DS & AI ENGINEERING

Artificial Intelligence

Informed search



Lecture No.- 03

Recap of Previous Lecture







Properties of Henristics - Admissible
- Consistent

Topics to be Covered











About Aditya Jain sir



- 1. Appeared for GATE during BTech and secured AIR 60 in GATE in very first attempt City topper
- 2. Represented college as the first Google DSC Ambassador.
- 3. The only student from the batch to secure an internship at Amazon. (9+ CGPA)
- 4. Had offer from IIT Bombay and IISc Bangalore to join the Masters program
- 5. Joined IIT Bombay for my 2 year Masters program, specialization in Data Science
- 6. Published multiple research papers in well known conferences along with the team
- 7. Received the prestigious excellence in Research award from IIT Bombay for my Masters thesis in ML
- 8. Completed my Masters with an overall GPA of 9.36/10
- 9. Joined Dream11 as a Data Scientist
- 10. Have mentored 15,000+ students & working professions in field of Data Science and Analytics
- 11. Have been mentoring & teaching GATE aspirants to secure a great rank in limited time
- 12. Have got around 27.5K followers on Linkedin where I share my insights and guide students and professionals.



Revision: Graph and Tree Search



- Graph Search \rightarrow We make a closed list and never visit closed node.
- Tree Search → We do not make closed list.



Revision: Graph and Tree Search



- Graph Search → Consist noptimal solution.
- Tree Search → Admissible hoptimal solution.



h=2

- 1) Admissible? 2) Consistent?

1) Admissible? — No

15-2 2 5 1+1

2) Consistent? — No $1 \le 2+1$ $h(n) \le c(n,n') + h(n')$

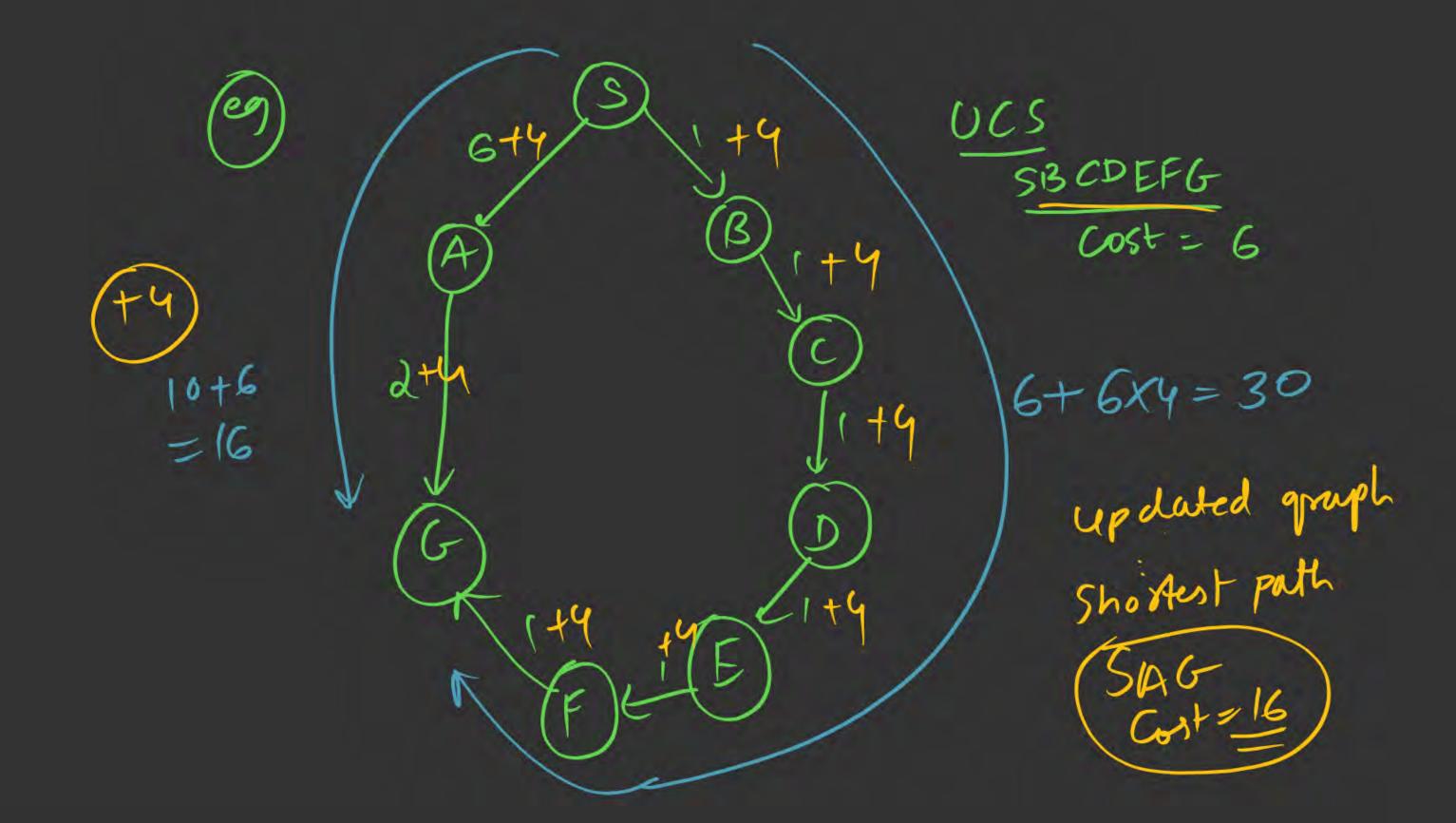
h=2 Bad Heuristics not optimal h=1 h=6 closed SBCG n=0 Cost = 6

h=2 1 h= G h=0

(2) F	Tree Search: CX
SX AB CG	Cost=6 Cost=6 Cost=6

Q.I Graph - UCS -, can path change?

(Yes) it may change



Graph ->UCS - (pun) **C (cray edge multiplied by some const c)

Graph' - UCS -, can path change? No, will not change



Revision: Time and Space Complexity



• $A^* \Rightarrow O(b^d) \rightarrow We might store all nodes (Visit all nodes)$



Revision: Heuristic - consistent and admissible



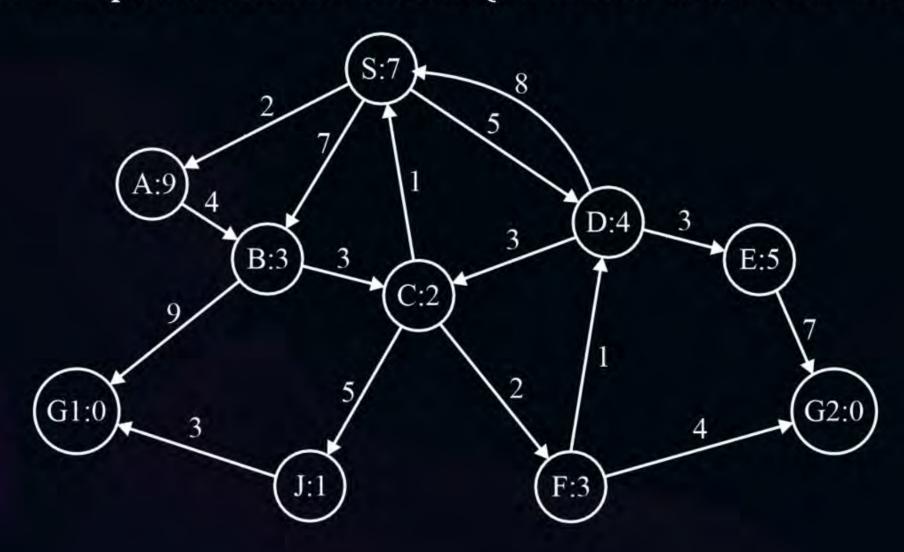
Summary:

	A* tree	A* graph
If C heuristic	✓	✓
If A heuristic	✓	×
If neither C nor A heuristics	×	*





#Q. Consider the search space below, where S is the start node and Gl and G2 satisfy the goal test. Arcs are labeled with the cost of traversing them and the heuristic cost to a goal is reported inside nodes (so lower scores are better).







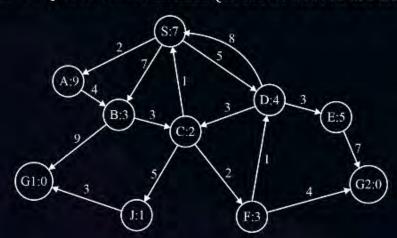
For A* search, indicate which goal state is reached at what cost and list, in order, all the states popped off of the OPEN list. You use a search graph to show your work.

Note: When all else is equal, nodes should be removed from OPEN in alphabetical order.





#Q. Consider the search space below, where S is the start node and Gl and G2 satisfy the goal test. Arcs are labeled with the cost of traversing them and the heuristic cost to a goal is reported inside nodes (so lower scores are better).







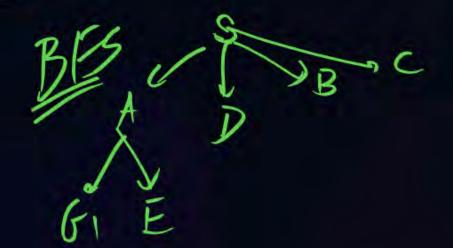


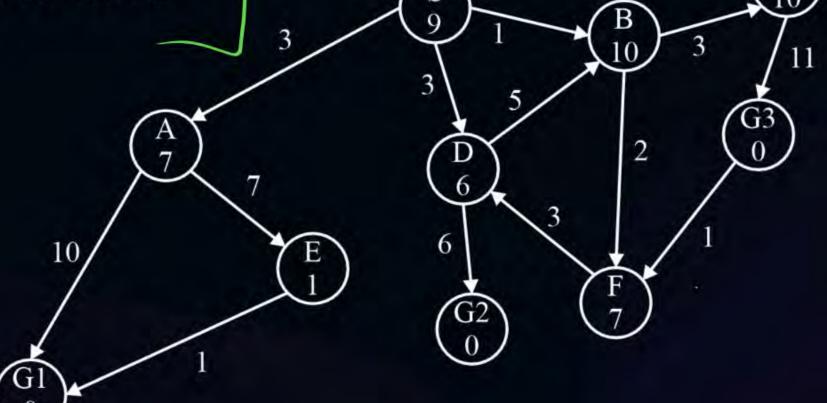
#Q. Consider the search graph below, where S is the start node and G1, G2, and G3 are goal states. Arcs are labeled with the cost of traversing them and the heuristic cost to a goal is shown inside the nodes. For each of the three search strategies below, indicate which of the goal states is reached:

(a) Breadth-first search. Goal reached: G1

(b) Uniform cost search. Goal reached: G2

(c) A* search. Goal reached: G2





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Advantages

- Optimality: A* is guaranteed to find the least-cost path to the goal if the heuristic used is <u>admissible</u> (never overestimates the true cost) and consistent (satisfies the triangle inequality).
- Efficiency: A* is generally more efficient than uninformed search algorithms like Dijkstra's algorithm because it uses heuristics to guide the search. This often results in exploring fewer nodes, reducing the overall computation.
- Flexibility: A* can be adapted to various problems by changing the heuristic function. This flexibility makes it applicable to a wide range of scenarios, from simple grid navigation to complex graph-based pathfinding.
- Combines Advantages of Greedy and Uniform-Cost Search: A* combines the strengths of greedy best-first search (which is fast but not always optimal) and uniform-cost search (which is optimal but can be slow) by balancing exploration and exploitation.





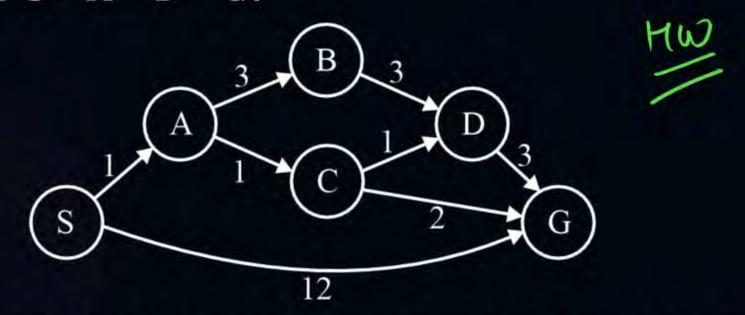
Disadvantages

- High Memory Usage: A* requires storing all generated nodes in memory, which can lead to high memory consumption, especially in large search spaces. This makes it less suitable for problems with very large state spaces.
- Computationally Intensive: In the worst case, A* can be computationally expensive, especially if the branching factor is high or the heuristic is not well-optimized. The algorithm's time complexity is exponential in the worst case.
- Heuristic Dependency: The performance of A* heavily depends on the quality of the heuristic. If the heuristic is poorly chosen, the algorithm can degrade to a less efficient search, potentially exploring many unnecessary nodes.
- Not Always Practical for Real-Time Applications: Due to its potential high time and space complexity, A* may not be suitable for real-time applications where quick decisions are necessary.





#Q. Answer the following questions about the search problem shown above. Break any ties alphabetically. For the questions that ask for a path, please give your answers in the form 'S - A - D - G'.



- (a) What path would breadth-first graph search return for this search problem?
- (b) What path would uniform cost graph search return for this search problem?
- (c) What path would depth-first graph search return for this search problem?
- (d) What path would A* graph search, using a consistent heuristic, return for this search problem?





#Q. Consider the heuristics for this problem shown in the table below.

(i) Is h₁ admissible? Yes (No)

(ii) Is h₁ consistent? Yes (No)

(iii) Is h₂ admissible? (Yes) No

(iv) Is h₂ consistent? Yes No

		1
State	$ ho_1$	h ₂
S	5	4
Α	3	2
ВН	6	6
С	2	1
D	3	3
G	0	0

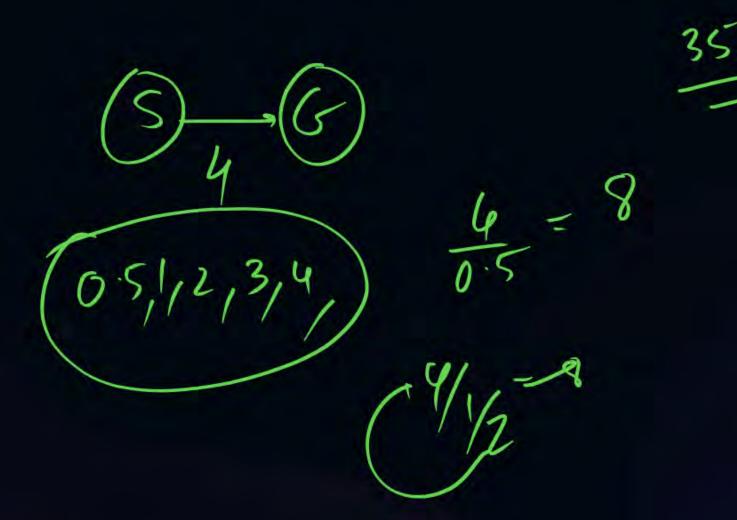




#Q. Let h_1 and h_2 be two <u>admissible</u> heuristics used in A* search. Which ONE of the following expressions is always an admissible heuristic?

- $h_1 + h_2$
- $h_1 \times h_2 \times$
- $h_1 / h_2, (h_2 \neq 0)$
- |h₁ h₂|









#Q. If h1 and h2 are two admissible heuristics, then which of the following are guaranteed to be admissible?

- Min(h1, h2)
- max(h1, h2)
- (h1 + h2) / 2
- (h1.h2)^0.5



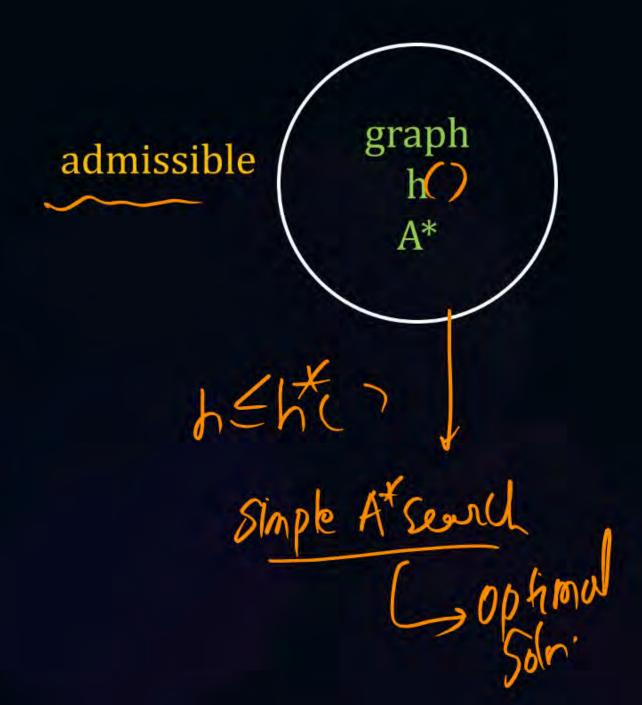


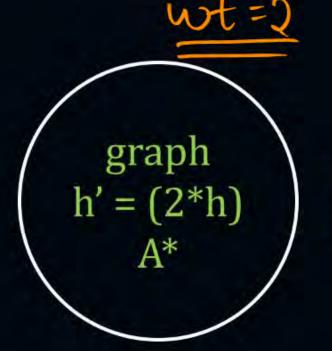
Weighted A* search Algorithm:

- Weighted A* is a variant of the classic A* algorithm used for faster pathfinding by introducing a weight factor to the heuristic. It trades optimality for speed, making it useful in time-critical applications.
- A* guarantees optimal path (if heuristic is admissible and consistent) but can be slow. Weighted A* speeds up the search by being more aggressive in exploring nodes closer to the goal, even if it sacrifices optimality.









h' = no more admissible So, no more optimal

Same as applying weighted A* on original graph with weight = 2





Weighted A* search Algorithm:

Classic A* uses:

$$f(n) = g(n) + h(n)$$

Weighted A* modifies this to:

$$f(n) = g(n) + w \cdot h(n)$$

Where:

- g(n) = cost from start to current node n
- h(n) = heuristic estimate from n to goal
- $w \ge 1 = weight factor$





Weighted A* search Algorithm:

If w = 1: behaves like classic A*, optimal and complete

If w > 1 : heuristic is over-emphasized, search is faster but not guaranteed to be optimal

• Higher the more greedy the algorithm becomes (closer to Greedy Best-First Search)



f(n) = g(n)



Weighted A* search Algorithm:

Uniform Cost

Algorithm	Modification	Notes
A*	f(n) = g(n) + h(n)	Optimal
Weight A*	$f(n) = g(n) + w \cdot h(n)$	Faster, suboptimal
Greedy BFS	f(n) = h(n)	Fastest, no optimality

Ignores heuristic, but optimal



