An assembler is a translator that translates an assembler into a conventional machine language program. Basically, the assembler goes through the program one line at a time, and generate machine code for instruction.



The intermediate file include each source statement, assigned address and error indicator.

Translate assemply language programs to object programs to machine code is called as assembler.
One simple way to eliminate this problem: require that all areas be defined before they are referenced It is possible, although inconvient, to do so for data items.

Assemble directives:
- Assembler directives are pseudo instructions.
  - They will not be translated into machine instructions.
  - They one provide instruction/direction/information to the assembler.

- Basic assembles directives :
  - START : Specify name and starting address for the program.
  - END: Indicate the end of the source program.
  - EQU: The EQU directive is used to replace a number by a symbol. For example: MAXIMUM EQU 98. After using this directive, every appearance for the label "Maximum" in the program will be interpreted by the assembler.

Three main Data structures
- Operation Code table (OPTAB)
- Location Counter (LOCCTR)
- Symbol Table (SYMTAB).

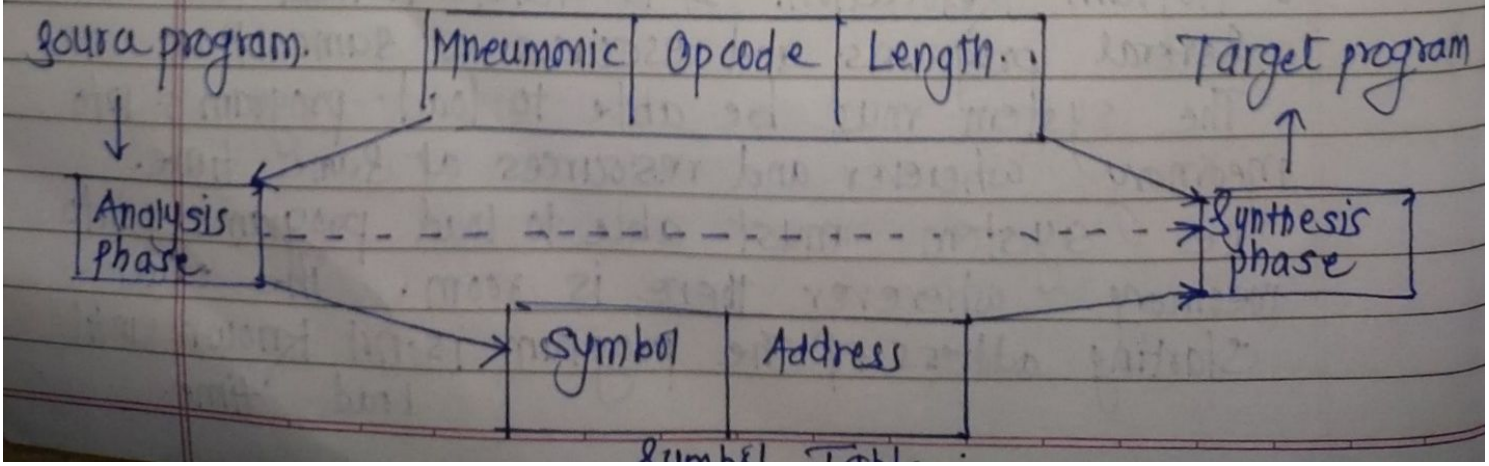
Instruction formats:
- Addressing modes: Direct addressing (address of operand is given in instruction itself). Register addressing (one of the operand is general purpose register) Register indirect addressing (address of operand is specified by register pair. Immediate addressing (operand - data is specified in the instruction itself) Implicit addressing (mostly the operation operates on the contents of accumulator).

- Program Relocation: It is desible to load and run several programs and resources at same time. The system must be able to load programs into memory wherever and resources at same time. The system must able to load programs into memory whereever there is room. The exact starting address of the program is not known until load time.

Q ~~22~~
~~Macro~~ name table

Literal :
it is convinient for the programmer to be able to write the value of a constant operand as a part of the instruction that uses it. Such an operand is called a literal.
The difference between literal operands and immediate operands
- For literal operand we use '=' as prefix, and with immediate operand we use '#' as prefix.
- During immediate addressing, the operand value is assembled as part of the machine instruction, and there is no memory refer
- With a literal, the assembler generates the specified value as a constant at some other memory location.

One-Pass assembler :
A one pass assembler passes over the soura file exactly once, in the same pass collecting the labels, resolving future references and doing the actual assembly.



Forward referena in one assembler..

- Omits the operand address if the symbol has not yet been defined.
- Enters this undefined symbol into SYMTAB and indicates that it is undefined
- Adds the address of this operand address to a list of forward references associated with the SYMTAB entry.
- When the defination for the symbol is encountered, scans the referena list and inserts the address.
- At the end defination for the symbol is encountered, scans the reference list and inserts the address.
- At the end of the program the error if there are still SYMTAB entries indicated undefined symbols.

Data structures for assembler:

OP code Table

Looked up for the translation of mnemonic code.

Key: mnemonic code.

Algorithm for pass 1 assembler:

```
begin
    if starting address is given
        LOCCTR = starting address;
    else
        LOCCTR = 0;
    while OPCODE != END do      ;; or EOF
        begin
        read a line from the code
        if there is a label
            if this label is in SYMTAB, then error
```

else insert (label, LOCCTR) into SYMTAB
Search OPTAB for the op code.
if found
LOCCTR += N ;; N is the length of this instruction (4 for MIPS)
else if this is an assembly directive
update LOCCTR as directed.
else error
write line to intermediate file
end
program size = LOCCTR - Starting address;
end.

Algorithm 4.1 (Assembler first Pass)
1. loc_cntr := 0 (default value)
pooltab_ptr := 1; POOLTAB[1] := 1;
littab_ptr := 1;
2: while next Statement is not an END Statement.
(a) if label is present then
this_label := symbol in label field;
Enter(this, label, loc_cntr) in SYMTAB.
(b) If an LTORG Statement then
(i) Process literals LITTAB[POOLTAB[POOLTAB_ptr]]... LITTAB(lit_tab_ptr-1) to allocate memory and put the address in the address field.
(ii) pooltab_ptr := pooltab_ptr + 1;
(iii) POOLTAB[pooltab_ptr] := littab_ptr,
(c) If a START or ORIGIN Statement then
loc_cntr := value Specified in operand field;

(d) If an EQU Statement then
(i) this_addr := value of <address spec>;
(ii) Correct the symtab entry for this label to (this_label, this_addr).
(e) If a declaration statement then
(i) code := code of the declaration statement;
(ii) size := size of memory area required by DC/DS.
(iii) loc_cntr := loc_cntr + size;
(iv) Generate IC (DL code)....'

(f) If an imperative statement then
(i) code := machine opcode from OPTAB;
(ii) loc_cntr := loc_cntr + instruction length from OPTAB;
(iii) If operated is a literal then
this_literal := literal in operand field;
LITTAB [littab_ptr] := this_literal;
littab_ptr := littab_ptr + 1;
else (i.e) operand is a symbol)
this_entry := SYMTAB entry number of operand;
Generate IC '(IS, code) (S, this_entry)';

3. (Processing of END Statement)
(a) Perform step 2(b).
(b) Generate IC '(AD, 02).
(c) Go to Pass II

Input:
START 200
MOVER AREG, ='4'

```
MOVEM AREG, A
MOVER BREG = '1'
LOOP MOVER CREG, B
LTORG
ADD GREG, ='6'
STOP
A DS 1
B DS 1
END
```

Expected Output : Symbol Table

A	208
LOOP	203
B	209

Intermediate Code

AD	01	C	200	
IS	04	1	L	1
IS	05	1	S	1
IS	04	2	L	2
IS	04	3	S	3
AD	05			
IS	01	3	L	3
IS	00			
DL	02	C	1	
DL	02	C	1	
AD	02			

Conclusion :
Thus we have implement PASS-1 Assembler using object oriented features.

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Assignment No. 01 [Pass 1 Assembler]

Problem Satement: Design suitable data structures and implement pass-I of a twopass assembler for pseudo-machine in Java using object oriented feature. Implementation should consist of a few instructions from each category and few assembler directives

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1. Pass 1 Program:

```
import
java.io.BufferedReader;
import java.io.*; import
java.io.IOException; import
java.util.*;

public class Pass1 { public static void
        main(String[] args) {

                BufferedReader br = null;
                FileReader fr = null;

                FileWriter fw = null;
                BufferedWriter bw = null;

                try {
                        String inputfilename = "/home/sagar-ravan/Desktop/Input.txt";
                        fr = new FileReader(inputfilename); br = new
                        BufferedReader(fr);

                        String OUTPUTFILENAME = "/home/sagar-ravan/Desktop/IC.txt";
                        fw = new FileWriter(OUTPUTFILENAME);
                        bw = new BufferedWriter(fw);

                        Hashtable<String, String> is = new Hashtable<String, String>();
                        is.put("STOP", "00"); is.put("ADD", "01"); is.put("SUB",
                        "02"); is.put("MULT", "03"); is.put("MOVER", "04");
                        is.put("MOVEM", "05"); is.put("COMP", "06"); is.put("BC",
                        "07"); is.put("DIV", "08"); is.put("READ", "09");
                        is.put("PRINT", "10");

                        Hashtable<String, String> dl = new Hashtable<String, String>();
                        dl.put("DC", "01"); dl.put("DS", "02");
                        Hashtable<String, String> ad = new Hashtable<String, String>();

                        ad.put("START", "01");
                        ad.put("END", "02");
                        ad.put("ORIGIN", "03");
```

```
ad.put("EQU", "04");
ad.put("LTORG", "05");

Hashtable<String, String> symtab = new Hashtable<String, String>();
Hashtable<String, String> littab = new Hashtable<String, String>();
ArrayList<Integer> pooltab = new ArrayList<Integer>();

String sCurrentLine; int
locptr = 0; int litptr = 1; int
symptr = 1; int pooltabptr =
1; sCurrentLine =
br.readLine();

String s1 = sCurrentLine.split(" ")[1];
if (s1.equals("START")) {
        bw.write("AD \t 01 \t");
        String s2 = sCurrentLine.split(" ")[2];
        bw.write("C \t" + s2 + "\n");
        locptr = Integer.parseInt(s2);
}

while ((sCurrentLine = br.readLine()) != null) { int mind_the_LC = 0;
        String type = null; int flag2 = 0; // checks whether addr is
        assigned to current symbol

        String s = sCurrentLine.split(" |\\,")[0]; // consider the first word in the
line

        for (Map.Entry m : symtab.entrySet()) { // allocating addr to arrived
symbols if (s.equals(m.getKey())) {
                m.setValue(locptr);
                flag2 = 1;
            }
        }
        if (s.length() != 0 && flag2 == 0) { // if current string is not " " or
addr is not assigned,

    // then the current string must be a new symbol.
          symtab.put(s, String.valueOf(locptr));
              symptr++;
        }
        int isOpcode = 0; // checks whether current word is an opcode or not

        s = sCurrentLine.split(" |\\,")[1]; // consider the second word in the
line

        for (Map.Entry m : is.entrySet()) { if (s.equals(m.getKey())) {
                bw.write("IS\t" + m.getValue() + "\t"); // if match found
in imperative stmt
```

```
                        type = "is";
                        isOpcode = 1;
                    }
                }

                for (Map.Entry m : ad.entrySet()) { if (s.equals(m.getKey())) {
                    bw.write("AD\t" + m.getValue() + "\t"); // if match
found in Assembler Directive type = "ad"; isOpcode = 1;
                    }
                }
                for (Map.Entry m : dl.entrySet()) { if (s.equals(m.getKey())) {
                    bw.write("DL\t" + m.getValue() + "\t"); // if match
found in declarative stmt type = "dl"; isOpcode = 1;
                    }
                }

                if (s.equals("LTORG")) {
                    pooltab.add(pooltabptr);
                    for (Map.Entry m : littab.entrySet()) { if (m.getValue() == "")
                        { // if addr is not assigned to the
literal
                            m.setValue(locptr);
                            locptr++;
                            pooltabptr++;
                            mind_the_LC = 1;
                            isOpcode = 1;
                        }
                    }
                }

                if (s.equals("END")) {
                    pooltab.add(pooltabptr);
                    for (Map.Entry m : littab.entrySet()) {
                        if (m.getValue() == "") {
                            m.setValue(locptr);
                            locptr++; mind_the_LC =
                            1;
                        }
                    }
                }

                if (s.equals("EQU")) { symtab.put("equ",
                    String.valueOf(locptr));
                }

                if (sCurrentLine.split(" |\\,").length > 2) { // if there are 3 words
                    s = sCurrentLine.split(" |\\,")[2]; // consider the 3rd word

                    // this is our first operand.
```

```
                        // it must be either a
                        Register/Declaration/Symbol if
                        (s.equals("AREG")) { bw.write("1\t"); isOpcode
                        = 1;
                        } else if (s.equals("BREG")) {
                                bw.write("2\t");
                                isOpcode = 1;
                        } else if (s.equals("CREG")) {
                                bw.write("3\t");
                                isOpcode = 1;
                        } else if (s.equals("DREG")) {
                                bw.write("4\t");
                                isOpcode = 1;
                        } else if (type == "dl") {
                                bw.write("C\t" + s + "\t");
                        } else { symtab.put(s, ""); // forward referenced
                        symbol }
                }

                if (sCurrentLine.split(" |\\,").length > 3) { // if there are 4 words

                        s = sCurrentLine.split(" |\\,")[3]; // consider 4th word.


        // this is our 2nd operand

        // it is either a literal, or a symbol if
                                (s.contains("=")) {
                                littab.put(s, "");
                               bw.write("L\t" + litptr + "\t");
                                isOpcode = 1;
                                litptr++;
                                } else { symtab.put(s, ""); // Doubt : what if the current
                                symbol
is already present in SYMTAB?
                                                                        // Overwrite?
                                bw.write("S\t" + symptr + "\t");
                                symptr++;
                                        }
                }

                    bw.write("\n"); // done with a line.

                if (mind_the_LC == 0)
                        locptr++;
        }

        String f1 = "/home/sagar-ravan/Desktop/SYMTAB.txt";
                FileWriter fw1 = new FileWriter(f1);
        BufferedWriter bw1 = new BufferedWriter(fw1); for
        (Map.Entry m : symtab.entrySet()) { bw1.write(m.getKey()
```

```
                + "\t" + m.getValue() + "\n");
                System.out.println(m.getKey() + " " + m.getValue());
                }

                String f2 = "/home/sagar-ravan/Desktop/LITTAB.txt";
                        FileWriter fw2 = new FileWriter(f2);
                BufferedWriter bw2 = new BufferedWriter(fw2); for
                (Map.Entry m : littab.entrySet()) { bw2.write(m.getKey() +
                "\t" + m.getValue() + "\n"); System.out.println(m.getKey()
                + " " + m.getValue());
                }

                String f3 = "/home/sagar-ravan/Desktop/POOLTAB.txt";
                        FileWriter fw3 = new FileWriter(f3);
                BufferedWriter bw3 = new BufferedWriter(fw3);
                for (Integer item : pooltab) {
                        bw3.write(item + "\n");
                        System.out.println(item);
                }

                bw.close();
                bw1.close();
                bw2.close();
                bw3.close();

        } catch (IOException e) {
                e.printStackTrace();
        }

    }

}
```

PASS 1 - ASSEMBLER OUTPUT:

Pritam-spos@pritam-HP:~/Desktop$ javac Pass1.java

Note: Pass1.java uses unchecked or unsafe operations.

Note: Recompile with –Xlint:unchecked for details.

Pritam-spos@-HP:~/Desktop$ java Pass1 Input.txt

A 8

LOOP 3

B 9

= ‘4’ 4

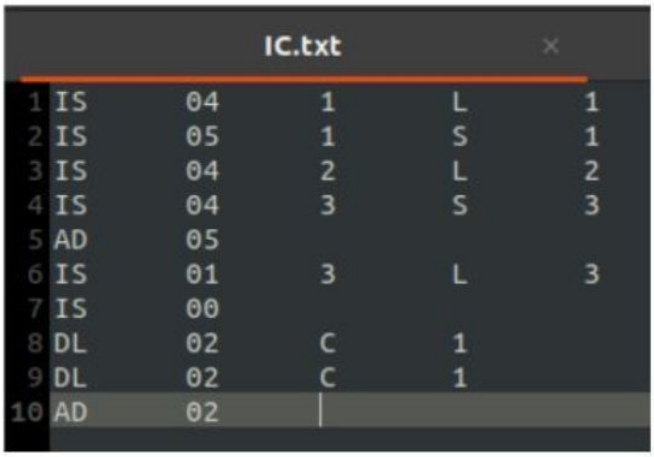
= ‘6’ 10

= ‘1’ 5

1

3

IC.txt



SYMTAB.txt

SYMTAB.txt

```
1 A       8
2 LOOP    3
3 B       9
```

LITTAB.txt
POOLTAB.txt

POOLTAB.txt

```
1 1
2 3
```

LITTAB.txt

```
1 ='4'    4
2 ='6'    10
3 ='1'    5
```