```
for (i = 0 to n)
                 if (over [i] == nation)

// element form 1
Q2. Iterative:
            - word insertion- sort (int A [], int n)

for (int i = 1; i < n; i++)
                                  x = A[i];
while (j > -1 && A[i] > x)
A[j+1] = A[j]
                               A [j+1]= x;
            uoid insertion-sort 1 int were I, int n)
                         if (n <= 1).
                             return;
                         insertion-spect ( arc, n-1);
                         int last = are [n-1];
                         nt j = n-a;
```

while
$$|j| = 0 \% \%$$
 area $[j] > last$)

{

over $[j+1] = aver [j];$
 $j--;$

over $[j+1] = last;$

Insertion sort is called online sort because it does not need to know conjething about what values it will sort and the information is requested while the algorithm is running

Other Sorting algorithm:

. Bullele Sort

· Quick Sout

· Merge Sort

· Seclection Sort

· Heap Sort.

| | Best | worst | Asserage |
|----------------|-------------|---------------------|----------------------|
| | 0(12) | 0 (n ²) | 0 (n²). |
| Selection Sort | 0(n) | 0 (n ²) | 0 (n ²). |
| Bulble Sort | O(n) | 0 (n²) | . 27 |
| Insertion Sout | | N(n(ean) | O (n log n). |
| Hear Sout | O (n log n) | 0(n²) | o (n log n) |
| vaice, as | n (n lean) | n(n log n) | O(nlogn) |
| Merge Sort | Ollings | 1000 | |

| Ouich Sort Mary Hoap Sort | Stable Sorting Menge Sort Bubble Insertion Count | Online Souting Insertion |
|------------------------------|--|--------------------------|
|------------------------------|--|--------------------------|

```
Q5. Iterative =
             livrary search (int are E.J., int 1, int se, int key).
                 while ( L = = +c)
                      int m = ((1+2)/2);
if (avel m] = = key).
                          return m;
                       else if ( key < arec [m])
                            2 = m-1 8;
                      else.
1=m+1;
            int binarysearch (int are [], int l, int u, int key).
     Recursive :
                         while (L <= n)
                              int m= (1L+2)/2);
                              it ( key == over [m])
                                   return m;
                              else if ( key < over [m]).
                                   ruturn lunaryslarch (aver, L, mid-1, key);
                                netwen linarysearch (are, mid +1, &, key);
                        return -1;
    Time Complexity: 3.
     · Linear Sewich - O(n)
    · Birary Search - O (log n)
```

- (1)

T(n) = 5(n/2) +1

A 6.

Q10. Worst Case (0 (n2)): The worst case occurs when the picked (5) pued is always un externe (smalls or large) element. This happens when input array is sorted or surverse sorted and either first or last element is picked as pinet. Best Case (0 (slogn)): The last case occurs when we will select purot element as a mean element. o (n log n). Q11. Merge Sout: Best case: - T(n) = 2T(n/2) + O(n) thoust case: T(n) = 2T/1/2) + 0(n). Dutck Sort: - Best Case: - T(n) = 2T(n/2) + O(n) -> O(n logn). Worst Case; - T(n)= T(n-1) + O(n) -> O(n2) Deuck Sort: The average of elements is divided into parts to repeatedly continue of it is not possible to divide it further It is not necessary to divide half. Merge Sort: The elements are split into two sub-array (1/2) again & again until one element to is left 超日12 for luti=0; i<n-1; i++) int min= i; for (int j = i+1; j = n; j++) { b (a E min] > a [j]) min = d', ort key - a [min]. while I min >i). { a [min] = 0 [min - j]; min - -; aci] = key;

```
hoid buildle (vit a CJ, vit n)

{

for (vit i=0; i < n; i++)

int stuaps=0;

for (vit j=0; j < n-i-j; j++)

{

it (a Ej ] > a Ej+17)

vit t = a Cj J;

a Ej ] = a Ej+17;

a Ej+13 = t;

swaps + +;

}

if (swaps=-0)

breake;

}
```