

Day 4

10 November 2022 20:31

Given an array `arr[]` and an integer `K` where `K` is smaller than size of array, the task is to find the K^{th} smallest element in the given array. It is given that all array elements are distinct.

Example 1:

```
Input:
N = 6
arr[] = 7 10 4 3 20 15
K = 3
Output : 7 ★
Explanation :
3rd smallest element in the given
array is 7.
```

$N=6$ $K=3$ ★

7	10	4	3	20	15
---	----	---	---	----	----

3rd smallest element in the array
1st smallest $\rightarrow 3$
2nd smallest $\rightarrow 4$
3rd smallest $\rightarrow \underline{\underline{7}}$ Ans

आइए हम देखेंगे $\frac{N}{2}$ ||

1st Approach

K^{th} (3rd smallest)

7	10	4	3	20	15
---	----	---	---	----	----

3	4	7	10	15	20
0	1	2			

 \uparrow
(7) O/P

T.C - $O(n \log n)$
S.C - $O(1)$

S.C - $O(1)$

Sorting $\rightarrow O(n \log n)$

कोई ऐसा method चाहिए जिससे हम
बिना sort किए k^{th} smallest element
find कर पाएँ॥

HEAP

k^{th} smallest
 k^{th} largest 2nd largest

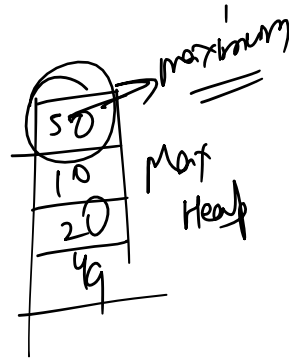
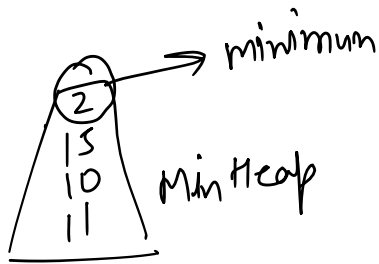
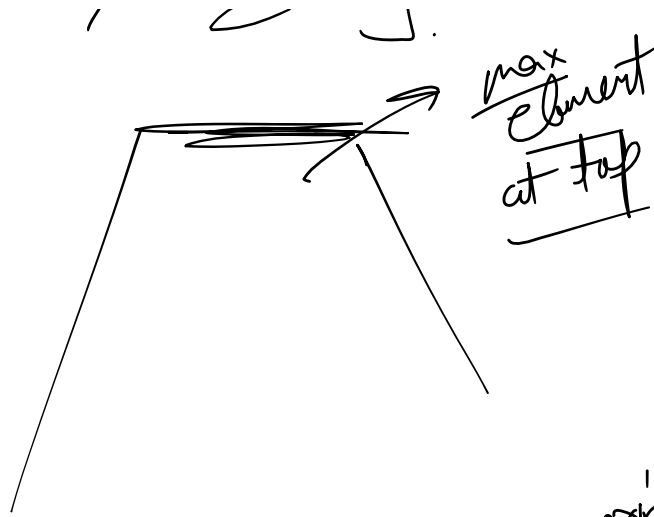
1	5	8	2	9
---	---	---	---	---

8

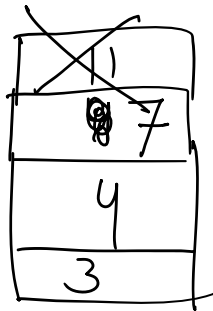
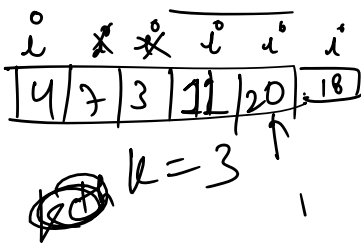
Min Heap k^{th} smallest
Max Heap k^{th} largest



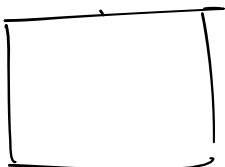
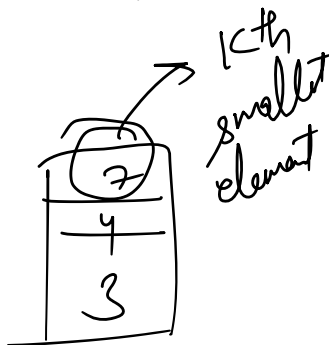
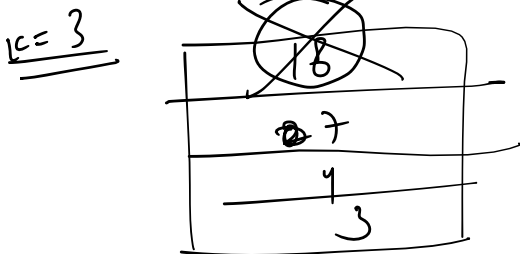
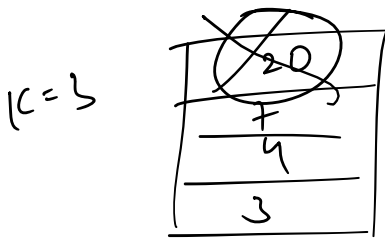
Max Heap



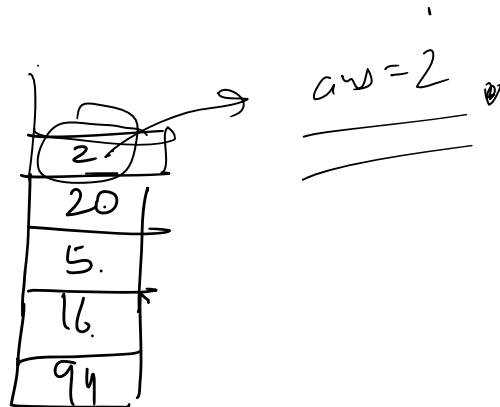
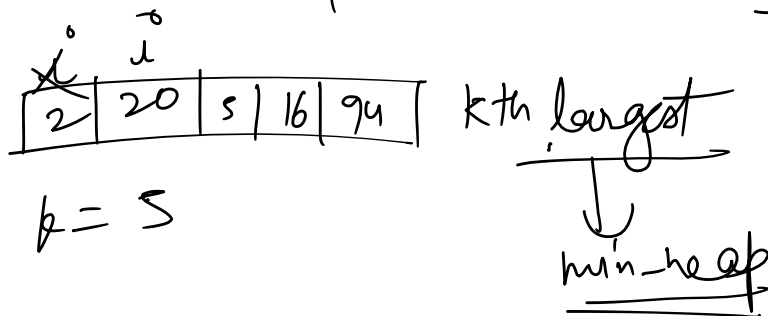
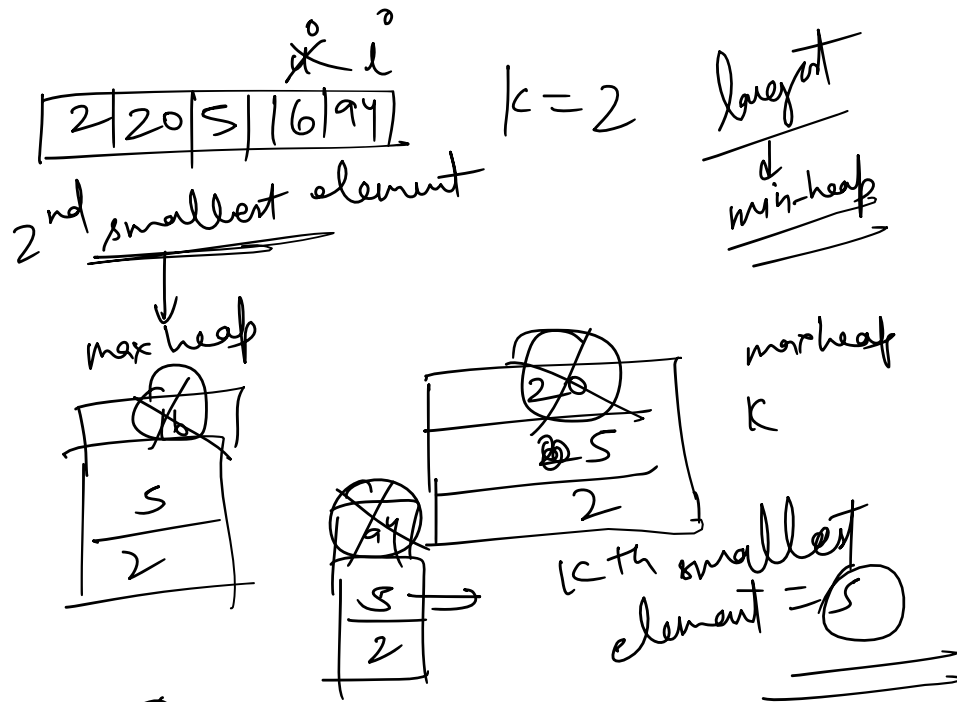
k^{th} smallest \rightarrow max heap



max heap
k value
direction
 \rightarrow 11



k element तक हमने रखेगा



Implementation

min heap \rightarrow
max heap \rightarrow

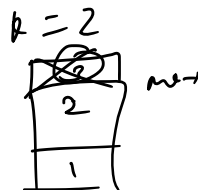
$\text{priority_queue} < \text{int} > \text{maxheap};$

min-heap $\rightarrow \text{priority_queue} < \text{int}, \text{vector} < \text{int} >, \text{greater} < \text{int} > > \text{min-heap};$

$\text{priority_queue} \langle \underline{\text{DT}} \rangle \text{max_heap};$
 $\text{priority_queue} \langle \text{DT}, \text{vector} \langle \text{DT} \rangle, \text{greater} \langle \text{DT} \rangle \rangle \text{min_heap};$
 comparator function

k^{th} smallest Element Code

1 2 3 4 5



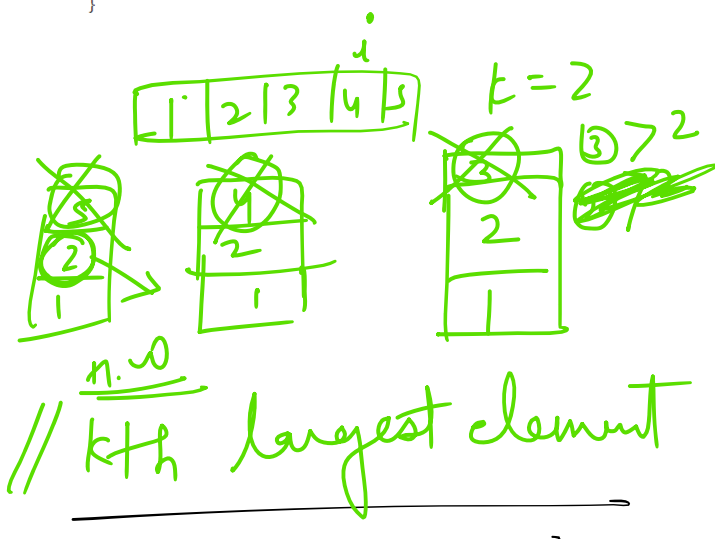
```

int kthSmallest(int arr[], int l, int r, int k) {
    // smallest -> max_heap
    priority_queue<int> max_heap;
    for(int i=l; i<=r; i++){
        max_heap.push(arr[i]);
        if(max_heap.size() > k){
            // pop hota h voh toh se hota h
            max_heap.pop();
        }
    }
    return max_heap.top();
}

```

T.C - $O(n \log k)$

$\log k$



Given an array of size N containing only 0s, 1s, and 2s; sort the array in ascending order.

Example 1:

Input:
 N = 5
 arr[] = {0 2 1 2 0}

Given an array of size N containing only 0s, 1s, and 2s; sort the array in ascending order.

Example 1:

Input:
N = 5
arr[] = {0 2 1 2 0}
Output:
0 0 1 2 2
Explanation:
0s 1s and 2s are segregated into ascending order.



1st Approach

Sorting Algorithm

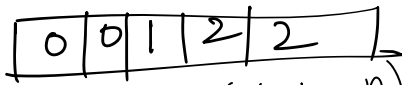
sort



T.C - $O(n \log n)$

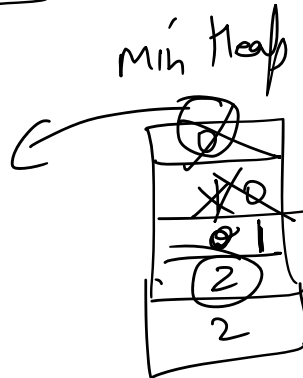
S.C - $O(1)$

2nd ~~Sorting~~ \Rightarrow Heap



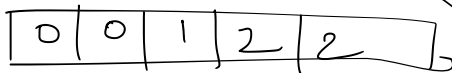
T.C - $O(n \log n)$

S.C - $O(n)$



3rd

no. of zeros = 2
ones = 1
twos = 2



O/p

Ans

T.C - $O(N)$

S.C - $O(N)$

4th Algorithm (Duth Plag)

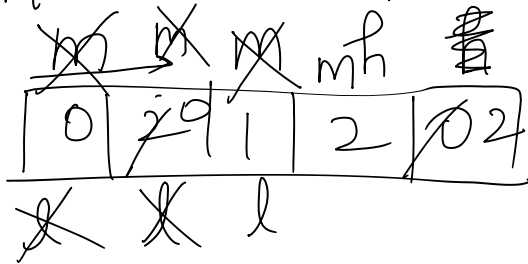
3 pointer

~~1~~
~~m~~
~~h~~

traverse

① $l, m \rightarrow 0$

② $h \rightarrow n-1$

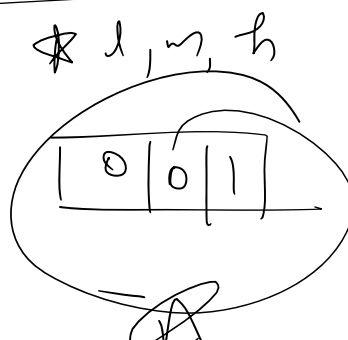
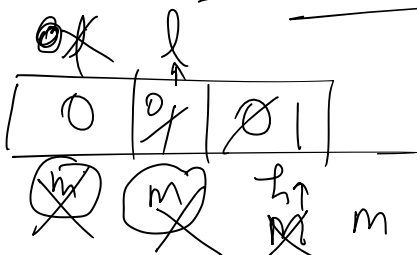


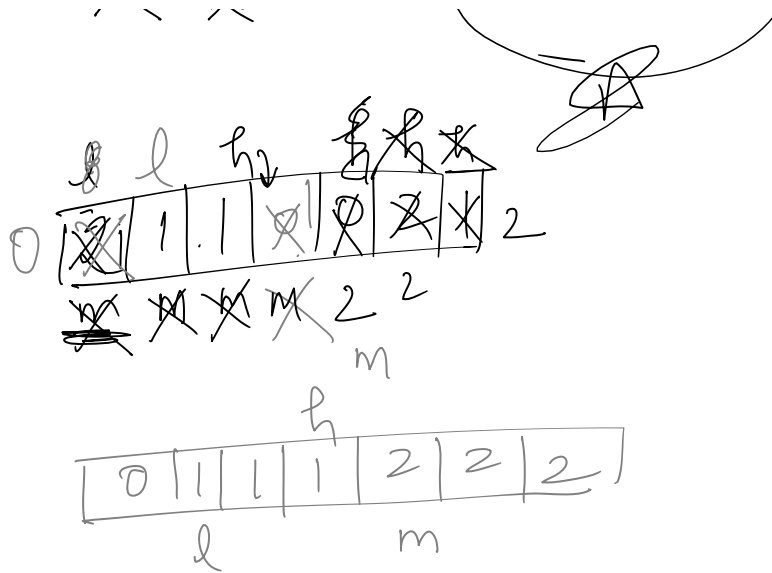
Case 1 element 0
 \rightarrow
 $\text{swap}(\text{arr}[l], \text{arr}[m])$
 $l++, m++;$

Case 2 element 2
 \rightarrow
 $\text{swap}(\text{arr}[m], \text{arr}[h])$
 $h--;$

Case 3 element 1
 \rightarrow
 $m++;$

while ($m \leq h$) // condition





\star l $\frac{1}{2}$ left (0)
 \star h $\frac{1}{2}$ right (2)

Implementation

int l=0,

m=0, h=n-1;

```
void sort012(int a[], int n)
{
    int l = 0, m = 0, h = n-1;
    while(m <= h){
        if(a[m] == 0){
            swap(a[m], a[l]);
            m++; l++;
        } else if(a[m] == 1){
            m++;
        } else{
            swap(a[m], a[h]);
            h--;
        }
    }
}
```

T.C - $O(N)$

S.C - $O(1)$

$\star\star$

Dutch Flag Algorithm

Move all negative numbers to beginning and positive to end with constant extra space

Difficulty Level : Easy • Last Updated : 14 Jun, 2022

Read

Discuss

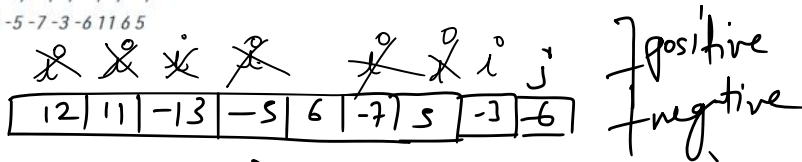


An array contains both positive and negative numbers in random order. Rearrange the array elements so that all negative numbers appear before all positive numbers.

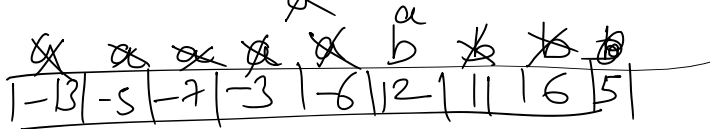
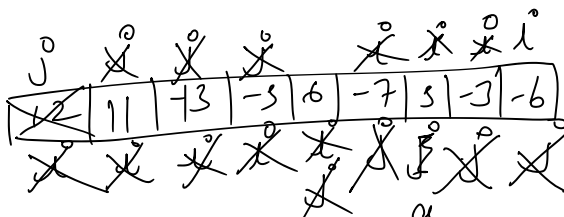
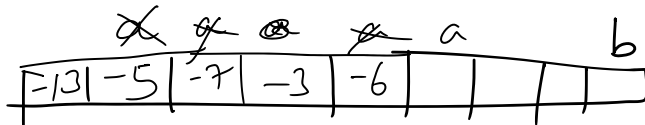
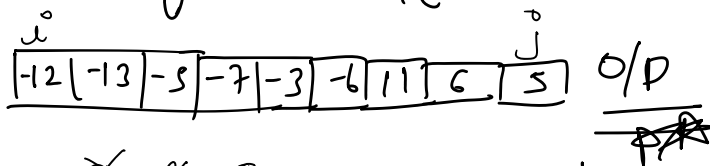
Examples :

Input: -12, 11, -13, -5, 6, -7, 5, -3, -6

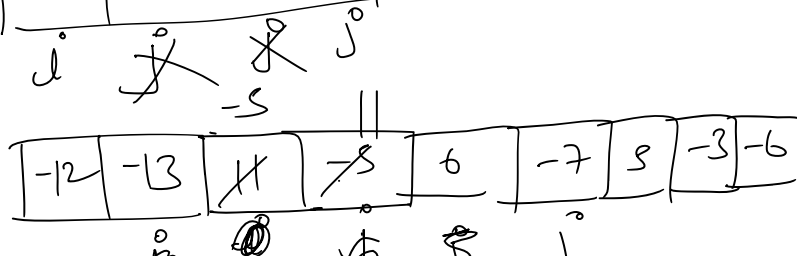
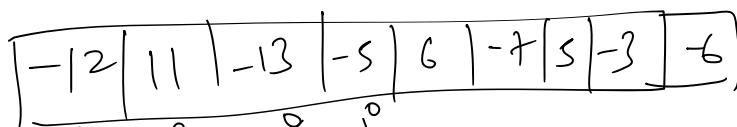
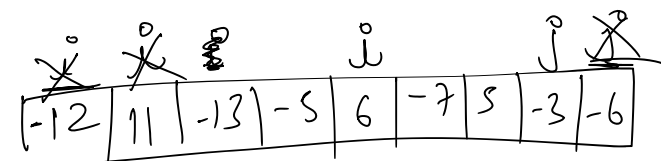
Output: -12 -13 -5 -7 -3 -6 11 6 5

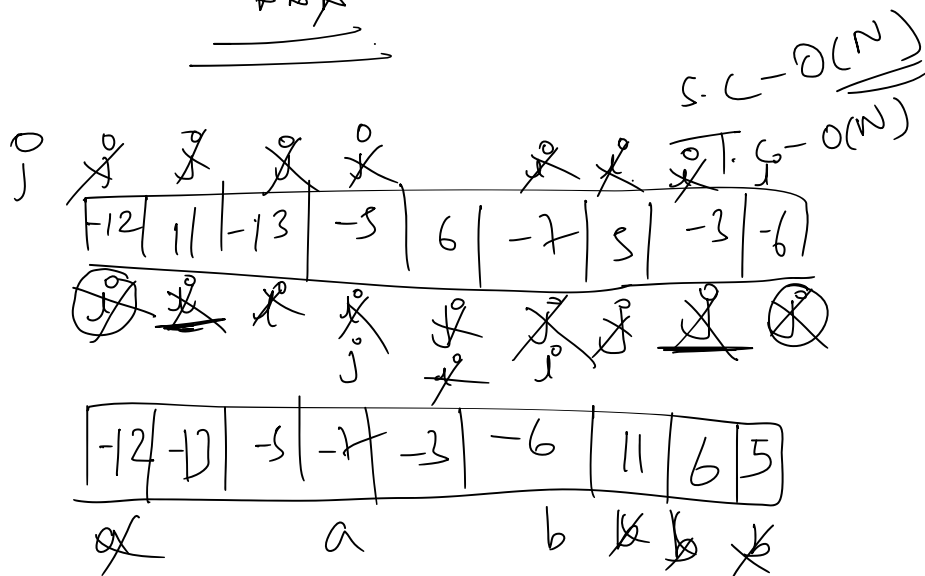
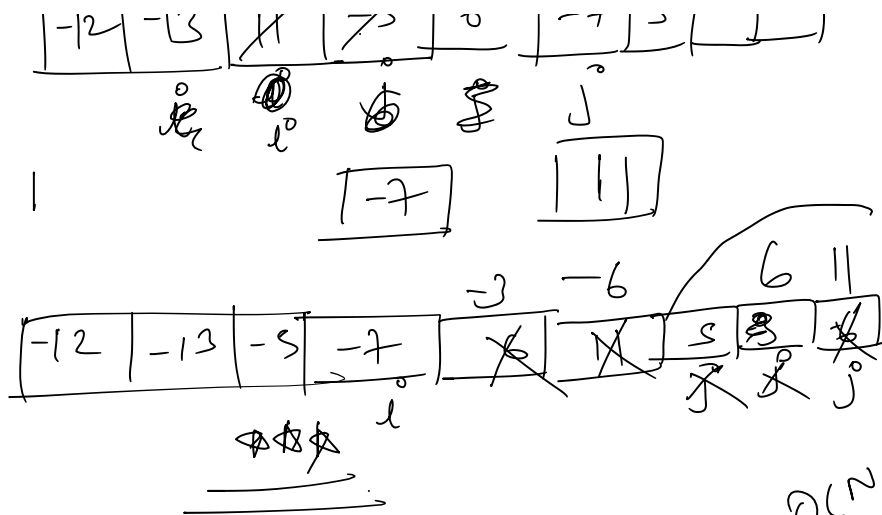


Order change ना हो जाए



Approach ✓ ✓ Answer





Code

```

int j = n-1;
vector<int> answer(n); int a=0, b=n-1;
for(int i=0; i<n; i++, j--){
    if(arr[i] < 0){
        answer[a] = arr[i]; a++;
    }
    if(arr[j] > 0){
        answer[b] = arr[j]; b--;
    }
}
return answer;

```

Complexity: $T.C - O(N)$, $S.C - O(N)$

Union of two sorted array

Given two arrays $a[]$ and $b[]$ of size n and m respectively. The task is to find union between these two arrays.

Union of the two arrays can be defined as the set containing distinct elements from both the arrays. If there are repetitions, then only one occurrence of element should be printed in the union.

Note : Elements are not necessarily distinct.

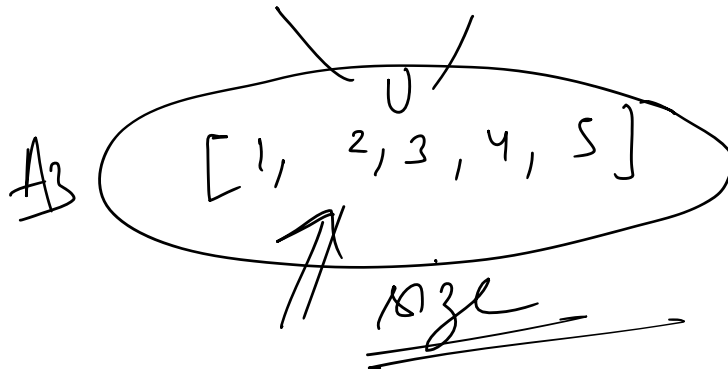
Example 1:

```

Input:
5 3
1 2 3 4 5
1 2 3
Output:
5
Explanation:
1, 2, 3, 4 and 5 are the
elements which comes in the union set
of both arrays. So count is 5.

```

$$A_1 = [1, 2, 3, 4, 5] \quad A_2 = [1, 2, 3]$$



1st Approach

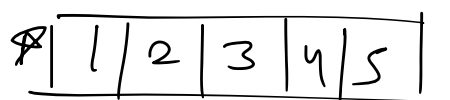
~~Set~~ $\rightarrow \{1, 2, 3, 4, 5\}$ (5)

A_1 ~~Read~~ $n \log n$ A_2 $m \log m$

T.C - $O(n \log n + m \log m)$

$(\min) \log(m+n)$

S.C - $O(m+n)$



n

65	25	1	32	54	6
----	----	---	----	----	---

m

85	2
----	---

① sort A₁ and A₂ Distinct

A₁

11	6	25	32	54	65
----	---	----	----	----	----

A₂

2	85
---	----

1	2	6	25	32	54	85
---	---	---	----	----	----	----

T.C - $(m+n) \lg(m+n)$
S.C - $\underline{\underline{1}}$

Elements Distinct not

11	6	25	25	32	32	54	85
----	---	----	----	----	----	----	----

2	85	96	96	100
---	----	----	----	-----

A₁

2

A₂

3	4	6
---	---	---

c = 1

```

int doUnion(int a[], int n, int b[], int m) {
    sort(a,a+n);
    sort(b,b+m);
    int i = 0, j = 0;
    int c = 0;
    while(i < n && j < m){
        while(i+1 < n && a[i] == a[i+1]){
            i = i+1;
        }
        while(j + 1 < m && b[j] == b[j+1]){
            j = j + 1;
        }
        if(a[i] < b[j]){
            i++;
        } else if(a[i] > b[j]){
            j++;
        } else{
            i++;j++;
        }
        c++;
    }
}

```

T.C - $O(\dots)$

$O(m+n)$

$n \log n + m \log m$

T.C - $O(m+n \log(n+m))$

S.C - $O(1)$

```

if(i >= n){
    while(j < m){
        while(j+1 < m && b[j] == b[j+1]){
            j++;
        }
        j++;
        c++;
    }
} else{
    while(i < n){
        while(i+1 < n && a[i] == a[i+1]){
            i++;
        }
        c++;
        i++;
    }
}
return c;

```