

Searching - 1



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PSP - 68.6%



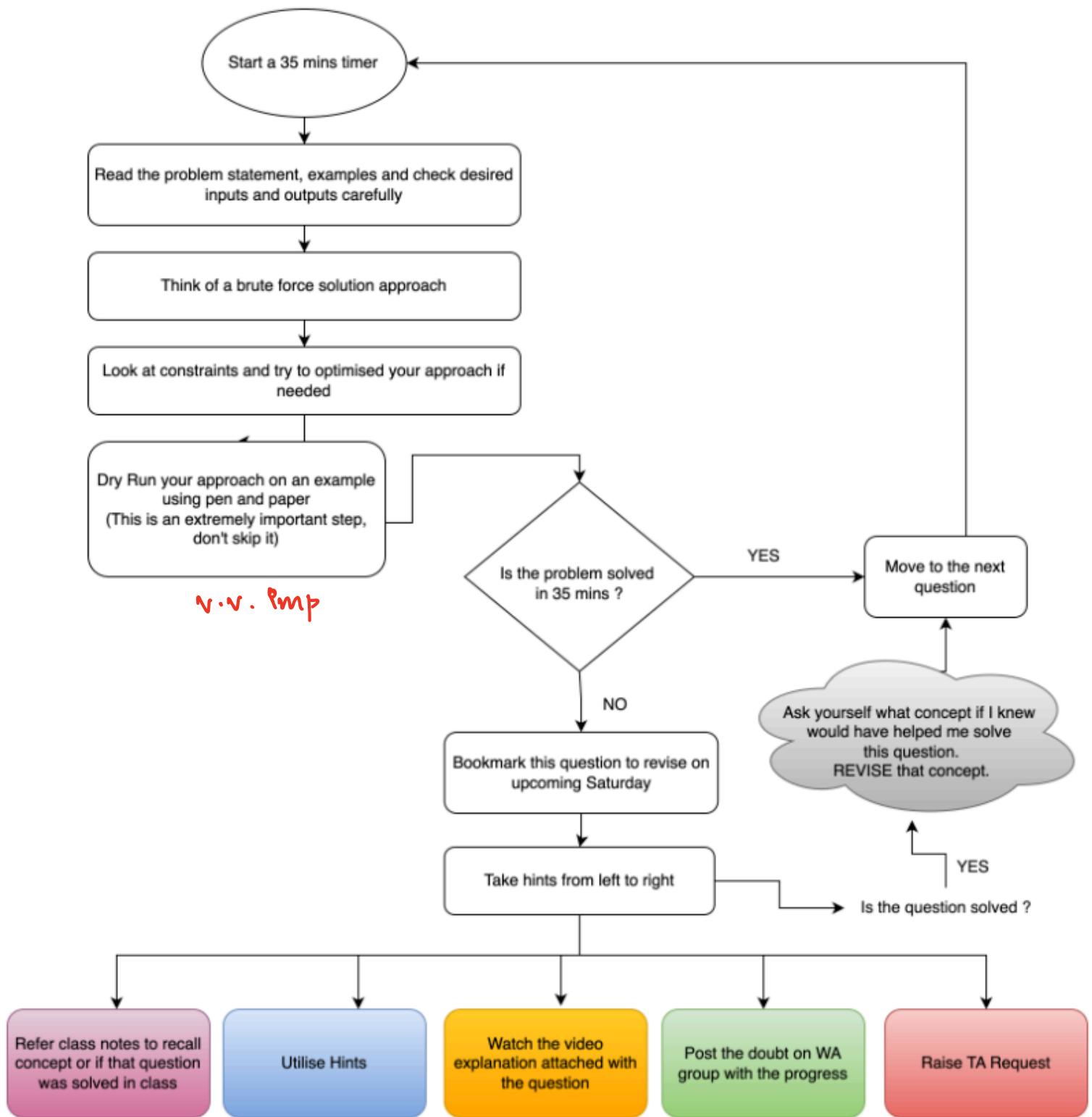
Challenges are what makes
life interesting.

Overcoming them is what
makes life meaningful.

Can we apply B.S only on sorted array?

No! :)

How to solve a problem ?



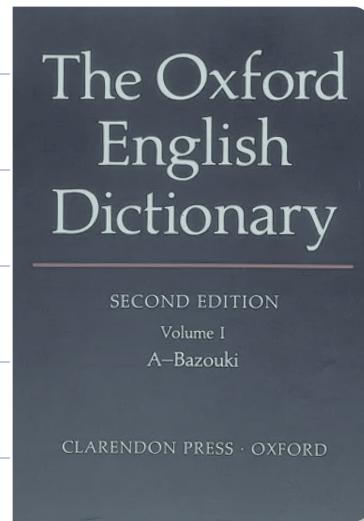


Searching

Searching a word in newspaper



Searching a word in dictionary



Binary Search

Search Space : The area to know the result exist
and we search there only.

Target : The item we are looking for.

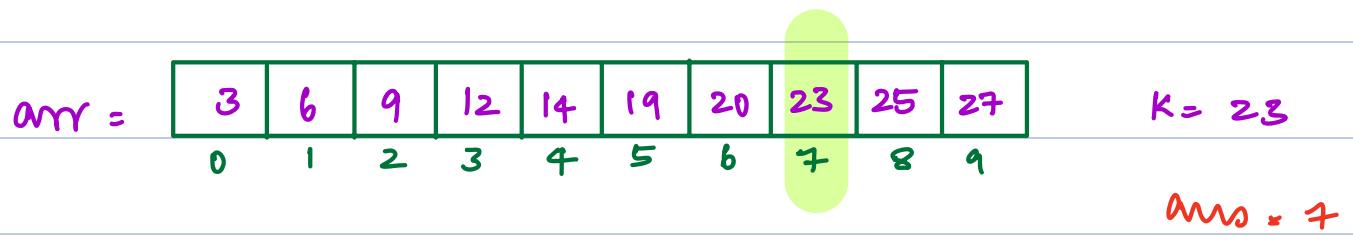
Condition : Condition to discard some part of search

space repeatedly

we discard one half of the array



Given a sorted arr[]. Search an element K. If K is present return its index otherwise return -1.



BF Approach

Int find K (Int arr[], Int K) {

 for (Int P=0; P<arr.length; P++) {

 if (arr[P] == K) return P;

 return -1;

y

T.C = O(N)

S.C = O(1)

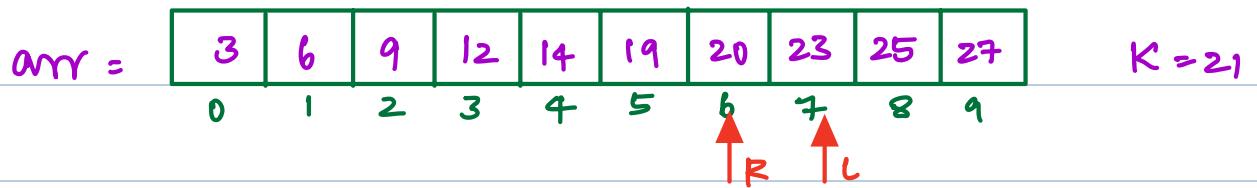
Idea -2

Binary Search

Search Space: The entire array

Target: Input K

Condition:



L	R	MID	Remarks
0	9	$(0+9)/2 = 4$	$14 < 21$ (go right) $I = mid + 1$
5	9	$(5+9)/2 = 7$	$23 > 21$ (go left)
5	6	$(5+6)/2 = 5$	$19 < 21$ (go right)
6	6	$(6+6)/2 = 6$	$20 < 21$ (go right)
7	6	—	stop

pseudo code

```

int binary-search (int A[], int B) {
    int l=0, R = A.length - 1;
    while (l <= R) {
        int mid = (l+R)/2;
        if (A[mid] == B) return mid;
        else if (A[mid] < B) l=mid+1;
        else R = mid - 1;
    }
    return -1;
}

```

$$\text{int } mid = (l+R)/2;$$

$$l + \frac{(R-l)}{2}$$

If $A[mid] == B$ return mid;

else if $A[mid] < B$ $l=mid+1$;

else $R = mid - 1$;

return -1;

T.C = $O(\log N)$

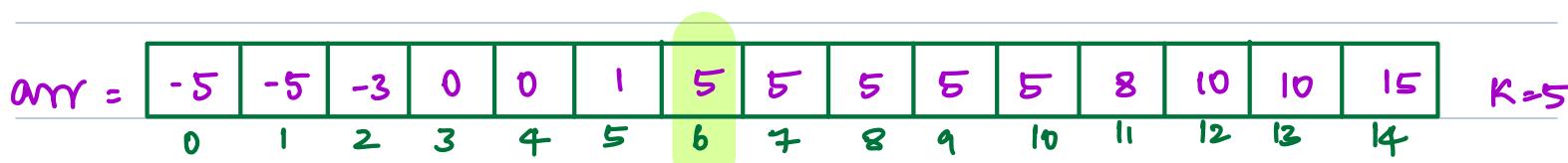
S.C = $O(1)$



Story Based Question

All emails in your mailbox are sorted chronologically. Can you find the first mail that you received in 2024?

Given a sorted arr[N]. Find first occurrence of K.



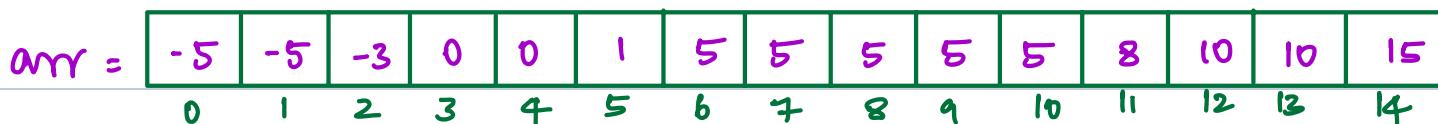
BF Approach

Linear Search

T.C = O(N) S.C = O(1)



Idea - 2



Ans =  b

L	R	MID	Remarks
0	14	7	$a[7] == 5$ (go left)
0	6	3	$a[3] < 5$ (go right)
4	6	5	$a[5] < 5$ (go right)
6	6	6	$a[6] == 5$ (go left)
6	5		Stop



pseudo code

```
int binary-search (int A[], int B) {  
    int l=0, R = A.length - 1, ans=-1;  
    while (l <= R) {  
        int mid = (l+R)/2;  
        ✓ If (A[mid] == B) { ans = mid; R = mid-1 }  
        else if (A[mid] < B) l = mid+1;  
        else R = mid-1;  
    }  
    return ans;  
}
```

✓
T.C = $O(\log n)$
S.C = $O(1)$

Quiz 3

When searching for first occurrence in sorted array

If $(\text{arr}[mid] == k)$ → store ans
proceed left



Every element occurs twice except for 1. Find that unique element.



Note : Duplicates are adjacent to each other and array isn't necessarily sorted.



XOR to get unique element

T.c = O(n) S.e = O(1)

9:55 → 10:05



Search Space : The entire array

Target : unique element

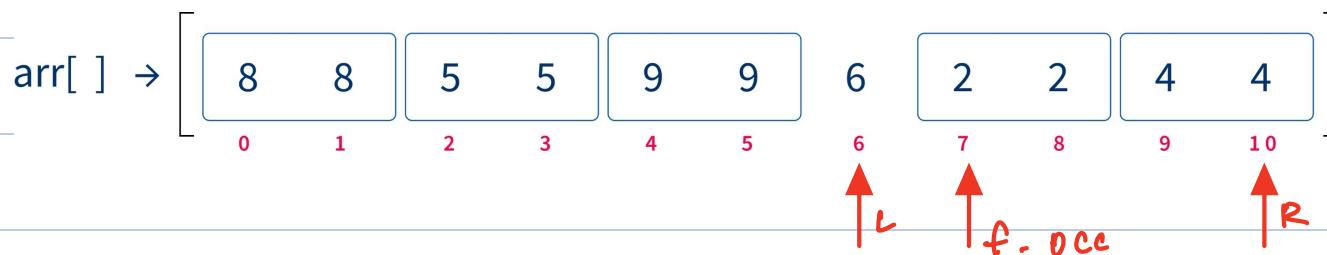
Condition :

First occurrence index

odd (move left)

even (move right)

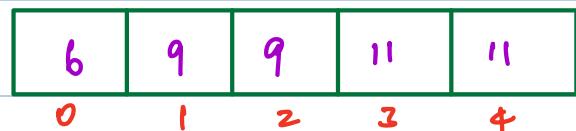
Here the concept is that the first_occ and my unique element will be on even index. When the first_occ began starting with odd index, my unique element would be on the left



L	R	MID	1st Occ.	Remark
0	10	5	4	$4 \cdot 1 \cdot 2 = 0$ f. occ + 2
6	10	8	7	$7 \cdot 1 \cdot 2 = 1$ f. occ - 1
6	6	6	—	ACB ↵

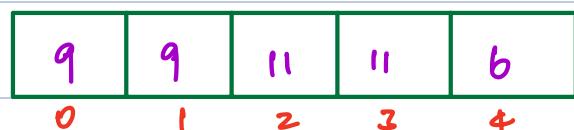
Edge cases

ans at 0th index



$\text{mid} == 0 \quad || \quad a[\text{mid} - 1].! = a[\text{mid}]$

ans at $n-1^{\text{th}}$ index



$\text{mid} == n-1 \quad || \quad a[\text{mid} + 1].! = a[\text{mid}]$

Combine to get unique element

Here the concept is that the first_occ and my unique element will be on even index. When the first_occ began starting with odd index, my unique element would be on the left

#pseudo code

Pnt unique_element (Pnt C A) {

Pnt L=0; R= A.length-1; N = A.length;

while (L <= R) { → Handle edge conditions

Pnt mid = (L+R) /2;

// correct ans conditions

If (C mid == 0 || (a[mid-1] != a[mid]))

&& (mid == N-1 || a[mid+1] != a[mid]))

return a[mid];

// Get the f-occ

Pnt f_occ = mid;

If (A[mid-1] == A[mid]) {

f_occ = mid-1;

// Basic conditions

If (f_occ % 2 == 0) { L = f_occ + 2; }

else { R = f_occ - 1; }

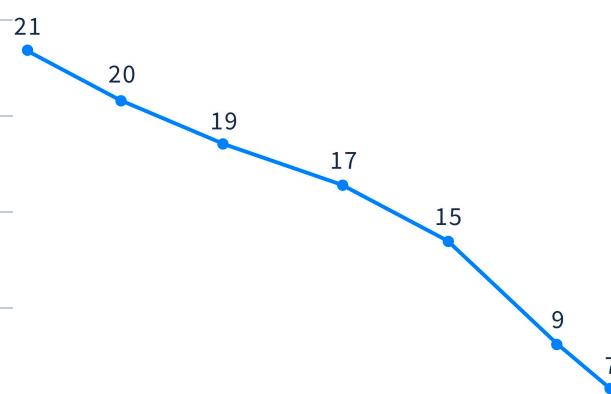
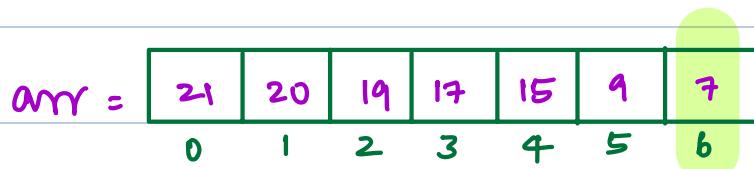
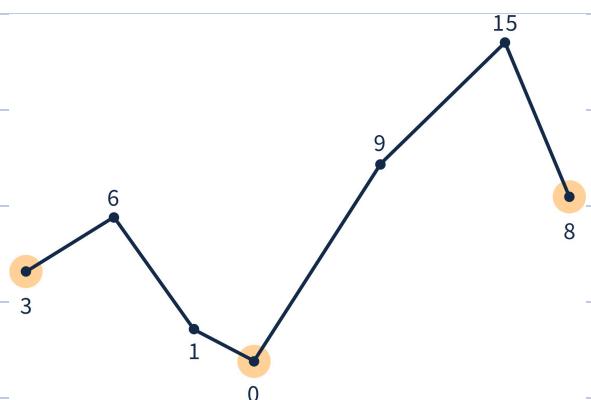
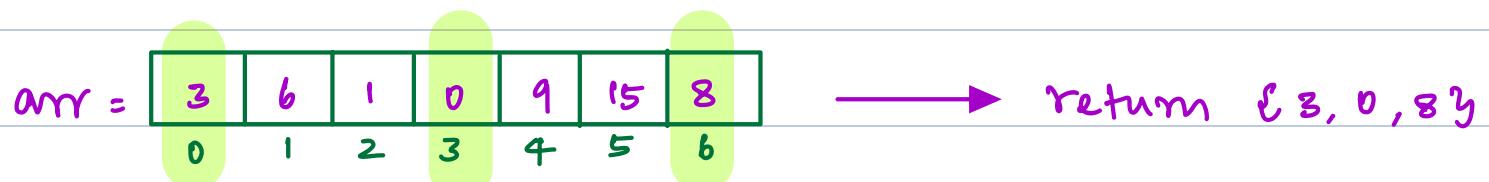
T.C = O(logn) S.C = O(1)



Given $\text{arr}[N]$. Find any one local minima.

Local Minima : An element which is less than or equals to it's adjacent elements.

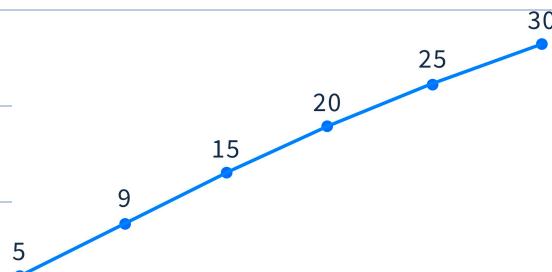
examples





arr =

5	9	15	20	25	30
0	1	2	3	4	5



arr =

5	8	12	3
0	1	2	3

any one of {5, 3}



BF Idea

Linear Search → $T.C = O(n)$ $S.C = O(1)$



Search Space: The entire array

Target: local minima

Condition: $a[\text{mid}] \leq a[\text{mid}-1] \text{ & } a[\text{mid}] \leq a[\text{mid}+1]$

$a[\text{mid}] \leq a[\text{mid}+1]$



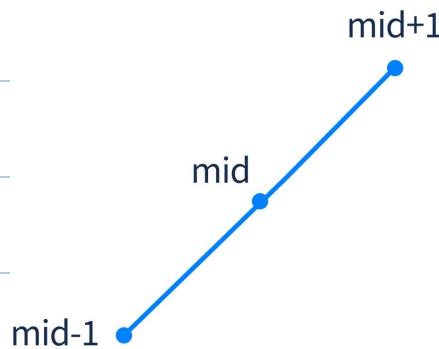
// True Condition

$a[\text{mid}] \leq a[\text{mid-1}] \&$
 $a[\text{mid}] \leq a[\text{mid+1}]$



// go Right

$a[\text{mid}] \leq a[\text{mid-1}] \&$
 $a[\text{mid}] > a[\text{mid+1}]$



// go Left

$a[\text{mid}] > a[\text{mid-1}] \&$
 $a[\text{mid}] \leq a[\text{mid+1}]$



Go either way



</> Code

Pnt localMinima (Pnt A) {

// only one element present

Pnt n = A.length;

If (n == 1) return A[0];

// local minima at 0th index

If (A[0] <= A[1]) return A[0];

// local minima at $n-1^{\text{th}}$ index

If (A[n-1] <= A[n-2]) return A[n-1];

Pnt L = 0; R = N - 1;

while (L <= R) {

Pnt mid = (L + R) / 2;

If (A[mid] <= A[mid - 1]) &&

A[mid] <= A[mid + 1]) {

return A[mid];

else If (A[mid] > A[mid + 1]) L = mid + 1;

else R = mid - 1;

T.C = $O(\log n)$ S.C = $O(1)$

