

Assignment - 2

CSA0389 - Data Structure for Stack
implementation.

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Topic : pseudocode for stack operation.

3) Insert 70
50 60 70

30 50 70

20 60
10 30 50 70

(3 children)

Describe the concept of Abstract data type (ADT) and how they differ from concrete data structures. Design an ADT for a stack and implement it using arrays and linked list in C. Include operations like Push, Pop, peek, is empty, is full and peek.

ABSTRACT DATA TYPE :

An Abstract data type is a theoretical model that defines a set of operations and the semantics (behaviour) of those operations on a data structure, without specifying how the data structure should be implemented. It provides a high level description of what operations can be performed on data and what constraints apply to those operations.

CHARACTERISTICS OF ADTS :

Operations : Defines a set of operations that can be performed on the data structure.

Semantics : Specifies the behaviour of each operation.

Encapsulation : Hides the implementation details, focusing on the interface provided to the user.

ADT FOR STACK :

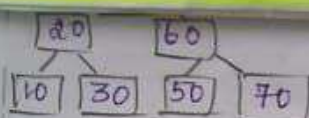
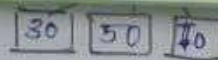
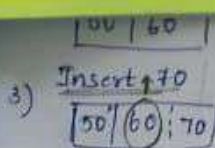
A stack is a fundamental data structure that follows the last in, first out (LIFO) principle. It supports the following operations.

Push : Adds an element to the top of the stack.

Pop : Removes and returns the element from the stack of top.

Peek : Returns the element from the top of the stack without removing it.

is empty : Check if the stack is empty.



(3 children)

is full : check if the stack is full.

Concrete data Structures :

The implementations using arrays and linked lists are specific ways of implementing the stack ADT in C.

how ADT differ from concrete data structure :

ADT focuses on the operations and their behaviour, while concrete data structures focus on how these operations are realised using specific programming constructs (arrays and linked list)

ADV OF ADT :

By separating from its implementation system, you achieve modularity, encapsulation, and flexibility in designing and using data structures in programs. This separation allows for easier maintenance, code reuse and abstraction of the complex operations.

Implementation in C using Arrays :

```
#include <stdio.h>
#define MAX_SIZE 100
typedef struct {
    int items[MAX_SIZE];
    int top;
} stackArray;

int main() {
    stackArray stack;
    stack.top = -1;
    stack.items[++stack.top] = 10;
    stack.items[++stack.top] = 20;
    stack.items[++stack.top] = 30;
```

```

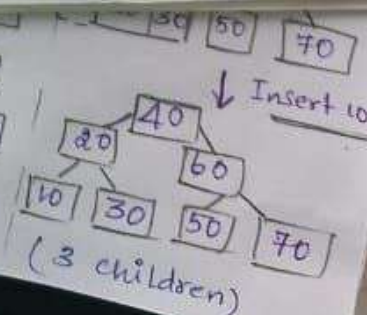
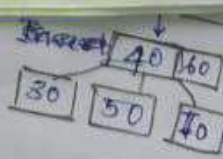
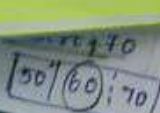
printf("memory allocation failed \n");
return 1;
}
newNode->data = 30;
newNode->next = top;
top = newNode;
if (top != NULL) {
    printf("top element : %d\n", top->data);
}
else {
    printf("stack is empty : \n");
}
if (top == NULL) {
    Node *temp = top;
    printf("popped element : %d\n", temp->data);
    printf("stack underflow : \n");
}
if (stack.top != -1) {
    printf("top element after pops : %d\n", stack.items[stack.top]);
}
return 0;
}

```

```

#include <stdio.h>
#include <stdlib.h>
typedef struct Node {
    int data;

```




```

if (stack.top != -1) {
    printf ("Top element : %d \n", stack.items[stack.top]);
} else {
    printf ("stack is empty : \n");
}
if (stack.top == -1) {
    printf ("popped element : %d \n", stack.items[stack.top]);
}
else {
    printf ("stack underflow! \n");
}
if (stack.top != -1) {
    printf ("popped element : %d \n", stack.items[stack.top]);
}
else {
    top = newNode;
    newnode = (Node*) malloc (sizeof (node));
    if (newnode == null) {
        printf ("memory allocation failed \n");
        return 1;
    }
    newnode->data = 20;
    newnode->next = top;
    top = newnode;
    newnode = (node*) malloc (sizeof (node));
    if (newnode == null) {

```

```

    struct Node *next;
} node;

int main() {
    node *top = null;
    node *newnode = (node *) malloc (size of (node));

    if (newnode == null) {
        printf ("memory allocation failed\n");
        return 1;
    }

    newnode->data = 10;
    newnode->next = top;
    top = top->next;
    free (temp);
}

else {
    printf ("stack underflow\n");
    while (top != null) {
        node *temp = top;
        top = top->next;
        free (temp);
    }
    return 0;
}

```

to reduce the number of disc access.

- * Time complexity to insert, delete - $O(\log n)$.
- * All leaves are at same value.
- * $m-1$ key values.

(i) 2-3 TREE : (each node has 2 or 3 children w/ atleast val)

1) Insert 50

```

    [50]

```

2) Insert 60

```

    [60]
   /  \
 [50] [70]

```

3) Insert 30

```

    [40 | 60]
   /  \
 [50] [70]

```

4) Insert 40

```

    [40 | 60]
   /  \
 [50] [70]

```

5) Insert 20

```

    [20 | 40 | 60]
   /  |  \
 [50] [70] [ ]

```

Types of

- 2-3 Tree
- 2-3-4
- 2-3-4-5
- 2-3-4-5-6

1) Initialize stack () :

the stack.

Initialize necessary variable or structure to present

2) push (elements) :

if stack is full :

print "stack overflow"

else :

add element to the top of the stack

increment top pointer

3) pop () :

if stack is empty :

print ("stack underflow")

return null (or appropriate error value)

else :

remove and return the element to the top of the stack

decrement end pointer

4) peek () :

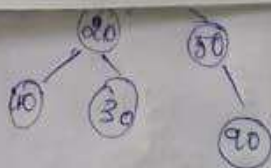
if stack is empty () :

print ("stack is empty")

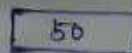
return null or appropriate error value

5) isEmpty () :

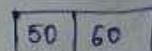
return true if Top is -1 (stack is empty)



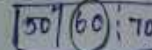
1) Insert 50



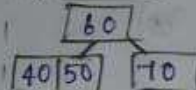
2) Insert 60



3) Insert 70



4) Insert 40



6) Is full:

return True, if top is equal to maxsize - 1 (stack is full)
otherwise, return False.

Initialize the necessary variables or data structure to represent a stack

Adds an element to the top of the stack, check if the stack is full before pushing

Removes and returns the element from the top of the stack, check if the stack is empty before popping.

Returns the element at the top of the stack without removing it, check if the stack is empty without peeking.

Checks if the stack is empty by inspecting the top pointer or equivalent variable

Checks if the stack is full by comparing the top pointer or equivalent variable to the maximum size of the stack.

Linear search works by checking each element in the list one by one until the desired element found in the end of the list that doesn't require any prior sorting of the data.

50 | 60 | 70

10 | 30 | 50 | 70

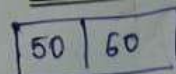
(3 children)


```

#include <stdio.h>
int main () {
    int regNumbers = {20142010};
    int target = 20142010;
    int n = sizeof(regNumbers) / sizeof(regNumbers[0]);
    int found = 0;
    int i;
    for (i = 0; i < n; i++) {
        if (regNumbers[i] == target) {
            printf ("Registration number id found at index: %d\n",
                    target, i);
            found = 1;
            break;
        }
    }
    if (!found) {
        printf ("Registration number id not found in
                list : \n", target);
    }
    return 0;
}

```

2) Insert 60



Insert 70

4) Insert 40

