

Day 5

11 November 2022 20:35

Easy



9.1K

12.6K



Companies

Given an integer array `nums` sorted in **non-decreasing order**, remove the duplicates **in-place** such that each unique element appears only **once**. The **relative order** of the elements should be kept the **same**.

Since it is impossible to change the length of the array in some languages, you must instead have the result be placed in the **first part** of the array `nums`. More formally, if there are `k` elements after removing the duplicates, then the first `k` elements of `nums` should hold the final result. It does not matter what you leave beyond the first `k` elements.

Return `k` after placing the final result in the first `k` slots of `nums`.

Do **not** allocate extra space for another array. You must do this by **modifying the input array in-place** with $O(1)$ extra memory.

Example 1:

Input: `nums = [1,1,2]`

Output: `2, nums = [1,2,_]`

Explanation: Your function should return `k = 2`, with the first two elements of `nums` being 1 and 2 respectively. It does not matter what you leave beyond the returned `k` (hence they are underscores).

Example 2:

Input: `nums = [0,0,1,1,1,2,2,3,3,4]`

Output: `5, nums = [0,1,2,3,4,_,_,_,_,_]`

Explanation: Your function should return `k = 5`, with the first five elements of `nums` being 0, 1, 2, 3, and 4 respectively. It does not matter what you leave beyond the returned `k` (hence they are underscores).

Ex1

1	1	2
---	---	---



1	2	...
---	---	-----

2

Ex2

`[0, 0, 1, 1, 1, 2, 2, 3, 3, 4]`

`[0, 1, 2, 3, 4, _, _, _, _, _]`

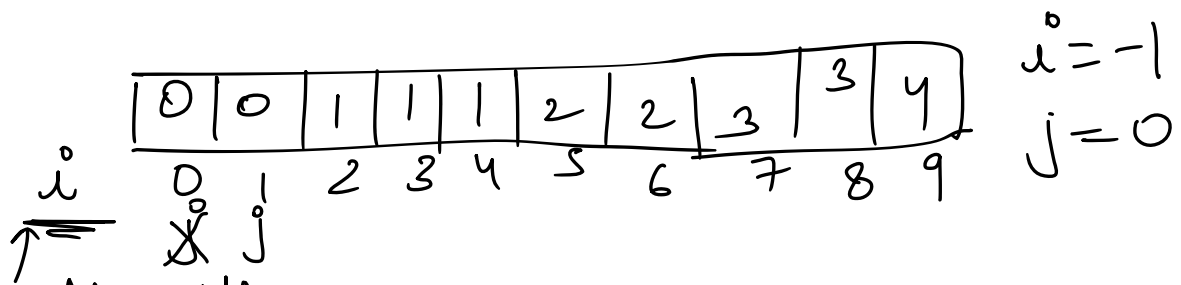
5
Ans

Contik

Inplace → No Extra

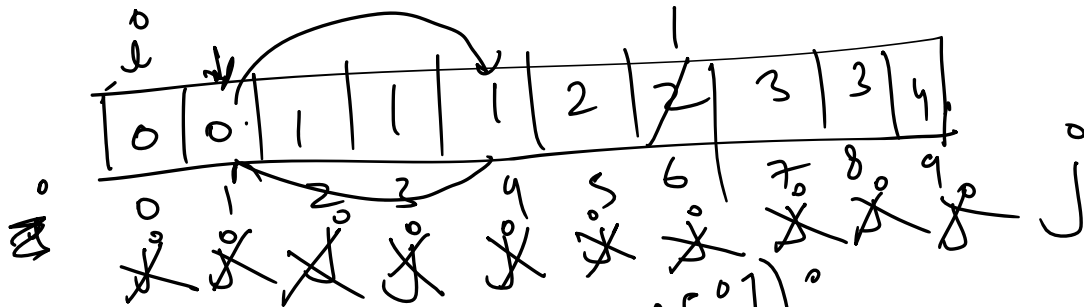
0	0	1	1	1	2	2	3	3	4
---	---	---	---	---	---	---	---	---	---

$i = -1$

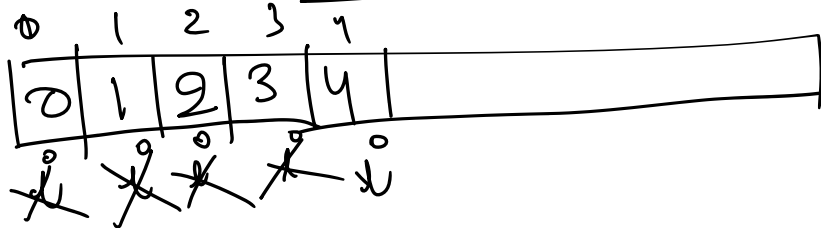


Algorithm

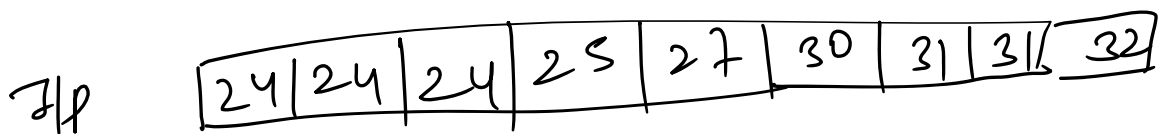
swap(arr[i+1], arr[j])



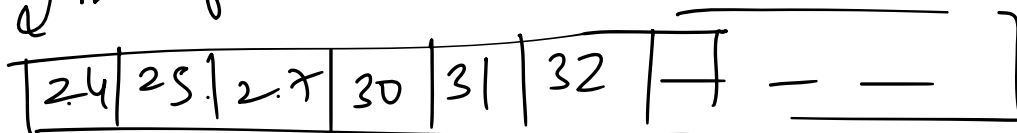
~~swap(arr[i+1], arr[j]);~~



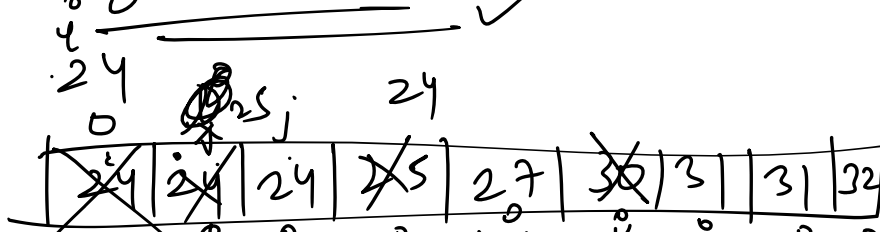
$i+1$ ✓



modified



Size = 6



$i = -1$
 $j = 0$

24	24	24	25	27	30	31	32
j	j	j	j	j	j	j	j

last occurrence

swap(nums[i+1], nums[j])

$i++$

j	j	j	j	j	j
24	25	27	30	31	32
0	1	2	3	4	5

$(i+1) \Rightarrow 6$ ✓

```

int removeDuplicates(vector<int>& nums) {
    int i = -1, j = 0;
    while(j < nums.size()){
        if(j + 1 < nums.size() && nums[j] == nums[j+1]){
            j++; continue;
        }
        // last occurrence
        swap(nums[i+1], nums[j]);
        i++;
        j++;
    }
    return i+1;
}

```

T.C - $O(N)$

S.C - $O(1)$

✓

Intersection of Two Arrays

A_1

1	2	3	4	5
---	---	---	---	---

A_1

1	2	3	4	5
---	---	---	---	---

A_2

1	2	3
---	---	---

* sort the A_1 and A_2

A_1

85	86	91	100	102
----	----	----	-----	-----

A_2

2	3	4	6	82	85	100
---	---	---	---	----	----	-----

O/p

85	100
----	-----

Algorithm

* sort A_1 and A_2

$i=0, j=0, c=0$

while ($i < n$ && $j < m$) {

if ($A_1[i] < A_2[j]$) {
 $i++$;

else if ($A_1[i] > A_2[j]$) {
 $j++$;

} else {
 $c++$; $i++$; $j++$;

3 cases
C++, i++, j++
y

Given an array, rotate the array by one position in clock-wise direction.

Example 1:

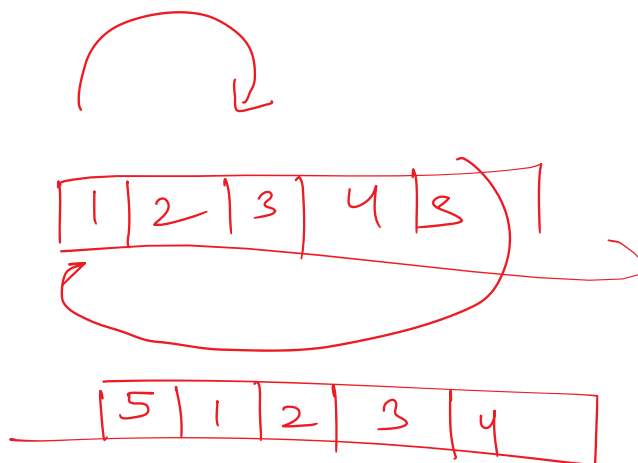
Input:

N = 5

A[] = {1, 2, 3, 4, 5}

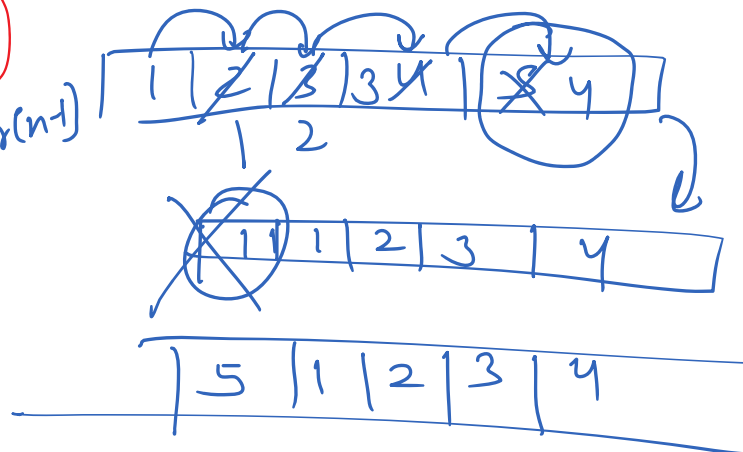
Output:

5 1 2 3 4



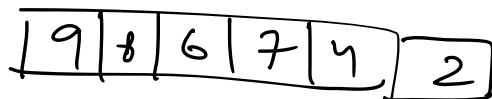
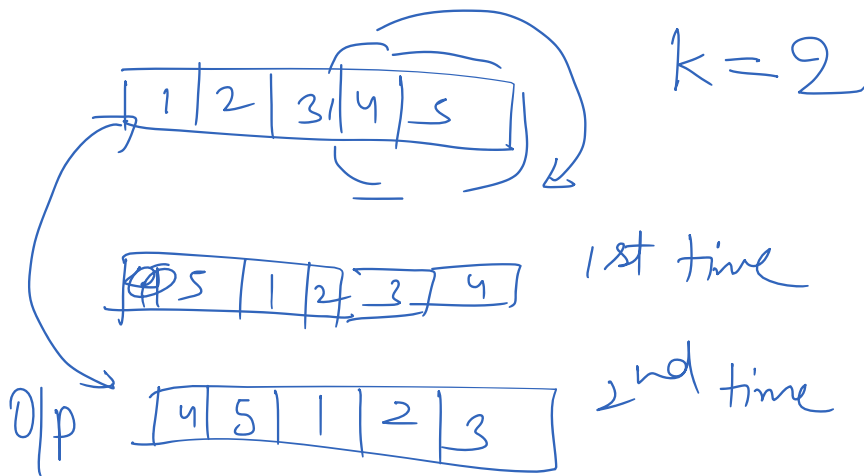
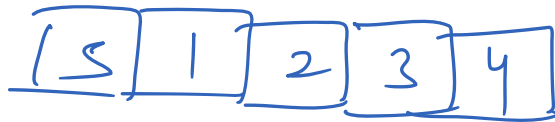
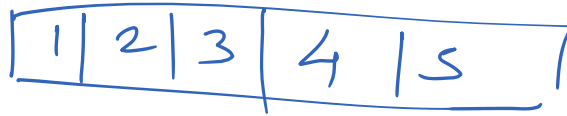
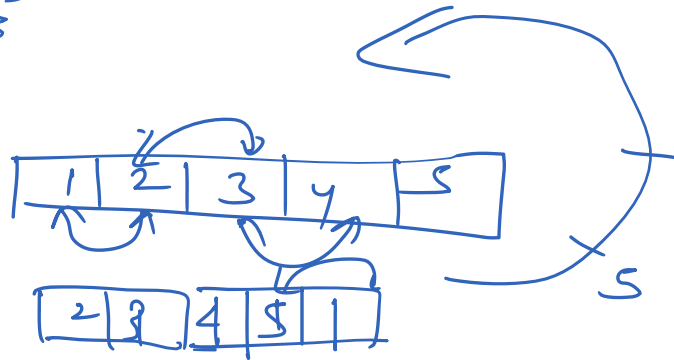
T.C - $O(N)$
S.C - $O(1)$

temp = arr[n-1]



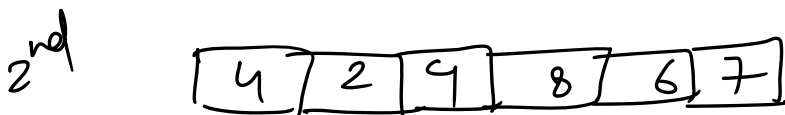
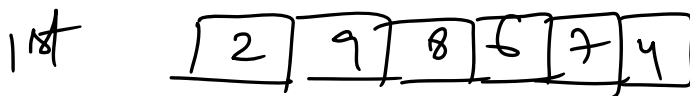
T.C - $O(N)$
S.C - $O(1)$

1. L
S.C - O(1)



$k=4$

4 times clockwise



3rd

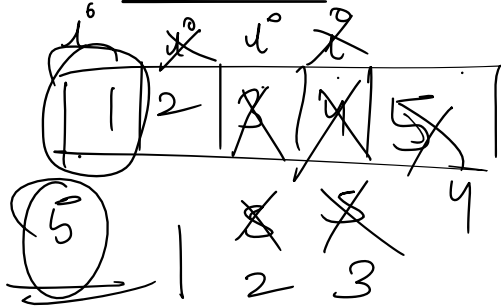
7	4	2	9	8	6
---	---	---	---	---	---

4th

6	7	4	2	9	8
---	---	---	---	---	---

0/p

$$O(n \times k)$$



for (int i = 0; i < n; i++) {

```

int temp = arr[n-1];
int i = n-2;
while(i >= 0){
    arr[i+1] = arr[i];
    i--;
}
arr[0] = temp;

```

$$O(N \times k)$$

TLE

Different Approach

9	8	6	7	4	2
---	---	---	---	---	---

$k=4$

0/p

6	7	4	2	9	8
---	---	---	---	---	---

Step 1 last k elements are reverse

for rest

for $n-1$ ||

9	8	2	4	7	6
---	---	---	---	---	---

Step 2 6 $\xrightarrow{1}$ ~~2~~ $\xrightarrow{n-k}$ elements \rightarrow reverse
 \rightarrow 9 8 ||

8	9	2	4	7	6
---	---	---	---	---	---

Step 3 All elements should be reversed

6	7	4	2	9	8
---	---	---	---	---	---

o/p

T.C = $O(N)$

--

Anticlockwise

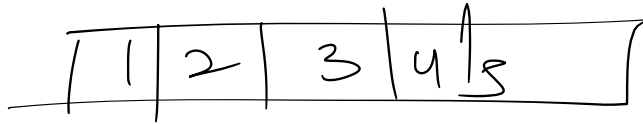
S-1 starting \rightarrow k elements reverse

S-2 $k+1$ to $n-1$ ||

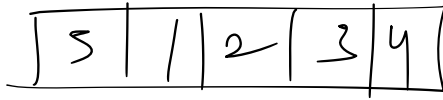
S-3 Array reverse

$k=3$ \rightarrow

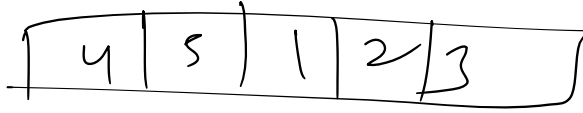
$k=3$ ↓



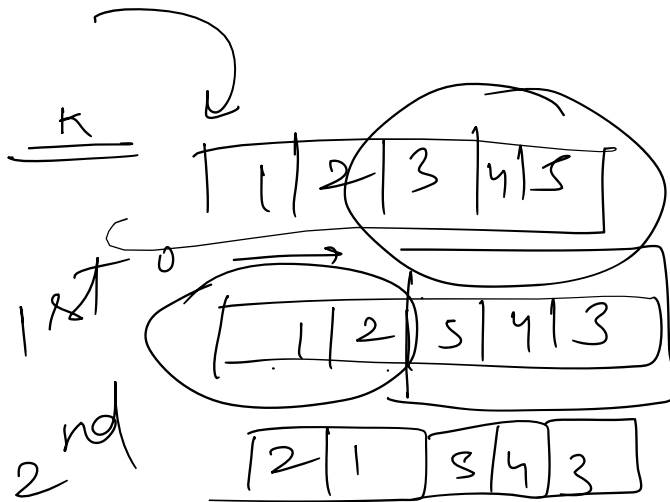
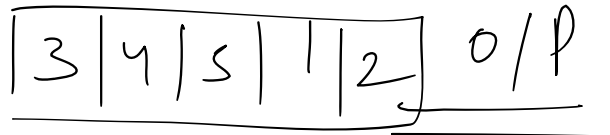
0/p 1st



2nd

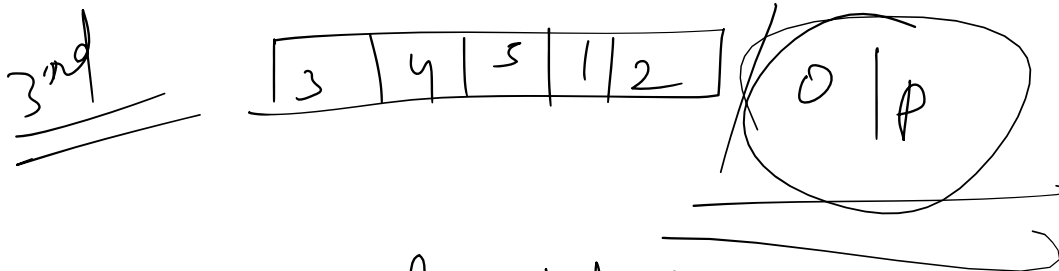


3rd



$$n-k = 5-3$$

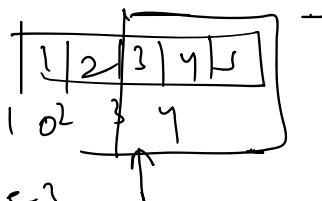
$$k=3$$



Clockwise Implementation

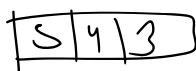
1st step reverse(arr + (n-k), arr + n);

$k=3$



$5-3$

2



$$arr + (5-3), arr + 5$$

$$reverse(arr + 2, arr + 5)$$

2

5	4	3
---	---	---

2nd step $reverse(arr, arr + (n - k))$

④ $\frac{5}{arr + 5}$ $reverse(arr, arr + 2)$

2	1
---	---

5	4	3
---	---	---

3rd step $reverse(arr, arr + n)$

3	4	5	1	2
---	---	---	---	---

 O/P

$k \rightarrow 0 \text{ to } n$

1	2	3	4	5
---	---	---	---	---

 $\xrightarrow{k=10000}$
 $\swarrow n=5$
 1st

5	1	2	3	4
---	---	---	---	---

2nd

4	5	1	2	3
---	---	---	---	---

3rd

3	4	5	1	2
---	---	---	---	---

4th

2	3	4	5	1
---	---	---	---	---

5th

1	2	3	4	5
---	---	---	---	---

 O/P
 $999 \% 5$

$k = k \% n$
 $100 \% 5 = 0$
 T.C - $O(N)$
 S.C - $O(1)$

Given an array **Arr[]** of **N** integers. Find the contiguous sub-array (containing at least one number) which has the maximum sum and return its sum.

Example 1:

Input:

$N = 5$

$Arr[] = \{1, 2, 3, -2, 5\}$

Output:

9

Explanation:

Max subarray sum is 9

of elements $\{1, 2, 3, -2, 5\}$ which is a contiguous subarray.

subarray

$$\frac{n(n+1)}{2}$$

1 1 2 3

$\{1\}, \{2\}, \{3\}, \{1, 2\}, \{2, 3\}, \{1, 2, 3\}$
 $\{1, 2, 3\}$

maximum sum

7 1 1 1

↓ Sum of subarray

1	2	3	4	5	6	8
1	2	3	4	5	6	8

1	2	3	4
---	---	---	---

CS = ~~1~~ ~~2~~ ~~3~~ ~~4~~ ~~5~~ ~~6~~ 8
 LS = ~~INT MIN~~ ~~1~~ ~~2~~ ~~3~~ ~~4~~ ~~5~~ 8

Ans 11

T.C - $O(N)$
 S.C - $O(1)$

Naive Algo

★ start subarray generate krte

★ Sum compare max sum kr return krte

Implement

1	2	3	4	5
---	---	---	---	---

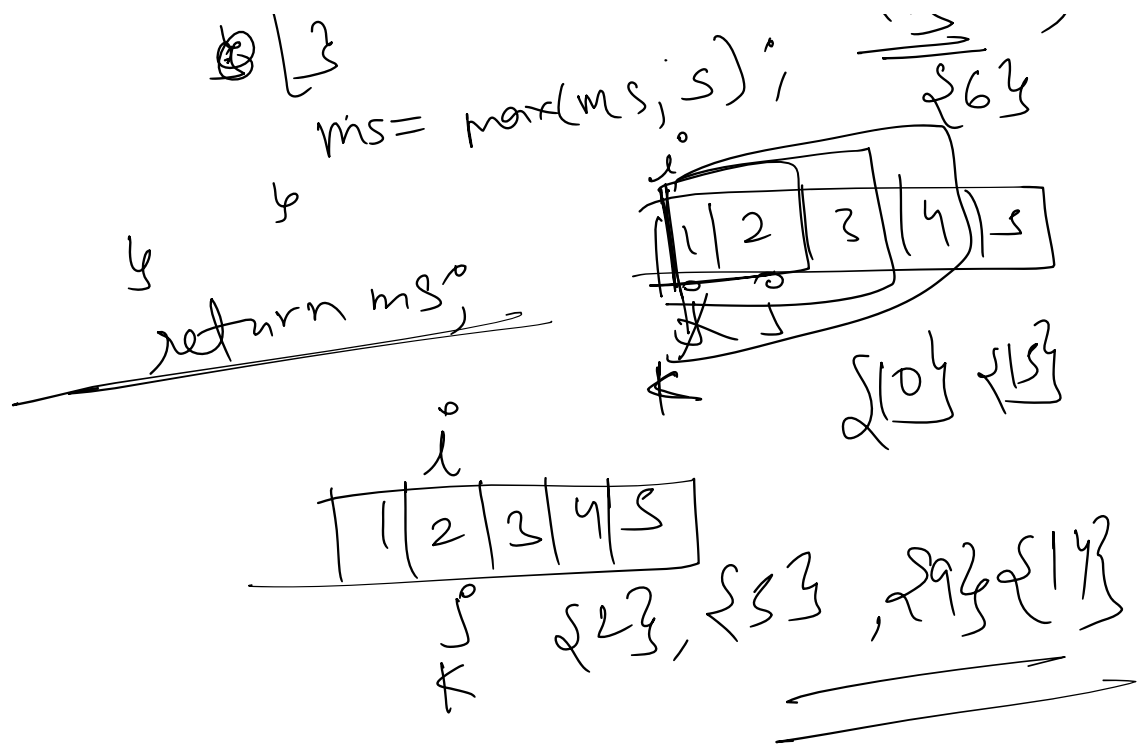
```

int ms = INT_MIN;
for(int i = 0; i < n; i++) {
    for(int j = i; j < n; j++) {
        int s = 0;
        for(int k = i; k <= j; k++) {
            s += arr[k];
        }
        ms = max(ms, s);
    }
}

```

3

2 3 4
 2 3 4
 2 3 4



T.C - $O(N^3)$
 S.C - $O(1)$

TUE

$\{1, 2, 3, 4, 5\}$
 $\{1\}$ $\{3\}$ $\{6\}$ $\{10\}$ $\{15\}$
 $\{2\}$ $\{5\}$ $\{9\}$ $\{14\}$
 $\{3\}$ $\{7\}$ $\{12\}$
 $\{4\}$ $\{8\}$

```

int ms = INT_MIN;
for (int i = 0; i < n; i++) {
    int s = 0;
    for (int j = i; j < n; j++) {
        s += a[j];
        ms = max(ms, s);
    }
}

```

$$MS = \max (ms, s)^0$$

return ms;

THE 

1	2	3	-2	5
---	---	---	----	---

```
for(int i=0; i<n; i++){
    cs += a[i];
    if(cs < a[i]) {
        cs = a[i];
    }
}
```

$$\text{ans} = \max(\text{ans}, \text{cs})$$

Ans = ~~2N~~ \times $\frac{2N}{2}$
CS = 0

(3)

⑥

10

15

44

10 - 1

CS = 1

1	2	3	4	5	6	7
---	---	---	---	---	--------------	---

(3) * (6) * (10) (15)

Handwritten work for Question 10:

Initial array: $2N = 1, 2, 3, 4, 5, 6, 7, 8, 9, 10$

Array after removing elements: $[1, 2, 3, -2, 5]$

Sum of elements: $1 + 2 + 3 + (-2) + 5 = 9$

Final result: 9

$x = 1\%$
 ~~$CS = -3$~~
 $y(-3 < -2)$
 $CS = -2$

$$a_m = -2$$

eg $(-5 < -3)$ ✓

$$\frac{z}{ans} = \frac{-1,73}{-1}$$
$$CS = -3 + -4 = -7$$

of $(-7 < -4)$

$$CS = -4$$

$$i = 4;$$

$$cs = -4 + 0 = -4$$

$$\text{if } (-4 < 0) \{$$

$$cs = 0;$$

$$\text{Ans} = 0$$

Algorithm

① Cumulative Sum Calculate krni krta

$$\text{If } (cs < a[i]) \{$$

$$cs = a[i]$$

}

$$\text{ans} = \max(\text{ans}, cs)$$

$$\text{T.C} - O(N)$$

$$\text{S.C} - \underline{O(1)}$$


```

// Function to find the sum of contiguous subarray with max.
long long maxSubarraySum(int arr[], int n){
    long long ans = INT_MIN;
    long long cs = 0;
    for(int i = 0; i < n; i++){
        cs += arr[i];
        if(cs < arr[i]){
            cs = arr[i];
        }
        ans = max(ans, cs);
    }
    return ans;
}

```

T.C — $O(N)$

S.C — $O(1)$
