

Problem Statement

There are N stones, numbered $1, 2, \dots, N$. For each i ($1 \leq i \leq N$), the height of Stone i is h_i .

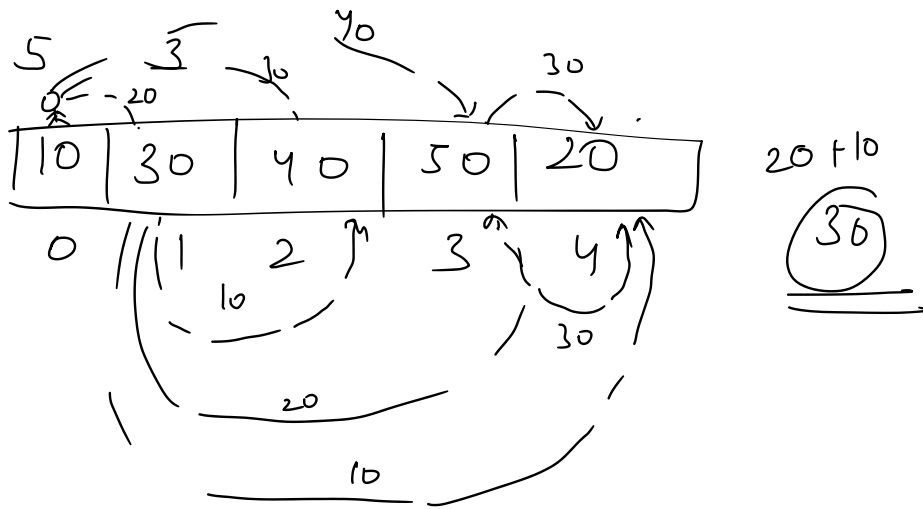
There is a frog who is initially on Stone 1. He will repeat the following action some number of times to reach Stone N :

- If the frog is currently on Stone i , jump to one of the following: Stone $i + 1, i + 2, \dots, i + K$. Here, a cost of $|h_i - h_j|$ is incurred, where j is the stone to land on.

Find the minimum possible total cost incurred before the frog reaches Stone N .

Constraints

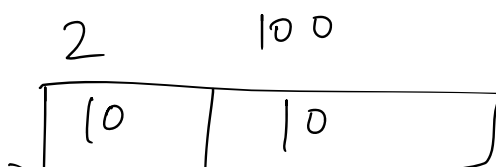
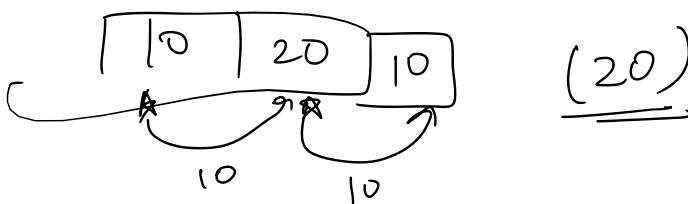
- All values in input are integers.
- $2 \leq N \leq 10^5$
- $1 \leq K \leq 100$
- $1 \leq h_i \leq 10^4$

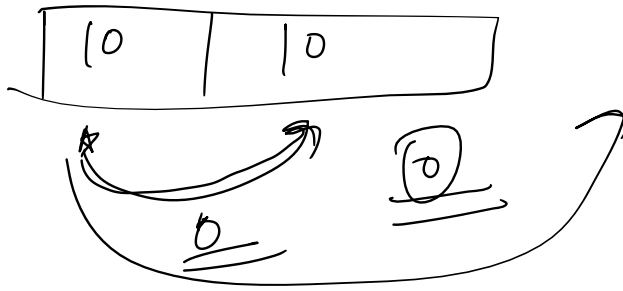


$1 \rightarrow 2 \rightarrow 5$ (30) ★ ★

$1 \rightarrow 2 \rightarrow 4 \rightarrow 5$ (20)

3 1

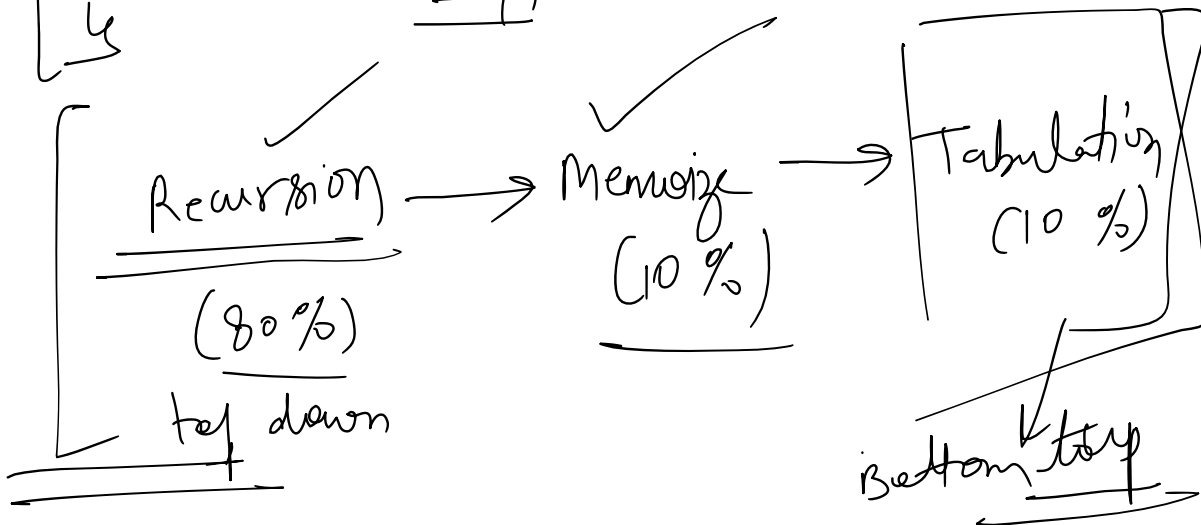




```

int solve(int i, vector<int> cost,
          int k)
{
    if (i == 0) return 0;
    // choice
    int mini = INT_MAX;
    for (int jump = 1; jump <= k; jump++) {
        if (i - jump >= 0) {
            mini = min(mini, solve(i - jump, cost, k)
                        + abs(cost[i]
                              - cost[i - jump]));
        }
    }
    return mini;
}

```



```

using namespace std;
int solve(int i, int k, vector<int> &cost, vector<int> &dp){
    if(i == 0) return 0;
    if(dp[i] != -1){
        return dp[i];
    }
    int mini = INT_MAX;
    for(int p = 1; p <= k; p++){
        if(i - p >= 0){
            mini = min(mini, solve(i - p, k, cost, dp) + abs(cost[i] - cost[i - p]));
        }
    }
    return dp[i] = mini;
}

int main(){
    int n, k;
    cin >> n >> k;
    vector<int> cost(n);
    for(int i = 0; i < n; i++){
        cin >> cost[i];
    }
    vector<int> dp(n+1, -1);
    cout << solve(n-1, k, cost, dp);
}

```

Recursion

T.C — $O(K^N)$

S.C — $O(N)$

Memorization

T.C — $O(K \times N)$

S.C — $O(N)$

N $+$ N
 \downarrow \downarrow
 Recursion DP

Tabulation

T.C — $O(K \times N)$

S.C — $O(N)$

DP

Buffer
overflow

198. House Robber

Medium



15.5K

303



Companies

You are a professional robber planning to rob houses along a street. Each house has a certain amount of money stashed, the only constraint stopping you from robbing each of them is that adjacent houses have security systems connected and **it will automatically contact the police if two adjacent houses were broken into on the same night.**

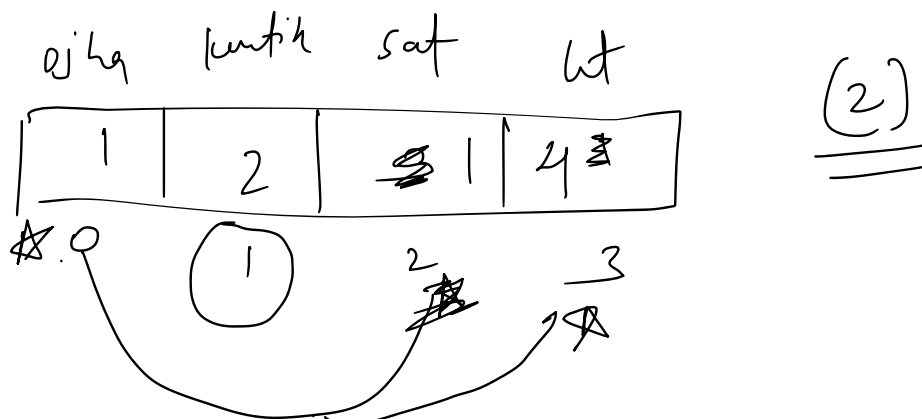
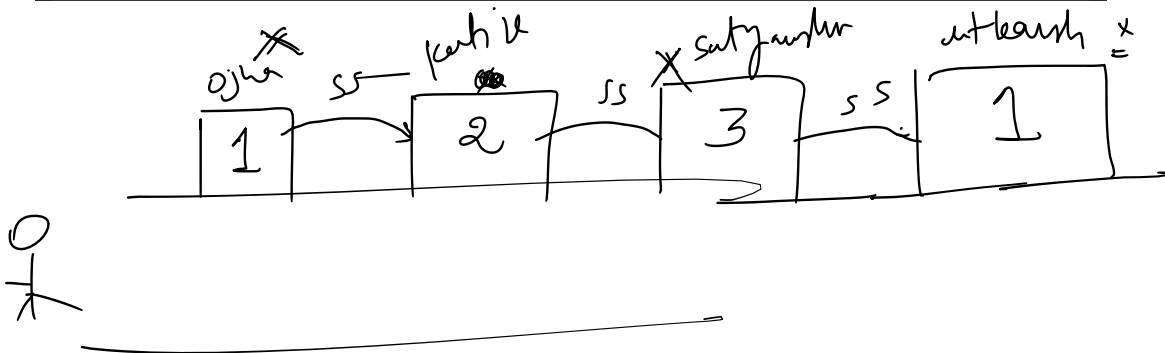
Given an integer array `nums` representing the amount of money of each house, return the *maximum* amount of money you can rob tonight **without alerting the police.**

Example 1:

Input: `nums = [1,2,3,1]`

Output: 4

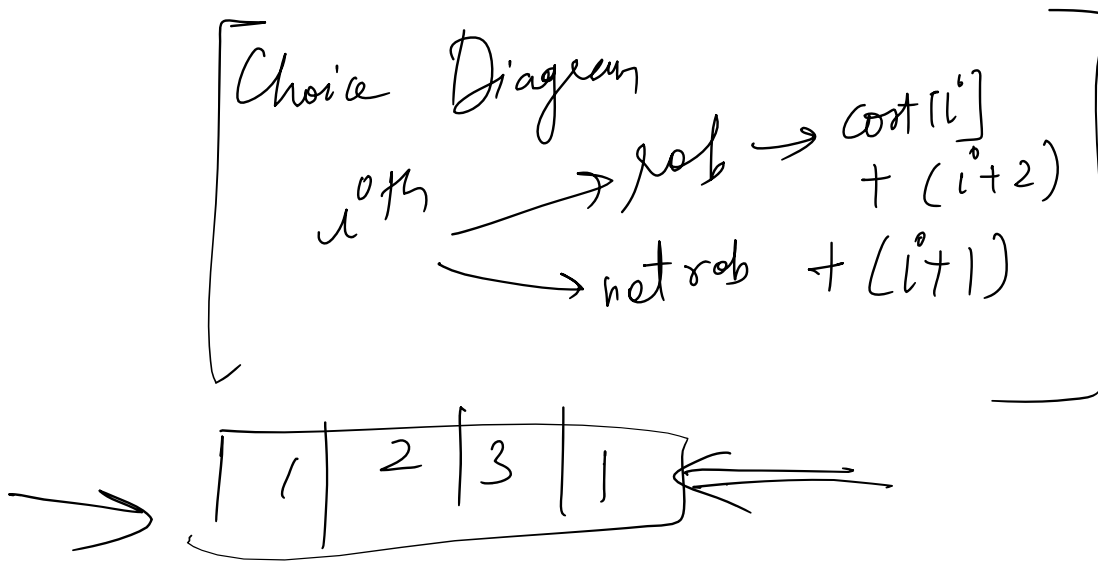
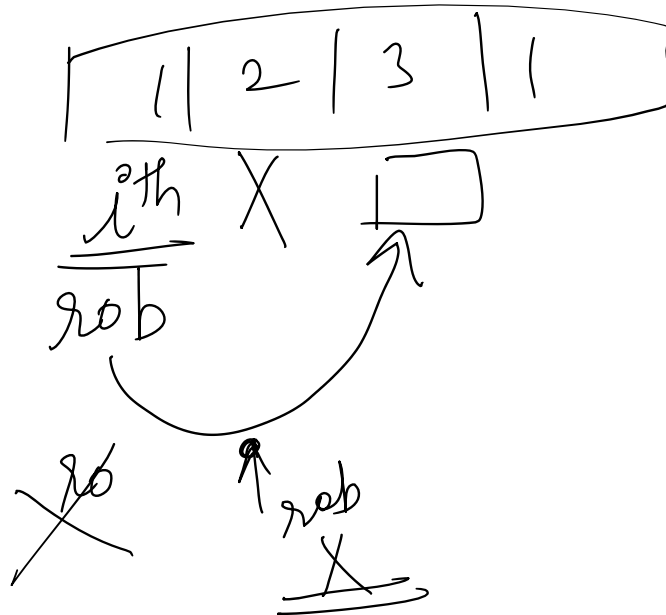
Explanation: Rob house 1 (money = 1) and then rob house 3 (money = 3).
Total amount you can rob = 1 + 3 = 4.



Recursion

Choice + Decision

Choice + Common



int solve(vector<int> &paisa,
int n)

if $(n == 0)$ return 0;
 // res int maxi = 0;

$$\text{maxi}^6 = \text{max}(\text{maxi}, \text{solve}(\text{paisa}, n-2) + \underline{\text{paisa}(n-1)});$$

rel X

$$\max_i = \max(\max_i, \text{solve}(\text{pairs}, n-1))$$

return \max_i^0

2

Recursive TLE

```
class Solution {
public:
    int solve(vector<int>& paisa, int n){
        if(n <= 0) return 0;
        int maxi = 0;
        // rob
        maxi = max(maxi, solve(paisa, n-2) + paisa[n-1]);
        // rob nahi krta
        maxi = max(maxi, solve(paisa, n-1));
        return maxi;
    }
    int rob(vector<int>& nums) {
        return solve(nums, nums.size());
    }
};
```

T.C $\leftarrow O(2^N)$

S.C — $O(N)$

Memoization

```
class Solution {
public:
    int solve(vector<int>& paisa, int n, vector<int> &dp){
        if(n <= 0) return 0;
        if(dp[n] != -1) return dp[n];
        int maxi = 0;
        // rob
        maxi = max(maxi, solve(paisa, n-2, dp) + paisa[n-1]);
        // rob nahi krta
        maxi = max(maxi, solve(paisa, n-1, dp));
        return dp[n] = maxi;
    }
    int rob(vector<int>& nums) {
        vector<int> dp(nums.size() + 1, -1);
        return solve(nums, nums.size(), dp);
    }
};
```

T.C - $O(N)$
 S.C - $O(N)$ \rightarrow $N \Rightarrow$ Recursion
 \rightarrow $N \Rightarrow$ DP Array

```
int rob(vector<int>& nums) {
    vector<int> dp(nums.size() + 1, 0);
    // base case
    dp[0] = 0;
    int n = nums.size();
    for(int i = 1; i <= n; i++){
        int maxi = 0;
        // rob
        maxi = max(maxi, (i-2 >= 0 ? dp[i-2] : 0) + nums[i-1]);
        // rob nahi krta
        maxi = max(maxi, dp[i-1]);
        dp[i] = maxi;
    }
    return dp[n];
}
```

$\left[\begin{array}{l} \text{T.C} - O(N) \\ \text{S.C} - O(N) \end{array} \right] \rightarrow \underline{\underline{dp}}$
 $[i-2] \Rightarrow dp[i]$

$i=3$
 $dp[3-2] \Rightarrow dp[1]$
 $dp[3-1] \Rightarrow dp[2]$
 $i=2$
 $dp[2-2] \Rightarrow dp[0]$
 $dp[2-1] \Rightarrow dp[1]$

$i=1$
 $dp[0]$

$i=4$
 $dp[4-2] \Rightarrow dp[2]$
 $dp[4-1] \Rightarrow dp[3]$

$i=n$
 $dp[n-2]$
 $dp[n-1]$
 $int\ prev = -1$
 $int\ curr = 0$
 $dp[0] = 0$

```

for(int i = 1; i <= n; i++){
    int maxi = 0;
    // rob
    maxi = max(maxi, (i-2 >= 0 ? prev : 0)
+ nums[i-1]);
    // rob nahi krta
    maxi = max(maxi, curr);
    Prev = curr; (0)
    curr = maxi;
}

```



```

    curr = maxi;
}

```

$dp[i-2] \iff prev$
 $, dp[i-1] \iff curr$

$O(1)$

~~$dp[i]$~~ = $curr$;
 $curr = maxi$;

~~prev~~
~~curr~~ ~~curr~~ ~~curr~~ ~~curr~~ ~~curr~~

1	2	3	1
---	---	---	---

T.C — $O(N)$
 S.C — $O(1)$

```

int rob(vector<int>& nums) {
    // base case
    int curr = 0; // dp[i-1]
    int prev = -1; // dp[i-2]
    int n = nums.size();
    for(int i = 1; i <= n; i++){
        int maxi = 0;
        // rob
        maxi = max(maxi, (i-2 >= 0 ? prev : 0) + nums[i-1]);
        // rob nahi krta
        maxi = max(maxi, curr);
        prev = curr;
        curr = maxi;
    }
    return curr;
}

```

T.C — $O(N)$
 S.C — $O(1)$

213. House Robber II

Hint

Medium



7.3K



108



Companies

You are a professional robber planning to rob houses along a street. Each house has a certain amount of money stashed. All houses at this place are **arranged in a circle**. That means the first house is the neighbor of the last one. Meanwhile, adjacent houses have a security system connected, and **it will automatically contact the police if two adjacent houses were broken into on the same night**.

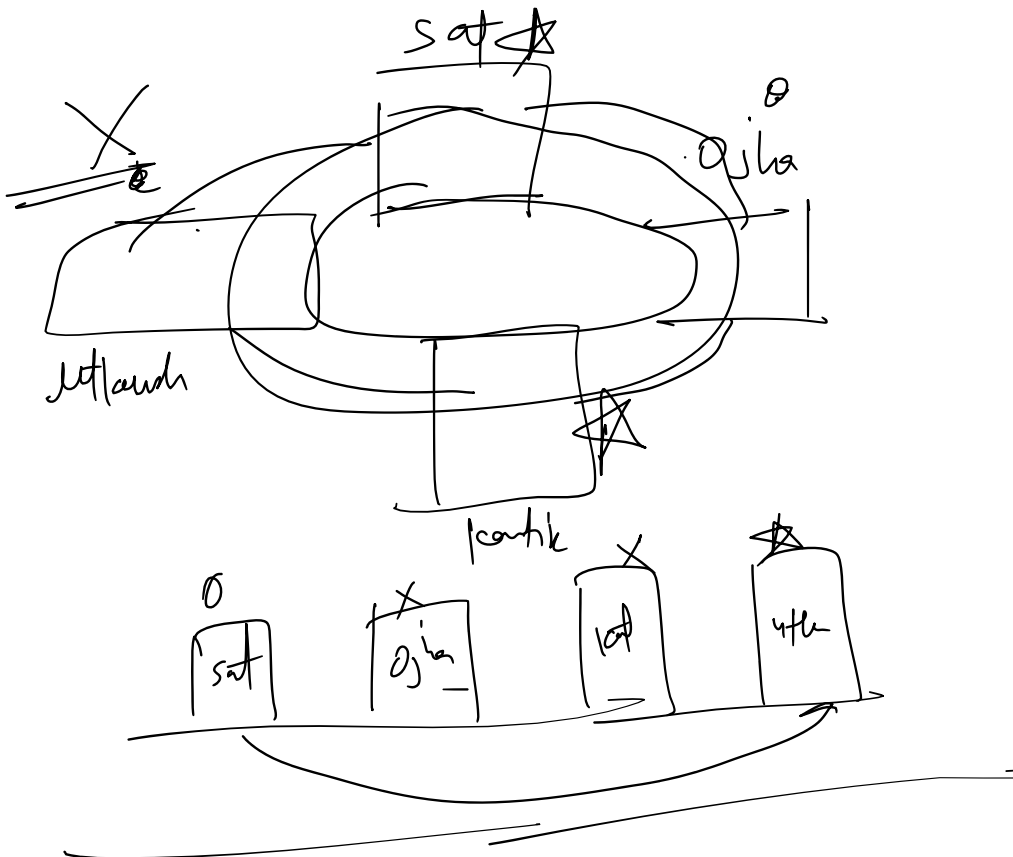
Given an integer array `nums` representing the amount of money of each house, return *the maximum amount of money you can rob tonight **without alerting the police***.

Example 1:

Input: `nums = [2,3,2]`

Output: 3

Explanation: You cannot rob house 1 (money = 2) and then rob house 3 (money = 2), because they are adjacent houses.



```

int simple_rob(vector<int>& nums) {
    // base case
    int curr = 0; // dp[i-1]
    int prev = -1; // dp[i-2]
    int n = nums.size();
    for(int i = 1; i <= n; i++){
        int maxi = 0;
        // rob
        maxi = max(maxi, (i-2 >= 0 ? prev : 0) + nums[i-1]);
        // rob nhi krta
        maxi = max(maxi, curr);
        prev = curr;
        curr = maxi;
    }
    return curr;
}

int rob(vector<int>& nums) {
    if(nums.size() == 1) return nums[0];
    if(nums.size() == 0) return 0;
    vector<int> first_skip, last_skip;
    int n = nums.size();
    for(int i = 0; i < n; i++){
        if(i != 0){
            first_skip.push_back(nums[i]);
        }
        if(i != n-1){
            last_skip.push_back(nums[i]);
        }
    }
    return max(simple_rob(first_skip), simple_rob(last_skip));
}

```

$T.C - O(2N) \approx O(N)$
 $S.C - O(2N) \approx O(N)$