

24/02/25

## 6. Searching

\* List is doubly linked list.

\* Iterator  $\text{vector<int>} :: \text{iterator} \quad \text{itr} = \text{a.begin();}$   
 $\hookrightarrow \text{advance(ptr, x)}$

#

### Binary Search

TC:  $O(\log n)$ SC:  $O(1) / O(\log n)$ 

Iterative

Recursive

#### Iterative

```
int fun(int a[], int n, int x){
    int l=0, h=n-1;
    while (l <= h) {
        int m = (l+h)/2;
        if (a[m] == x)
            return m;
        if (a[m] < x)
            l = m+1;
        if (a[m] > x)
            h = m-1;
    }
    return -1;
}
```

#### Recursive

```
int fun(int a[], int l, int h, int x) {
    if (l > h)
        return -1;
    int m = (l+h)/2;
    if (a[m] == x)
        return m;
    else if (a[m] > x)
        return fun(a, l, m-1, x);
    else if (a[m] < x)
        return fun(a, m+1, h, x);
}
```

#

### 1<sup>st</sup> occurrence of a number

```
int firstocc (int a[], int n, int x) {
    int l=0, h=n-1;
    while (l <= h) {
        int m = (l+h)/2;
        if (a[m] > x)
            h = m-1;
    }
}
```

$Tc: O(\log n)$   
 $Sc: O(1)$

```

else if ( $a[m] < x$ )
     $l = m + 1$ ;
else {
    if ( $m = 0$  or  $a[m] \neq a[m - 1]$ )
        return  $m$ ;
    else
         $h = m - 1$ ;
}

```

## # Last occurrence of a number

$Tc: O(\log n)$   
 $Sc: O(1)$

```
int lastocc (int a[], int n, int x) {
```

```
    int l=0, h=n-1;
```

```
    while ( $l \leq h$ ) {
```

```
        int m = ( $l+h$ )/2;
```

```
        if ( $a[m] > x$ )
```

```
            h = m-1;
```

```
        else if ( $a[m] < x$ )
```

```
            l = m+1;
```

```
        else {
```

```
            if ( $m == n-1$  or  $a[m] \neq a[m+1]$ )
```

```
                return m;
```

```
            else
```

```
                l = m+1;
```

```
}
```

```
3
```

## ## Square root of given number.

```

int squareRoot ( int n ) {
    int l = 0, h = n ;
    int ans ;
    while ( l <= h ) {
        int m = ( l + h ) / 2 ;
        if ( ( m * m ) == n )
            return m ;
        else if ( ( m * m ) < n ) {
            l = m + 1 ;
            ans = m ;
        }
        else
            h = m - 1 ;
    }
    return ans ;
}

```

TC:  $O(\log n)$   
SC:  $O(1)$

## ## Search in Infinite array.

```

int fun ( int x, int a[] ) {
    if ( a[0] == x )
        return 0 ;
    int i = 1 ;
    while ( a[i] < x )
        i *= 2 ;
    if ( a[i] == x )
        return i ;
    return binarySearch ( a, i / 2 + 1, i, x ) ;
}

```

TC:  $O(\log n)$   
SC:  $O(1)$

```

int binarySearch ( int a[], int l, int h, int x) {
    while ( l <= h ) {
        int m = (l+h)/2;
        if ( a[m] == x )
            return m;
        if ( a[m] > x )
            h = m-1;
        if ( a[m] < x )
            l = m+1;
    }
    return -1;
}

```

### #### Search in a sorted rotated array.

$a[] = 3 4 5 1 2$	
$x = 4$	

TC:  $O(\log n)$   
SC:  $O(1)$

```

int fun ( int a[], int n, int x) {
    int l=0, h=n-1;
    while ( l <= h ) {
        int m=(l+h)/2;
        if ( a[m] == x )
            return m;
        if ( a[m] >= a[l] ) {
            // sorted increasing part
            if ( a[l] <= x and a[m] >= x )
                h = m-1;
            else
                l = m+1;
        }
        else {

```

```

else if // rotated part here
if (a[h] >= x and a[m] <= x)
    l = m + 1;
else
    h = m - 1;
}

```

1  
return -1;

}

#### find Peak element in a given array.  
 ↳ Peak element:  $\Delta$  element on left & right  $< \text{element}$

$a = [1, 2, 4, 5, 3, 6]$

and = 5 or 6

```

int fun(int a[], int n) {
    int l = 0, h = n - 1;
    if (a[0] > a[1])
        return 0;
    if (a[n - 1] > a[n - 2])
        return n - 2;
}

```

while (l < h) {

int m = (l + h) / 2;

~~if (a[m] > a[m - 1] & a[m] > a[m + 1])~~

if (a[m] > a[m - 1] or a[m] > a[m + 1]))

return m;

else if (m != 0 and a[m] < a[m - 1])

h = m - 1;

else

l = m + 1;

3.

## find pair with given sum 's' in unsorted array

```
bool fun ( int a[], int n, int s ) {
    map < int, int > hash;
    for ( int i=0; i<n; i++ ) {
        if ( hash [ s - a[i] ] != 0 )
            return true;
        hash [ a[i] ]++;
    }
    return false;
}
```

TC: O(n)

SC: O(n)

## find pair with given sum 's' in sorted array .

```
bool fun ( int a[], int n, int s ) {
    int l=0, h=n-1;
    while ( l < h ) {
        if ( a[l] + a[h] == s )
            return true;
        if ( a[l] + a[h] < s )
            l++;
        if ( a[l] + a[h] > s )
            h--;
    }
    return false;
}
```

TC: O(n)

SC: O(1)

~~## find pair with given sum 's' in sorted array~~

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#### find triplet with given sum's in a sorted array

```
bool triplet ( int a[], int n, int s ) {
    for( int i = 0; i < n; i++ ) {
        if( pair(a, i+1, n-1, s - a[i]) )
            return true;
    }
}
```

3.

```
return false;
```

$Tc: O(n^2)$
$Sc: O(1)$

```
bool pair ( int a[], int l, int h, int s ) {
```

```
    while( l < h ) {
```

```
        if( a[l] + a[h] == s )
```

```
            return true;
```

```
        else if( a[l] + a[h] < s )
```

```
            l++;
        
```

```
        else
```

```
            h--;
        
```

3

```
return false;
```

3.

★ If unsorted array is given, we can first sort array in  $O(n \log n)$  then apply above method.  $\rightarrow Tc: O(n^2)$   $Sc: O(1)$ .

#### Count no. of pairs with a given sum.

① Unsorted array

$Tc: O(n)$   $Sc: O(n)$

```
int fun( int a[], int n, int s ) {
    int ans = 0;
    map<int, int> hash;
```

② Sorted array

$Tc: O(n)$   $Sc: O(1)$

```
int fun( int a[], int n, int s ) {
    int l = 0, h = n - 1;
    int ans = 0;
```

```

for (int i=0; i<n; i++) {
    if (hash[s-a[i]] != 0)
        ans++;
    hash[a[i]]++;
}
return ans;

```

3.

```

while (l < h) {
    if (a[l] + a[h] == s)
        ans++;

```

```

    l++; h--;
}

```

```

else if (a[l] + a[h] < s) {
    l++;
}

```

```

else if (a[l] + a[h] > s) {
    h--;
}

```

```

return ans;

```

3.

## ## Count triplet with given Sum 's'

### ① Sorted array

TC: O(n<sup>2</sup>) SC: O(1)

```

int func(int a[], int n, int s) {
    int ans = 0;
    for (int i=0; i<n; i++) {
        if (pair(a, i+1, n-1, s-a[i]))
            ans++;
    }
    return ans;
}

```

```

bool pair(int a[], int l, int h, int s) {
    while (l < h) {
        if (a[l] + a[h] == s)
            return true;
        else if (a[l] + a[h] > s)
            h--;
        else

```

### ② Unsorted array

TC: O(n<sup>2</sup>) SC: O(1)

→ sort array → O(n log n)

do same like this

l++;

}

return false;

3.

##

find triplet a, b, c such that  $a^2 + b^2 = c^2$  [Pythagorean theorem]

\* If unsorted, sort the array first.

```
bool
fun( int a[], int n ) {
    for( int i=0; i<n; i++ ) {
        a[i] *= a[i];
    }
    sort( a, a+n );
    for( int i=n-1; i>=0; i-- ) {
        if( ispair( a, 0, i-1, a[i] ) )
            return true;
    }
}
```

TC: O(n<sup>2</sup>)

SC: O(1)

3

return false;

3.

```
bool ispair( int a[], int l, int h, int s ) {
    while( l < h ) {
        if( a[l] + a[h] == s )
            return true;
        else if( a[l] + a[h] < s )
            l++;
        else
            h--;
    }
}
```

return false;

3.

Video

## # Medium of 2 sorted array

here  $a = \{ \dots \} \rightarrow n_1$  $b = \{ \dots \} \rightarrow n_2$  $n_1 \leq n_2$ 

```

int l = 0, h = n_1;
while (l <= h) {
    int i_1 = (l+h)/2;
    int i_2 = (n_1+n_2+1)/2 - i_1;

```

max <sub>1</sub>	min <sub>1</sub>
max <sub>2</sub>	min <sub>2</sub>

first point is

$$m = (l+h)/2$$

second point is.

$$m = \frac{(n_1+n_2+1)}{2} - m$$

```
int max1 = (i1 == 0) ? INT_MIN, a[i1-1];
```

```
int min1 = (i1 == n1) ? INT_MAX, a[i1];
```

```
int max2 = (i2 == 0) ? INT_MIN, b[i2-1];
```

```
int min2 = (i2 == n2) ? INT_MAX, b[i2];
```

```
if (max1 <= min2 and max2 <= min1) {
```

```
    if ((n1+n2) % 2)
```

```
        return max(max1, max2);
```

else

```
    return [max(max1, max2) + min(min1, min2)] / 2;
```

```
else if (max1 > min2) {
```

```
    h = i1 - 1;
```

```
    else if (max2 > min1) {
```

```
        l = i1 + 1;
```

J.

J. of next meet:

TC : O(log n)

SC : O(1)

## ## Repeating element. → video

```

int sol (vector<int> a, int n) {
    int slow = a[0] + 1, fast = a[0] + 1;
    do {
        slow = a[slow] + 1;
        fast = a[a[fast] + 1] + 1;
    } while (slow != fast);
    slow = a[0] + 1;
    while (slow != fast) {
        slow = a[slow] + 1;
        fast = a[fast] + 1;
    }
    return slow - 1;
}

```

## ## Allocate minimum pages

$TC: O(n \log(\text{sum}))$   
 $SC: O(1)$

### ① Naive Recursive Approach.

```

int fmin (int a[], int n, int k) {
    if (n == 1)
        return a[0];
    if (k == 1)
        return sum (a, 0, n-1);
    int res = INT-MAX;
    for (int i = 1; i < n; i++) {
        res = min (res, max (fmin(a, i, k-1), sum(a, i, n-1)));
    }
    return res;
}

```

## ② Efficient approach

ans should lie b/w [max<sup>m</sup> element, sum of all element]

then apply binary search

$\text{if } m < \text{sum} \rightarrow \text{if mid element is feasible} \Rightarrow \text{check for value} \leq \text{mid}$   
 if there is any feasible value.

else find feasible in right half {value  $>$  mid}.

```
int findmin ( int a[], int n, int k ) {
```

// here n → size of array, k → no. of people to divide

int maxim = INT\_MIN;

int tot = 0;

for( int i=0 ; i<n ; i++ ) {

    tot += a[i];

    maxim = max( maxim, a[i] );

}

int l = maxim, h = tot, res;

while( l <= h ) {

    int m = (l+h)/2;

    if( isfeasible (a, n, k, m) ) {

        h = m-1;

    } else {

        l = m+1;

    }

    3

    return res;

3

```

bool isfeasible ( int a[], int n, int k, int m ) {
    int count = 1, cur = 0;
    for ( int i = 0; i < n; i++ ) {
        if ( cur + a[i] > m ) {
            count++;
            cur = a[i];
        } else {
            cur += a[i];
        }
    }
    return ( count <= k );
}

```

- # binary-search ( start\_ptr , end\_ptr , element ) → true / false .
- # upper-bound ( start\_ptr , end\_ptr , element ) → ptr to element just greater than element .
- # lower\_bound ( start\_ptr , end\_ptr , element ) → ptr to element if present / element just small than element .

### ##### Minimum number in a sorted rotated array .

```

int fun ( int a[], int l, int h ) {
    while ( l <= h ) {
        int m = ( l + h ) / 2;
        if ( a[m] == a[h] )
            h--;
        if ( a[m] > a[h] )
            l = m + 1;
        else
            h = m;
    }
    return a[h];
}

```

$T.C: O(\log n)$   
 $S.C: O(1)$

## two selected element

```
int fun( int a[], int n) {
    vector<int> ans;
    for( int i = 0; i < N + 2; i++) {
        a[abs(a[i]) - 1] *= -1;
    }
}
```

3. return ans;

```
for( int i = 0; i < N; i++) {
    if( a[i] > 0) {
        ans.push_back( i);
    }
}
```

3.

return ans;

## Maxm water b/w 2 building.

↓

delete (n-2) building → water trapped = (gap b/w building) \* min(a[l], a[h])

```
int fun( int a[], int n) {
    if( n <= 2)
```

return 0;

int ans = 0, maxm = 0;

int l = 0, h = n - 1;

while( l <= h) {

ans = max( ans, (h - l - 1) \* min(a[l], a[h]));

maxm = max( maxm, ans);

if( a[l] < a[h])

l++;

else

h--;

TC: O(n)

SC: O(1)

return maxm; // step 4 out ##

3. tri. (No tri) with two

two < n > values

##

Smallest positive missing number.

Time complexity O(n)

int missingno (int a[], int n) {

    int p = 0;

    for (int i = 0; i < n; i++) {

        if (a[i] == i)

            p++;

    }

    if (p == 0)

        return 1;

    for (int i = 0; i < n; i++) {

        if (a[i] <= 0 || a[i] > n), M #77

            a[i] = 1;

    }

    for (int i = 0; i < n; i++) {

        if (a[i] <= (n+1)).

            return i+1;

    }

    return n+1;

    3. no ans;

##

Count only repeated

int fun (int a[], int n) {

    int l = 0, h = n-1;

    while (l <= h) {

        int m = (l+h)/2;

        if (a[m] >= a[0] + m)

            l = m+1;

        else

            h = m;

$f(m, n)$  if ( $n - a[n-1] + a[0] == 1$ )  
 $m = m - 1$  return  $\{ -1, -1 \}$ ;  
 $1 - m == n$  else.

return  $\{ a[l], n - a[n-1] + a[0] \}$ ;

3.  $l + m = 1$

### ## Count more than $n/k$ occurrences

int fun( int a[], int n, int k) {

map<int, int> m;

for( int i=0; i < n; i++) {

if (m[a[i]] > 0) m[a[i]]++;

else

3.

int p = n/k; int count = 0;

for( auto x : m) {

if (x.second > p)

count++;

3.

return count;

3.

### ## Allocate minimum number of pages

int findPages( int a[], int n, int k) {

int l, h, ans = -1;

l = INT\_MIN;

h = 0;

for( int i=0; i < n; i++) {

if (l > a[i])

l = a[i];

h += a[i];

3.

while (l <= h) {

m = (l+h)/2;

