Lab report no: 08

Lab report name: Implementation of SJF scheduling algorithm.

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Objective:- Shortest Job First (SJF) scheduling algorithm technique in detail.

To understand the advantage and disadvantage of Shortest Job First (SJF) scheduling algorithm technique.

Shortest Job First:- Shortest job first can be either preemptive or non-preemptive. Owing to its simple nature, shortest job first is considered optimal. It also reduces the average waiting time for other processes awaiting execution.

Algorithm:

- 1. Sort all the process according to the arrival time.
- 2. Then select that process which has minimum arrival time and minimum Burst time.
- 3. After completion of process make a pool of process which after till the completion of previous process and select that process among the pool which is having minimum Burst time.
- 1. Completion Time: Time at which process completes its execution.
- 2. Turn Around Time: Time Difference between completion time and arrival time. Turn Around Time = Completion Time — Arrival Time
- 3. Waiting Time (W.T): Time Difference between turn-around time and burst time.

Waiting Time = Turn Around Time – Burst Time In this post, we have assumed arrival times as 0, so turn around and completion times are same.

Gantt chart:-

P1		P3	P2		P4	
0	7	8		12		16

process	Arrival time	Burst time	Completion	Turn-around	Waiting time
			time	time	
P1	0	7	7	7	0
P2	2	4	12	10	10
Р3	4	1	8	4	3
P4	5	4	16	11	7

Average waiting time = (0+6+3+7)/4

=4ms

Average turn-around time = (7+10+8+11)/4 =8ms

```
Code:-
```

*a = *b;

```
// C++ program to implement Shortest Job first with Arrival Time
#include<iostream>
using namespace std;
int mat[10][6];

void swap(int *a, int *b)
{
  int temp = *a;
```

```
*b = temp;
void arrangeArrival(int num, int mat[][6])
{
for(int i=0; i<num; i++)
{
for(int j=0; j<num-i-1; j++)</pre>
 {
 if(mat[j][1] > mat[j+1][1])
  for(int k=0; k<5; k++)
  swap(mat[j][k], mat[j+1][k]);
void completionTime(int num, int mat[][6])
{
```

```
int temp, val;
mat[0][3] = mat[0][1] + mat[0][2];
mat[0][5] = mat[0][3] - mat[0][1];
mat[0][4] = mat[0][5] - mat[0][2];
for(int i=1; i<num; i++)
{
temp = mat[i-1][3];
int low = mat[i][2];
for(int j=i; j<num; j++)</pre>
{
 if(temp >= mat[j][1] && low >= mat[j][2])
 low = mat[j][2];
 val = j;
 }
}
mat[val][3] = temp + mat[val][2];
mat[val][5] = mat[val][3] - mat[val][1];
mat[val][4] = mat[val][5] - mat[val][2];
for(int k=0; k<6; k++)
```

```
swap(mat[val][k], mat[i][k]);
 }
int main()
{
int num, temp;
cout<<"Enter number of Process: ";</pre>
cin>>num;
cout<<"...Enter the process ID...\n";
for(int i=0; i<num; i++)</pre>
 cout<<"...Process "<<i+1<<"...\n";
 cout<<"Enter Process Id: ";</pre>
 cin>>mat[i][0];
 cout<<"Enter Arrival Time: ";</pre>
 cin>>mat[i][1];
 cout<<"Enter Burst Time: ";</pre>
 cin>>mat[i][2];
```

```
arrangeArrival(num, mat);
    completionTime(num, mat);
     cout<<"Process ID\tArrival Time\tBurst Time\tWaiting Time\tTurnaround Time\n";</pre>
   for(int i=0; i<num; i++)
     {
    cout << mat[i][0] << "\t\t" << mat[i][1] << "\t\t" << mat[i][2] << "\t\t" << mat[i][4] << mat[
 "\t\t"<<mat[i][5]<<"\n";
     }
Output:
                        .Enter the process ID...
```

}

```
Enter number of Process: 4
...Enter the process ID...
...Process 1...
Enter Process Id: 4
Enter Process Id: 4
Enter Burst Time: 5
Enter Burst Time: 4
...Process 2...
Enter Process Id: 3
Enter Arrival Time: 4
Enter Burst Time: 1
...Process 3...
Enter Process Id: 2
Enter Arrival Time: 2
Enter Arrival Time: 2
Enter Burst Time: 4
...Process 4...
Enter Process Id: 1
Enter Arrival Time: 0
Enter Burst Time: 7
Process ID Arrival Time Burst Time Waiting Time Turnaround Time 1
0 7 0 7
3 4 1 3 4
4 5 4 3 7
2 2 4 10 14

Process returned 0 (0x0) execution time: 47.147 s
Press any key to continue.
```

Conclusion:- Regarding the average and waiting time is pretty advantage then FCFS .This because the execution of small processes in start brings a little time waiting for long processes. One of the major disadvantage of this algorithm is: a process must wait in the queue if he has a large time of execution although he may be in the queue for a long time if on the queue come processes with short execution time .