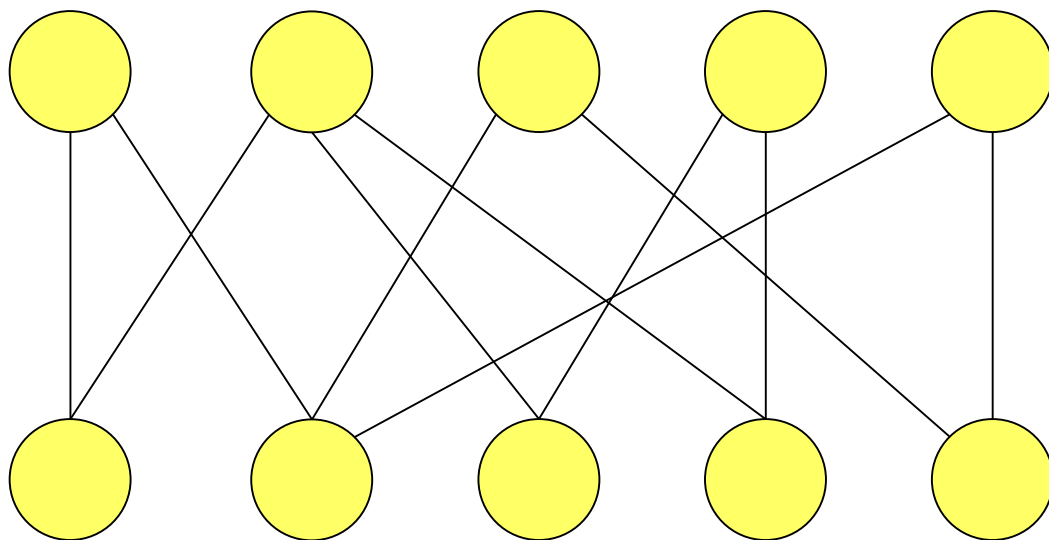
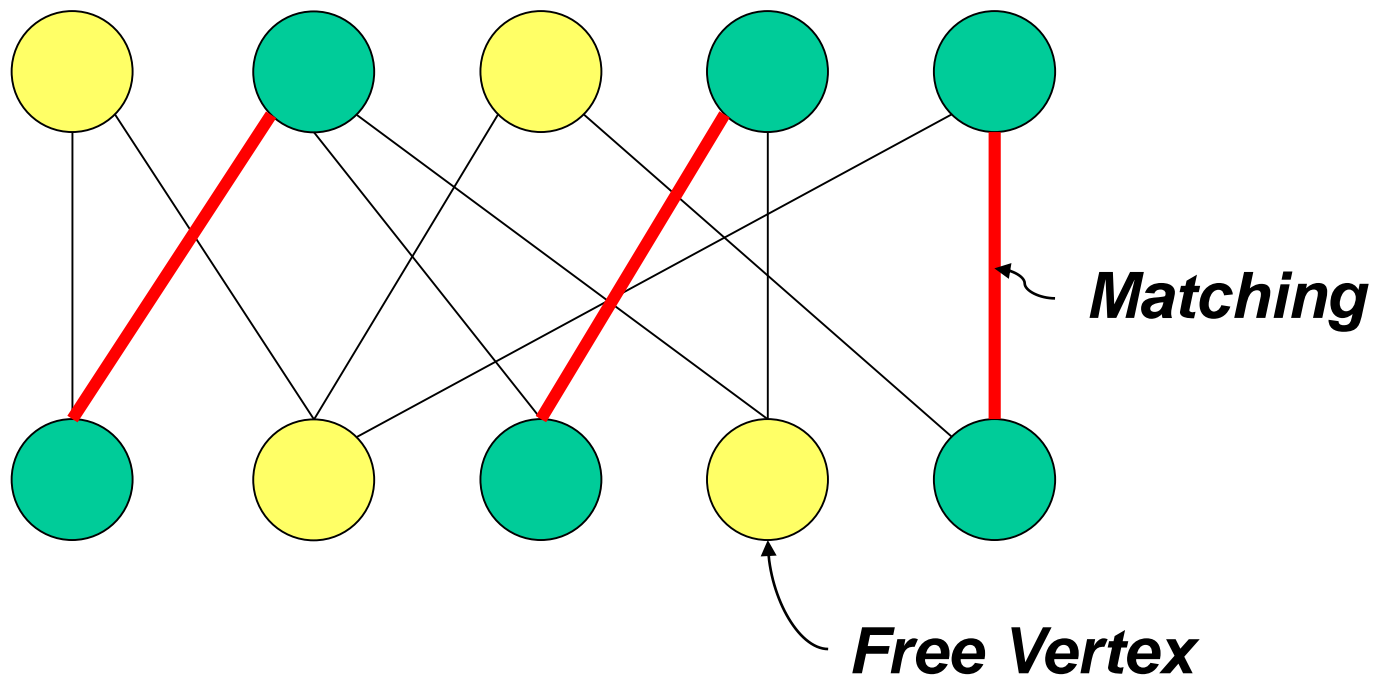


# Maximum Bipartite Matching

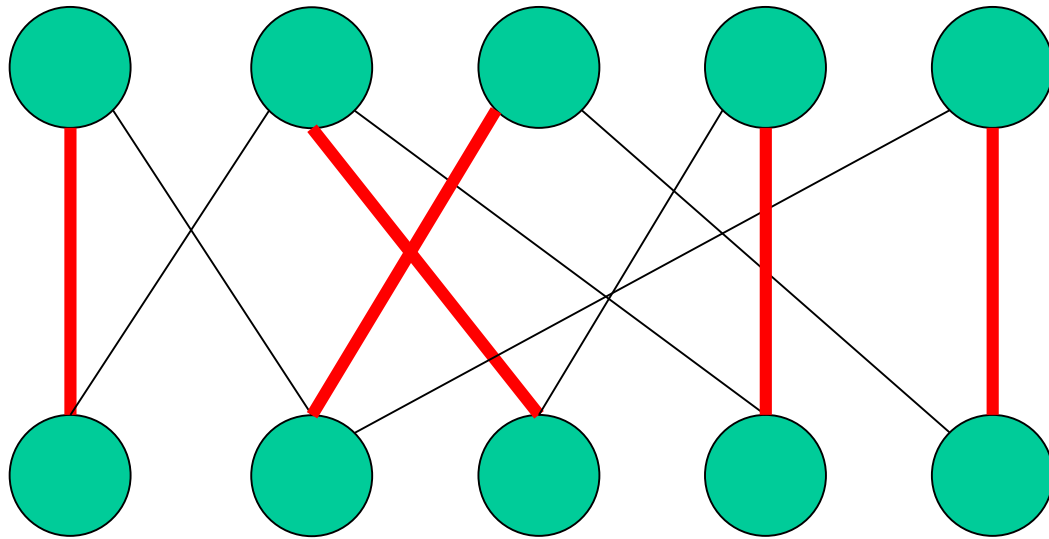
# Bipartite Matching



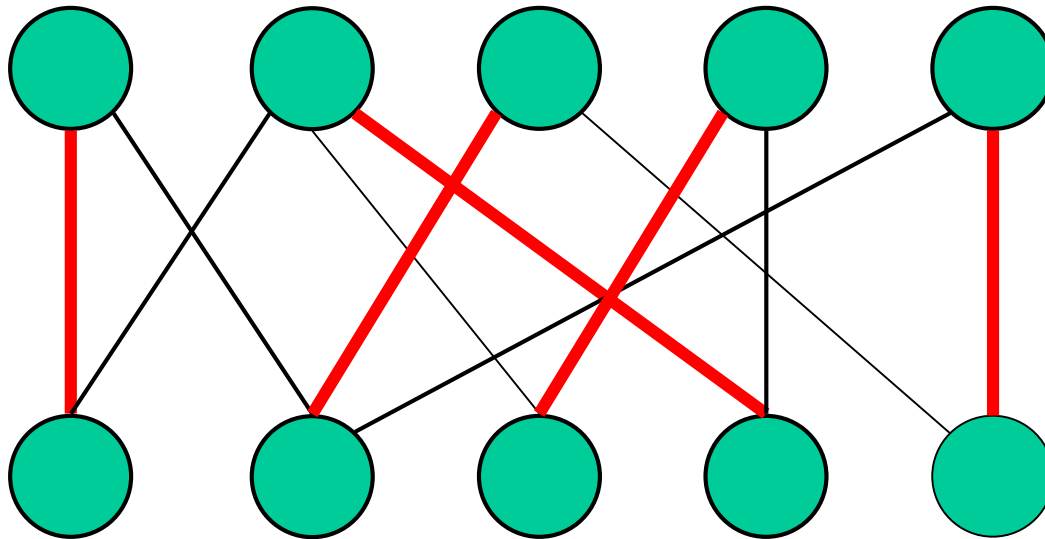
# Matching



- Maximum Matching: matching with the largest number of edges



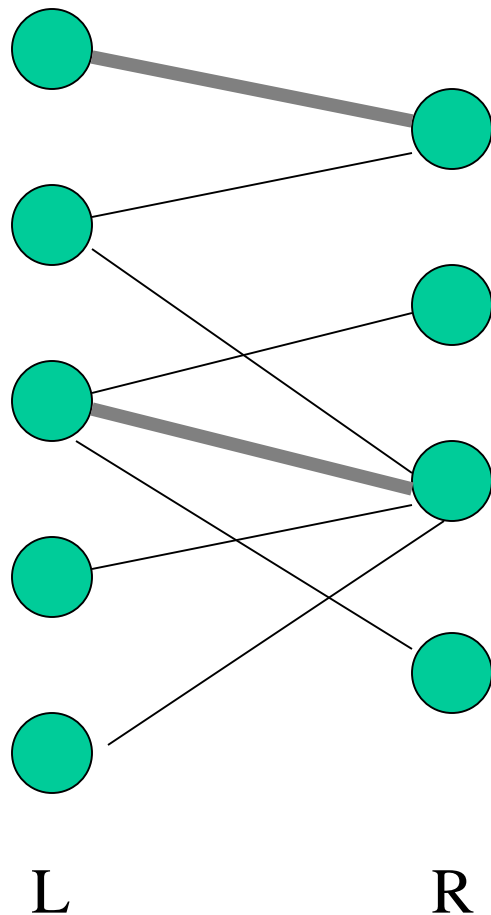
- Note that maximum matching is not unique.



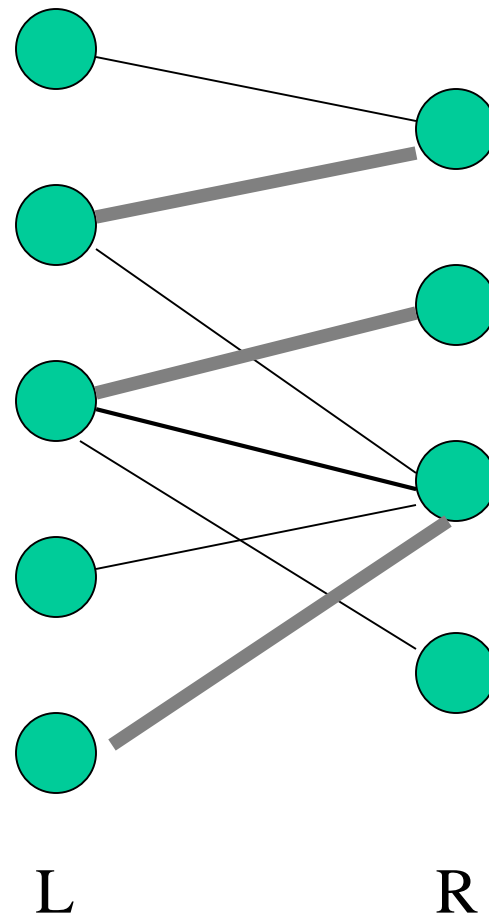
# Maximum bipartite matching:

- Bipartite graph: a graph  $(V, E)$ , where  $V=L\cup R$ ,  $L\cap R=\text{empty}$ , and for every  $(u, v)\in E$ ,  $u \in L$  and  $v \in R$ .
- Given an undirected graph  $G=(V,E)$ , a **matching** is a subset of edges  $M\subseteq E$  such that for all vertices  $v\in V$ , at most one edge of  $M$  is incident on  $v$ . We say that a vertex  $v \in V$  is **matched** by matching  $M$  if some edge in  $M$  is incident on  $v$ ; otherwise,  $v$  is **unmatched**. A **maximum matching** is a matching of maximum cardinality, that is, a matching  $M$  such that for any matching  $M'$ , we have

$$|M| \geq |M'|$$



(a)



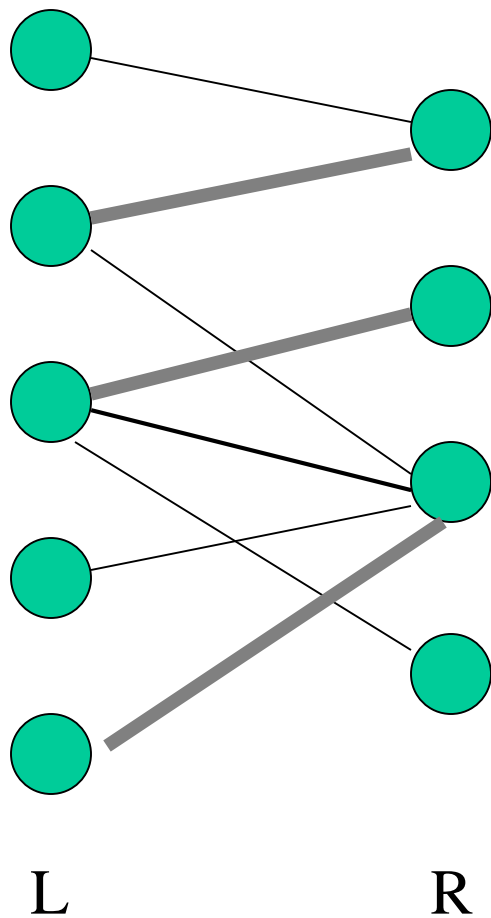
(b)

A bipartite graph  $G=(V,E)$  with vertex partition  $V=L\cup R$ .(a) A matching with cardinality 2.(b) A maximum matching with cardinality 3.

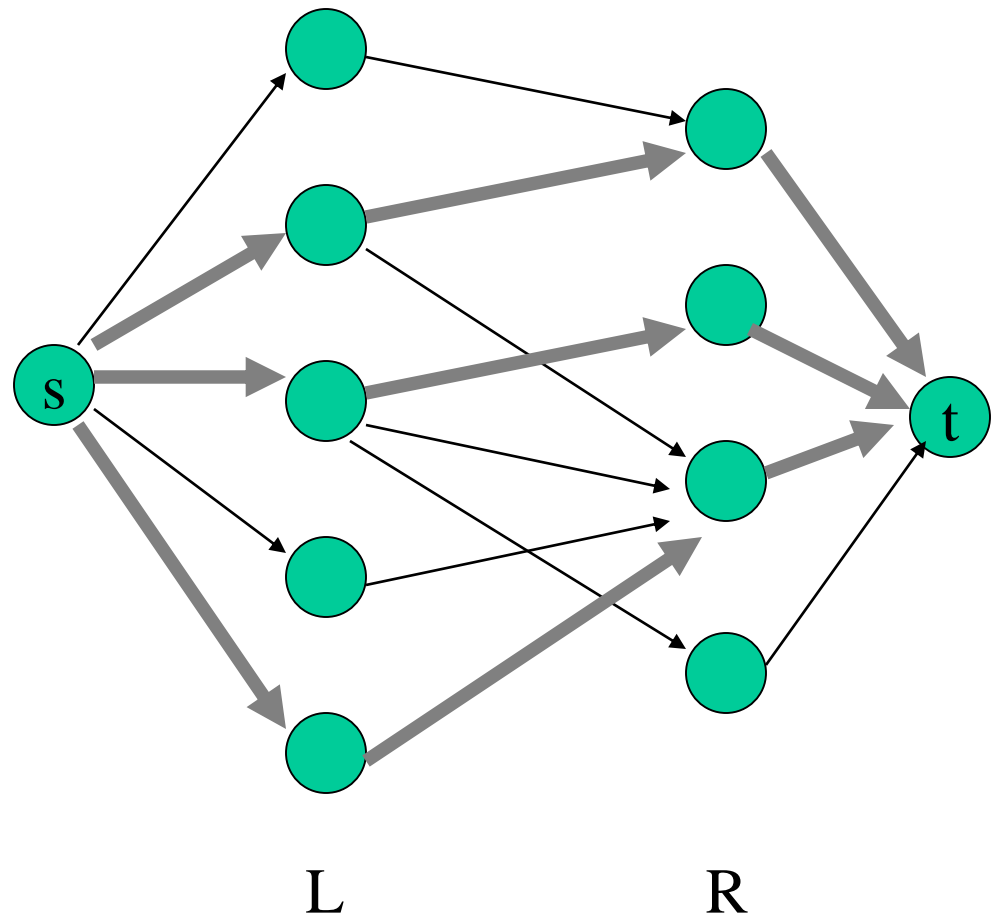
# Finding a maximum bipartite matching:

- We define the **corresponding flow network**  $G'=(V',E')$  for the bipartite graph  $G$  as follows. Let the source  $s$  and sink  $t$  be new vertices not in  $V$ , and let  $V'=V\cup\{s,t\}$ . If the vertex partition of  $G$  is  $V=L\cup R$ , the directed edges of  $G'$  are given by  $E'=\{(s,u):u\in L\}\cup\{(u,v):u\in L,v\in R,\text{ and } (u,v)\in E\}\cup\{(v,t):v\in R\}$ . Finally, we assign unit capacity to each edge in  $E'$ .
- We will show that a matching in  $G$  corresponds directly to a flow in  $G'$ 's corresponding flow network  $G'$ . We say that a flow  $f$  on a flow network  $G=(V,E)$  is **integer-valued** if  $f(u,v)$  is an integer for all  $(u,v)\in V\times V$ .





(a)



(b)

(a) The bipartite graph  $G=(V,E)$  with vertex partition  $V=L \cup R$ . A maximum matching is shown by shaded edges. (b) The corresponding flow network. Each edge has unit capacity. Shaded edges have a flow of 1, and all other edges carry no flow.

# Continue:

- Lemma .
- Let  $G=(V,E)$  be a bipartite graph with vertex partition  $V=L\cup R$ , and let  $G'=(V',E')$  be its corresponding flow network. If  $M$  is a matching in  $G$ , then there is an integer-valued flow  $f$  in  $G'$  with value  $|f|=|M|$ . Conversely, if  $f$  is an integer-valued flow in  $G'$ , then there is a matching  $M$  in  $G$  with cardinality  $|M|=|f|$ .
- Reason: The edges incident to  $s$  and  $t$  ensures this.
  - Each node in the first column has in-degree 1
  - Each node in the second column has out-degree 1.
  - So each node in the bipartite graph can be involved once in the flow.