

Binary Search

Criteria

1. **Search Space:** Range in which we are performing our search
2. **Target:** Key of search
3. **Condition:** For answer; To eliminate left; To eliminate right

1.Problem Statement – Search in rotated array

Given a sorted array of integers **A** of size **N** and an integer **B**, where array **A** is rotated at **some pivot** unknown beforehand.

For example, the array [0, 1, 2, 4, 5, 6, 7] might become [4, 5, 6, 7, 0, 1, 2].

Your task is to search for the target value **B** in the array. If **found**, return its **index**; **otherwise**, return **-1**.

You can assume that **no duplicates** exist in the array.

NOTE: You are expected to solve this problem with a time complexity of **$O(\log(N))$** .

Problem Constraints

$1 \leq N \leq 1000000$

$1 \leq A[i] \leq 10^9$

All elements in **A** are **Distinct**.

Output Format: Return index of **B** in array **A**, otherwise return **-1**

Example Input

A = [4, 5, 6, 7, 0, 1, 2, 3]

B = 4

Example Output: 0

Example Explanation: Target 4 is found at index 0 in A.

A = [1, 2, 3, 4, 5, 6, 7, 8]

Rotated Array

A = [2, 3, 4, 5, 6, 7, 8, 1]

A = [3, 4, 5, 6, 7, 8, 1, 2]

A=[4, 5, 6, 7, 8, 1, 2, 3]

Brute force Idea

Do a linear search

TC = $O(N)$

SC = $O(1)$

Idea -2

Observation 1:

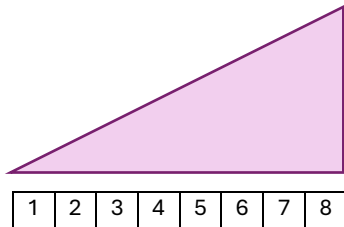
A=[4, 5, 6, 7, 8, 1, 2, 3]

In rotated sorted array, we have 2 sorted subarray

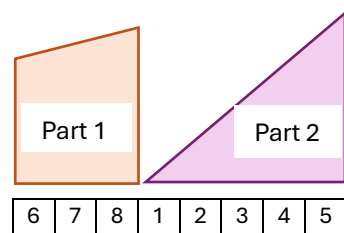
Observation 2:

1. How can we identify if $A[i]$ is part of P1 or P2?

Compare with $A[0]$, if $(A[i] \geq A[0])$ then it belongs to P1 else It belongs to P2.



2. If any element = 8, if $A[0] < 8$, then it will surely come in P1 and not in P2.



Binary Search

1. Search Space: The entire array
2. Target: Key
3. Condition: $mid = (l+r)/2$

Check if $A[mid]$ and key are in same part → If, Yes: Apply binary search

→ If No: Move to the actual part by changing mid value

Actual Code:

```
public class Solution {
    public int search(final int[] A, int B) {
        boolean targetinP1 = true;
        if(B < A[0]) targetinP1 = false;
        int l = 0, r = A.length-1;
        while(l <= r){
            int mid = (l+r)/2;
            //identify part of mid
            boolean midinP1 = true;
            if(A[mid] < A[0]) midinP1 = false;
            //if both are in same part
            if(midinP1 == targetinP1){
                if(A[mid] == B) return mid;
                else if(A[mid] < B) l = mid+1;
                else r = mid-1;
            }
            else if(midinP1){
                //if mid is in P1 but target is in P2 → move my mid to P2.
                l = mid+1;
            }
            else{
                r = mid-1;
            }
        }
    }
}
```

```

        r = mid-1;
    }
}
return -1;
}

```

Dry Run:

6	7	8	1	2	3	4	5
0	1	2	3	4	5	6	7

K = 11

targetinFirstHalf = True

L	R	Mid	A[mid]	A[mid] part	Both k and A[mid] -- Same part	Condition
0	7	$(0+7)/2 = 3$	A[3] = 1	$1 < 6$; Part 2	No since $11 > 6$	R = mid-1
0	2	$(0+2)/2 = 1$	A[1] = 7	$7 \geq 6$; Part 1	Yes	L = mid + 1
2	2	$(2+2)/2 = 2$	A[2] = 8	$8 \geq 6$; Part 1	Yes	L = mid + 1
3	2	$(3+2)/2 = 2$	-----	-----	-----	Return -1

Stop: since $l > r \rightarrow$ return -1

k = 5

targetinFirstHalf = False as $5 < A[0]$

L	R	Mid	A[mid]	A[mid] part	Same part	Condition
0	7	$(0+7)/2 = 3$	A[3] = 1	$1 < 6$; Part 2	Yes	So basic condition will be binary search L = mid+1
4	7	$(4+7)/2 = 5$	A[5] = 3	$3 < 6$; Part 2	Yes	Since $3 < 5$ eliminate left side; L = mid + 1
6	7	$(6+7)/2 = 6$	A[6] = 4	$4 < 6$; Part 2	Yes	L = mid + 1
7	7	$(7+7)/2 = 7$	-----	-----	-----	Return -1

2.Problem Statement – Square Root of Integer

Given an integer A. Compute and return the square root of A. If A is not a perfect square, return floor(sqrt(A)).

NOTE:

The value of A*A can cross the range of Integer.

Do not use the sqrt function from the standard library.

Users are expected to solve this in $O(\log(A))$ time.

Output Format: Return floor(sqrt(A))

Example Input: 11

Example Output: 3

Example Explanation

When $A = 11$, square root of $A = 3.316$. It is not a perfect square so we return the floor which is 3.

N	Floor(sqrt(N))
9	sqrt(9) = 3
12	sqrt(12) = 3
16	sqrt(16) = 4
24	sqrt(24) = 4

Brute Force(n =50)

i	i * i	ans
1	1*1 = 1	1
2	2*2 = 4	2
3	3*3 = 9 (9<50)	3 → potential answer
4	4*4 = 16	4
..
..
7	7*7 = 49	7 → ans
8	8*8 = 64	since (64 > 50) → stop

Brute force code:

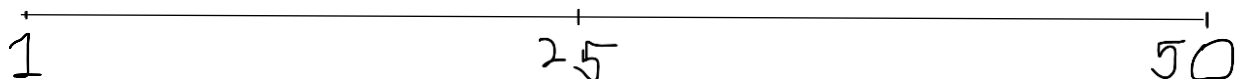
```
int Floorsqrt(int N){
    int ans = 1;
    for(int i =1; i*i <=N; i++){
        ans = i;
    }
    return ans;
}
```

TC = $O(\sqrt{N})$

SC = $O(1)$

Binary Search:

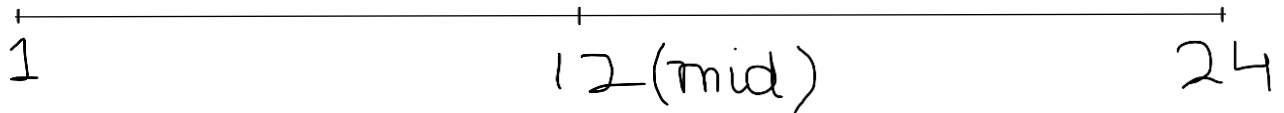
1. **Search Space:** smallest ans =1; largest (possible) ans = N
2. **Target:** floor(sqrt(N))
3. **Condition:** if($\text{mid} * \text{mid} \leq N$) $L = \text{mid} + 1$; //potential ans
else { $r = \text{mid} - 1$ }



for $n = 50$

$\text{mid} = 25$, $25 * 25 = 625$

Since $625 > 50$, discard the elements on the right.



mid = 12, $12 * 12 = 144$

Since $144 > 50$, discard the elements on the right.

mid = 5, $5 * 5 = 25$

Since, $25 < 50$, so '5' could be my potential ans right now but will look for better answer so $l = \text{mid} + 1$.

Actual Code

```
public class Solution {
    public int sqrt(int A) {
        long l = 1;
        long r = A, ans = 0;
        while(l <= r){
            long mid = l + (r - l) / 2;
            if(mid * mid <= A) { //potential ans
                ans = mid;
                l = mid + 1;
            }
            else{
                r = mid - 1;
            }
        }
        return (int)ans;
    }
}
```

Dry Run

n = 50, ans = ~~6~~ 7

L	R	Mid	mid*mid	where next?
1	50	25	$25 * 25 = 625$	$625 > 50$; r = mid-1
1	24	12	$12 * 12 = 144$	$144 > 50$; r = mid-1
1	11	6	$6 * 6 = 36$	$36 < 50$; l = mid + 1; one of my potential ans
7	11	9	$9 * 9 = 81$	$81 > 50$; r = mid-1
7	8	7	$7 * 7 = 49$	$49 < 50$; next potential ans
8	8	8	$8 * 8 = 64$	$64 > 50$; r = mid-1
8	7	-----	-----	return ans = 7

3.Problem Statement – Median of two sorted array

Given two sorted arrays A and B of size M and N respectively, return the median of the two sorted arrays.
Round of the value to the floor integer [$2.6=2$, $2.2=2$]

Output Format: Return an integer.

Example Input:

A = [1, 3]

B = [2]

Example Output: 3

Median → is middle element in a sorted array

A =

10	20	30	40	45	46	50	60	70	80
0	1	2	3	4	5	6	7	8	9

 Median = (45+46)/2

Median of two sorted arrays → Odd elements

A =

1	4	5
---	---	---

 C =

1	2	3	4	5
---	---	---	---	---

B =

2	3
---	---

Brute Force

1. Merge into a single sorted array
2. Find the median array
3. return median

TC = O(N+M)**SC = O(N+M)****Idea -2 : Binary Search**

1. **Search Space:** Apply binary search in a combined array – 2 arrays A & B
2. **Target:** Median of A & B
3. **Condition:** if(L1 <= R2 && L2 <= R1) → ans
 else if (L2 > R1) l = mid + 1 → eliminate left
 else r = mid - 1 → eliminate right

A =

10	30	45	46	60	80
----	----	----	----	----	----

B =

20	40	50	70
----	----	----	----

C =

10	20	30	40	45	46	50	60	70	80
----	----	----	----	----	----	----	----	----	----

Observation1: Both parts have some elements from Array A & some elements from Array B**Observation2:** All elements in Part1 (P1) < All elements in Part2(P2)**TRIAL & ERROR**

A =

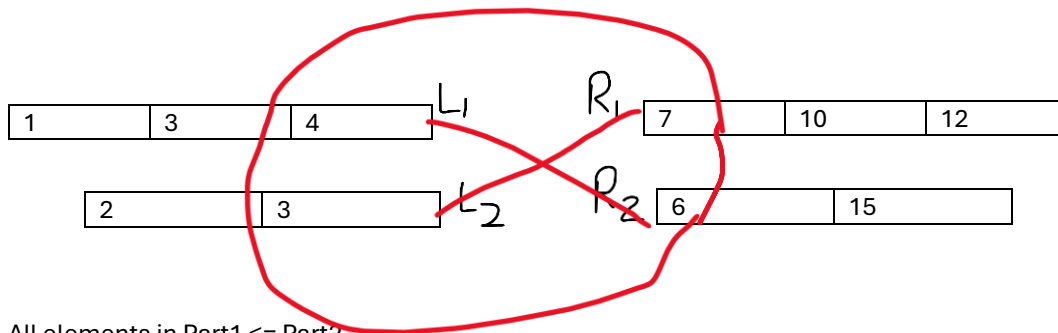
1	3	4	7	10	12
---	---	---	---	----	----

B =

2	3	6	15
---	---	---	----

Try 1:

Let me pick 3 elements from A in Part1



All elements in Part1 \leq Part2

Here $L1 \leq R2$ and $L2 \leq R1$

Part1:

1	2	3	3	4
---	---	---	---	---

Part2:

6	7	10	12	15
---	---	----	----	----

Median = $(\text{Max}(L1, L2) + \text{Min}(R1, R2))/2$ in case of even elements.

Total elements in Part1 = $(N+M)/2$

Dry Run:

A =

1	2	3	4	9	11
---	---	---	---	---	----

0 1 2 3 4 5

B =

7	12	14	15
---	----	----	----

0 1 2 3

Total elements in Part1 = $(N+M)/2 = (6+4)/2 = 5$

L	R	mid	L1	R1	L2	R2	is valid?
0	5	$0+5/2 = 2$ (this is the total #of elements taken in P1)	$A[1] = 2$	3	14	15	Since $L2 > R1$; $L = \text{mid}1+1$
3	5	$3+5/2 = 4$	4	9 since $L2 < R1$ & $L1 < R2$ so this is a good split	7	12	$4 \leq 12$; $7 \leq 9$

Elements taken from A in Part1 = 2 (mid element)

$L1 = A[\text{mid}1-1]$

$R1 = A[\text{mid}1]$

$L2 = B[\text{mid}2-1]$

$R2 = B[\text{mid}2]$

$\text{mid}2 = \text{total elements in PartA} - \text{mid} \rightarrow 5 - 2 = 3$

$\text{ans} = (\text{Max}(L1, L2) + \text{Min}(R1, R2))/2$

$$\text{Max}(4,7) + \text{Min}(9, 12)/2 = 7+9 /2 = 8$$

If $L2 < R1 \rightarrow$ select more elements from A in Part1

Note: Time Complexity depends on length of array A. So among both array, whosoever would have less length, I could take lesser length array as array A.

How to handle odd length array

Total elements in P1 = $(n+m+1)/2$  Take one extra element in Part 1

Median = $\text{Max}\{L1, L2\}$

Edge Cases

if (mid1 == 0) \rightarrow L1 will go out of bound L1 = - infinity	if (mid1 == n) \rightarrow R1 will go out of bound R1 = + infinity
if (mid2 == 0) L2 = -infinity	if (mid2 == m) R2 = + infinity

Actual Code

```
public class Solution {
    public int solve(int[] A, int[] B) {
        //always consider smaller array as A
        if(A.length > B.length){
            int [] temp = A;
            A = B;
            B = temp;
        }
        //initialize all values
        int n = A.length, m = B.length;
        int totalinPartA = (m+n+1)/2;
        int l = 0, r = n;
        double median = 0;
        //Binary Search
        while(l <= r){
            int mid1 = (l+r)/2;
            int mid2 = totalinPartA - mid1;
            //initialize L1, L2, R1, R2

            int L1 = (mid1 > 0) ? A[mid1 - 1]: Integer.MIN_VALUE;
            int R1 = (mid1 < n) ? A[mid1]: Integer.MAX_VALUE;
            int L2 = (mid2 > 0) ? B[mid2 - 1]: Integer.MIN_VALUE;
            int R2 = (mid2 < m) ? B[mid2]: Integer.MAX_VALUE;

            //condition for ans
            if(L1 <= R2 && L2 <= R1){
                if((n+m) % 2 == 0){
                    median = (Math.max(L1, L2) + Math.min(R1, R2))/2;
                    return (int) median;
                }else{
                    median = Math.max(L1, L2);
                    return (int) median;
                }
            }
            if(L1 < R2){
                l = mid1 + 1;
            }else{
                r = mid1 - 1;
            }
        }
    }
}
```



```

    }
}
else if (L2 > R1){
    l = mid1 + 1;
}else{
    r = mid1 - 1;
}
}
return (int) median;
}
}

```

Problem Statement 4 – Matrix Median

Problem Description

Given a matrix of integers **A** of size $N \times M$ in which each row is sorted.

Find and return the overall median of matrix A.

NOTE: No extra memory is allowed.

NOTE: Rows are numbered from top to bottom and columns are numbered from left to right.

Problem Constraints

$1 \leq N, M \leq 10^5$

$1 \leq N \times M \leq 10^6$

$1 \leq A[i] \leq 10^9$

$N \times M$ is odd

Input Format

The first and only argument given is the integer matrix A.

Output Format

Return the overall median of matrix A.

Example Input

Input 1:

```

A = [ [1, 3, 5],
      [2, 6, 9],
      [3, 6, 9] ]

```

Input 2:

```

A = [ [5, 17, 100] ]

```

Example Output

Output 1:

5

Output 2:

17

Example Explanation

Explanation 1:

$A = [1, 2, 3, 3, 5, 6, 6, 9, 9]$

Median is 5. So, we return 5.

Explanation 2:

Median is 17.

Actual Code

```
public class Solution {
    public int findMedian(ArrayList<ArrayList<Integer>> A) {
        int l = Integer.MAX_VALUE, r = Integer.MIN_VALUE, ans = 0;
        int n = A.size();
        int m = A.get(0).size();
        //int count = 0;
        for(int i = 0; i < n; i++){
            l = Math.min(l, A.get(i).get(0));
            r = Math.max(r, A.get(i).get(m-1));
        }
        int req_cnt = (n*m + 1)/2; // required count of numbers less than the median // median would
        be req_cnt+1th number
        while(l < r){
            int mid = l + (r-l)/2;
            int count = 0;
            for (int i = 0; i < n; i++){
                count += countlessthanorequalto(A.get(i), mid);
            }
            if(count < req_cnt){
                l = mid+1; // Median is larger
            }else{
                r = mid; // Median is smaller or equal
            }
        }
        return l;
    }
    private int countlessthanorequalto(ArrayList<Integer> row, int mid){
        int l = 0, r = row.size() - 1;
        while(l <= r){
            int mid2 = l + (r-l)/2;
            if(row.get(mid2) <= mid){
                l = mid2 + 1;
            }else{
                r = mid2 - 1;
            }
        }
        return l;
    }
}
```

Problem Statement 5 - Ath Magical Number

Problem Description

You are given three positive integers, **A**, **B**, and **C**.

Any positive integer is magical if divisible by either **B** or **C**.

Return the **Ath** smallest magical number. Since the answer may be very large, return modulo **$10^9 + 7$** .

Note: Ensure to prevent **integer overflow** while calculating.

Problem Constraints

$1 \leq A \leq 10^9$

$2 \leq B, C \leq 40000$

Input Format

The first argument given is an integer **A**.

The second argument given is an integer **B**.

The third argument given is an integer **C**.

Output Format

Return the **Ath** smallest magical number. Since the answer may be very large, return modulo **$10^9 + 7$** .

Example Input

Input 1:

A = 1

B = 2

C = 3

Input 2:

A = 4

B = 2

C = 3

Example Output

Output 1: 2

Output 2: 6

Example Explanation

Explanation 1:

1st magical number is 2.

Explanation 2:

First four magical numbers are 2, 3, 4, 6 so the 4th magical number is 6.

Actual Code

```
public class Solution {  
    public int solve(int A, int B, int C) {  
        long l = 1, ans = 0;  
        long r = (long) A * Math.min(B, C);  
        long mod = 1000000007L;  
    }  
}
```

```

long gcd = findgcd(B, C);
long lcm = (long) B * (long) C / gcd;
while(l <= r){
    long mid = l + (r-l)/2;
    long count = isPossible(mid, B, C, lcm);
    if(count >= A){
        ans = mid;
        r = mid-1;
    }else{
        l = mid+1;
    }
}
return (int) (ans % mod);
}

private long isPossible(long mid, long B, long C, long lcm){ //count magical numbers
    return mid/B + mid/C - mid/lcm;
}

private long findgcd(long a, long b){
    if (b==0) return a;
    return findgcd(b, a % b);
}
}

```

$T.C. :- O(\log(\min(B, C) \times A))$
 $S.C. :- O(1)$

Problem Statement 6 - Find Smallest Again

Problem Description

Given an integer array **A** of size **N**.

If we store the sum of each triplet of the array **A** in a new list, then find the **Bth** smallest element among the list.

NOTE: A triplet consists of three elements from the array. Let's say if **A[i]**, **A[j]**, **A[k]** are the elements of the triplet then **i < j < k**.

Problem Constraints

$3 \leq N \leq 500$

$1 \leq A[i] \leq 10^8$

$1 \leq B \leq (N*(N-1)*(N-2))/6$

Input Format

The first argument is an integer array A.

The second argument is an integer B.

Output Format

Return an integer denoting the B^{th} element of the list.

Example Input

Input 1:

A = [2, 4, 3, 2]

B = 3

Input 2:

A = [1, 5, 7, 3, 2]

B = 9

Example Output

Output 1:

9

Output 2:

14

Example Explanation

Explanation 1:

All the triplets of the array A are:

(2, 4, 3) = 9

(2, 4, 2) = 8

(2, 3, 2) = 7

(4, 3, 2) = 9

[7,8,9,9]

So the 3rd smallest element is 9.

Actual Code

```
import java.util.*;

public class Solution {
    public int solve(ArrayList<Integer> A, int B) {
        Collections.sort(A); // Step 1: Sort the array
        int n = A.size();
        long low = (long) A.get(0) + A.get(1) + A.get(2);
        long high = (long) A.get(n - 1) + A.get(n - 2) + A.get(n - 3);
        int ans = 0;

        while (low <= high) {
```

```

        long mid = low + (high - low) / 2;
        long count = countTriplets(A, mid); // Step 3: Count triplets

        if (count >= B) {
            ans = (int) mid; // Update answer
            high = mid - 1; // Search for smaller values
        } else {
            low = mid + 1; // Search for larger values
        }
    }

    return ans;
}

private long countTriplets(ArrayList<Integer> A, long mid) {
    int n = A.size();
    long count = 0;

    for (int i = 0; i < n - 2; i++) {
        int j = i + 1, k = n - 1;

        while (j < k) {
            long sum = (long) A.get(i) + A.get(j) + A.get(k);
            if (sum <= mid) {
                count += (k - j); // All pairs from j to k are valid
                j++;
            } else {
                k--; // Decrease the upper pointer
            }
        }
    }

    return count;
}
}

```

Problem Statement 7 - ADD OR NOT - **Unsolved**

Problem Description

Given an array of integers **A** of size **N** and an integer **B**.

In a single operation, any one element of the array can be increased by 1. You are allowed to do at most **B** such operations.

Find the number with the **maximum** number of occurrences and return an array **C** of size 2, where **C[0]** is the number of occurrences, and **C[1]** is the number with maximum occurrence. If there are several such numbers, your task is to find the **minimum** one.

Problem Constraints

$1 \leq N \leq 10^5$

$-10^9 \leq A[i] \leq 10^9$

$0 \leq B \leq 10^9$

Input Format

The first argument given is the integer array A.

The second argument given is the integer B.

Output Format

Return an array C of size 2, where C[0] is number of occurrence and C[1] is the number with maximum occurrence.

Example Input

Input 1:

A = [3, 1, 2, 2, 1]

B = 3

Input 2:

A = [5, 5, 5]

B = 3

Example Output

Output 1: [4, 2]

Output 2: [3, 5]

Example Explanation

Explanation 1:

Apply operations on A[2] and A[4]

A = [3, 2, 2, 2, 2]

Maximum occurrence = 4

Minimum value of element with maximum occurrence = 2

Explanation 2:

A = [5, 5, 5]

Maximum occurrence = 3

Minimum value of element with maximum occurrence = 5

Actual Code

```
import java.util.*;

public class Solution {
    public ArrayList<Integer> solve(ArrayList<Integer> A, int B) {
        Collections.sort(A); // Sort the array
        int n = A.size();
        int maxOccurrences = 1;
        int minValue = A.get(0);

        int left = 0, right = 0, totalCost = 0;
```

```

// Sliding window to calculate max occurrences
while (right < n) {
    // Calculate the cost of making A[left...right] equal to A[right]
    totalCost += (right - left) * (A.get(right) - A.get(right - 1));

    // If cost exceeds B, move the left pointer
    while (totalCost > B && left < right) {
        totalCost -= (A.get(right) - A.get(left));
        left++;
    }

    // Update the maximum occurrences and the minimum value
    int currentOccurrences = right - left + 1;
    if (currentOccurrences > maxOccurrences) {
        maxOccurrences = currentOccurrences;
        minValue = A.get(right);
    } else if (currentOccurrences == maxOccurrences) {
        minValue = Math.min(minValue, A.get(right));
    }

    right++;
}

ArrayList<Integer> result = new ArrayList<>();
result.add(maxOccurrences);
result.add(minValue);
return result;
}
}

```