

# Time Complexity Problems

Lecture-20

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## Let's jump to some starters first

```
int c = 0;

for(int i = 1; i <=n; i += i) {

    c++;

}
```

$$i = 1, 2, 4, 8, 16 \dots n$$
  
 $i = 1, 2^1, 2^2, 2^3, 2^4, \dots 2^n$   
 $n+1$  terms

No. of iterations = No. of values

of i

$$= n+1$$

$$T. C. = O(n+1) = O(n)$$

$$2^{n} = n$$

$$T \cdot C \cdot = O(\log_{2} n)$$

$$T \cdot C \cdot = O(\log n)$$

$$a^{n} = b$$

$$\log_{a} b = x$$

$$= 0 \left(\log_{2} n\right) = 0 \left(\log_{2} n\right)$$

$$\log_{2} 2 \cdot \log_{2} n$$

$$\Rightarrow 2^n = n = 0 (K \cdot \log_2 n)$$

$$= \mathcal{O}(\log_2 n)$$



```
int c = 0;
for(int i = 1; i <=n; i += i) {
    for(int j = 0; j < i; j++) {
        c++;
    }
}</pre>
```

$$i = 1, 2, 2^{2}, 2^{3}, \dots n$$
  
 $\sim n = log_{2}n$   
 $i = 1, j = 0 \rightarrow 1$   
 $i = 2, j = 0, 1 \rightarrow 2$   
 $i = 4, j = 0, 1, 2, 3 \rightarrow 4$   
 $i = 8, j = 0, 1, 2, 3, 4, 5, 6, 7 \rightarrow 8$ 

Total No. of Sterations: (M-1) 2x=n

$$= 1 + 2' + 2^2 + 2^3 + 2^4 \cdot \cdot \cdot \cdot \cdot 2^{x}$$
2+1 terms

$$S_{\tau} = a[\tau^{\tau} - 1]$$

$$= \chi \left[ \frac{2^{n+1}-1}{2} \right] = 2^{n+1}-1$$

$$= T \cdot C \cdot = O(2^{n+1}-1) = O(2^{n}) = O(Kn)$$

$$7 \cdot C \cdot = O(n)$$

Total No. of Sterations: 
$$(M-2)$$

$$\Rightarrow 1+2+4+8+16\cdots+\frac{n}{2}+n$$

$$\Rightarrow (2+2)+4+8+16\cdots \frac{n}{2}+n-1$$

$$T \cdot C \cdot = O(2n-1)$$

$$T \cdot C = O(n)$$

R SKILLS 2x - (n = 16) 1+1+2+4+8+16-1 4+4+8+16-1 8 + 8 + 16 - 1

2n-1



## Now see the main course



```
int c = 0;

for(int i = 1; i <=n; i += i) { \longrightarrow T·C· = O(log n)

for(int j = n; j >=0; j--) { \longrightarrow T·C· = O(n)
             C++; >in+1' times
             T.C. = O(n·logn)
```

```
0(logn) i=1 - j=0,4-, 1
int c = 0;
                                  i=2, -j=0,1 - 2
for(int i = 1; i < n; i *= 2) {
                                   i=4: 7j=0,1,2,3 -34
   for(int j = 0; j < i; j++) {
      C++; >i' times > 0(i) d
            0(n)
```

```
i=1,2,4,8,16...22
               ri²<n → i<√n
int c = 0;
for(int i = 1; i * i <=n; i *= 2) {
                                      \rightarrow 2^n \cdot 2^n = n
   for(int j = 0; j < i; j++) {
                                       a 22 = \n
   Total iterations = 1+2+3+4+.... 2x
                  = 22+1-
```

$$\neg T \cdot c \cdot = O(2^n) \Rightarrow T \cdot c \cdot = O(\sqrt{n})$$

```
i = 1, 2, 4, 8. \dots 2^n
                  i<In
                                          \rightarrow 2^n = 5n \rightarrow n = loq_2 n
int c = 0;
                                        i=1, j=n, n-1, \dots 3, 2: n-1
for(int i = 1; i * i < n; i += i) {
                                        i=2, j=n, n-1, ... 4, 3: n-2
    for(int j = n; j > i; j--)
                                        i=4, j=n,n-1,...6,5 : n-4
       C++; - 'n-i'
                                                             : n-8
Total No = (n-1) + (n-2) + (n-4) + (n-8) i= In - n-In
               +(n-16)+...(n-1<sup>n</sup>)
of iterations
```

Total No = (n-1)+(n-2)+(n-4)+(n-8)+(n-16)+...(n-2) (skills

of iterations
$$= (n + n + n + \cdots) - (1 + 2 + 4 + 8 + \cdots + 2^{\times})$$

$$= n(x+1) - [2^{x+1} - 1] = n \cdot x + n - 2 \cdot 2^{\times} + 1$$

$$T \cdot C \cdot = 0 (n \cdot \log_2 \ln + n - \ln) = 0 (n \cdot \log_2 \ln)$$

$$T \cdot C \cdot = 0 (n \cdot \log_2 \ln)$$

$$T \cdot C \cdot = 0 (n \cdot \log_2 \ln)$$

# Thodi si Moths $\log_a b^m = m \cdot \log_a b$ $\Rightarrow \log_2 \ln m = \log_2 n^{1/2} = \frac{1}{2} \log_2 n$

\*Ques : Calculate the time complexity for the following code snippet.

i = 2, 4, 16, 256, (256)

$$\frac{1}{2}$$
,  $2^{2}$ ,  $2^{4}$ ,  $2^{8}$ ,  $2^{16}$ ,  $2^{32}$ ,  $2^{64}$ ...

$$i = 2^{1}, 2^{2^{1}}, 2^{2^{2}}, 2^{2^{3}}, 2^{2^{4}}, 2^{2^{5}}, 2^{2^{6}} \cdot \dots \cdot 2^{2^{5}}$$

$$T \cdot C \cdot = O(x) \Rightarrow 2^{n} = n$$

$$\Rightarrow 2^{t} = n [t = 2^{x}]$$

$$\Rightarrow log_2 n = 2^{\times}$$

$$2^{x} = K \left[ K = log_{2}n \right]$$
 $log_{2}k = X$ 

$$T \cdot C = O(log(logn))$$

```
1 i <= In
int c = 0;
for(int i = 2; i * i <-n; i *= i) {
i = 2, 4, 16, 256, 256*256
i = 2, 2^2, 2^4, 2^8, 2^{16}
i = 2^1, 2^2, 2^2, 2^3, 2^4 \dots 2^{2^{\chi}}
                X+1
```

$$T \cdot C = O(x)$$

$$2^{(2^{k})} = \ln$$

$$\Rightarrow x = \log_{2}(\log_{2} \ln)$$

$$\Rightarrow x = \log_{2}(\log_{2} n^{1/2})$$

$$\Rightarrow x = \log_{2}(\frac{1}{2} \cdot \log_{2} n)$$

$$\Rightarrow x = \log_{2}(\frac{1}{2} \cdot \log_{2} n)$$

$$\Rightarrow T \cdot C = O(\log_{2}(\log_{2} n) + \log_{2}(\frac{1}{2})$$

$$\Rightarrow T \cdot C = O(\log(\log_{2} n))$$

Exponential  $\rightarrow O(2^n) \rightarrow recur sine calls$ 

Sorting ??
Binary Search -> O(log\_n)

### Thank you!

-> 4 Lectures