

Longest common subsequence

INPUT: two strings

OUTPUT: longest common subsequence

ACTGAACTCTGTGCACT

TGACTCAGCACAAAAC

Longest common subsequence

INPUT: two strings

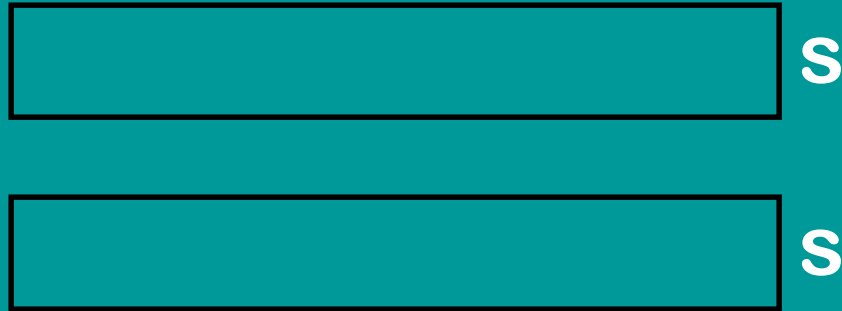
OUTPUT: longest common subsequence

ACTGAACTCTGTGCACT

TGACTCAGCACAAAAC

Longest common subsequence

If the sequences end with the same symbol s , then LCS ends with s .



Longest common subsequence

Sequences x_1, \dots, x_n , and y_1, \dots, y_m

$\text{LCS}(i,j)$ = length of a longest common
subsequence of x_1, \dots, x_i and y_1, \dots, y_j

Longest common subsequence

Sequences x_1, \dots, x_n , and y_1, \dots, y_m

$\text{LCS}(i, j)$ = length of a longest common
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if $x_i = y_j$ then

$\text{LCS}(i, j) =$

Longest common subsequence

Sequences x_1, \dots, x_n , and y_1, \dots, y_m

$\text{LCS}(i,j)$ = length of a longest common subsequence of x_1, \dots, x_i and y_1, \dots, y_j

if $x_i = y_j$ then

$$\text{LCS}(i,j) = 1 + \text{LCS}(i-1,j-1)$$

Longest common subsequence

Sequences x_1, \dots, x_n , and y_1, \dots, y_m

$\text{LCS}(i,j)$ = length of a longest common subsequence of x_1, \dots, x_i and y_1, \dots, y_j

if $x_i \neq y_j$ then

$$\text{LCS}(i,j) = \max(\text{LCS}(i-1,j), \text{LCS}(i,j-1))$$

x_i and y_j cannot both be in LCS

Longest common subsequence

Sequences x_1, \dots, x_n , and y_1, \dots, y_m

$\text{LCS}(i, j)$ = length of a longest common subsequence of x_1, \dots, x_i and y_1, \dots, y_j

if $x_i = y_j$ then

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if $x_i \neq y_j$ then

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Longest common subsequence

Running time ?

Sequences x_1, \dots, x_n , and y_1, \dots, y_m

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Longest common subsequence

Running time = $O(mn)$

Sequences x_1, \dots, x_n , and y_1, \dots, y_m

$LCS(i,j)$ = length of a longest common
subsequence of x_1, \dots, x_i and y_1, \dots, y_j

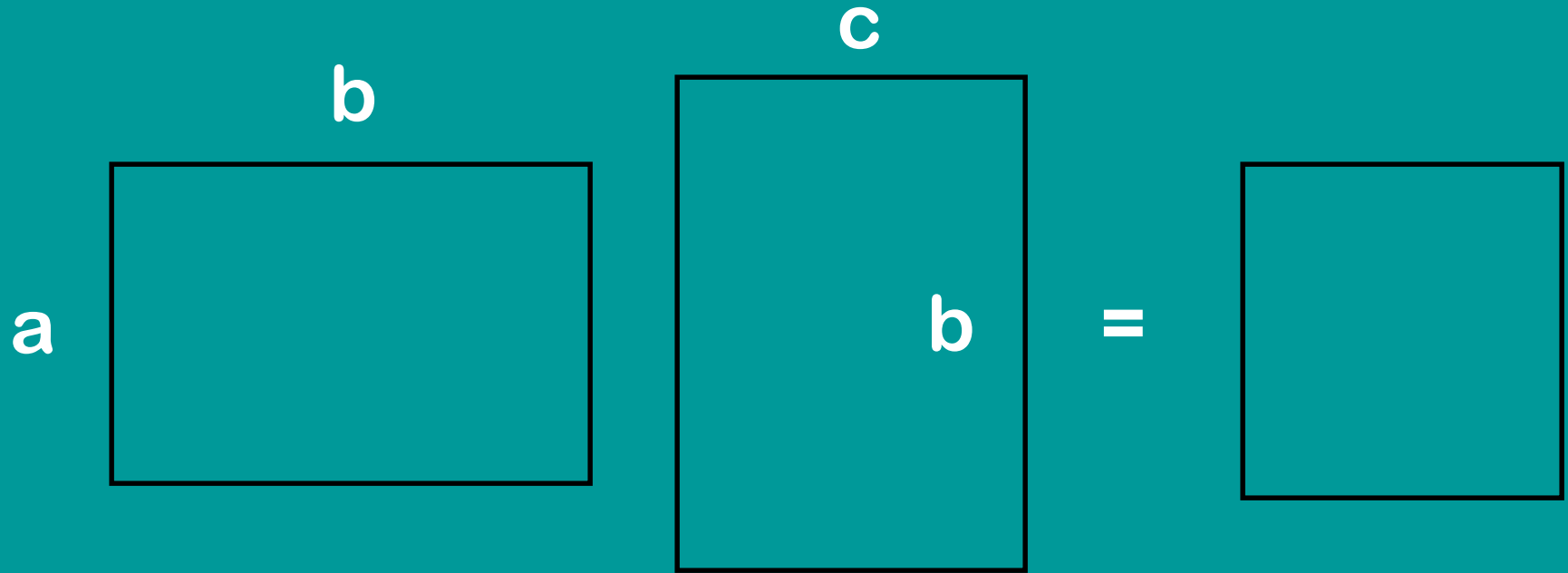
if $x_i = y_j$ then

$$LCS(i,j) = 1 + LCS(i-1,j-1)$$

if $x_i \neq y_j$ then

$$LCS(i,j) = \max (LCS(i-1,j), LCS(i,j-1))$$

Optimal matrix multiplication

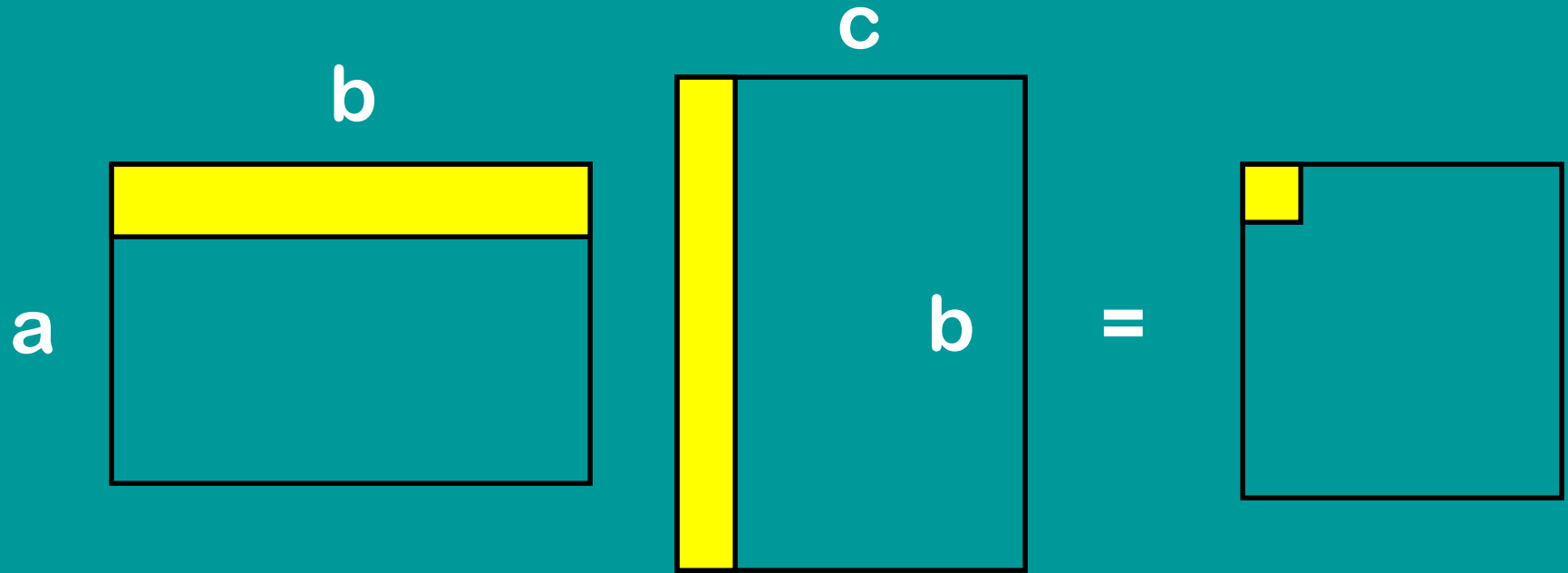


$A = a \times b$ matrix

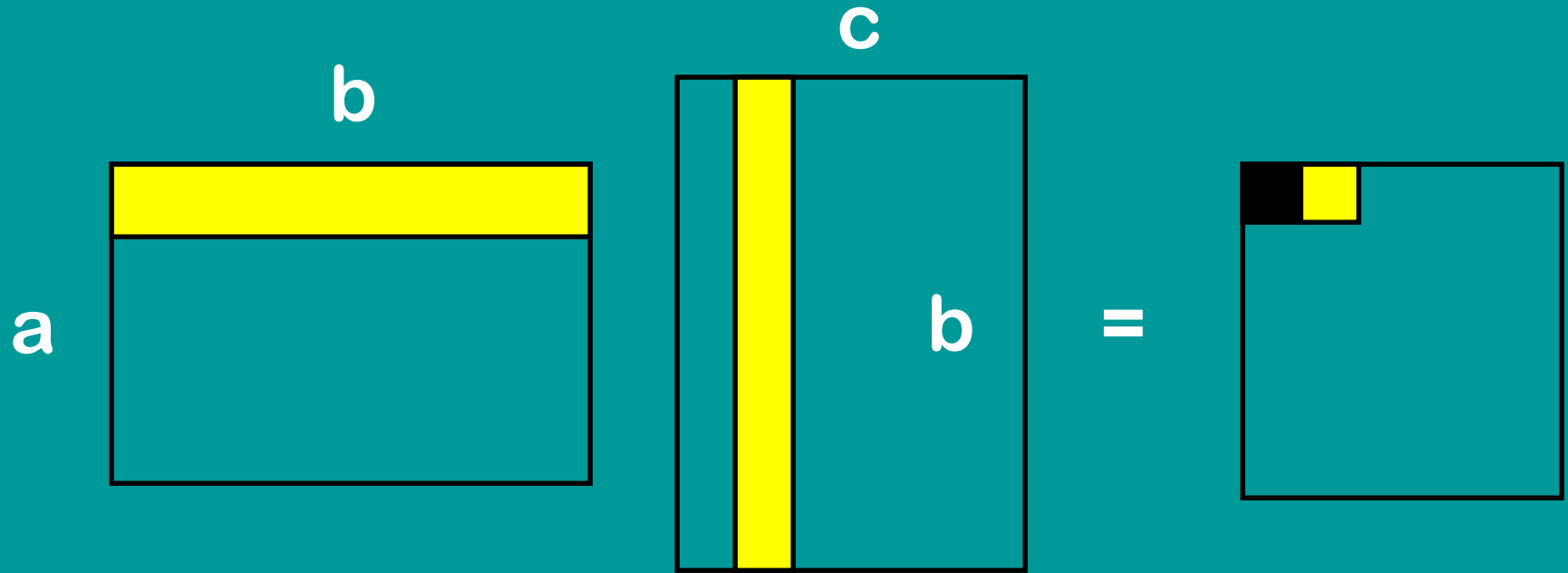
$B = b \times c$ matrix

How many operations to compute AB ?

Optimal matrix multiplication



Optimal matrix multiplication



each entry of $A * B$ takes $\Theta(b)$ time

need to compute ac entries $\Rightarrow \Theta(abc)$ time
total

Optimal matrix multiplication



A



B

Compute AB , time = ?

Optimal matrix multiplication



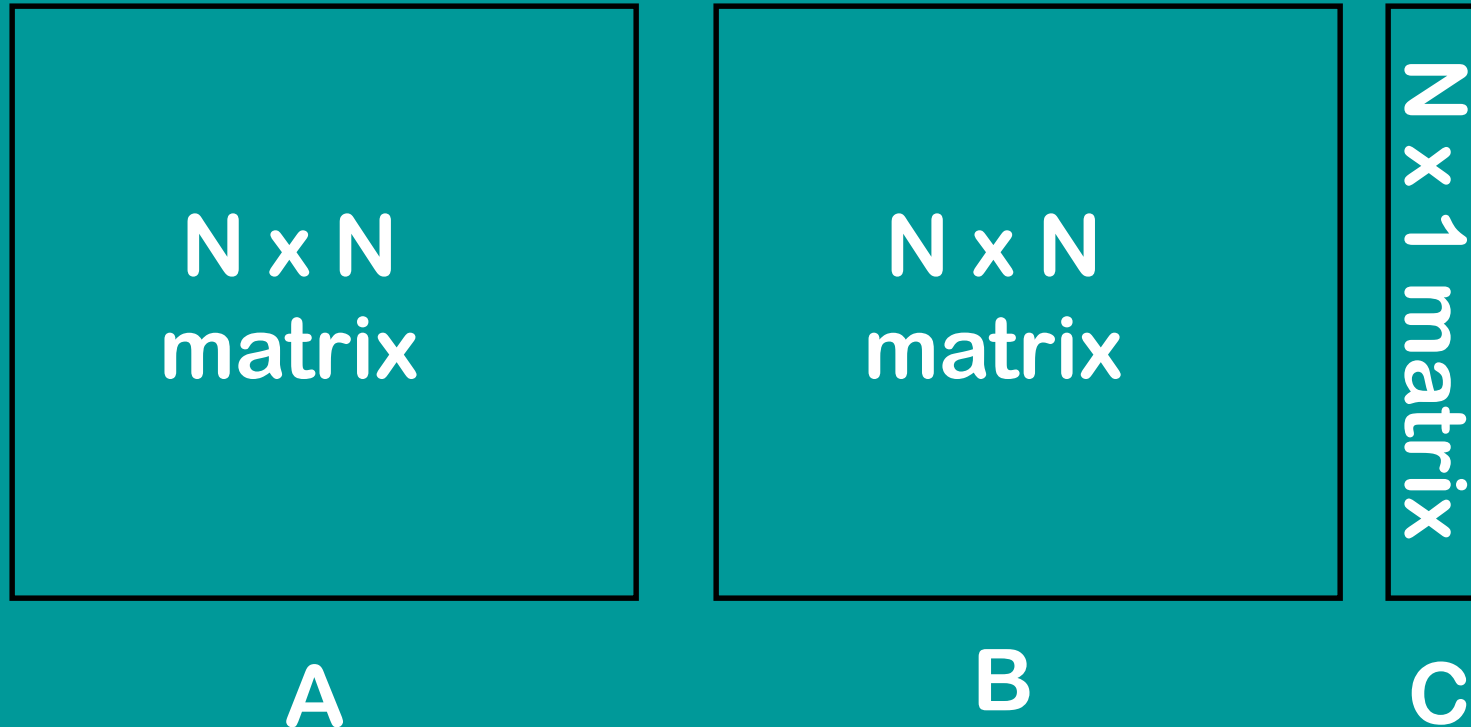
A



B

Compute AB , time = $\Theta(N^3)$

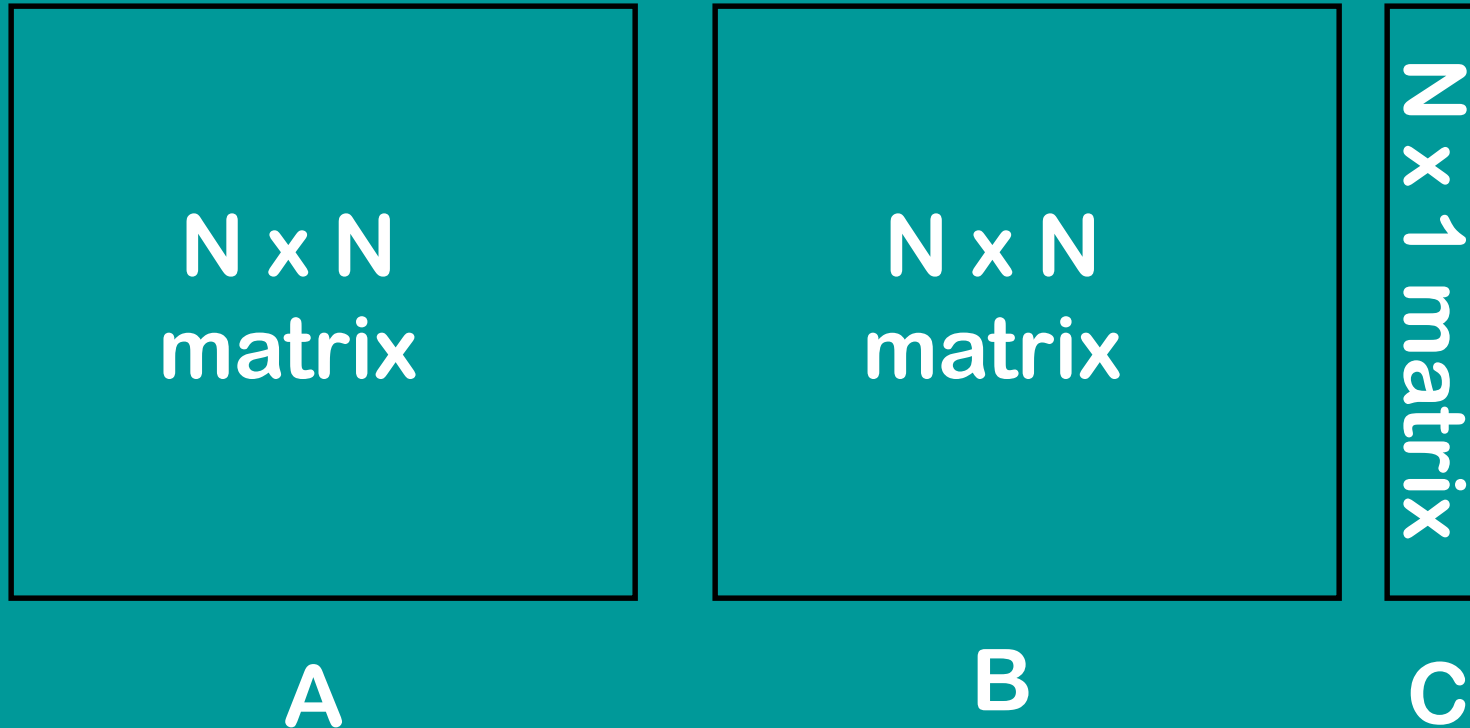
Optimal matrix multiplication



Compute AB , time = $\Theta(N^3)$

Compute ABC , time = ?

Optimal matrix multiplication

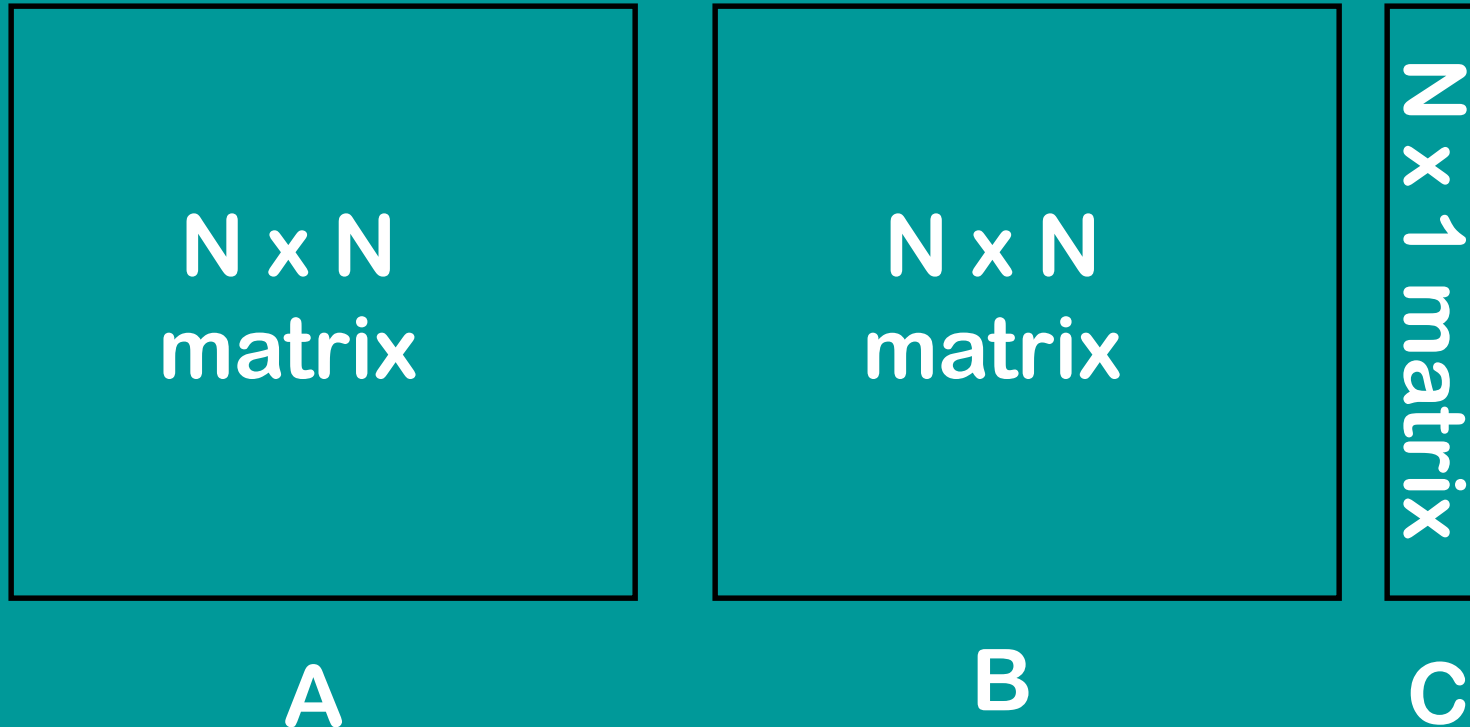


Compute $D=BC$, time = ?

Compute AD , time = ?

Compute ABC , time = ?

Optimal matrix multiplication

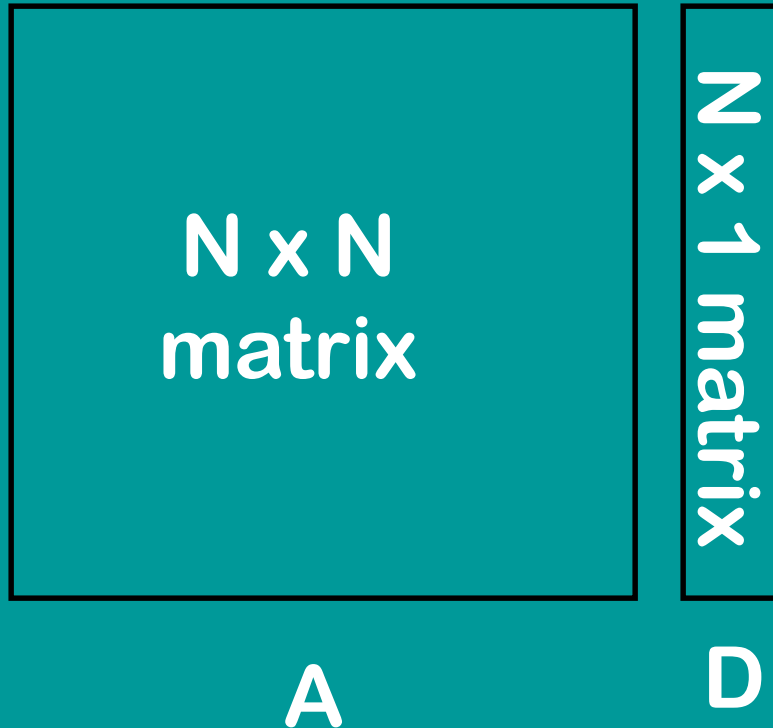


Compute $D=BC$, time = $\Theta(N^2)$

Compute AD , time = ?

Compute ABC , time = ?

Optimal matrix multiplication

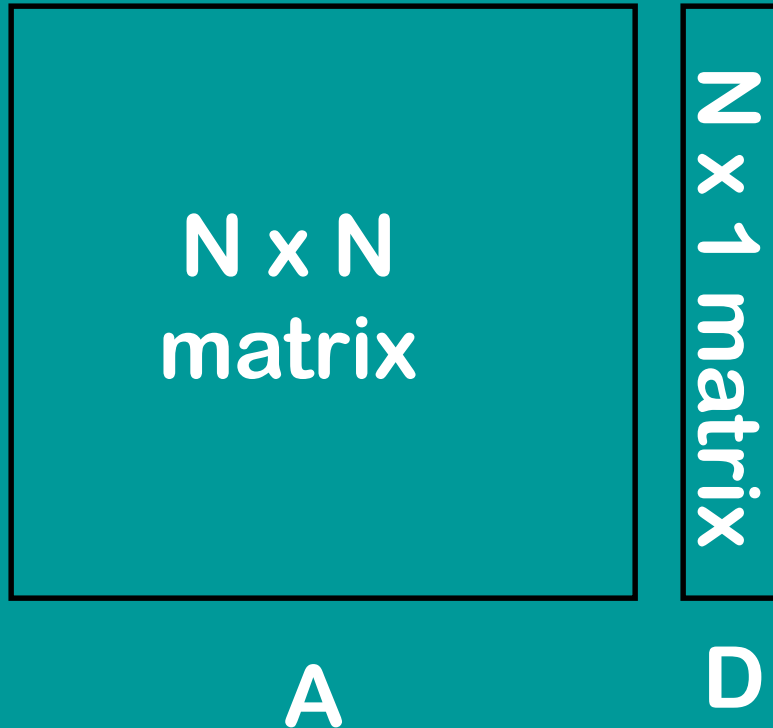


Compute $D=BC$, time = $\Theta(N^2)$

Compute AD , time = ?

Compute ABC , time = ?

Optimal matrix multiplication

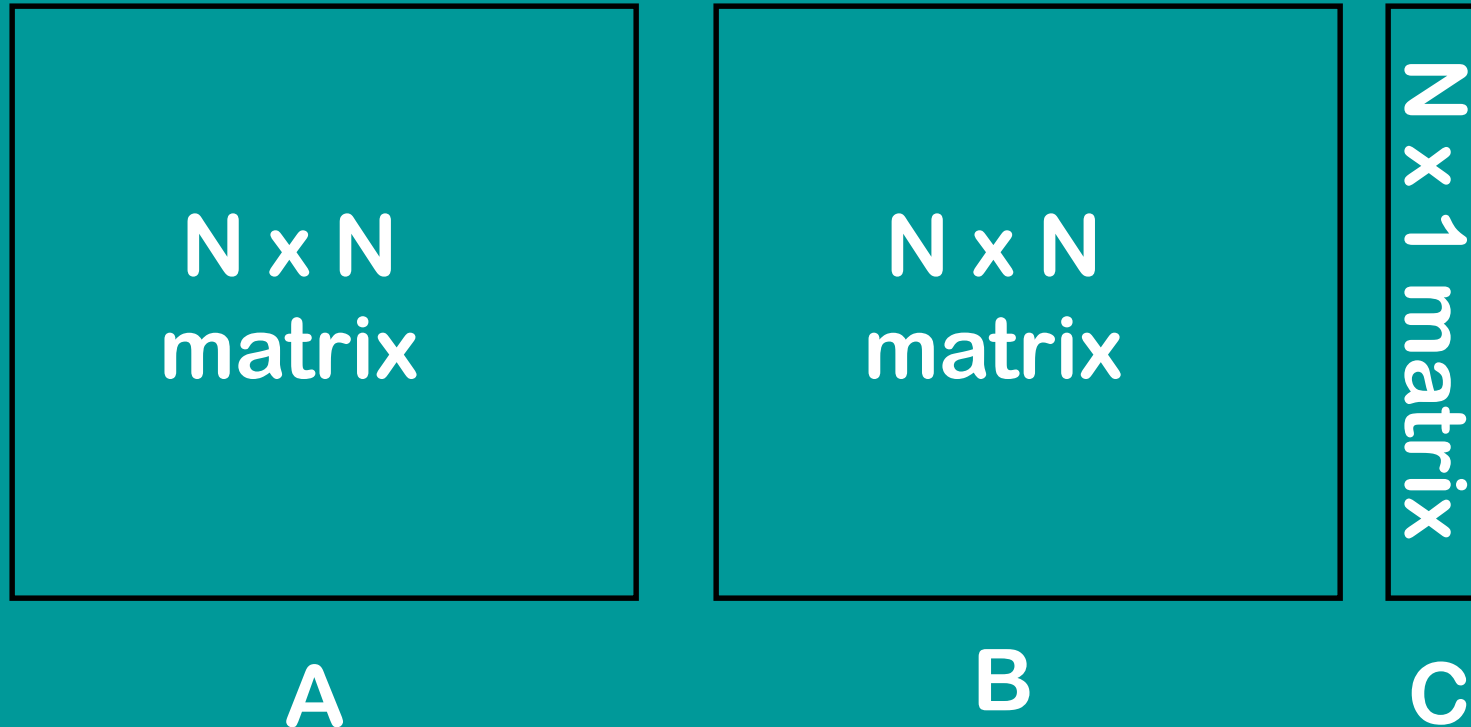


Compute $D=BC$, time = $\Theta(N^2)$

Compute AD , time = $\Theta(N^2)$

Compute ABC , time = $\Theta(N^2)$

Optimal matrix multiplication



$$(AB)C = ABC = A(BC)$$

The order of evaluation
does not change the result
can change the amount of work needed

Optimal matrix multiplication

$a_1, a_2, a_3, \dots, a_n$

$n-1$ matrices of sizes

$a_1 \times a_2$ B_1

$a_2 \times a_3$ B_2

$a_3 \times a_4$ B_3

\dots

$a_{n-1} \times a_n$ B_{n-1}

What order should we multiply them in?

Optimal matrix multiplication

$$B_1 B_2 B_3 B_4 \dots B_{n-1}$$

$$B_1 (B_2 B_3 B_4 \dots B_{n-1})$$

$$(B_1 B_2) (B_3 B_4 \dots B_{n-1})$$

$$(B_1 B_2 B_3) (B_4 \dots B_{n-1})$$

...

$$(B_1 B_2 B_3 B_4 \dots) B_{n-1}$$

Optimal matrix multiplication

$$B_1 B_2 B_3 B_4 \dots B_{n-1}$$

$K[i,j]$ = the minimal number of operations
needed to multiply $B_i \dots B_j$

| | |
|-------------------------------------|---|
| $B_i (B_{i+1} B_{i+2} \dots B_j)$ | $K[i,i] + K[i+1,j] + a_i a_{i+1} a_{j+1}$ |
| $(B_i B_{i+1}) (B_{i+2} \dots B_j)$ | $K[i,i+1] + K[i+2,j] + a_i a_{i+2} a_{j+1}$ |
| $(B_i B_{i+1} B_{i+2}) (\dots B_j)$ | $K[i,i+2] + K[i+3,j] + a_i a_{i+3} a_{j+1}$ |
| \dots | |
| $(B_1 B_2 B_3 \dots) B_j$ | $K[i,j-1] + K[j,j] + a_i a_j a_{j+1}$ |

Optimal matrix multiplication

$$B_1 B_2 B_3 B_4 \dots B_{n-1}$$

$K[i,j]$ = the minimal number of operations
needed to multiply $B_i \dots B_j$

$$K[i,i]=0$$

$$K[i,j] = \min_{i \leq w < j} K[i,w] + K[w+1,j] + a_i a_{w+1} a_j$$

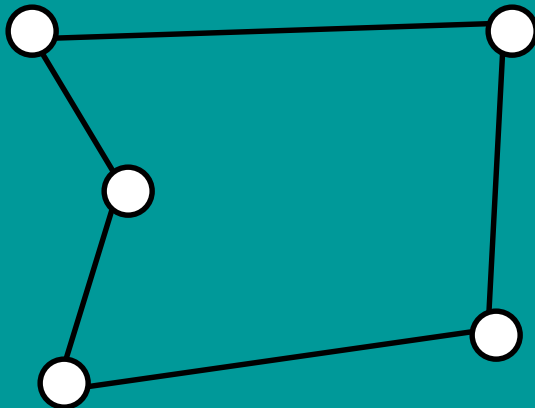
Travelling Salesman Problem

INPUT:

N cities, $N \times N$ symmetric matrix D ,
 $D(i,j)$ = distance between city i and j

OUTPUT:

the shortest tour visiting all the cities



Travelling Salesman Problem

Algorithm 1 – try all possibilities

for each permutation π of $\{1, \dots, n\}$
visit the cities in the order π ,
compute distance travelled,

pick the best solution

running time = ?

Travelling Salesman Problem

Algorithm 1 – try all possibilities

for each permutation π of $\{1, \dots, n\}$
visit the cities in the order π ,
compute distance travelled,

pick the best solution

running time $\approx n!$

is $(n+1)! = O(n!)$?

Travelling Salesman Problem

for each subset S of the cities with $|S| \geq 2$ and each $u, v \in S$

$K[S, u, v]$ the length of the shortest path that

- * starts at u
- * ends at v
- * visits all cities in S

How large is K ?

Travelling Salesman Problem

for each subset S of the cities with $|S| \geq 2$ and each $u, v \in S$

$K[S, u, v]$ the length of the shortest path that

- * starts at u

- * ends at v

- * visits all cities in S

How large is K ?

$$\approx 2^n n^2$$

Travelling Salesman Problem

$K[S, u, v]$

some vertex $w \in S - \{u, v\}$
must be visited first

$d(u, w)$ = we get to w

$K[S - u, w, v]$ = we need to get from w to v
and visit all vertices in $S - u$

Travelling Salesman Problem

$K[S,u,v]$ the length of the shortest path that

- * starts at u

- * ends at v

- * visits all cities in S

if $S=\{u,v\}$ then $K[S,u,v]=d(u,v)$

if $|S|>2$ then

$$K[S,u,v] = \min_{w \in S - \{u,v\}} K[S - u, w, v] + d(u, w)$$

Travelling Salesman Problem

if $S=\{u,v\}$ then $K[S,u,v]=d(u,v)$

if $|S|>2$ then

$$K[S,u,v] = \min_{w \in S - \{u,v\}} K[S - u, w, v] + d(u, w)$$

Running time = ?

$$K \approx 2^n n^2$$

Travelling Salesman Problem

if $S=\{u,v\}$ then $K[S,u,v]=d(u,v)$

if $|S|>2$ then

$$K[S,u,v] = \min_{w \in S - \{u,v\}} K[S - u, w, v] + d(u, w)$$

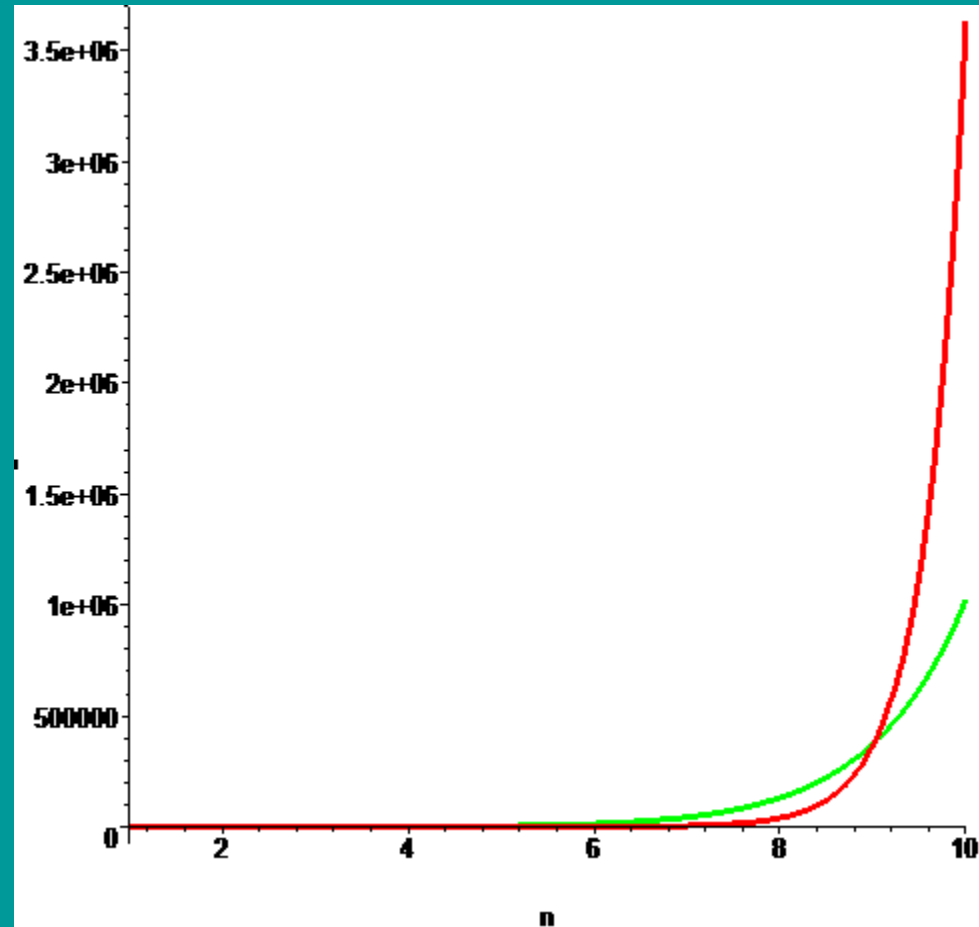
Running time = $O(n^3 2^n)$

$$K \approx 2^n n^2$$

Travelling Salesman Problem

dynamic programming = $O(n^3 2^n)$

brute force = $O(n!)$



Longest increasing subsequence

INPUT: numbers a_1, a_2, \dots, a_n

OUTPUT: longest increasing subsequence

1,9,2,4,7,5,6



1,9,2,4,7,5,6

Longest increasing subsequence

INPUT: numbers a_1, a_2, \dots, a_n

OUTPUT: longest increasing subsequence

reduce to a problem that we saw today

Longest increasing subsequence

INPUT: numbers a_1, a_2, \dots, a_n

OUTPUT: longest increasing subsequence

Longest increasing subsequence

INPUT: numbers a_1, a_2, \dots, a_n

OUTPUT: longest increasing subsequence

$K[0..n, 0..n]$

$K[i, j]$ = the minimum last element of an increasing sequence in a_1, \dots, a_i of length j (if no sequence $\Rightarrow \infty$)

Longest increasing subsequence

$K[0..n, 0..n]$

$K[i, j]$ = the minimum last element of an increasing sequence in a_1, \dots, a_i of length j (if no sequence $\Rightarrow \infty$)

true/false: $K[i, j] \leq K[i, j+1]$?

Longest increasing subsequence

$K[0..n, 0..n]$

$K[i, j]$ = the minimum last element of an increasing sequence in a_1, \dots, a_i of length j (if no sequence $\Rightarrow \infty$)

$K[0, j] = ?$ for $j \geq 1$

$K[0, 0] = ?$

Longest increasing subsequence

$K[0..n, 0..n]$

$K[i, j]$ = the minimum last element of an increasing sequence in a_1, \dots, a_i of length j (if no sequence $\Rightarrow \infty$)

$$K[0, j] = \infty \quad \text{for } j \geq 1$$

$$K[0, 0] = -\infty$$

Longest increasing subsequence

$K[0..n, 0..n]$

$K[i, j]$ = the minimum last element of an increasing sequence in a_1, \dots, a_i of length j (if no sequence $\Rightarrow \infty$)

$K[i, j] = ?$

Longest increasing subsequence

K[0..n,0..n]

$K[i,j]$ = the minimum last element of an increasing sequence in a_1, \dots, a_i of length j (if no sequence $\Rightarrow \infty$)

$$K[i,j] = a_i \quad \text{if} \quad \begin{array}{l} a_i < K[i-1,j] \\ \text{and} \\ a_i > K[i-1,j-1] \end{array}$$

K[i,j] = K[i-1,j] otherwise

Longest increasing subsequence

$K[0..n, 0..n]$

$K[i, j]$ = the minimum last element of an increasing sequence in a_1, \dots, a_i of length j (if no sequence $\Rightarrow \infty$)

$$K[i, 0] = -\infty$$

$$K[i, 1] =$$

$$K[i, 2] =$$

...

$$K[i, j] =$$

$$K[i, j+1] = \infty$$

$$a_i < K[i-1, j]$$

and

$$a_i > K[i-1, j-1]$$

Longest increasing subsequence

$$K[0,0] = -\infty$$

$$K[0,1] = \infty$$


$$K[0,2] = \infty$$

$$K[0,3] = \infty$$

$$K[0,4] = \infty$$

$$K[0,5] = \infty$$

$$K[0,6] = \infty$$

1,9,2,4,7,5,6



Longest increasing subsequence

$$K[1,0] = -\infty$$

$$K[1,1] = 1$$

$$K[1,2] = \infty$$

$$K[1,3] = \infty$$

$$K[1,4] = \infty$$

$$K[1,5] = \infty$$

$$K[1,6] = \infty$$

1,9,2,4,7,5,6



Longest increasing subsequence

$$K[1,0] = -\infty$$

$$K[1,1] = 1$$

$$K[1,2] = \infty$$

$$K[1,3] = \infty$$

$$K[1,4] = \infty$$

$$K[1,5] = \infty$$

$$K[1,6] = \infty$$

1,9,2,4,7,5,6



Longest increasing subsequence

$$K[2,0] = -\infty$$

$$K[2,1] = 1$$

$$K[2,2] = 9$$

$$K[2,3] = \infty$$

$$K[2,4] = \infty$$

$$K[2,5] = \infty$$

$$K[2,6] = \infty$$

1,9,2,4,7,5,6



Longest increasing subsequence

$$K[2,0] = -\infty$$

$$K[2,1] = 1$$



$$K[2,2] = 9$$

$$K[2,3] = \infty$$

$$K[2,4] = \infty$$

$$K[2,5] = \infty$$

$$K[2,6] = \infty$$

1,9,2,4,7,5,6



Longest increasing subsequence

$$K[3,0] = -\infty$$

$$K[3,1] = 1$$



$$K[3,2] = 2$$

$$K[3,3] = \infty$$

$$K[3,4] = \infty$$

$$K[3,5] = \infty$$

$$K[3,6] = \infty$$

1,9,2,4,7,5,6



Longest increasing subsequence

$$K[3,0] = -\infty$$

$$K[3,1] = 1$$

$$K[3,2] = 2$$

$$K[3,3] = \infty$$

$$K[3,4] = \infty$$

$$K[3,5] = \infty$$

$$K[3,6] = \infty$$



1,9,2,4,7,5,6



Longest increasing subsequence

$$K[4,0] = -\infty$$

$$K[4,1] = 1$$

$$K[4,2] = 2$$

$$K[4,3] = 4$$

$$K[4,4] = \infty$$

$$K[4,5] = \infty$$

$$K[4,6] = \infty$$



1,9,2,4,7,5,6



Longest increasing subsequence

$$K[4,0] = -\infty$$

$$K[4,1] = 1$$

$$K[4,2] = 2$$

$$K[4,3] = 4$$

$$K[4,4] = \infty$$

$$K[4,5] = \infty$$

$$K[4,6] = \infty$$



1,9,2,4,7,5,6



Longest increasing subsequence

$$K[5,0] = -\infty$$

$$K[5,1] = 1$$

$$K[5,2] = 2$$

$$K[5,3] = 4$$

$$K[5,4] = 7$$

$$K[5,5] = \infty$$

$$K[5,6] = \infty$$



1,9,2,4,7,5,6



Longest increasing subsequence

$$K[5,0] = -\infty$$

$$K[5,1] = 1$$

$$K[5,2] = 2$$

$$K[5,3] = 4$$

$$K[5,4] = 7$$

$$K[5,5] = \infty$$

$$K[5,6] = \infty$$



1,9,2,4,7,5,6



Longest increasing subsequence

$$K[6,0] = -\infty$$

$$K[6,1] = 1$$

$$K[6,2] = 2$$

$$K[6,3] = 4$$

$$K[6,4] = 5$$

$$K[6,5] = \infty$$

$$K[6,6] = \infty$$



1,9,2,4,7,5,6



Longest increasing subsequence

$$K[6,0] = -\infty$$

$$K[6,1] = 1$$

$$K[6,2] = 2$$

$$K[6,3] = 4$$

$$K[6,4] = 5$$

$$K[6,5] = \infty$$

$$K[6,6] = \infty$$



1,9,2,4,7,5,6



Longest increasing subsequence

$$K[7,0] = -\infty$$

$$K[7,1] = 1$$

$$K[7,2] = 2$$

$$K[7,3] = 4$$

$$K[7,4] = 5$$

$$K[7,5] = 6$$

$$K[7,6] = \infty$$

answer = 5

1,9,2,4,7,5,6