

Consider the right recursive grammar $E \rightarrow T + E \mid T$ $T \rightarrow V * T \mid V$ $V \rightarrow \text{<id>}$	Top-down Parsing with Back tracking																
<u>Left Factoring</u>  Common prefix requires applying left factoring $E \rightarrow T + E \mid T$ $T \rightarrow V * T \mid V$ Left factor E $E \rightarrow T + E \mid T$ $E \rightarrow TE''$ $E'' \rightarrow + E \mid \epsilon$ Left factor T $T \rightarrow V * T \mid V$ $T \rightarrow VT''$ $T'' \rightarrow * T \mid \epsilon$	Top-down parsing without back tracking																
Modified Left Factored Grammar $E ::= T \{ +T \}^*$ $T ::= V \{ * V \}^*$ $V ::= \text{<id>}$																	
Consider the Left Recursive Grammar $E \rightarrow E + T \mid T$ $T \rightarrow T * V \mid V$ $V \rightarrow \text{<id>}$	RD parser																
<u>Eliminating left-recursion</u> Apply left recursion elimination $E \rightarrow E + T \mid T$ $T \rightarrow T * V \mid V$  Eliminate left recursion of E $E \rightarrow TE'$ $E' \rightarrow \epsilon \mid + TE'$ Eliminating left recursion of T $T \rightarrow VT'$ $T' \rightarrow \epsilon \mid + VT'$	Not Used as it is for Top-Down Parsing																
Operator Grammar using (OPM)	LL1 / Table driven / Top-down Parsing																
	Operator Precedence Parser (Bottom Up Parser)																
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