

Array / String
Fixed size window (window
size)

Computation return

maximum / minimum
subarray / substring

Max Sum Subarray of size K

Basic Accuracy: 49.6% Submissions: 55120 Points: 1

Given an array of integers Arr of size N and a number K . Return the maximum sum of a subarray of size K .

Example 1:

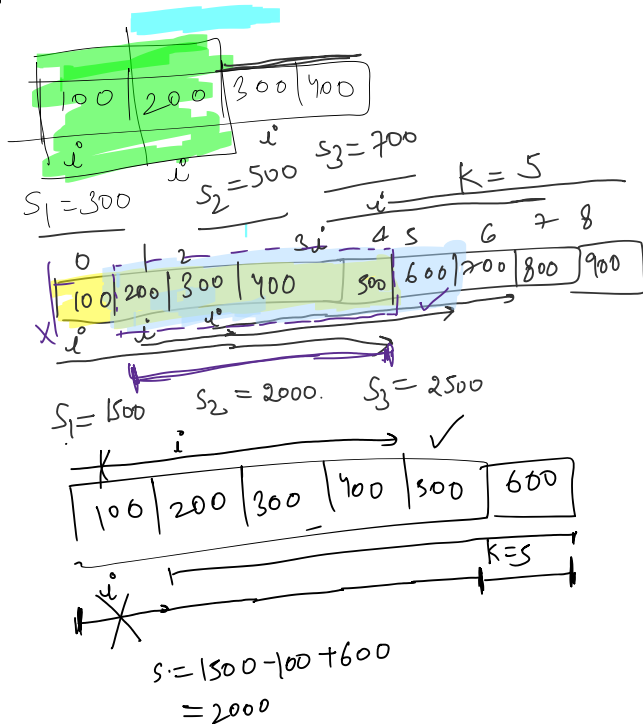
Input:
 $N = 4, K = 2$
 $Arr = [100, 200, 300, 400]$
Output:
700
Explanation:
 $Arr_3 + Arr_4 = 700$,
which is maximum.

Array (N)

100	200	300	400
-----	-----	-----	-----

$K = 2$

maximum sum



Array / String
subArray / substring

Array / String
SubArray / SubString
K / window size
99%

$k=3$

0	1	2	3	4
100	200	300	400	500

$\underline{\underline{x}}$ $\underline{\underline{x}}$ $\underline{\underline{i}}$

Normal

T.C - $O(N \times K)$ $\rightarrow O(N^2)$

100	200	300	400	50
-----	-----	-----	-----	----

~~$\underline{\underline{x}}$ $\underline{\underline{x}}$~~

T.C - $O(N)$

Answer compute

$k=2$ $N=4$

$\underline{\underline{x}}$	$\underline{\underline{ws}}$	$\underline{\underline{we}}$		
100	200	300	400	
0	1	2	3	

$ws \rightarrow$ window start
 $we \rightarrow$ window end

$\star \underline{we - ws + 1} \rightarrow \underline{\underline{window_length}}$

$we = 0$

$ws = 0$

$$\text{window_length} = 0 - 0 + 1 = \underline{\underline{1}}$$

$$= 1 - 0 + 1 = \underline{\underline{2}}$$

$$ws_size = we - ws + 1 = \underline{\underline{1}}$$

$$ws_size = we - ws + 1 = \underline{\underline{2}}$$

// $ws = 0$ $\int \text{maxi} = \text{INT_MIN}$
// $we = 0$ $\int \text{sum} = 0$

while () {

$\text{sum} += \underline{\underline{A[we]}}; //$

if ($we - ws + 1 < k$) {
 $we++;$

} else if ($we - ws + 1 == k$) {
// compute output
// window slide
}

$\text{sum} = 0$

$\underline{\underline{x}}$	$\underline{\underline{ws}}$	$\underline{\underline{we}}$	$\underline{\underline{2}}$	
100	200	300	400	
0x	1	2	3	

$$\begin{aligned} \text{sum} &= 100 + 200 \\ &= 300 \\ k &= 2 \end{aligned}$$

$$\begin{aligned} &0 - 0 + 1 \\ &1 - 0 + 1 = \underline{\underline{2}} \end{aligned}$$

// output compute
 $\text{maxi} = \max(\text{maxi}, \text{sum});$

// slide the window
 $\text{sum} -= A[\text{ws}];$

$\text{ws}++;$

$\text{we}++;$

ws	ws	ws	ws	we
100	200	300	400	
0	1	2	3	

$$\text{sum} = 100 + 200 + 300 + 400 = 1000$$

while ($\text{we} < N$)
 $\text{sum} += A[\text{we}];$
 $\text{we}++;$

expnd // if ($\text{we} - \text{ws} + 1 < k$)
 $\text{we}++;$

else if ($\text{we} - \text{ws} + 1 == k$)
 $\text{maxi} = \max(\text{maxi}, \text{sum});$

$\text{sum} -= A[\text{ws}];$

$\text{ws}++;$

$\text{we}++;$

ws	ws	ws	ws	ws	ws	ws	we
100	200	300	400	500	600	100	
0	1	2	3	4	5	6	

// Initial $\text{ws} = 0, \text{we} = 0$
 $0 - 0 + 1 = 1$
 $1 - 0 + 1 = 2$
 $\text{sum} = 100 + 200 + 300 + 400 + 500 + 600 + 100 = 2200$
 $\text{sum} = 2200 - 100 = 2100$
 $\text{sum} = 2100 - 200 = 1900$
 $\text{sum} = 1900 - 300 = 1600$
 $\text{sum} = 1600 - 400 = 1200$
 $\text{sum} = 1200 - 500 = 700$

```
int we = 0, ws = 0;
while(we < N){
    // cumulative computation
    if(we - ws + 1 < K){
        // window size of aage badjhana h
    }
    else if(we - ws + 1 == K){
        // output compute
        // window slide
    }
}
```

```

long maximumSumSubarray(int K, vector<int> &arr, int N){
    int ws = 0, we = 0;
    long maxi = INT_MIN;
    long sum = 0;
    while(we < N){
        // cumulative computation
        sum += arr[we];
        if(we - ws + 1 < K){
            // window size of aage badhna h
            we++;
        }
        else if(we - ws + 1 == K){
            // output compute
            maxi = max(maxi, sum);
            // window slide
            sum -= arr[ws];
            ws++;
            we++;
        }
    }
    return maxi;
}

```

$T.C - O(N)$
 $S.C - O(1)$

1876. Substrings of Size Three with Distinct Characters

A string is **good** if there are no repeated characters.

Given a string `s`, return the number of **good substrings** of length **three** in `s`.

Note that if there are multiple occurrences of the same substring, every occurrence should be counted.

A **substring** is a contiguous sequence of characters in a string.

Example 1:

Input: `s = "xyzzaz"`
 Output: 1
 Explanation: There are 4 substrings of size 3: "xyz", "yzz", "zza", and "zaz".
 The only good substring of length 3 is "xyz".

sliding window

String Substring K=3

xyzzaaz

1

~~aa~~ ~~ab~~ ~~ba~~ ~~ab~~ ~~bc~~ ~~ca~~ ~~bc~~

~~1~~
~~2~~
~~3~~
4

K=3

~~ws~~ ~~ws~~ ~~ws~~ ~~ws~~ ~~ws~~ ~~ws~~ ~~ws~~ ~~ws~~

~~we~~ ~~we~~ ~~we~~ ~~we~~ ~~we~~ ~~we~~ ~~we~~ ~~we~~

~~a~~ ~~a~~ ~~b~~ ~~a~~ ~~b~~ ~~c~~ ~~a~~ ~~b~~ ~~c~~

0 1 2 3 4 5 6 7 8

0-0+1=1 1-0+1=2 2-0+1=3

```

int ws = 0, we = 0;
while(we < N){
    // cumulative computation
    if(we - ws + 1 < K){
        // window size of aage badhna h
        we++;
    }
    else if(we - ws + 1 == K){
        // output compute
        // window slide
    }
}

```

char count

~~a: 0~~
~~b: 0~~
~~c: 0~~

b: 1
c: 1

count = ~~1~~ ~~2~~ ~~3~~ 4

map

Give me count 0 1 2 3 4 5 6 7 8

delete the character

```

int countGoodSubstrings(string s) {
    int ws = 0, we = 0, n = s.size();
}

```

```

int countGoodSubstrings(string s) {
    int ws = 0, we = 0, n = s.size();
    unordered_map<char, int> counter;
    int count = 0;
    while(ws < n){
        counter[s[ws]]++;
        if(we - ws + 1 < 3){
            we++;
        }
        else if(we - ws + 1 == 3){
            // answer found
            if(counter.size() == 3){
                count++;
            }
            // window slide
            // remove window start value part
            counter[s[ws]]--;
            if(counter[s[ws]] == 0){
                counter.erase(s[ws]);
            }
            ws++;
            we++;
        }
    }
    return count;
}

```

T.C - $O(N)$
S.C - $O(K)$

2379. Minimum Recolors to Get K Consecutive Black Blocks

Easy 384 12 12 12 12

Companies

You are given a 0-indexed string `blocks` of length `n`, where `blocks[i]` is either `'W'` or `'B'`, representing the color of the `i`th block. The characters `'W'` and `'B'` denote the colors white and black, respectively.

You are also given an integer `k`, which is the desired number of **consecutive** black blocks.

In one operation, you can **recolor** a white block such that it becomes a black block.

Return the **minimum** number of operations needed such that there is at least **one** occurrence of `k` consecutive black blocks.

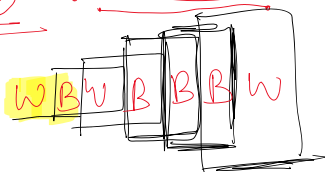
Example 1:

Input: `blocks = "WBBWWBBWBW"`, `k = 7`
Output: 3
Explanation:
 One way to achieve 7 consecutive black blocks is to recolor the 0th, 3rd, and 4th blocks so that `blocks = "BBBBBBBWBW"`.
 It can be shown that there is no way to achieve 7 consecutive black blocks in less than 3 operations.
 Therefore, we return 3.

T.C - $O(N)$
S.C - $O(1)$

blocks = "WBBWWBBWBW" k = 7

W.C = 3



k = 2

438. Find All Anagrams in a String

Medium 9.3K 287 12 12 12 12

Companies

Given two strings `s` and `p`, return an array of all the start indices of `p`'s anagrams in `s`. You may return the answer in **any order**.

An **Anagram** is a word or phrase formed by rearranging the letters of a different word or phrase, typically using all the original letters exactly once.

Example 1:

Input: `s = "cbaebabacd"`, `p = "abc"`
Output: `[0,6]`
Explanation:
 The substring with start index = 0 is "cba", which is an anagram of "abc".
 The substring with start index = 6 is "bac", which is an anagram of "abc".

Example 2:

Input: `s = "abab"`, `p = "ab"`
Output: `[0,1,2]`
Explanation:
 The substring with start index = 0 is "ab", which is an anagram of "ab".
 The substring with start index = 1 is "ba", which is an anagram of "ab".

Anagram

fat \Leftrightarrow atf \Leftrightarrow taf

p & anagram s at start starting index

c d a e b a b a c d a b c
 0 1 2 3 4 5 6 7 8 9
 [6]

a b a b
 0 1 2 3

a b
 [0, 1, 2]

"abc"
 window size = p.size()
 count = 3

pattern

a	:	1
b	:	1
c	:	1

we we we
 c b a e b a b a c d
 0 1 2 3 4 5 6 7 8 9

cumulative map

c	:	1
b	:	1
a	:	1

we we we we
 a b c a z

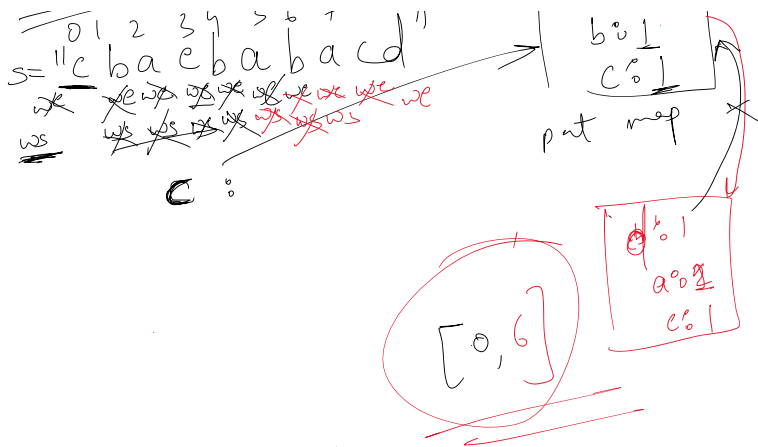
(4) p = abc
 count = 3

a	:	0
b	:	0
c	:	0

 [0]

HashMap travel k=3
 0 1 2 3 4 5 6 7 8 9
 s = "c b a e b a b a c d"
 we we we we we

a	:	1
b	:	1
c	:	1



// pattern
 k=3 count=3
 0 1 2 3 4 5 6 7 8
 " c b a c b a b a c d "
~~we~~ ~~we~~ ~~we~~ ~~we~~ ~~we~~ ~~we~~ ~~we~~ ~~we~~
 count = ~~we~~ ~~we~~ ~~we~~ ~~we~~ ~~we~~ ~~we~~ ~~we~~ ~~we~~
 // frequent decr can
 // frequency is 0
 count--
 *2
 [0, 5]

// hide
 ws \rightarrow map present
 // frequency is 0
 count++
 frequency++
 s[we] \rightarrow map

frequency--
 if frequency == 0
 count--
 if count == 0
 ws \rightarrow array

ws \rightarrow map
 if (frequency == 0)

if (frequency == 0)
 count++
frequency++

Window in array
 we use a, b, c, d

c = ~~2~~ ~~1~~ ~~1~~
~~a: 0~~
~~b: 0~~
~~c: 0~~

M
 T.C - $O(M+N)$
 S.C - $O(M)$

T.C - $O(M+N \times S \times M)$
 S.C - $O(M)$

1 7 4 9 2 5

1	7	6
1	4	3
1	9	8
1	2	1
1	5	4

7	4	-3
7	9	2
7	2	-5
7	5	-2

4	9	5
4	2	-2
4	5	1

9	2	-7
9	5	-4

2	5	3
---	---	---

1 → 7 → 4 → 9 → 2 → 5

1 17 5 10 13 15 10 5 16 8
 0 1 2 3 4 5 6 7 8 9

1	17	16
1	5	4
1	10	9
1	13	12
1	15	14
1	10	9
1	5	4
1	16	15
1	8	7

17	5	-ive
17	10	-ive
17	13	-ive
17	15	-ive
14	10	-ive
17	5	-ive
17	16	-ive
17	8	-ive

5	10	time
5	13	time
5	15	time
5	10	time
5	5	time
5	16	time
5	8	time

10	13	+ive
10	15	+ive
10	10	-ive
10	5	-ive
10	16	time
10	8	-ive

13	15	time
13	10	-ive
13	5	-ive
13	16	time
13	8	-ive

15	10	-ive
15	5	-ive
15	16	time
15	8	-ive

10	5	-ive
10	16	time
10	8	-ive

5	16	time
5	8	-ive

16	8	-ive
----	---	------

1 → 17 → 5 → 10 → 5 → 16 → 8



1 | 17 | 5 | 10 | 13 | 15 | 10 | 5 | 16 | 8

bottom to top

s d

1	17	+
1	5	+
1	10	+
1	13	+
1	15	+
1	10	+
1	5	+
1	16	+
1	8	+

17	5	-
17	10	-
17	13	-
17	15	-
17	10	-
17	5	-
17	16	-
17	8	-

5	10	+
5	13	+
5	15	+
5	10	+
5	5	X
5	16	+
5	8	+

10	13	+
10	15	+
10	10	X
10	5	-
10	16	+
10	8	-

13	15	+
13	10	-
13	5	-
13	16	+
13	8	-

15	10	-
15	5	-
15	16	+
15	8	-

10	5	-
10	16	+
10	8	-

5	16	+
5	8	+

16	8	-
----	---	---

1 → 17 → 5 → 10 → 5 → 16 → 8

1 | 17 | 5 | 10 | 13 | 15 | 10 | 5 | 16 | 8

top to bottom

8	16	+
8	5	-
8	10	+
8	13	+
8	15	+
8	10	+
8	5	-
8	17	+
8	1	-

16	5	-
16	10	-
16	13	-
16	15	-
16	10	-
16	5	-
16	17	+
16	1	-

5	10	+
5	13	+
5	15	+
5	10	+
5	5	X
5	17	+
5	1	-

10	15	+
10	13	+
10	10	X
10	5	-
10	17	+
10	1	-

15	13	-
15	10	-
15	5	-
15	17	+
15	1	-

13	10	-
13	5	-
13	17	+
13	1	-

10	5	-
10	17	+
10	1	-

5	17	+
5	1	-

17	1	-
----	---	---

8 → 16 → 5 → 10 → 5 → 17 → 1

8 → 5 → 10 → 5 → 17 → 1

i

	0	1	2
0	10	40	70
1	20	50	80
2	30	60	90