

Eg:- $m \times (1111 \overset{3}{1} \overset{2}{0} \overset{1}{0} \overset{0}{0})$

$$= mx(-2^3 + 2')$$

$$= m \times (-2^3 + 2^2 - 2^1)$$

For the initial block of only, difference arises. R just like case 1.

~~Here~~


Hamacher: Booth's Algorithm

Bit 1	Bit -1	Partial Product (Summand)
0	0	0xM
0	1	1xM
1	0	-1xM
1	1	-0xM

Ex:- $(+13) \times (-6)$

$$+13 = \quad 0 \quad 1 \quad 1 \quad 0 \quad 1$$

$$-6 = 11010 \boxed{0}$$

Booth's 
Recording

$$0 - 1 + 1 - 1 + 0$$

0 0 0 0 0 0 0 0 0 0

1 1 1 1 1 0 0 1 1 x

0 0 0 0 1 1 0 1 X

1 0 1 0 0 1 1 1 x

○ ○ ○ ○ ○ ○ ×

110110010

↓ 2's compl.

$$= -(000100110)_2 = -(78)_{10}$$

Eg:-

$$-7 = 1001$$

$$-3 = 1101 \boxed{0}$$

$$\begin{array}{l} 0 - 1 + 1 - 1 \\ -4 + 2 - 0 \end{array}$$

$$1001 \xrightarrow{2^{\text{ls Comp}}} 0111$$

$$0 - 1 + 1 - 1$$

$$\begin{array}{ccccccc} 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 0 & 0 & 1 & X \\ 0 & 0 & 0 & 1 & 1 & 1 & X & X \\ 0 & 0 & 0 & 0 & 0 & X & & \end{array}$$

$$\begin{array}{ccccccc} 1 & 1 & 1 & 1 & 0 & 0 & 1 & X \\ 0 & 0 & 0 & 1 & 1 & 1 & X & X \\ 0 & 0 & 0 & 0 & 0 & X & & \end{array}$$

$$\begin{array}{ccccccc} 0 & 0 & 0 & 1 & 1 & 1 & X & X \\ 0 & 0 & 0 & 0 & 0 & X & & \end{array}$$

$$\begin{array}{ccccccc} 0 & 0 & 0 & 0 & 0 & X & & \end{array}$$

Discard

Sign
bit

$$(21)_{10} \rightarrow (011001000)_2$$

Eg:-

$$-7 = 1001 \boxed{0}$$

$$+3 = 0011 \boxed{0}$$

$$0 + 10 - 1$$

$$1001 \rightarrow 0111$$

$$0 + 10 - 1$$

$$\begin{array}{ccccccc} 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & X \\ 1 & 1 & 1 & 0 & 0 & 1 & X & \\ 0 & 0 & 0 & 0 & 0 & X & & \end{array}$$

$$\begin{array}{ccccccc} 0 & 0 & 0 & 0 & 0 & 0 & X \\ 1 & 1 & 1 & 0 & 0 & 1 & X \\ 0 & 0 & 0 & 0 & 0 & X & & \end{array}$$

$$\begin{array}{ccccccc} 1 & 1 & 1 & 0 & 0 & 1 & X \\ 0 & 0 & 0 & 0 & 0 & X & & \end{array}$$

$$\begin{array}{ccccccc} 1 & 1 & 1 & 0 & 0 & 1 & X \\ 0 & 0 & 0 & 0 & 0 & X & & \end{array}$$

$$\begin{array}{ccccccc} 1 & 1 & 1 & 0 & 0 & 1 & X \\ 0 & 0 & 0 & 0 & 0 & X & & \end{array}$$

$$\begin{array}{ccccccc} 1 & 1 & 1 & 0 & 0 & 1 & X \\ 0 & 0 & 0 & 0 & 0 & X & & \end{array}$$

$$\begin{array}{ccccccc} 1 & 1 & 1 & 0 & 0 & 1 & X \\ 0 & 0 & 0 & 0 & 0 & X & & \end{array}$$

$$= -(00010101)_2 = (-21)_{10}$$

Fast Multiplication:-

• Reduction of the number of summands,

• CSA

Bit-pair of Booth's Recoding,

$$-6 = 11010 \rightarrow \begin{array}{ccccc} 0 & -1 & +1 & -1 & 0 \\ \hline 0 & -1 & -2 & & \end{array}$$

$$\begin{bmatrix} -1 & +1 \\ = -2m + m \\ = -m \end{bmatrix}$$

$$+13 \rightarrow 01101 \rightarrow 10011$$

$$0 \quad 111+2$$

$$\begin{array}{cccccccc} 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 \\ 1 & 1 & 1 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & x & x \end{array}$$

$$\square 111010010$$

↓

$$-(000100110)_2 = (-78)_{10}$$

*No. of Summands are reduced by $\lceil \frac{n}{2} \rceil$ if $n = no.$

Eg:-

$$-7 \rightarrow 1001$$

$$-3 \rightarrow 1101 \quad \square$$

$$\begin{array}{cc} 0-1+1-1 \\ \square \quad \square \\ -1 \quad +1 \end{array}$$

$$\begin{array}{cccccccc} 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ x & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 \end{array} \rightarrow 0111$$

$$\begin{array}{cccccccc} x & 1 & 0 & 0 & 1 & -1 & +1 & \end{array}$$

$$11010101001$$

$$898111xx$$

$$\square 000100101$$

$$= (2)_{10}$$

Eg:-

$$+6 \rightarrow 00110$$

$$-13 \rightarrow 10011 \quad \square$$

$$\begin{array}{ccccc} -1 & 0 & +1 & 0 & -1 \\ \square & \square & \square & \square & \square \end{array}$$

$$-1 \quad +1 \quad -1$$

$$00110 \rightarrow 11010$$

$$\begin{array}{ccccc} -1 & +1 & -1 & \end{array}$$

$$111111010$$

$$00000110xx$$

$$111010xx$$

$$\square 111011010$$

$$= -(000100110)_2 = (-78)_{10}$$