

# Competitive Algorithm Design and Practice

## Longest Increasing Sub-sequence

### 2014/03/19

---

Yi Long, Lu (mike199250)

*mike199250@gmail.com*

[http://myweb.ncku.edu.tw/~f74991073/20140319\\_DP.zip](http://myweb.ncku.edu.tw/~f74991073/20140319_DP.zip)

Department of Computer Science and Information Engineering  
National Cheng Kung University  
Tainan, Taiwan



# Longest Increasing Sub-sequence

---



# LIS

- Find a **sub-sequence** of a given sequence in which the **sub-sequence's elements** are **in sorted order**, lowest to highest.
- And the sub-sequence is **as long as possible**.

1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15



# LIS

1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15

- An increasing sub-sequence:

1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15

- The longest one:

1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15



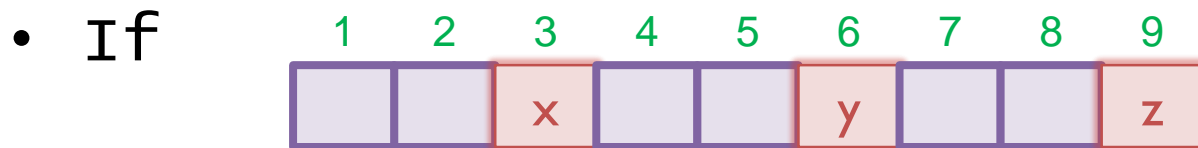
# LIS

- Naïve solutions:
  - For every sub-sequence, check if it is LIS.
- Time-complexity:
  - Every sub-sequence,  $O(2^N)$ .
  - For each sub-sequence checking needs  $O(N)$ .
  - Total:  $O(N * 2^N)$

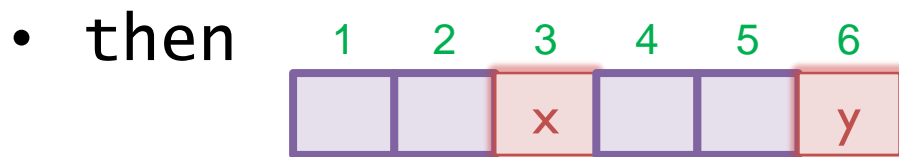


# LIS

- Observation:



- is LIS ended at z,



- must be LIS ended at y



# LIS

- what do we want to know?
  - The LIS ended at index  $k$ , i.e.  $LIS[k]$
- How can we get that?
  - Find the previous number with longest LIS.
  - $LIS[k] =$   
 $\max( 1, LIS[i]+1 )$  for  $i < k$  &  $num[i] < num[k]$



# LIS

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS																





# LIS

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1															



# LIS

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1															



# LIS

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	1														



# LIS

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2														



# LIS

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2														



# LIS

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	1													



# LIS

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2													



# LIS

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2													





# LIS

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	1												



# LIS

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	2												



# LIS

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3												



# LIS

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3												



# LIS

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3												



# LIS

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6



# LIS

- Time-complexity:
  - For each element find the previous longest LIS.
  - Every element,  $O(N)$
  - Each element, check previous elements,  $O(N)$
  - Total:  $O(N*N) \Rightarrow O(N^2)$
- Space-complexity:
  - Additional LIS array,  $O(N)$



# LIS

```

1  /* file name: LIS.c */
2  #include <stdio.h>
3
4  int num[17]={0,0,8,4,12,2,10,6,14,1,9,5,13,3,11,7,15};
5  int LIS[17];
6  void Find_LIS()
7  {
8      int i,j;
9      for(i=1;i<=16;i++)
10     {
11         LIS[i]=1;
12         for(j=1;j<i;j++)
13             if(num[j]<num[i] && LIS[j]+1>LIS[i])
14                 LIS[i]=LIS[j]+1;
15     }
16 }
17 int main()
18 {
19     int i;
20     Find_LIS();
21     printf("num:");
22     for(i=1;i<=16;i++)printf("%3d",num[i]);
23     printf("\nLIS:");
24     for(i=1;i<=16;i++)printf("%3d",LIS[i]);
25     putchar('\n');
26     return 0;
27 }
28

```

```

num:  0  8  4 12  2 10  6 14  1  9  5 13  3 11  7 15
LIS:  1  2  2  3  2  3  3  4  2  4  3  5  3  5  4  6

Process returned 0 (0x0)   execution time : 0.049 s
Press any key to continue.

```





# POJ 2533

---



# Learn more!

- Sometimes we need to **output** a **solution** too, but **how**?
- 1. Find the number **backwards**.
  - Assume  $\text{ans} = \text{length of LIS}$
  - Find  $k$  such that  $\text{LIS}[k] = \text{ans}$ .
  - Then from index  $1 \sim k$  find  $i$  such that  $\text{LIS}[i] = \text{LIS}[k] - 1$  &&  $\text{num}[i] < \text{num}[k]$ , and so on
- 2. Additional array **pre**.
  - $\text{pre}[k] = i$  such that  $\text{LIS}[i] = \text{LIS}[k] - 1$  &&  $\text{num}[i] < \text{num}[k]$



# Learn more!

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6
pre	1	1	1	2	1	2	5	4	11	7	3	10	5	10	7	14



# Learn more!

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6
pre	1	1	1	2	1	2	5	4	11	7	3	10	5	10	7	14



# Learn more!

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6
pre	1	1	1	2	1	2	5	4	11	7	3	10	5	10	7	14



# Learn more!

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6
pre	1	1	1	2	1	2	5	4	11	7	3	10	5	10	7	14



# Learn more!

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6
pre	1	1	1	2	1	2	5	4	11	7	3	10	5	10	7	14



# Learn more!

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6
pre	1	1	1	2	1	2	5	4	11	7	3	10	5	10	7	14





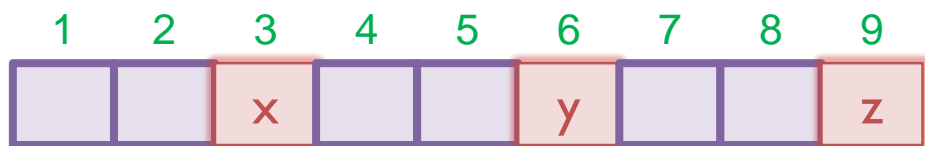
# Learn more!

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6
pre	1	1	1	2	1	2	5	4	11	7	3	10	5	10	7	14



# Learn more!

- Sometimes  $O(N^2)$  is too slow.....



- If  $LIS[x] = LIS[y]$  (mean-while  $x \geq y$ )
- For some  $z$  that  $LIS[z] = LIS[x] + 1$   
We can append  $z$  after either  $x$  or  $y$  to form LIS, by the way  $y$  might be smaller!
- That is, for  $LIS[] = k$  we can memorize just one number, the smaller one,  $y$ .



# Learn more!

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6



# Learn more!

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6

LIS = 1

0



# Learn more!

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6

LIS =2

LIS =1

8
0



# Learn more!

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6

LIS =2

LIS =1

4
0



# Learn more!

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6

LIS =3

LIS =2

LIS =1

12
4
0



# Learn more!

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6

LIS =3

LIS =2

LIS =1

12
2
0





# Learn more!

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6

LIS =3

LIS =2

LIS =1

10
2
0



# Learn more!

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6

LIS =3

LIS =2

LIS =1

6
2
0



# Learn more!

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6

LIS =4

LIS =3

LIS =2

LIS =1

14
6
2
0



# Learn more!

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6

LIS =4

LIS =3

LIS =2

LIS =1

14
6
1
0



# Learn more!

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6

LIS =4	9
LIS =3	6
LIS =2	1
LIS =1	0



# Learn more!

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6

LIS =4

9

LIS =3

5

LIS =2

1

LIS =1

0



# Learn more!

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6

LIS =5

13

LIS =4

9

LIS =3

5

LIS =2

1

LIS =1

0



# Learn more!

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6

LIS =5

13

LIS =4

9

LIS =3

3

LIS =2

1

LIS =1

0





# Learn more!

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6

LIS =5

11
9
3
1
0

LIS =4

LIS =3

LIS =2

LIS =1



# Learn more!

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6

LIS =5

11
7
3
1
0

LIS =4

LIS =3

LIS =2

LIS =1



# Learn more!

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
num	0	8	4	12	2	10	6	14	1	9	5	13	3	11	7	15
LIS	1	2	2	3	2	3	3	4	2	4	3	5	3	5	4	6

LIS =6	15
LIS =5	11
LIS =4	7
LIS =3	3
LIS =2	1
LIS =1	0



# Learn more!

---

- There will be some **ordering**, so we can use **binary-search** to find  $LIS[k]$ , then update the array.
- Find the best place it can be, and update
- It`s greedy!

