

Solution Details:

- Reading input file as per the input specification in the sample input file.
- As we read the input file, we create separate list (bank_details_may and bank_details_june) for month of May and June. We also create another list (bank_details_new) for new records.
- Display bank details (bank name and denomination count) of each list. The time complexity for this operation is **$O(n)$** .
- Calculate the total count of Rs. 2000 for the months of May and June. The time complexity for this operation is **$O(n)$** .
- - a. Append new records to the lists – bank_details_may and bank_details_june according to the months. The time complexity for this operation is **$O(n)$** .
 - b. Create the max heap for the month of May and June using max_heap_build method. The time complexity for this operation is **$O(n\log(n))$** .

Heapification:

Function max_heapify(bank_details, start_index, key_func){

 Left = 2*start_index + 1

 Right = 2 * start_index + 2

 Max_index = start_index

 If bank_details is not empty

 Size = len(bank_details)

 If Left < size and

 key_func(bank_details[Left]) > key_func(bank_details[Max_index])

 Max_index = Left

 If Right < size and

 key_func(bank_details[Right]) > key_func(bank_details[Max_index])

 Max_index = Right

 If Max_index != start_index then swap(bank_details, Max_index, start_index)

 If (Max_index == Right) then max_heapify(bank_details, Right, key_func)

 Else max_heapify (bank_details, Left, key_func)

 Else return null

}

- c. Maximum denomination count for each is the first element for that month (bank_details_may[0] and bank_details_june[0]). The time complexity for this operation is **$O(1)$** .
- Sum the denomination count for all the banks from both the lists. The time complexity for this operation is **$O(n)$** .

- Remove the first element from the may month maxheap (maxheap_pop function), and heapify the maxheap. This is done twice to remove the largest 2 elements. The time complexity for this operation is $O(\log(n))$.

Deletion:

```
Function Max_heap_pop(bank_details, key_func){
    If length(bank_details) > 1
        Element = bank_details[0]
        bank_details[0] = bank_details.pop()
        max_heapify( bank_deatils,0,key_func)
    Else Element = bank_details.pop()
    Return Element;
}
```

- Display bank details (bank name and denomination count) of each maxheap. The time complexity for this operation is $O(n)$.
- [Instruction 2]** Errors handling for invalid input has been implemented. These are logged to console. Underflow or empty list/heap is also handled in the code.

```
c:\Users\l... \bits\sem-1\BITS\DSAD\assignments>python ps07_group27.py
Record - 'Bank12,, ' was skipped as denomination count is expected to be positive integer.
Record - 'Bank13, , may' was skipped as denomination count is expected to be positive integer.
Record - 'Bank14, -12, may' was skipped as denomination count is expected to be positive integer.
Record - 'Bank15, -12, June' was skipped as denomination count is expected to be positive integer.
```

- [Instruction 8.a]** We have used 2 separate maxheaps to store the bank details for each month. This ensure that the data is easy to access the details of largest denomination count (time complexity is $O(1)$). Also, when such element is removed, it is quite cheap to find the next largest denomination count and maintain the structure of the heap (time complexity is $O(n\log(n))$).
- [Instruction 8.b]** Each operation and its time complexity.
 - **build-maxheap:** It iterates through non-leaf nodes starting the last one and call maxheapification. Choosing maxheap reduces the time complexity to $O(n\log(n))$, while with other structures, the sorting would be quite expensive with time complexity of $O(n^2)$.
 - **maxheapification:** maxheapification maintains the heap structure with largest element being at the top. This maxheapification process is of time complexity $O(n\log(n))$.
 - **maxheap-pop:** It removes the largest element from the heap and heapifies the remain structure. The time complexity of the process is $O(\log(n))$.
- [Instruction 8.c]** The *alternate approach* would be storing the bank details a *stack* based on the descending order of denomination count.
 - *Building* the sorted stack will be expensive with time complexity of $O(n^2)$.
 - *Traversing* this stack will be time complexity $O(n)$.
 - *Popping* the largest element will be computationally cheap and have time complexity of $O(1)$.