

Subset Sum Problem



Medium

Accuracy: 32.0%

Submissions: 103K+

Points: 4

Given an array of non-negative integers, and a value *sum*, determine if there is a subset of the given set with sum equal to given *sum*.

Example 1:

Input:

N = 6

arr[] = {3, 34, 4, 12, 5, 2}

sum = 9

Output: 1

Explanation: Here there exists a subset with sum = 9, 4+3+2 = 9.

non-negative integer
 $0 \leq \text{+ive}$

Sum

T.C |

3	34	4	12	5	2
0	1	2	3	4	5

 N=6

sum = 9

⑨

⑨

subset {5, 4} {4, 3, 2}

yes

3	34	4	12	5	2
---	----	---	----	---	---

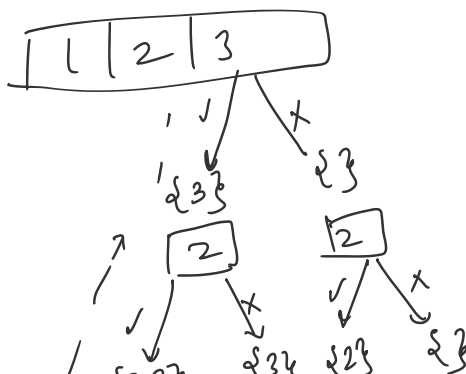
sum=30

{3, 4, 12, 5, 2}

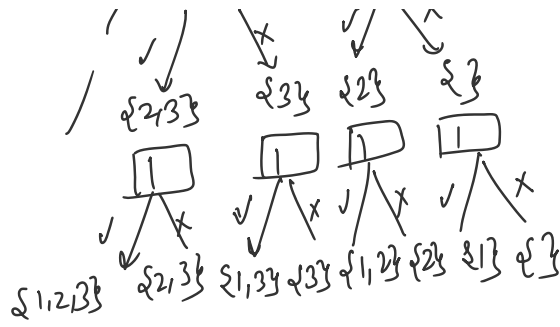
{26}

False

sum → True



$O(2^N \times N)$
 \Rightarrow $[1, 2, 3]$
 $O(2^N \times N + 2^N)$
 subset find



Loop and find total
 subset find

weight array

3	34	4	12	5	2
0	1	2	3	4	5

value array

1	1	1	1	1	1
---	---	---	---	---	---

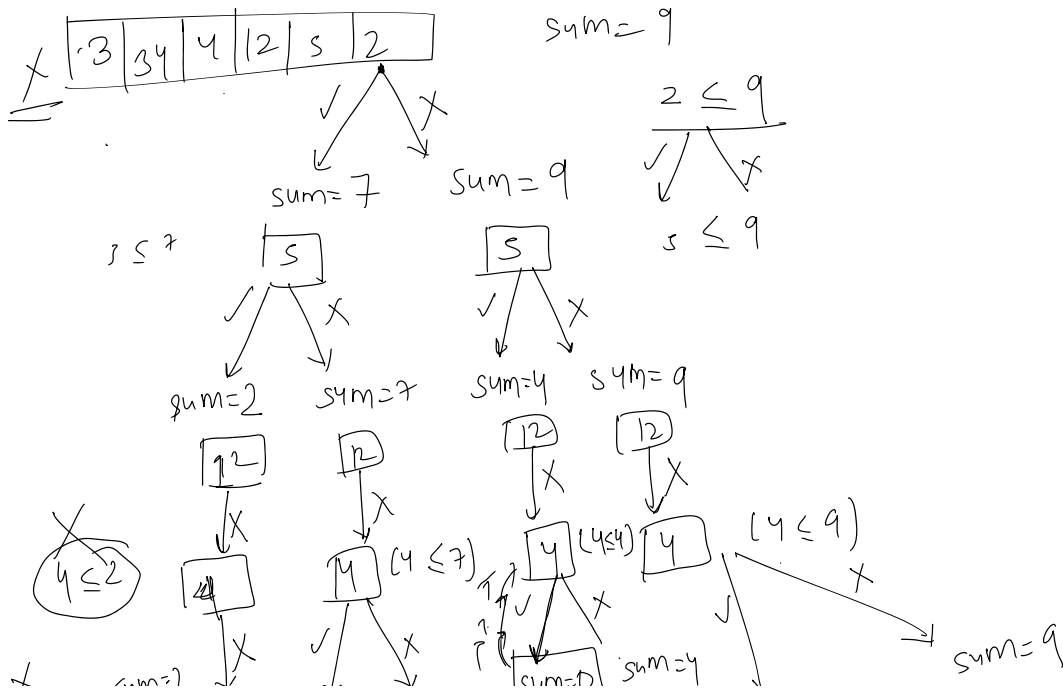
sum = 9

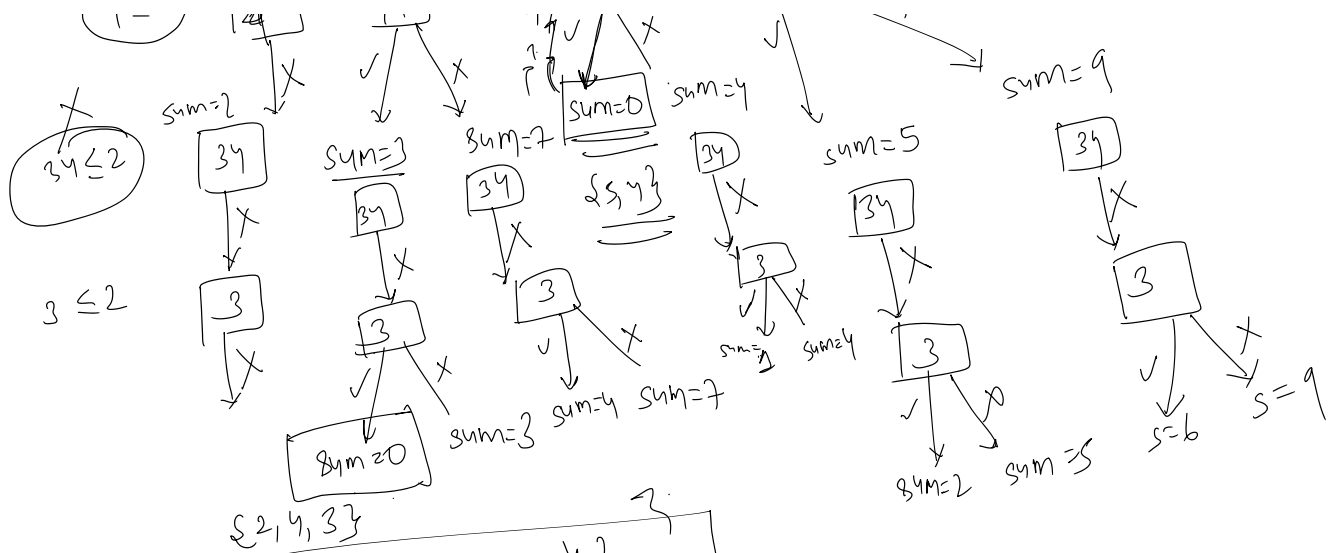
$w = 9$
 weight array
 value array

② weight $\leq Kc$
 Include / Not Include

if (weight item $\leq Kc$)
 { Included knapsack
 Not Included
 }

T.C.
 sum = 9
 2 ans





$\{2, 4, 3\}$
 $\{2, 4, 3\}, \{5, 4\}$

Choice + Decision (Recursion)

```

class Solution{
public:
    bool solve(vector<int> &arr, int n, int sum){
        // base case
        if(sum == 0) return true;
        if(n == 0) return false;
        if(arr[n-1] <= sum){
            // include
            return solve(arr, n-1, sum-arr[n-1]) || solve(arr, n-1, sum);
        }else{
            // not include
            return solve(arr, n-1, sum);
        }
    }
    bool isSubsetSum(vector<int> &arr, int sum){
        int n = arr.size();
        return solve(arr, n, sum);
    }
};

```

Time limited Exceed.

T.C - $O(2^N)$
 S.C - $O(N)$

① change variable

2- $(2D)$

3- $(3D)$

$n - (nD)$

nt | sum + 1

```

class Solution{
public:
    bool solve(vector<int> &arr, int n, int sum, vector<vector<int>>& dp){
        // base case
        if(sum == 0) return true; ✓
        if(n == 0) return false; ✓
        if(dp[n][sum] != -1) return dp[n][sum];
        if(arr[n-1] <= sum){
            // include                not include
            return dp[n][sum] = solve(arr, n-1, sum-arr[n-1], dp) || solve(arr, n-1, sum, dp);
        }else{
            return dp[n][sum] = solve(arr, n-1, sum, dp);
        }
    }
    bool isSubsetSum(vector<int> arr, int sum){
        int n = arr.size();
        vector<vector<int>> dp(n+1, vector<int>(sum+1, -1)); ✓
        return solve(arr, n, sum, dp);
    }
};
// Driver Code Ends

```

Iterative
(Tabulation)
(Bottom to Top)

T.C - $O(n \times \text{sum})$
S.C - $O(n \times \text{sum})$

$(n+1) \times (\text{sum}+1)$

memoization

recursion
stack overflow

	0	1	2	3	4	5
0	F	F	F	F	F	F
1						
2						
3						
4						
5						

	0	1	2	3	4	5
0	3	3	4	4	12	3
1						
2						
3						
4						
5						

```

bool isSubsetSum(vector<int> arr, int sum){
    int n = arr.size();
    vector<vector<int>> dp(n+1, vector<int>(sum+1));
    // return solve(arr, n, sum, dp);
    // dp[i][j]
    // i -> n
    // j -> sum
    for(int j = 0; j < sum+1; j++){
        dp[0][j] = false;
    }
    for(int i = 0; i < n+1; i++){
        dp[i][0] = true;
    }
    for(int i = 1; i < n+1; i++){
        for(int j = 1; j < sum+1; j++){
            if(arr[i-1] <= sum){
                // include                not include
                dp[i][j] = dp[i-1][j-arr[i-1]] || dp[i-1][j];
            }else{
                dp[i][j] = dp[i-1][j];
            }
        }
    }
    return dp[n][sum];
}

```

T.C - $O(n \times \text{sum})$
S.C - $O(n \times \text{sum})$

h = 6, sum = 9, 3 ≤ 9

	0	1	2	3	4	5	6	7	8	9
0	T	F	F	F	F	F	F	F	F	F
1	T	F	F	F	F	F	F	F	F	F
2	T	F	F	F	F	F	F	F	F	F
3	T	F	F	F	F	F	F	F	F	F
4	T	F	F	F	F	F	F	F	F	F
5	T	F	F	F	F	F	F	F	F	F
6	T	F	F	F	F	F	F	F	F	F

$0 \leq 3$
 $dp[s][1] || dp[s][3]$
 $dp[s][2] || dp[s][4]$
 $dp[s][9]$

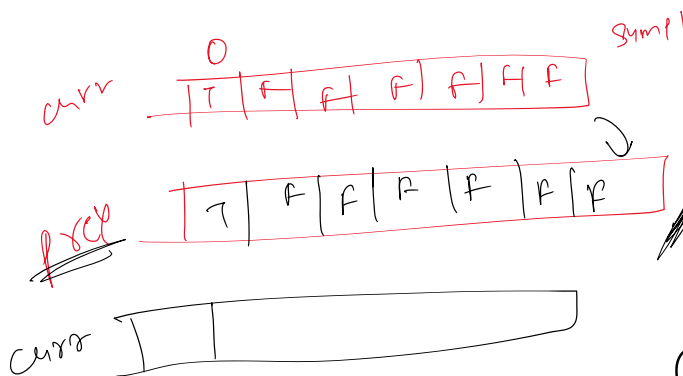
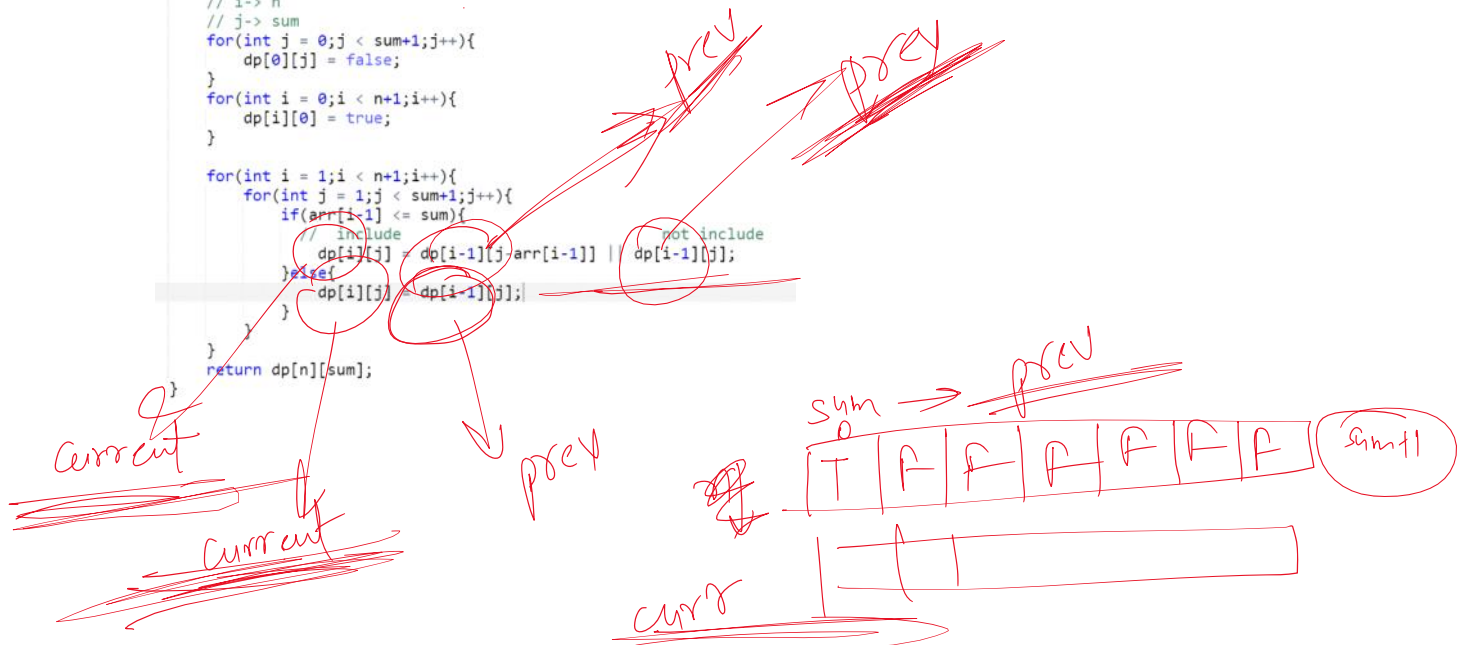
Recursion 80% → Memoization (10%) → Tabulation (10%) → Space Optimization

```

bool isSubsetSum(vector<int>arr, int sum){
    int n = arr.size();
    vector<vector<int>> dp(n+1,vector<int>(sum+1));
    // return solve(arr,n,sum,dp);
    // dp[i][j]
    // i-> n
    // j-> sum
    for(int j = 0; j < sum+1; j++){
        dp[0][j] = false;
    }
    for(int i = 0; i < n+1; i++){
        dp[i][0] = true;
    }

    for(int i = 1; i < n+1; i++){
        for(int j = 1; j < sum+1; j++){
            if(arr[i-1] <= sum){
                // include
                dp[i][j] = dp[i-1][j-arr[i-1]] || dp[i-1][j];
            } else {
                // not include
                dp[i][j] = dp[i-1][j];
            }
        }
    }
    return dp[n][sum];
}

```



```

vector<int> curr(sum+1,0);
curr[0] = 1;

for(int i = 1; i < n+1; i++){
    auto prev = curr;
    for(int j = 1; j < sum+1; j++){
        if(arr[i-1] <= sum){
            // include
            curr[j] = prev[j-arr[i-1]] || prev[j];
        } else {
            // not include
            curr[j] = prev[j];
        }
    }
}
return curr[sum];

```

T.C - $O(n \times \text{sum})$
 S.C - $O(2 \times \text{sum})$

2-D DP Array
 1-D DP Array

Partition Equal Subset Sum



Medium

Accuracy: 30.24%

Submissions: 149K+

Points: 4



This problem is part of GFG SDE Sheet. Click here to view more.



Given an array `arr[]` of size `N`, check if it can be partitioned into two parts such that the sum of elements in both parts is the same.

Example 1:

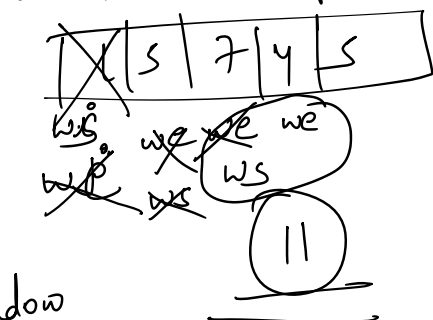
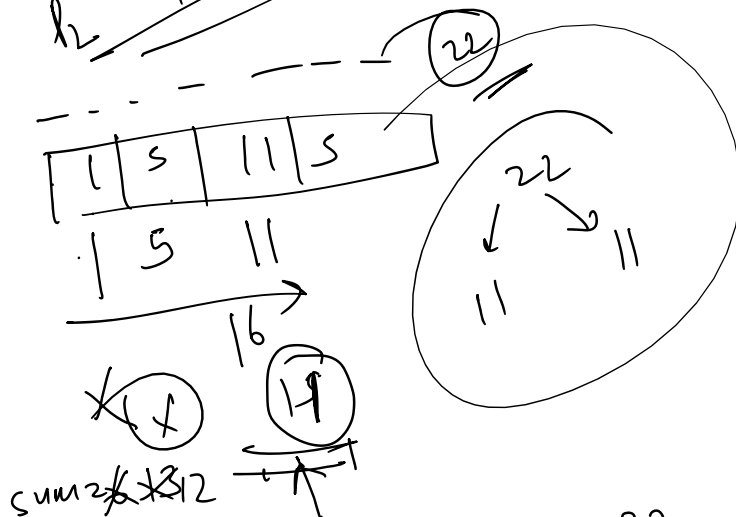
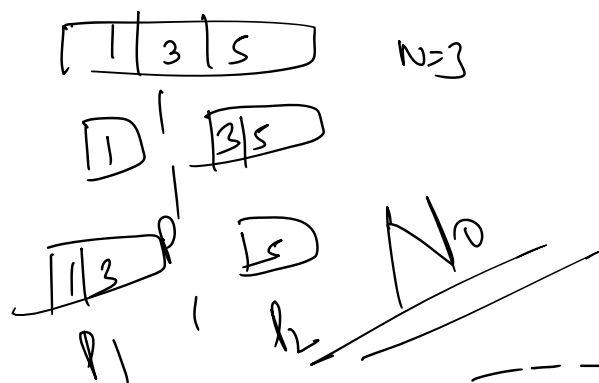
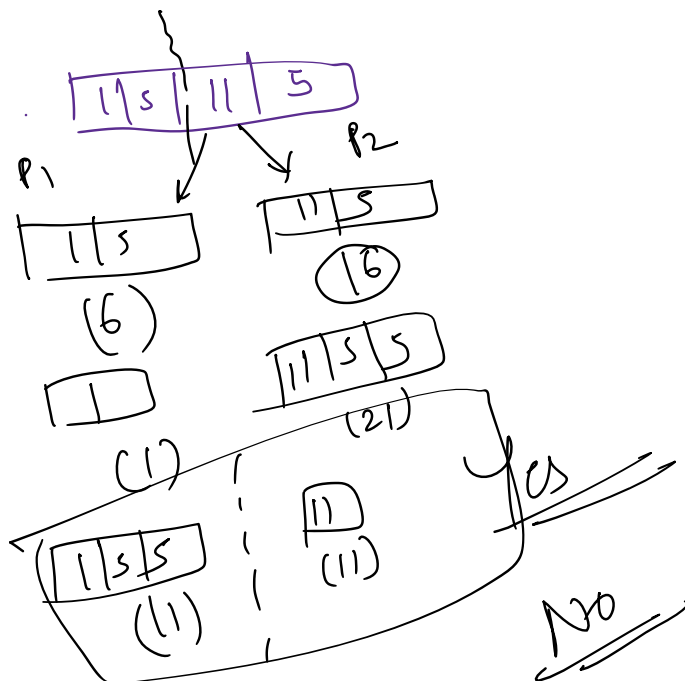
Input: `N = 4`

`arr = {1, 5, 11, 5}`

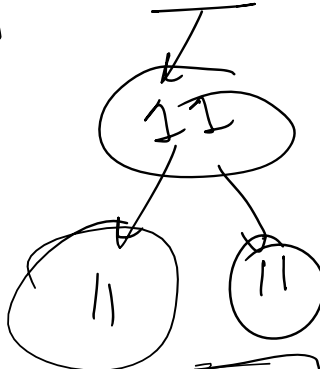
Output: YES

Explanation:

The two parts are {1, 5, 5} and {11}.



sum = 22



Sliding window

Sliding window

Array sum (S)

even

p1

p2

S/2

S/2

S/2

① Sum of Array

② Sum is even

p1

(S/2)

Subset sum problem

478 504 371 608 197 75 723 612 20

From: <https://practice.geeksforgeeks.org/problems/subset-sum-problem2014/1>

1792

20
612

9x8
509
982
3x1
1353
608
1961
197
2158
75
2233

2233
723
2956
612
3568
20
3588
2
1792

p1 = p2

Sum

sum % 2 == 0

p1

p2

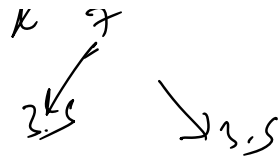
p1

p2

sum is odd

p1

p2



Perfect Sum Problem

Medium

Accuracy: 20.58%

Submissions: 111K+

Points: 4

Given an array `arr[]` of non-negative integers and an integer `sum`, the task is to count all subsets of the given array with a sum equal to a given `sum`.

Note: Answer can be very large, so, output answer modulo 10^9+7

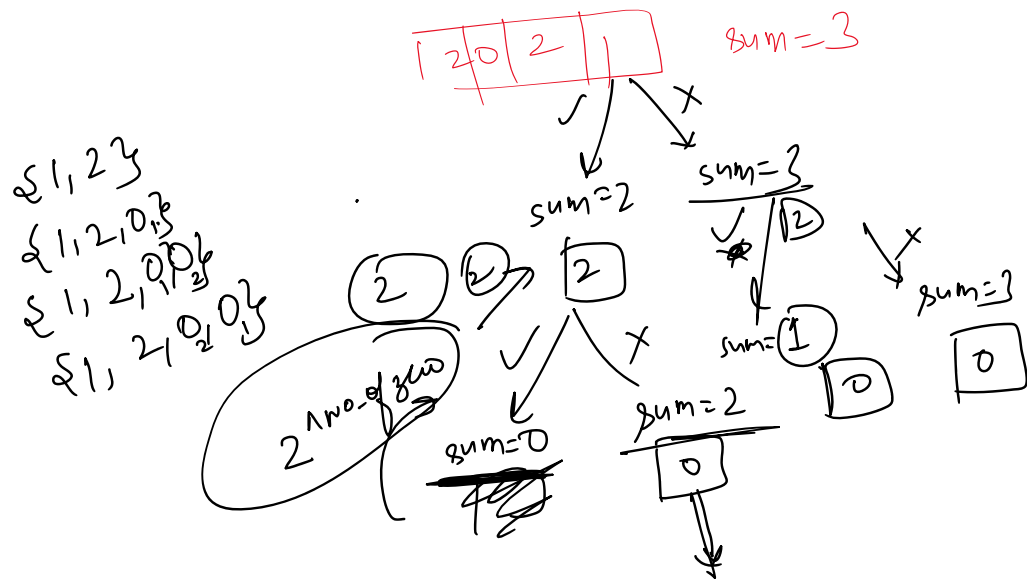
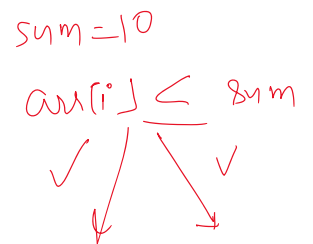
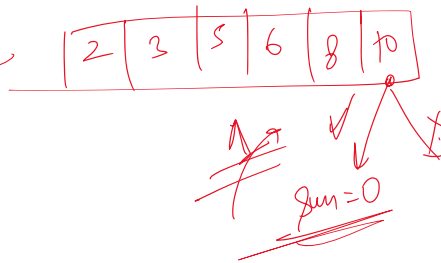
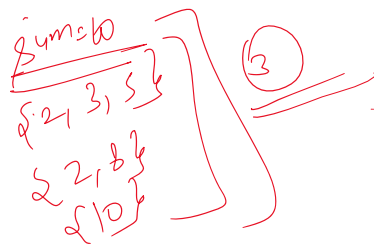
Example 1:

Input: `N = 6, arr[] = {2, 3, 5, 6, 8, 10}`

`sum = 10`

Output: 3

Explanation: {2, 3, 5}, {2, 8}, {10}




```

int mod = int(pow(10,9)) + 7;
int solve(int arr[], int n, int sum, int nz){
    // base case
    if(sum == 0) return (1 << nz);
    if(n == 0) return 0;
    if(arr[n-1] == 0){
        return solve(arr,n-1,sum,nz)%mod;
    }
    if(arr[n-1] <= sum){
        // include
        return (solve(arr,n-1,sum-arr[n-1],nz) + solve(arr,n-1,sum,nz))%mod;
    }else{
        // not include
        return solve(arr,n-1,sum,nz)%mod;
    }
}

int perfectSum(int arr[], int n, int sum)
{
    int no_of_zeros = 0;
    for(int i = 0; i < n; i++){
        if(arr[i] == 0){
            no_of_zeros++;
        }
    }
    return solve(arr,n,sum,no_of_zeros);
}

```

$N = \max(\text{length of array, sum})$

T.C - $O(2^N)$
S.C - $O(N)$

```

int perfectSum(int arr[], int n, int sum)
{
    int no_of_zeros = 0;
    for(int i = 0; i < n; i++){
        if(arr[i] == 0){
            no_of_zeros++;
        }
    }
    vector<int> curr(sum+1,0);
    curr[0] = (1 << no_of_zeros);

    for(int i = 1; i < n+1; i++){
        auto prev = curr;
        for(int j = 1; j < sum+1; j++){
            if(arr[i-1] == 0){
                curr[j] = prev[j]%mod;
            }
            else if(arr[i-1] <= j){
                // include
                curr[j] = (prev[j-arr[i-1]] + prev[j])%mod;
            }else{
                // not include
                curr[j] = prev[j]%mod;
            }
        }
    }
    return curr[sum]%mod;
}

```

T.C - $O(N \times \text{sum})$
S.C - $O(\text{sum})$