

# DATA STRUCTURES & ALGORITHMS

## 11: GRAPH SEARCHING - BFS

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# GRAPH REPRESENTATIONS

# GRAPH SEARCHING - BFS

Given graph  $\mathbf{G} = (\mathbf{V}, \mathbf{E})$ .

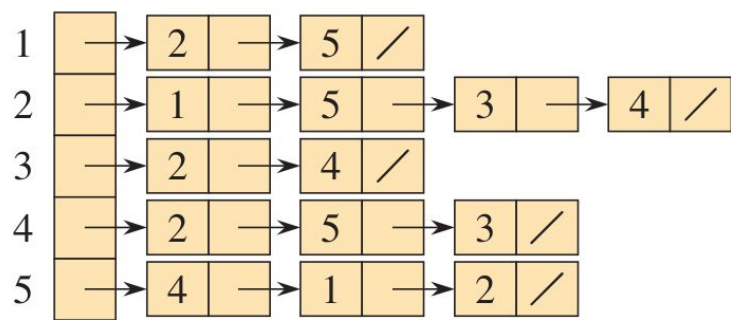
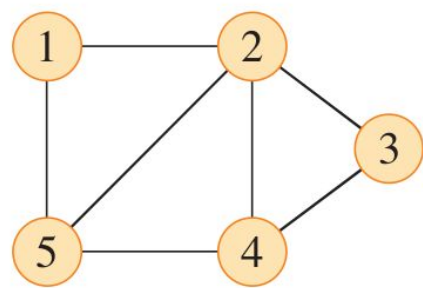
In pseudocode, represent vertex set by  $\mathbf{G.V}$  and edge set by  $\mathbf{G.E}$ .

G may be either **directed** or **undirected**.

Two common ways to represent graphs for algorithms:

1. Adjacency Lists.
2. Adjacency Matrix.

# GRAPH SEARCHING - BFS

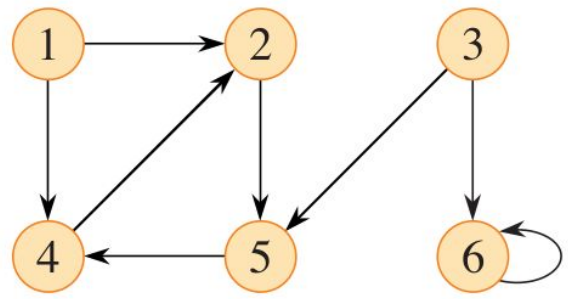


Adjacency List

	1	2	3	4	5
1	0	1	0	0	1
2	1	0	1	1	1
3	0	1	0	1	0
4	0	1	1	0	1
5	1	1	0	1	0

Adjacency Matrix

# GRAPH SEARCHING - BFS



1	→	2	→	4	/
2	→	5	/		
3	→	6	→	5	/
4	→	2	/		
5	→	4	/		
6	→	6	/		

Adjacency List

	1	2	3	4	5	6
1	0	1	0	1	0	0
2	0	0	0	0	1	0
3	0	0	0	0	1	1
4	0	1	0	0	0	0
5	0	0	0	1	0	0
6	0	0	0	0	0	1

Adjacency Matrix

# GRAPH SEARCHING - BFS

## Adjacency lists

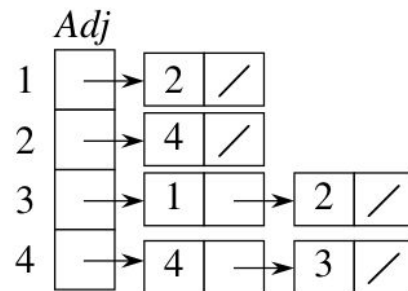
Array ***Adj*** of  $|V|$  lists, one per *vertex*.

Vertex ***u***'s list has all vertices ***v*** such that  $(u, v) \in E$ .

(Works for both directed and undirected graphs.)

In pseudocode, denote the array as attribute ***G.Adj***.

We will see notation such as ***G.Adj[u]***



## Adjacency matrix

$|V| \times |V|$  matrix  $A = (a_{ij})$

$$a_{ij} = \begin{cases} 1 & \text{if } (i, j) \in E, \\ 0 & \text{otherwise.} \end{cases}$$

	1	2	3	4
1	0	1	0	0
2	0	0	0	1
3	1	1	0	0
4	0	0	1	1

# GRAPH SEARCHING

# GRAPH SEARCHING - BFS

## Graph Searching

Searching a graph means systematically following the edges of the graph so as to visit the vertices of the graph.

*Graph searching algorithms discovers the structure of the graph.*

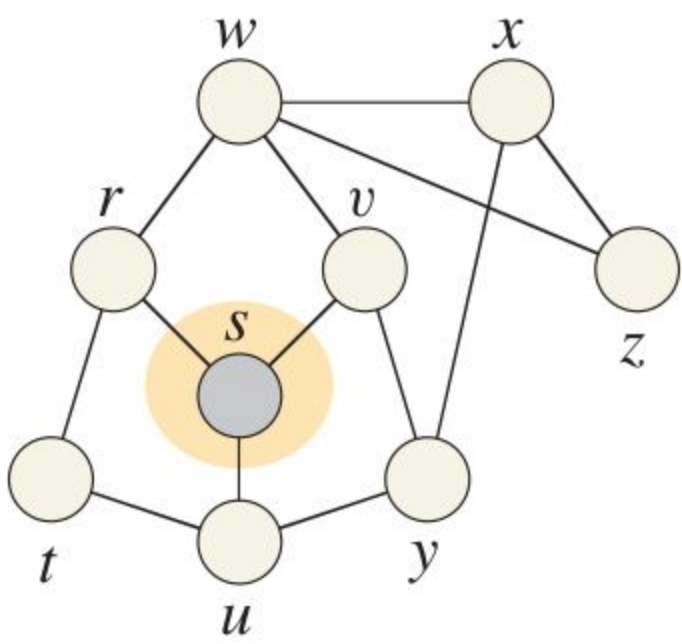
## Two prominent Graph Searching

- Breadth First Search(BFS)
- Depth First Search (DFS)



# BREADTH FIRST SEARCH (BFS)

# GRAPH SEARCHING - BFS



# GRAPH SEARCHING - BFS

## Breadth First Search (BFS)

BFS, from a given source vertex  $s$ , finds the shortest simple path to vertices.

$\Rightarrow$  discovers every vertex  $v$  reachable from  $s$  through a *shortest path*.

$\rightarrow$  **shortest path** is the path containing *smallest number of edges*.

BFS on a graph with source vertex  $s$ , generate a **Breadth First Tree** with  $s$  as root.

**Breadth First**  $\Rightarrow$  A wave (also mentioned as frontiers) emanating from source  $s$ , visiting vertices at **distance 1**, then **distance 2** etc until it has discovered every vertex reachable from  $s$ .

# GRAPH SEARCHING - BFS

## Breadth First Search (BFS)

To keep track of the wave, BFS uses a Queue.

⇒ Queue contains **some vertices** at distance **k** and,  
possibly **some vertices** at distance **k+1**.

Queue contains vertices belonging to *portions* of **two consecutive waves** at any time.

# GRAPH SEARCHING - BFS

## Breadth First Search (BFS)

### Colouring scheme in BFS:

BFS colors vertices **WHITE**, **GRAY** and **BLACK**.

- **WHITE** → not visited, and not reachable from  $s$
- **GRAY** → discovered first time a vertex is reachable from  $s$ . (*contained in queue*)
- **BLACK** → once all of the vertex's edges are explored.

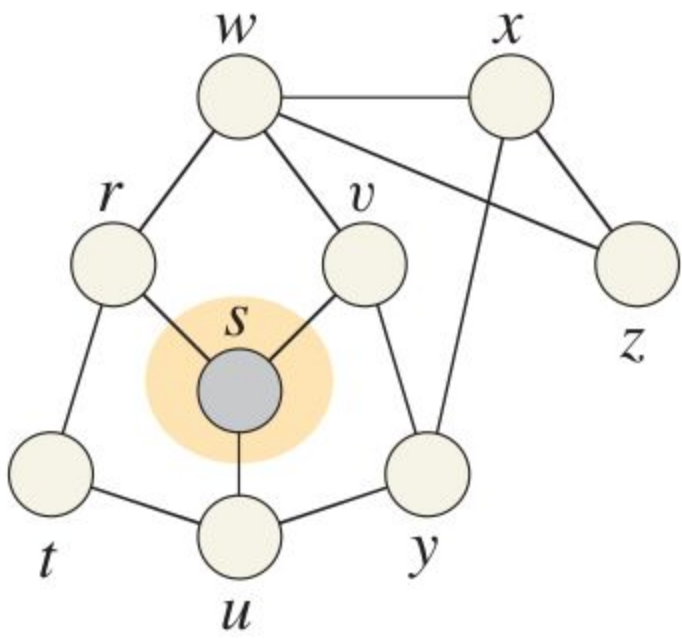
# GRAPH SEARCHING - BFS

## Breadth First Search (BFS)

**Additional attributes in to each vertex  $v$ :**

- $v.color$  is the color of  $v$ : WHITE, GRAY, or BLACK.
- $v.d$  holds the distance from the source vertex  $s$  to  $v$ , as computed by the algorithm.
- $v.\pi$  is  $v$ 's predecessor in the breadth-first tree. If  $v$  has no predecessor because it is the source vertex or is undiscovered, then  $v.\pi = \text{NIL}$ .

# GRAPH SEARCHING - BFS



# GRAPH SEARCHING - BFS

**BFS**( $G, s$ )

**for** each vertex  $u \in G.V - \{s\}$

$u.color = \text{WHITE}$

$u.d = \infty$

$u.\pi = \text{NIL}$

$s.color = \text{GRAY}$

$s.d = 0$

$s.\pi = \text{NIL}$

$Q = \emptyset$

**ENQUEUE**( $Q, s$ )

**while**  $Q \neq \emptyset$

$u = \text{DEQUEUE}(Q)$

**for** each vertex  $v$  in  $G.Adj[u]$

**if**  $v.color == \text{WHITE}$

$v.color = \text{GRAY}$

$v.d = u.d + 1$

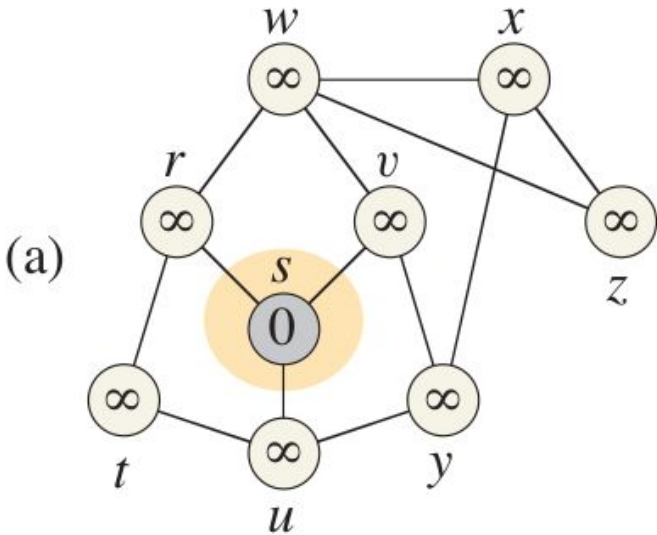
$v.\pi = u$

**ENQUEUE**( $Q, v$ )

$u.color = \text{BLACK}$



# GRAPH SEARCHING - BFS



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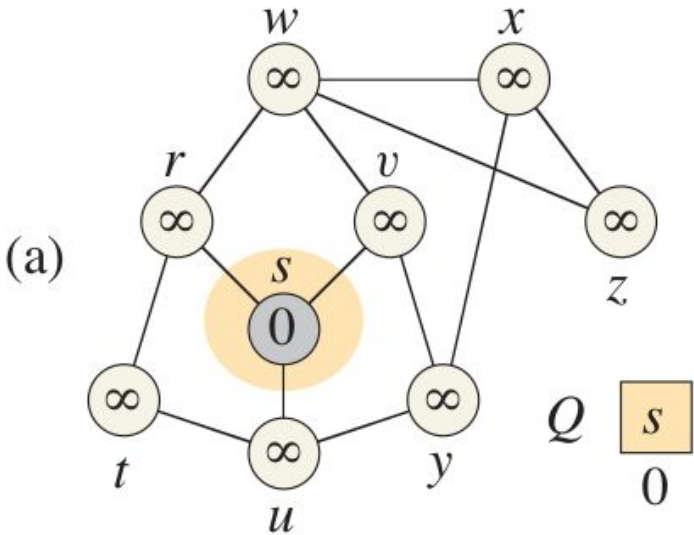
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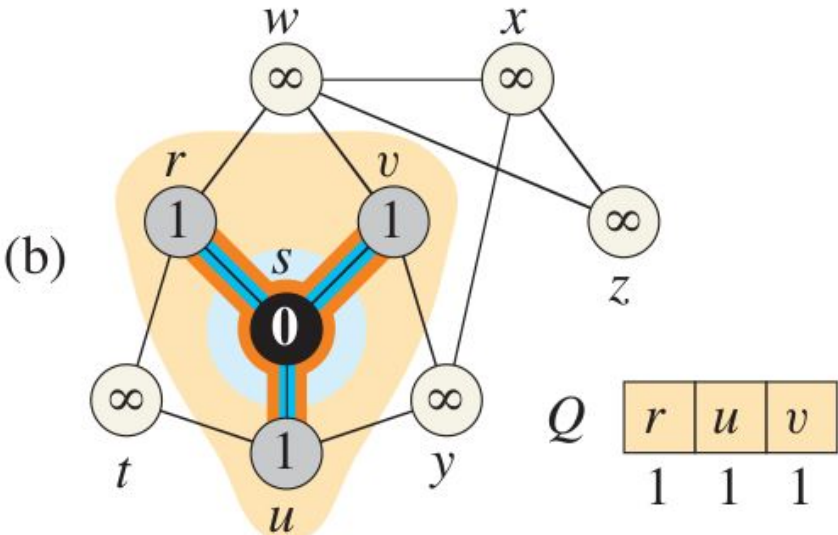
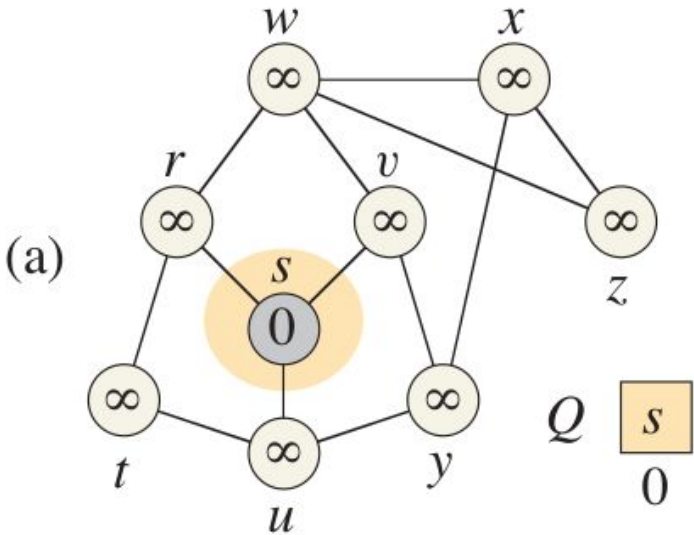
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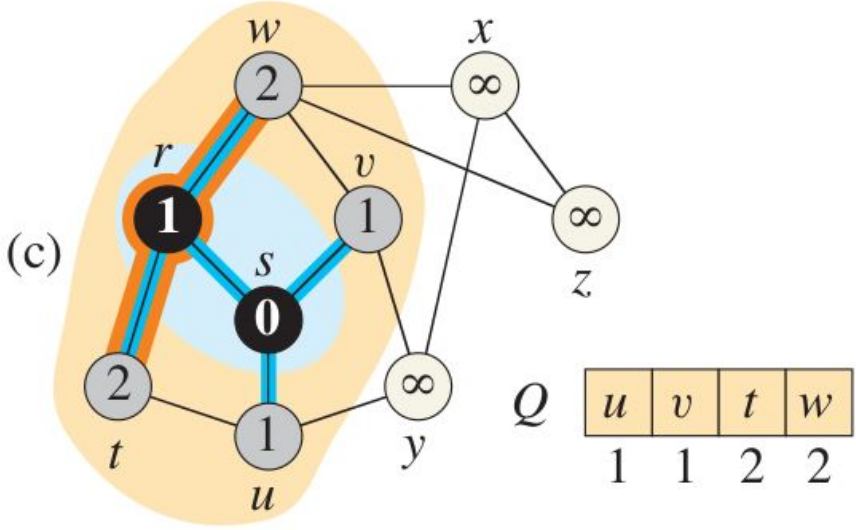
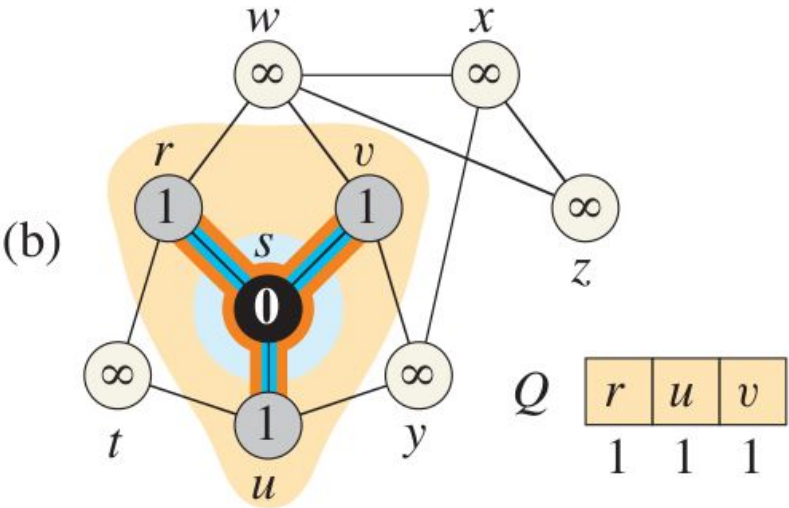
            ENQUEUE( $Q, v$ )

$u.color = \text{BLACK}$

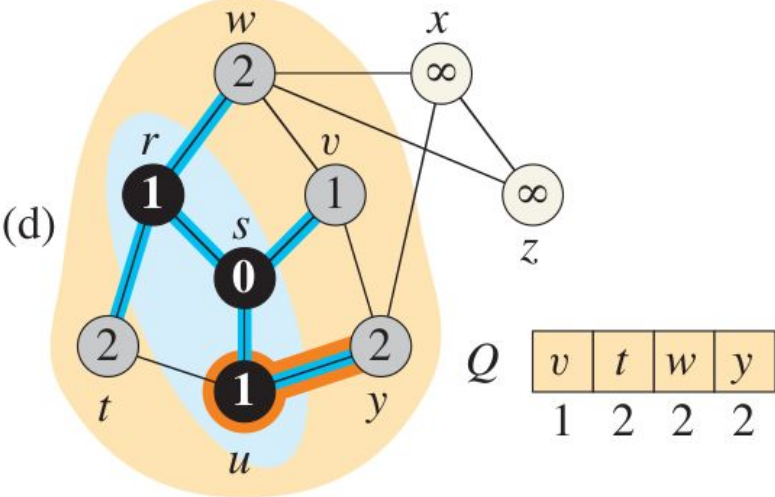
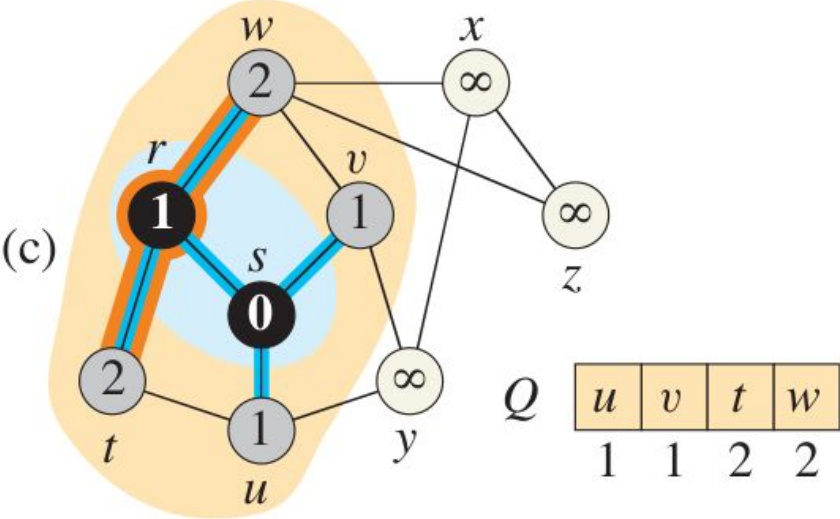
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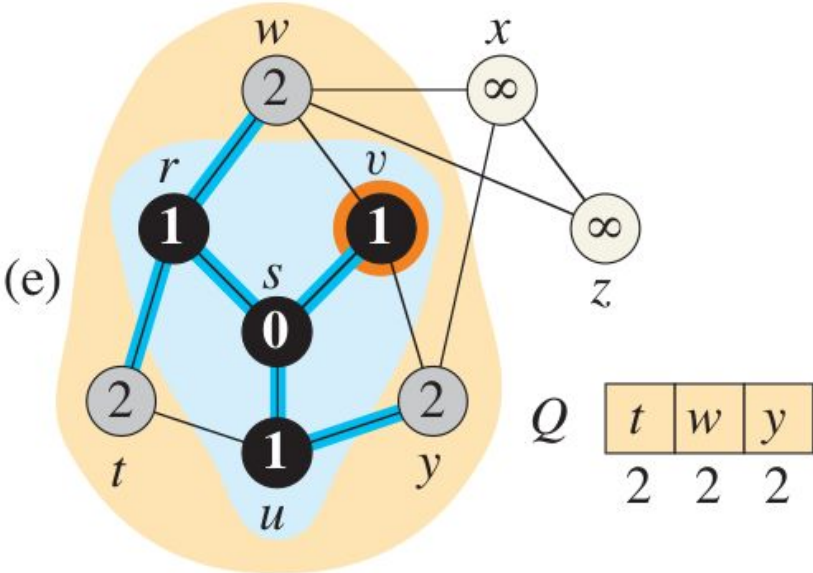
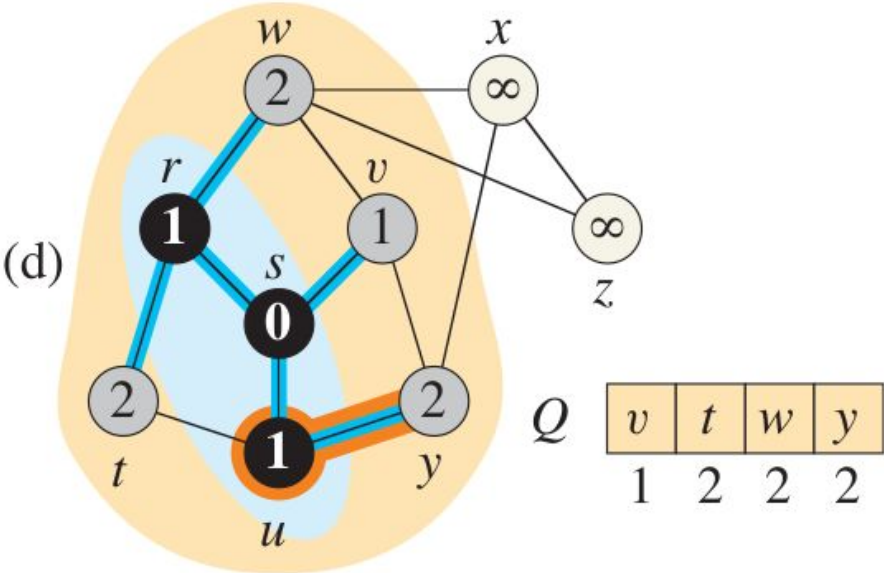
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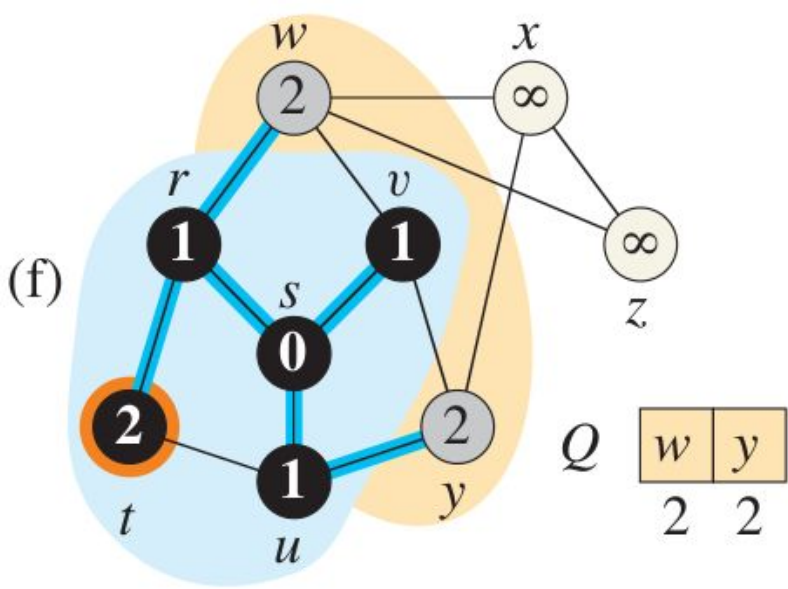
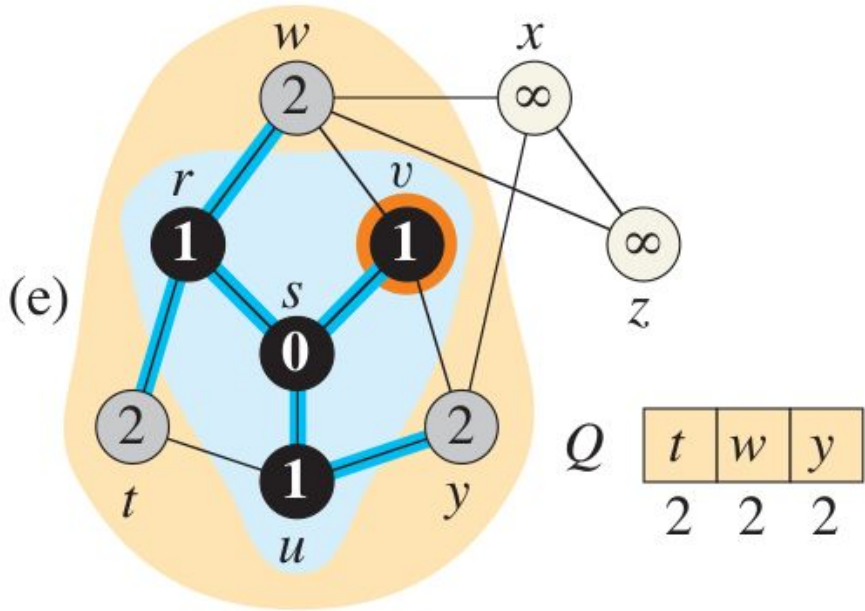


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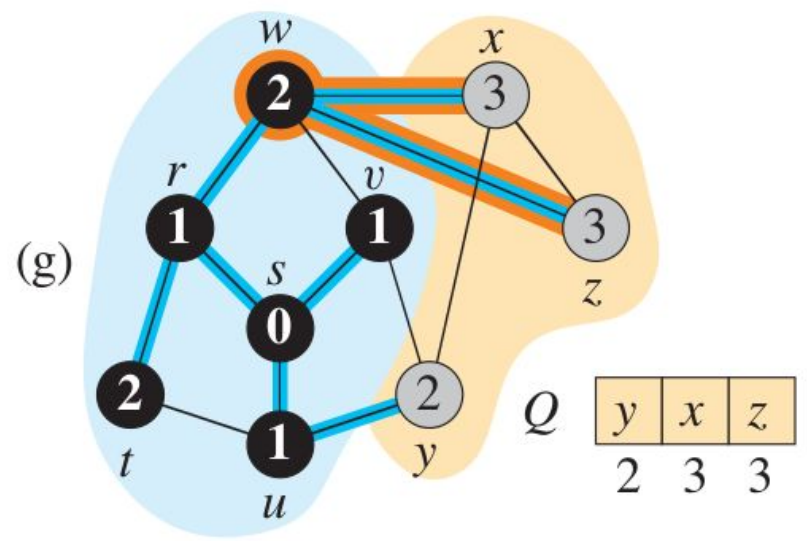
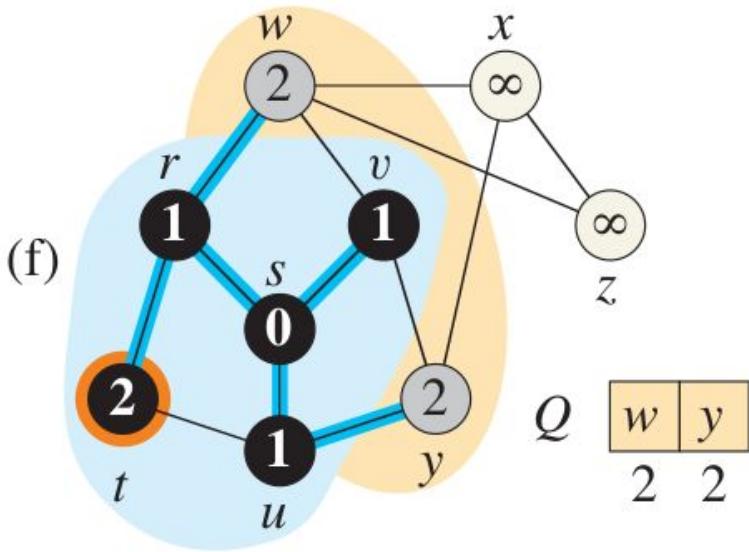




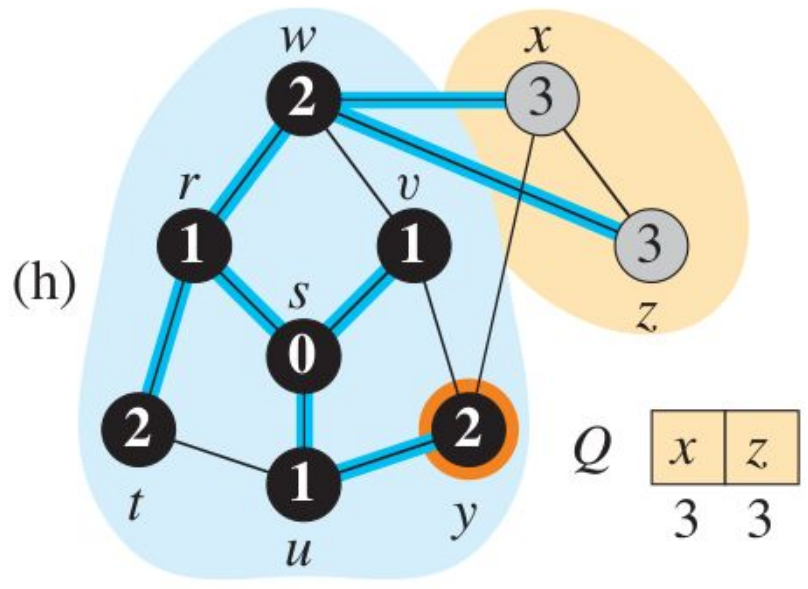
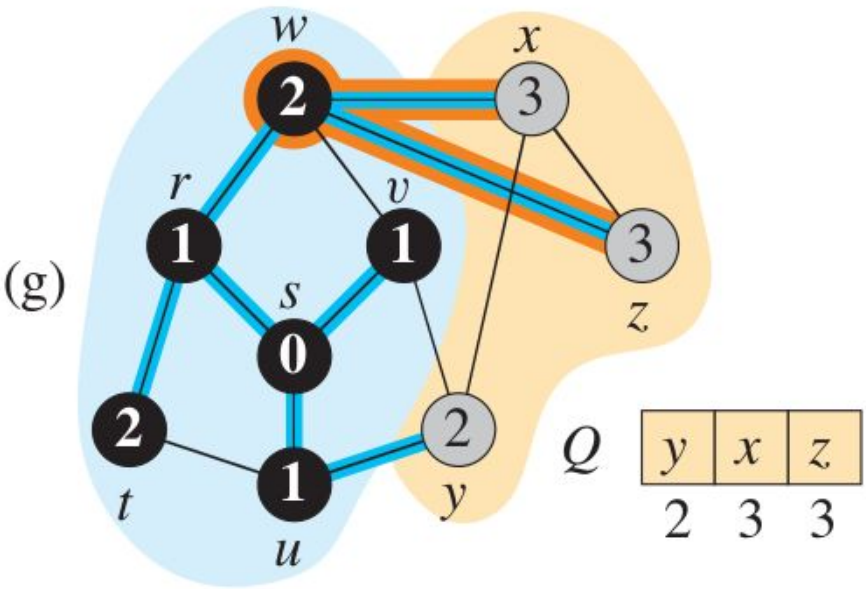
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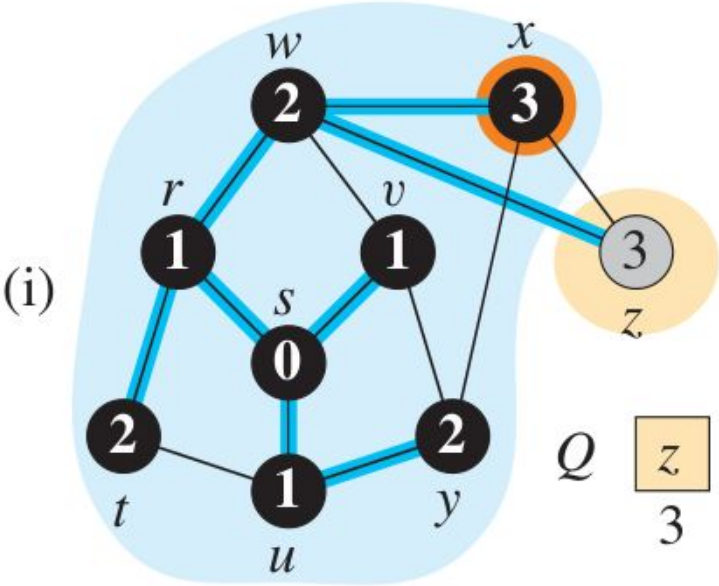
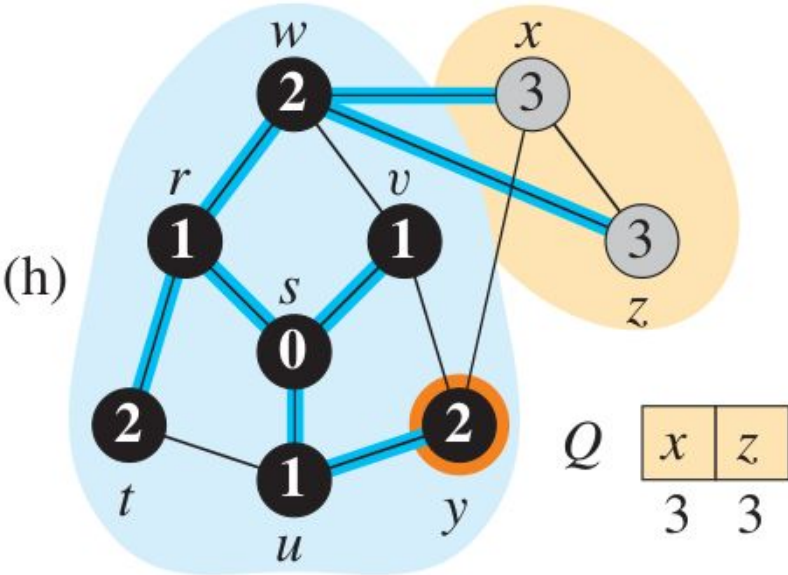
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