DATA STRUCTURES & ALGORITHMS

25: HASH TABLES

(COLLISION RESOLUTION)



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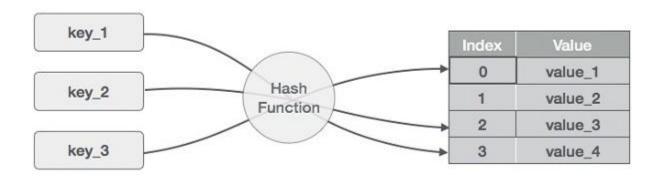
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Hash Tables

- A hash table is an effective data structure for implementing dictionaries.
- The built-in dictionaries of Python are implemented with hash tables.
- Indices can be integers, floats, chars, strings.

Hashing Function

- Function which transforms a key to an index.
- If **h** is a hash function that takes key as a parameter, it will compute index where value is placed.
- It is important to note that hash function returns the same index value every time for same key.



DIVISION METHOD

Division Method for hashing

- Hashing an integer **k** is to divide **k** by **M** and then to use the remainder **modulo M**.
 - This is called the **division method** of hashing.
 - - **k** is the key value
 - **M** is the hash table size

Division Method for hashing

- Hashing an integer **k** is to divide **k** by **M** and then to use the remainder **modulo M**.
 - This is called the **division method** of hashing.
 - - **k** is the key value
 - **M** is the hash table size

Example:

```
Hash table size M = 10

Key value, k = 112

Hash function = h(k) = k \mod M = 112 \mod 10 = 2
```

COLLISION

Collisions

• Collision occurs when one key value stores more than one value. Hash function produces the same position for two values.

Hash table cannot ensure to have single value corresponding to a key. This is because a
hash table has fixed number of keys.

• Any new entry in hash table, after all the slots are filled, causes collision. However, good hash tables use some mechanism to avoid the situations of collisions or to handle them efficiently.

Resolving Collisions

• Closed addressing Chaining

• Open addressing

Quadratic Probing

Double Hashing

Linear Probing

```
C --> 67
code = ['CS', 'CDS', 'MATH', 'STAT']
                                                          S --> 83
                                                          ASCII sum for CS is 150
for elem in code:
                                                          C --> 67
                                                          D --> 68
                                                          S --> 83
                                                          ASCII sum for CDS is 218
 val = 0
 for ch in elem:
                                                          M --> 77
                                                          A --> 65
   print (f'{ch} --> {ord (ch)}')
                                                          T --> 84
   val = val + ord(ch)
                                                          H --> 72
                                                          ASCII sum for MATH is 298
                                                          S --> 83
 print(f'ASCII sum for {elem} is {val}\n')
                                                          T --> 84
                                                          A --> 65
                                                          Т --> 84
                                                          ASCII sum for STAT is 316
```

CHAINING

Closed addressing – Chaining

$$keys = 5, 1, 6, 11$$

$$h(k) = 3k + 1$$
 $M = 6$

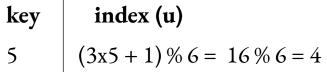
$$M = 6$$

index (u) key

$$keys = 5, 1, 6, 11$$

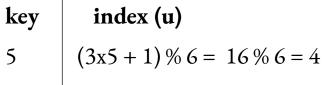
$$h(k) = 3k + 1$$
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$$keys = 5, 1, 6, 11$$

$$h(k) = 3k + 1$$
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$$keys = 5, 1, 6, 11$$

$$h(k) = 3k + 1$$
 $M = 6$

$$M = 6$$

key index (u)
5
$$(3x5+1)\%6 = 16\%6 = 4$$

1 $(3x1+1)\%6 = 4\%6 = 4$

Closed addressing – Chaining

$$keys = 5, 1, 6, 11$$

$$h(k) = 3k + 1$$
 $M = 6$

key

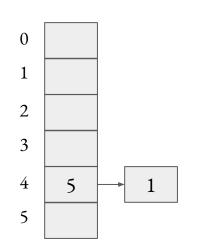
index (u)

(3x5 + 1)% 6 = 16% 6 = 4(3x1 + 1)% 6 = 4% 6 = 4

Closed addressing – Chaining

$$keys = 5, 1, 6, 11$$

$$h(k) = 3k + 1$$
 $M = 6$



index (u) key

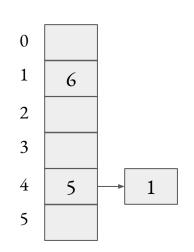
(3x5 + 1)% 6 = 16% 6 = 4 (3x1 + 1)% 6 = 4% 6 = 4 (3x6 + 1)% 6 = 19% 6 = 1

Closed addressing – Chaining

$$keys = 5, 1, 6, 11$$

$$h(k) = 3k + 1$$
 $M = 6$

$$M = 6$$



index (u) key

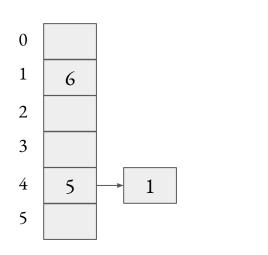
- (3x5 + 1)% 6 = 16% 6 = 4 (3x1 + 1)% 6 = 4% 6 = 4 (3x6 + 1)% 6 = 19% 6 = 1

Closed addressing – Chaining

$$keys = 5, 1, 6, 11$$

$$h(k) = 3k + 1$$
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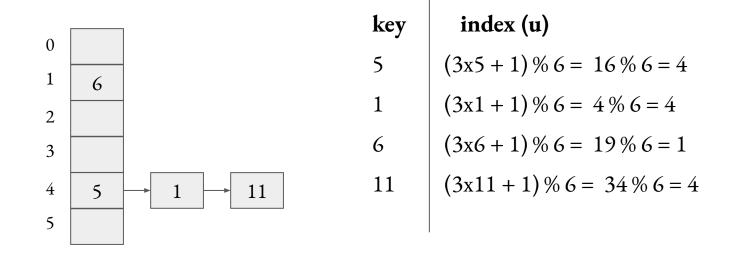


index (u) key

$$keys = 5, 1, 6, 11$$

$$h(k) = 3k + 1$$
 $M = 6$

$$M = 6$$



LINEAR PROBE

Open addressing – Linear Probe

$$keys = 5, 1, 6, 11$$

$$h(k) = 3k + 1$$
 $M = 6$

$$M = 6$$

index (u) (3x5 + 1) % 6 = 16 % 6 = 4key

Open addressing – Linear Probe

$$keys = 5, 1, 6, 11$$

$$h(k) = 3k + 1$$
 $M = 6$

key	index (u)	probe
5	(3x5 + 1)%6 = 16%6 = 4	1
1	(3x1+1)%6 = 4%6 = 4	2

Open addressing – Linear Probe

$$keys = 5, 1, 6, 11$$

$$h(k) = 3k + 1$$
 $M = 6$

key index (u) probe 5 (3x5 + 1)% 6 = 16% 6 = 4 1 1 (3x1 + 1)% 6 = 4% 6 = 4 2 6 (3x6 + 1)% 6 = 19% 6 = 1 1

Open addressing – Linear Probe

keys = 5, 1, 6, 11
$$h(k) = 3k + 1$$
 $M = 6$

$$h(k) = 3k + 1$$

$$M = 6$$

0	11	
1	6	
2		\
3		
4	5	
5	1	

key	index (u)	probe
5	(3x5 + 1)%6 = 16%6 = 4	1
1	(3x1+1)%6 = 4%6 = 4	2
6	(3x6+1)%6 = 19%6 = 1	1
11	(3x11 + 1)%6 = 34%6 = 4	3

Open addressing – Linear Probe

keys = 5, 1, 6, 11
$$h(k) = 3k + 1$$
 $M = 6$

$$h(k) = 3k + 1$$

$$M = 6$$

0				key	index (u)	probe
0	11		probing	5	(3x5 + 1)% 6 = 16% 6 = 4	1
1	6		$(\mathbf{u}+\mathbf{i})\%\mathbf{M}$	1	(3x1+1)%6 = 4%6 = 4	2
3			i = 0 to M-1	6	(3x6+1)%6 = 19%6 = 1	1
4	5			11	(3x11 + 1)% 6 = 34% 6 = 4	3
5	1					

Open addressing – Linear Probe

$$keys = 5, 1, 6, 11$$

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$$h(k) = 3k + 1$$
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0			key	index (u)	probe
0	11	probing	5	(3x5 + 1)% 6 = 16% 6 = 4	1
2	6	$(\mathbf{u} + \mathbf{i}) \% \mathbf{M}$	1	(3x1 + 1)% 6 = 4% 6 = 4	2
3		i = 0 to M-1	6	(3x6+1)%6 = 19%6 = 1	1
4	5		11	(3x11 + 1)% 6 = 34% 6 = 4	3
5	1				

Order of elements in hash table: 11, 6, ___, __, 5, 1

QUADRATIC PROBE

Open addressing - Quadratic Probe

$$keys = 5, 1, 6, 11$$

$$h(k) = 3k + 1$$
 $M = 6$

key index (u)
$$(3x5 + 1) \% 6 = 16 \% 6 = 4$$

$$(u + i^{2}) \% M$$

$$i = 0 \text{ to } M-1$$

$$(3x1 + 1) \% 6 = 4 \% 6 = 4$$

$$(3x6 + 1) \% 6 = 19 \% 6 = 1$$

$$11 \qquad (3x11 + 1) \% 6 = 34 \% 6 = 4$$

probe

Open addressing - Quadratic Probe

$$keys = 5, 1, 6, 11$$

$$h(k) = 3k + 1$$
 $M = 6$

key index (u) probe
$$5 \quad (3x5+1)\% 6 = 16\% 6 = 4$$

$$1 \quad (3x1+1)\% 6 = 4\% 6 = 4$$

$$6 \quad (3x6+1)\% 6 = 19\% 6 = 1$$

$$11 \quad (3x11+1)\% 6 = 34\% 6 = 4$$

Open addressing - Quadratic Probe

$$keys = 5, 1, 6, 11$$

$$h(k) = 3k + 1$$
 $M = 6$

key index (u) probe
$$5 \quad (3x5+1)\% 6 = 16\% 6 = 4$$

$$1 \quad (3x1+1)\% 6 = 4\% 6 = 4$$

$$6 \quad (3x6+1)\% 6 = 19\% 6 = 1$$

$$11 \quad (3x11+1)\% 6 = 34\% 6 = 4$$

Open addressing - Quadratic Probe

$$keys = 5, 1, 6, 11$$

$$h(k) = 3k + 1$$
 $M = 6$

$$M = 6$$

Open addressing - Quadratic Probe

$$keys = 5, 1, 6, 11$$

$$h(k) = 3k + 1 \qquad N$$

$$M = 6$$

Open addressing - Quadratic Probe

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Open addressing – Quadratic Probe

$$keys = 5, 1, 6, 11$$

keys = 5, 1, 6, 11
$$h(k) = 3k + 1$$
 $M = 6$

$$M = 6$$

0		$(\mathbf{u}+\mathbf{i}^2)~\%~\mathbf{M}$	key	index (u)	probe
0		i = 0 to M-1	5	(3x5 + 1)%6 = 16%6 = 4	1
2	6		1	(3x1 + 1)% 6 = 4% 6 = 4	2
3			6	(3x6 + 1)%6 = 19%6 = 1	1
4	5		11	(3x11 + 1)% 6 = 34% 6 = 4	
5	1				

Open addressing - Quadratic Probe

$$keys = 5, 1, 6, 11$$

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Open addressing – Quadratic Probe

$$keys = 5, 1, 6, 11$$

$$h(k) = 3k + 1$$
 $M = 6$

$$M = 6$$

0		$(\mathbf{u} + \mathbf{i}^2) \% \mathbf{M}$	key	index (u)	probe
0		i = 0 to M-1	5	(3x5 + 1)%6 = 16%6 = 4	1
2	6	$(4 + 0^2) \% 6 = 4$ $(4 + 1^2) \% 6 = 5$	1	(3x1 + 1)% 6 = 4% 6 = 4	2
3		(4+1)% 6 = 3 $(4+2^2)\% 6 = 2$	6	(3x6+1)%6 = 19%6 = 1	1
4	5		11	(3x11 + 1)%6 = 34%6 = 4	
5	1				

Open addressing – Quadratic Probe

keys = 5, 1, 6, 11
$$h(k) = 3k + 1$$
 $M = 6$

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$$M = 6$$

0		$(\mathbf{u} + \mathbf{i}^2) \% \mathbf{M}$	key	index (u)	probe
0		i = 0 to M-1	5	(3x5 + 1)% 6 = 16% 6 = 4	1
2	11	$(4 + 0^2) \% 6 = 4$ $(4 + 1^2) \% 6 = 5$	1	(3x1 + 1)% 6 = 4% 6 = 4	2
3	11	(4+1)% 6 = 3 $(4+2^2)\% 6 = 2$	6	(3x6 + 1)%6 = 19%6 = 1	1
4	5		11	(3x11 + 1)%6 = 34%6 = 4	3
5	1				

Open addressing - Quadratic Probe

$$keys = 5, 1, 6, 11$$

$$h(k) = 3k + 1$$
 $M = 6$

Order of elements in hash table: ___, 6, 11, ___, 5, 1

DOUBLE HASH

Open addressing – Double Hash

$$keys = 5, 1, 6, 11$$

$$h_1(k) = 3k + 1$$
 $M = 6$

$$M = 6$$

$$h_2(k) = 2k + 1$$

index (v)

key index (u)

$$(u + v^*i) \% M$$

$$i = 0 \text{ to } M-1$$

$$(3x5 + 1) \% 6 = 16 \% 6 = 4$$

$$(3x1 + 1) \% 6 = 4 \% 6 = 4$$

$$(3x6 + 1) \% 6 = 19 \% 6 = 1$$

$$(3x11 + 1) \% 6 = 34 \% 6 = 4$$

Open addressing – Double Hash

$$keys = 5, 1, 6, 11$$

$$h_1(k) = 3k + 1$$
 $M = 6$

$$M = 6$$

$$h_2(k) = 2k + 1$$

key index (u)
$$(u + v^*i) \% M$$

$$i = 0 \text{ to } M-1$$

$$(3x5 + 1) \% 6 = 16 \% 6 = 4$$

$$(3x1 + 1) \% 6 = 4 \% 6 = 1$$

$$(3x6 + 1) \% 6 = 19 \% 6 = 1$$

$$(3x11 + 1) \% 6 = 34 \% 6 = 4$$

index (v)

Open addressing – Double Hash

$$keys = 5, 1, 6, 11$$

$$h_1(k) = 3k + 1$$
 $M = 6$

$$M = 6$$

$$h_2(k) = 2k + 1$$

index (v)

Open addressing – Double Hash

$$keys = 5, 1, 6, 11$$

$$h_1(k) = 3k + 1$$
 $M = 6$

$$M = 6$$

$$h_2(k) = 2k + 1$$

index (v)

key index (u)

$$i = 0 \text{ to } M-1$$
 5
 $(3x5 + 1) \% 6 = 16 \% 6 = 4$

$$1$$
 $(3x1 + 1) \% 6 = 4 \% 6 = 4$

$$6$$
 $(3x6 + 1) \% 6 = 19 \% 6 = 1$

$$11$$
 $(3x11 + 1) \% 6 = 34 \% 6 = 4$

Open addressing – Double Hash

$$keys = 5, 1, 6, 11$$

$$h_1(k) = 3k + 1$$
 $M = 6$

$$M = 6$$

key

6

11

$$h_2(k) = 2k + 1$$

index (v)

(2x1+1)%6=3

index (u)

$$(3x5 + 1) \% 6 = 16 \% 6 = 4$$

 $(3x1 + 1) \% 6 = 4 \% 6 = 4$
 $(3x6 + 1) \% 6 = 19 \% 6 = 1$
 $(3x11 + 1) \% 6 = 34 \% 6 = 4$

Open addressing – Double Hash

$$keys = 5, 1, 6, 11$$

$$h_1(k) = 3k + 1$$
 $M = 6$

key

6

11

$$M = 6$$

$$h_2(k) = 2k + 1$$

$$(\mathbf{u} + \mathbf{v}^*\mathbf{i}) \% \mathbf{M}$$

 $\mathbf{i} = 0 \text{ to } \mathbf{M} - 1$
 $(4+3^*0)\%6 = 4$

index (u)

$$(3x5 + 1) \% 6 = 16 \% 6 = 4$$

 $(3x1 + 1) \% 6 = 4 \% 6 = 4$
 $(3x6 + 1) \% 6 = 19 \% 6 = 1$
 $(3x11 + 1) \% 6 = 34 \% 6 = 4$

index (v)

(2x1+1)%6=3

Open addressing – Double Hash

$$keys = 5, 1, 6, 11$$

$$h_1(k) = 3k + 1$$
 $M = 6$

$$M = 6$$

$$h_2(k) = 2k + 1$$

$$(\mathbf{u} + \mathbf{v}^*\mathbf{i}) \% \mathbf{M}$$

 $\mathbf{i} = 0 \text{ to } \mathbf{M} - 1$
 $(4+3^*0)\%6 = 4$
 $(4+3^*1)\%6 = 1$

key index (u)

$$5 (3x5 + 1) \% 6 = 16 \% 6 = 4$$

$$(3x1 + 1) \% 6 = 4 \% 6 = 4$$

$$(3x6+1)\% 6 = 19\% 6 = 1$$

 $(3x11+1)\% 6 = 34\% 6 = 4$

6

11

index (v)

(2x1+1)%6=3

Open addressing – Double Hash

$$keys = 5, 1, 6, 11$$

$$h_1(k) = 3k + 1$$
 $M = 6$

$$M = 6$$

$$h_2(k) = 2k + 1$$

key index (u)
$$(3x5+1)\% 6 = 16\% 6 = 4$$

$$(3x1+1)\% 6 = 4\% 6 = 4$$

$$(3x6+1)\% 6 = 19\% 6 = 1$$

$$(3x6+1)\% 6 = 19\% 6 = 1$$

 $(3x11+1)\% 6 = 34\% 6 = 4$

11

index (v)

(2x1+1)%6=3

Open addressing – Double Hash

$$keys = 5, 1, 6, 11$$

$$h_1(k) = 3k + 1$$
 $M = 6$

$$M = \epsilon$$

$$h_2(k) = 2k + 1$$

6

11

$$(3x5 + 1)\% 6 = 16\% 6 = 4$$

 $(3x1 + 1)\% 6 = 4\% 6 = 4$

$$(3x1+1)\%6 = 4\%6 = 4$$

$$(3x6+1)\%6 = 19\%6 = 1$$

 $(3x11+1)\%6 = 34\%6 = 4$

(2x1 + 1)%6 = 3

index (v)

Open addressing – Double Hash

$$keys = 5, 1, 6, 11$$

$$h_1(k) = 3k + 1$$
 $M = 6$

$$M = 6$$

key

6

11

$$h_2(k) = 2k + 1$$

index (v)

(2x1 + 1)%6 = 3

(2x6+1)%6=1

index (u)

$$(3x5+1)\% 6 = 16\% 6 = 4$$

 $(3x1+1)\% 6 = 4\% 6 = 4$
 $(3x6+1)\% 6 = 19\% 6 = 1$
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Open addressing - Double Hash

$$keys = 5, 1, 6, 11$$

$$h_1(k) = 3k + 1$$

$$M = 6$$

$$h_2(k) = 2k + 1$$

index (u)

$$(3x5 + 1)\% 6 = 16\% 6 = 4$$

 $(3x1 + 1)\% 6 = 4\% 6 = 4$
 $(3x6 + 1)\% 6 = 19\% 6 = 1$
 $(3x11 + 1)\% 6 = 34\% 6 = 4$

(2x1 + 1)%6 = 3(2x6 + 1)%6 = 1

index (v)

Open addressing – Double Hash

keys = 5, 1, 6, 11
$$h_1(k) = 3k + 1$$
 $M = 6$ $h_2(k) = 2k + 1$

Open addressing – Double Hash

keys = 5, 1, 6, 11
$$h_1(k) = 3k + 1$$
 $M = 6$ $h_2(k) = 2k + 1$

5

Open addressing – Double Hash

Open addressing – Double Hash

keys = 5, 1, 6, 11
$$h_1(k) = 3k + 1$$
 $M = 6$ $h_2(k) = 2k + 1$

Open addressing – Double Hash

keys = 5, 1, 6, 11
$$h_1(k) = 3k + 1$$
 $M = 6$ $h_2(k) = 2k + 1$

0		(u + v*i) % M key	index (u)	index (v)
0	1	i = 0 to M-1	(3x5 + 1)% 6 = 16% 6 = 4	
2.	6	(4+5*0)%6=4 1	(3x1 + 1)% 6 = 4% 6 = 4	(2x1+1)%6=3
3	11	$(4+5^*1)\%6 = 3$	(3x6+1)%6 = 19%6 = 1	(2x6+1)%6=1
4	5	11	(3x11+1)%6 = 34%6 = 4	(2x11 + 1)%6 = 5
5				

Open addressing - Double Hash

keys = 5, 1, 6, 11
$$h_1(k) = 3k + 1$$
 $M = 6$ $h_2(k) = 2k + 1$

Order of elements in hash table: __, 1, 6, 11, 5, __