

Day-17

Insertion Sort

⇒ Partially sorting the array

Eg: $\begin{matrix} 0 & 1 & 2 & 3 & 4 \\ 5 & 3 & 4 & 1 & 2 \end{matrix}$

1st pass: $\begin{matrix} 3 & 5 & 4 & 1 & 2 \\ p=0 \end{matrix}$

2nd pass: $\begin{matrix} 3 & 4 & 5 & 1 & 2 \\ p=1 \end{matrix}$

3rd pass: $\begin{matrix} 3 & 4 & 5 & 1 & 2 \\ p=2 \end{matrix}$

4th pass: $\begin{matrix} 1 & 2 & 3 & 4 & 5 \\ p=3 \end{matrix}$
Sorted array

Working:

$\begin{matrix} 0 & 1 & 2 & 3 & 4 \\ 5 & 3 & 4 & 1 & 2 \\ \text{Swap } 3 \text{ } 4 & \\ 3 & 5 & 4 & 1 & 2 \end{matrix}$

$\begin{matrix} 3 & 5 & 4 & 1 & 2 \\ 4 < 5 \text{ Swap} & \\ 3 & 4 & 5 & 1 & 2 \end{matrix}$

$\begin{matrix} 3 & 4 & 5 & 1 & 2 \\ \text{Swap } 1 < 5 & \\ 3 & 4 & 1 & 5 & 2 \\ \downarrow 1 < 4 & \\ 3 & 1 & 4 & 5 & 2 \\ \downarrow 1 < 3 & \\ 1 & 3 & 4 & 5 & 2 \end{matrix}$

$\begin{matrix} 1 & 3 & 4 & 5 & 2 \\ \text{Swap } 2 < 5 & \\ 1 & 3 & 4 & 2 & 5 \\ \downarrow 2 < 4 & \\ 1 & 3 & 2 & 4 & 5 \\ \downarrow 2 < 3 & \\ 1 & 2 & 3 & 4 & 5 \end{matrix}$

Sorted

1, 2, 3, 4, 5

Sorted Break

For every index:

Put that index element at the correct position of LHS

Outer loop:

Index	Pass	Value
0	1	5
1	2	3
2	3	4
3	4	1
4	5	2

$p < (n-2)$ | $j > 0$

0

1

1

2

2

3

3

4

Time Complexity:

Best case: Array is already sorted

1 2 3 4 5

(N-1) Linear

$O(N)$

Worst case: Desc sorted

5 4 3 2 1

$\frac{N(N+1)}{2}$ comparisons

$$\frac{(N-1)(N+1)}{2} = \frac{N^2 - N}{2} \Rightarrow O(N^2)$$

Why we use?

- Adaptive: No. of swaps reduced compared to bubble sort
- Stable:
- Used for smaller values of N works good when array is partially sorted. In Hybrid Sorting algo

Program:

```
import java.util.Arrays;
```

```
public class InsertionSort {
```

```
    public static void main (String [] args) {
```

```
        int [] arr = {3, 1, 4, 5, 2};
```

```
        insertion (arr);
```

```
        System.out.println (Arrays.toString (arr));
```

```
    } static void insertion (int [] arr) {
```

```
        for (int i = 0; i < arr.length - 1; i++) {
```

```
            for (int j = i + 1; j > 0; j--) {
```

```
                ✓ if (arr[j] < arr[j-1]) {
```

```
                    int temp = arr[j];
```

```
                    arr[j] = arr[j-1];
```

```
                    arr[j-1] = temp;
```

```
                } else {
```

```
                    break;
```

```
            }
```

```
        }
```

```
    }
```

```
}
```

4. (arr[j] > arr[j-1])

→ [5, 4, 3, 2, 1]

[1, 2, 3, 4, 5]

Selection Sort: Select element & put it on its correct index

(n-1) 4, 5, 1, 2, 3
 Swap

(n-2) 4, 3, 1, 2, 5
 swap

(n-3) 2, 3, 1, 4, 5
 swap

(n-4) 2, 1, 3, 4, 5
 swap

0 1, 2, 3, 4, 5
 Sorted array!

* Check the largest element & put it in correct index place
 (p.e) index = 4

* Check 2nd largest element

Also select minimum element & put it in correct index

0 1 2 3 4
 4, 5, 1, 2, 3
 Swap

⇒ 1, 5, 4, 2, 3

1, 2, 4, 5, 3

1, 2, 3, 5, 4

1, 2, 3, 4, 5 Sorted

Total comparison: $0 + 1 + 2 + \dots + (n-1)$

$$\frac{(n-1) \times (n+1-1)}{2}$$

$$= \frac{n(n-1)}{2} = \frac{n^2 - n}{2}$$

$O(n^2)$

Constant removed

Time complexity: Worst case: $O(n^2)$ Stable = No

Performs well on Small list/array. Best case: $O(n)$

⇒ In this sort, Largest number move to large index place by swapping method.

⇒ In best & worst case, both are same $O(n^2)$

Program:

```
import java.util.Arrays;
public class selection {
    public static void main (String [] args) {
        int [] arr = {64, 25, 12, 22, 11};
        selectionsort (arr);
        System.out.println (Arrays.toString(arr));
    }
```

```
    static void selectionsort (int [] arr) {
        for (int i = 0; i < arr.length; i++) {
            int min = i;
            for (int j = i + 1; j < arr.length; j++) {
                if (arr[j] < arr[min]) {
                    min = j;
                }
            }
            int temp = arr[min];
            arr[min] = arr[i];
            arr[i] = temp;
        }
    }
}
```

$i < min$
Asc order

[11, 12, 22, 25, 64]

$i > min$
Desc order

[64, 25, 22, 12, 11]