

# Probabilistic Method and Random Graphs

## Lecture 3. Chernoff bounds: behind and beyond

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Questions, comments, or suggestions?

# A brief review

## Moments

Expectation,  $k$ -moment, variance

## Inequalities

Universal: Union bound

1-moment: Markov's inequality

2-moment: Chebychev's inequality

## Chernoff bounds: independent sum

Let  $X = \sum_{i=1}^n X_i$ , where  $X_i$ 's are **independent** Poisson trials. Let  $\mu = \mathbb{E}[X]$ . Then

1. For  $\delta > 0$ ,  $\Pr(X \geq (1 + \delta)\mu) \leq \left(\frac{e^\delta}{(1+\delta)^{(1+\delta)}}\right)^\mu \leq e^{-\frac{\delta^2}{2+\delta}\mu}$ .
2. For  $1 > \delta > 0$ ,  $\Pr(X \leq (1 - \delta)\mu) \leq \left(\frac{e^{-\delta}}{(1-\delta)^{(1-\delta)}}\right)^\mu \leq e^{-\frac{\delta^2}{2}\mu}$ .

# General bounds for independent sums

Each  $X_i \in \{0, a_i\}$  where  $a_i \leq 1$

Basic Chernoff bounds remain valid, by Homework 2 of Week 2.

Each  $X_i \in [0, 1]$  but is not necessarily a Poisson trial

Basic Chernoff bounds remain valid (by  $e^{\lambda x} \leq xe^{1\lambda} + (1-x)e^{0\lambda}$ ).

The domains  $(a_i, b_i)$  of  $X_i$ 's differ

Hoeffding's Inequality:  $\Pr(|X - \mathbb{E}[X]| \geq t) \leq 2e^{-\frac{2t^2}{\sum_i (b_i - a_i)^2}}$ .  
Proposed in 1963.

Remarks of Hoeffding's Inequality

1. It considers the absolute, rather than relative, deviation.  
Particularly useful if  $\mu = 0$ .
2. When each  $X_i \in [0, s]$ , it is tighter than the simplified basic Chernoff bounds if  $\delta$  is big, and looser otherwise.

# Hoeffding's Inequality

Let  $X = \sum_{i=1}^n X_i$ , where  $X_i \in [a_i, b_i]$  are independent r.v. Then  
 $\Pr(|X - \mathbb{E}[X]| \geq t) \leq 2e^{-\frac{2t^2}{\sum_i (b_i - a_i)^2}}$  for any  $t > 0$

## Idea of the proof

1. Given r.v.  $Z \in [a, b]$  with  $\mathbb{E}[Z] = 0$ ,  $\mathbb{E}[e^{\lambda Z}] \leq e^{\frac{\lambda^2(b-a)^2}{8}}$   
-Hoeffding's Lemma
- 2.

$$\begin{aligned}\Pr(X - \mathbb{E}[X] \geq t) &\leq \frac{\prod_i \mathbb{E}[e^{\lambda(X_i - \mathbb{E}[X_i])}]}{e^{\lambda t}} \\ &\leq e^{\lambda^2 \sum_i \frac{(b_i - a_i)^2}{8} - \lambda t}\end{aligned}$$

3. Choose  $\lambda$  to minimize RHS. Likewise for  $\Pr(X - \mathbb{E}[X] \leq -t)$ .

# Proof of Hoeffding's Lemma

**Lemma:** Given r.v.  $Z \in [a, b]$  with  $\mathbb{E}[Z] = 0$ ,  $\mathbb{E}[e^{\lambda Z}] \leq e^{\frac{\lambda^2(b-a)^2}{8}}$

$$e^{\lambda z} \leq \frac{z-a}{b-a}e^{\lambda b} + \frac{b-z}{b-a}e^{\lambda a}, \text{ for } z \in [a, b]$$

$$\begin{aligned}\mathbb{E}[e^{\lambda Z}] &\leq \frac{be^{\lambda a}}{b-a} - \frac{ae^{\lambda b}}{b-a} \\ &= (1 - \theta + \theta e^u)e^{-\theta u} \quad \text{where } \theta = \frac{-a}{b-a}, u = \lambda(b-a) \\ &= e^{\phi(u)} \quad \text{where } \phi(u) \triangleq -\theta u + \ln(1 - \theta + \theta e^u)\end{aligned}$$

Taylor expansion  $\phi(u) = \phi(0) + \phi'(0) + \frac{\phi''(\xi)}{2}u^2$ .

Then  $\phi(u) \leq \frac{u^2}{8}$  since  $\phi(0) = \phi'(0) = 0, \phi''(\xi) \leq \frac{1}{4}$

# Example: Hoeffding's Inequality + Union bound

## Set balancing

Given a matrix  $A \in \{0, 1\}^{n \times m}$ , find  $b \in \{-1, 1\}^m$  s.t.  $\|Ab\|_\infty$  is minimized.

## Motivation

feature 1:  $\begin{bmatrix} a_{11} & a_{12} & \cdots & a_{1m} \end{bmatrix}$   
feature 2:  $\begin{bmatrix} a_{21} & a_{22} & \cdots & a_{2m} \end{bmatrix}$   
 $\vdots$   
feature  $n$ :  $\begin{bmatrix} \vdots & \vdots & \ddots & \vdots \end{bmatrix}$   
feature  $n$ :  $\begin{bmatrix} a_{n1} & a_{n2} & \cdots & a_{nm} \end{bmatrix}$ , each column is an object.  
Want to partition the objects so that every feature is balanced.

# Example: Hoeffding's Inequality + Union bound

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Want to partition the objects so that every feature is balanced.

## Algorithm

Uniformly randomly sample  $b$ .



# Performance analysis

## Performance

$$\Pr(\|Ab\|_\infty \geq \sqrt{4m \ln n}) \leq \frac{2}{n}$$

## Proof

For any  $1 \leq i \leq n$ ,  $Z_i = \sum_j a_{ij} b_j$  is the  $i$ th entry of  $Ab$ . By union bound, it suffices to prove  $\Pr(|Z_i| \geq \sqrt{4m \ln n}) \leq \frac{2}{n^2}$  for each  $i$ .

Fix  $i$ . W.l.o.g, assume  $a_{ij} = 1$  iff  $1 \leq j \leq k$  for some  $k \leq m$ . Then  $Z_i = b_1 + \dots + b_k$ .

Note that  $b_j$ 's are independent over  $\{-1, 1\}$  with  $\mathbb{E}[b_j] = 0$ .

By Hoeffding's Inequality,  $\Pr(|Z_i| \geq \sqrt{4m \ln n}) \leq 2e^{-\frac{8m \ln n}{4k}} \leq \frac{2}{n^2}$

# Reflection on moments and Chernoff bounds

## Chernoff Bounds

Why is it so good?

Can it be improved by non-exponential functions?

Anything to do with moments?

## Moments

Do moments uniquely determine the distribution?

The story begins with generating functions.

# Generating functions

## Informal definition

A power series whose coefficients encode information about a sequence of numbers.

## Example: Probability generating function

Given a discrete random variable  $X$  whose values are non-negative integers,  $G_X(t) \triangleq \sum_{n \geq 0} \Pr(X = n)t^n = \mathbb{E}[t^X]$ .

Example: Bernoulli and binomial random variables.

## Properties

**Convergence:** It converges if  $|t| < 1$ .

**Uniqueness:**  $G_X(\cdot) \equiv G_Y(\cdot)$  implies the same distribution.

## Application

Toy: Use uniqueness to show that the summation of independent identical binomial distribution is binomial.

Deriving Moments:  $G_X^{(k)}(1) = \mathbb{E}[X(X-1) \cdots (X-k+1)]$ .

# Moment generating functions

## Shortcoming of probability generating functions

Only valid for non-negative integer random variables.

## Moment generating functions

$$M_X(t) \triangleq \sum_x \Pr(X = x)e^{tx} = \mathbb{E}[e^{tX}].$$

Example of Bernoulli and binomial distributions.

## Properties

- If  $M_X(t)$  converges around 0,  $M_X^{(k)}(0) = \mathbb{E}[X^k]$ , meaning the moments are exactly the coefficients of the Taylor's expansion.
- **Convergence:**  $M_X(t)$  converges when  $X$  is bounded.
- If independent,  $M_{X+Y} = M_X M_Y$ .
- **Uniqueness:** If  $M_X(t)$  converges around 0, the distribution is uniquely determined by the moments. (Why? See later)

## Moments generating function may not converge

Cauchy distribution: density function  $f(x) = \frac{1}{\pi(1+x^2)}$  does not have moments for any order.

## An example of non-uniqueness of moments

Log-Normal-like distribution:

density function  $f_{X_n}(x) = \frac{e^{-\frac{1}{2}(\ln x)^2}}{\sqrt{2\pi x}} (1 + \sin(2n\pi \ln x))$ .

$k$ -Moments  $\mathbb{E}[X_n^k] = e^{k^2/2}$  for non-negative integers  $k$ .

# Characteristic functions

## Definition

$\varphi_X(t) \triangleq \int_{\mathbb{R}} e^{itx} dF_X(x)$  where  $i = \sqrt{-1}$  and  $t$  is real.

## Properties

**Convergence:** It always exists.

**Uniqueness:** It uniquely determines the distribution.

Rationale of the uniqueness.

## Uniqueness of convergent moments generating functions

Suppose  $M_X(t) = M_Y(t)$  converges around 0.

- $\phi_X(t)$  and  $\phi_Y(t)$  can be extended to the belt with small imaginary part (since formally,  $M_X(t) = \phi_X(it)$ )
- $\phi_X(t) = \phi_Y(t)$  when  $t$  is purely imaginary in this belt
- By the unique continuation theorem of analytic complex functions, the characteristic functions are equal

## Moments

Do moments uniquely determine the distribution?

Yes, but conditionally.

## Chernoff Bounds

- Why is it so good?
- Can it be improved by non-exponential functions?
- Anything to do with moments?

What's your answer?

# A story of generating function

Introduced in 1730 by Abraham de Moivre, to solve the general linear recurrence problem

Wisdom: A generating function is a clothesline on which we hang up a sequence of numbers for display. -Herbert Wilf

Application to Fibonacci numbers (by courtesy of de Moivre):

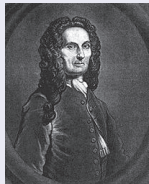
$$F(x) = \sum_{n=0}^{\infty} F_n x^n = x + \sum_{n=2}^{\infty} (F_{n-1} + F_{n-2}) x^n = x + xF(x) + x^2 F(x)$$

$$\Rightarrow F(x) = \frac{x}{1-x-x^2} = \frac{1}{\sqrt{5}} \left( \frac{\phi}{x+\phi} - \frac{\psi}{x+\psi} \right) = \sum_{n=0}^{\infty} \frac{1}{\sqrt{5}} (\phi^n - \psi^n) x^n$$

$$\Rightarrow F_n = \frac{1}{\sqrt{5}} (\phi^n - \psi^n).$$



# Brief introduction to Abraham de Moivre



- May 26, 1667-  
Nov. 27, 1754
- A French  
mathematician
- de Moivre's formula
- Binet's formula
- Central limit theorem
- Stirling's formula

## Legend

- Friends: **Isaac Newton**, Edmond Halley, and James Stirling
- Struggled for a living and lived for mathematics
- The Doctrine of Chances was prized by gamblers
  - 2nd probability textbook in history
- Predicted the exact date of his death

# Chernoff bound in a big picture

## Fundamental laws of probability theory

**Law of large numbers** (Cardano, Jacob Bernoulli 1713, Poisson 1837): The sample average converges to the expected value.

**Central limit theorem** (Abraham de Moivre 1733, Laplace 1812, Lyapunov 1901, Pólya 1920): The arithmetic mean of independent random variables is approximately normally distributed.

$$\lim_{n \rightarrow \infty} \Pr \left( \sqrt{n} \left( \left( \frac{1}{n} \sum_{i=1}^n X_i \right) - \mu \right) \leq x \right) = \Phi \left( \frac{x}{\sigma} \right)$$

## Marvelous but ...

Say nothing about the rate of convergence

## Large deviation theory

How fast does it converge? Beyond central limit theorem

# A glance at large deviation theory

## Motivation

$X_n$ : the number of heads in  $n$  flips of a fair coin.

By the central limit theorem,  $\Pr(X_n \geq \frac{n}{2} + \sqrt{n}) \rightarrow 1 - \Phi(1)$ .

What about  $\Pr(X_n \geq \frac{n}{2} + \frac{n}{3})$ ? Nothing but converging to 0.

## Chernoff bounds say...

$$\Pr(X_n \geq \frac{n}{2} + \frac{n}{3}) \leq \left( \frac{e^{\frac{2}{3}}}{(\frac{5}{3})^{\frac{5}{3}}} \right)^{\frac{n}{2}} \approx e^{-0.092n}.$$

## Actually

Direct calculation shows that

$$\Pr(X_n \geq \frac{n}{2} + \frac{n}{3}) \approx e^{-0.2426n + o(n)} \ll \text{Chernoff bound}.$$

Oh, no!

# Mission of Large Deviation Theory

Find the asymptotic probabilities of *rare* events - how do they decay to 0 as  $n \rightarrow \infty$ ?

*Rare* events mean large deviation.

So large that CLT is almost useless (deviation up to  $\sqrt{n}$ ).

## Intuition

Inspired by Chernoff bounds, conjecture that probabilities of rare events will be exponentially small in  $n$ :  $e^{-cn}$  for some  $c$ .

Q: Does  $\lim_{n \rightarrow \infty} \frac{1}{n} \ln \Pr(\mathcal{E}_n^{\text{rare}})$  exist? If so, what's it?

# Large Deviation Principle

Simple form (By courtesy of Cramer, 1938)

Let  $X_1, \dots, X_n, \dots \in \mathbb{R}$  be i.i.d. r.v. which satisfy  $\mathbb{E}[e^{tX_1}] < \infty$  for  $t \in \mathbb{R}$ . Then for any  $t > \mathbb{E}[X_1]$ , we have

$$\lim_{n \rightarrow \infty} \frac{1}{n} \ln \Pr \left( \sum_{i=1}^n X_i \geq tn \right) = -I(t),$$

where

$$I(t) \triangleq \sup_{\lambda > 0} \lambda t - \ln \mathbb{E}[e^{\lambda X_1}].$$

Remark

$I(\cdot)$ : rate function.

Many variants: the factor  $\frac{1}{n}$ , random variables

# Large Deviation Principle: Proof

## Large Deviation Principle

$$\lim_{n \rightarrow \infty} \frac{1}{n} \ln \Pr(\sum_{i=1}^n X_i \geq tn) = -(\sup_{\lambda > 0} \lambda t - \ln \mathbb{E}[e^{\lambda X_1}]).$$

## Proof: Upper bound

Let  $Y_n = \frac{\sum_{i=1}^n X_i}{n}$ ,  $M(\lambda) = \mathbb{E}[e^{\lambda X_1}]$ , and  $\psi(\lambda) = \ln M(\lambda)$ .

$$\Pr(Y_n \geq t) \leq e^{-\lambda nt} (M(\lambda))^n \text{ for any } \lambda \geq 0.$$

$$\frac{1}{n} \ln \Pr(Y_n \geq t) \leq -\lambda t + \psi(\lambda).$$

$$\frac{1}{n} \ln \Pr(Y_n \geq t) \leq -\sup_{\lambda \geq 0} (\lambda t - \psi(\lambda)).$$

# Large Deviation Principle: Proof

## Lower bound

The maximizer  $\lambda_0$  of  $\lambda t - \psi(\lambda)$  satisfies  $t = \int \frac{x e^{\lambda_0 x}}{M(\lambda_0)} d\mu(x)$ .

Let  $d\mu_0(x) = \frac{e^{\lambda_0 x}}{M(\lambda_0)} d\mu(x)$ . Its expectation  $\int x d\mu_0(x) = t$ .

Let  $A = \{Y_n \geq t\} \subseteq \mathbb{R}^n$ ,  $A_\delta = \{Y_n \in [t, t + \delta]\} \subseteq \mathbb{R}^n$ .

$$\begin{aligned} \Pr_\mu(A) &\geq \Pr_\mu(A_\delta) = \int_{A_\delta} \prod_{i=1}^n d\mu(x_i) \\ &= \int_{A_\delta} (M(\lambda_0))^n e^{-\lambda_0 \sum_{i=1}^n x_i} \prod_{i=1}^n d\mu_0(x_i) \\ &\geq \left( M(\lambda_0) e^{-\lambda_0(t+\delta)} \right)^n \Pr_{\mu_0}(A_\delta). \end{aligned}$$

Applying CLT to  $\mu_0$ , we have  $\lim_{n \rightarrow \infty} \Pr_{\mu_0}(A_\delta) = \frac{1}{2}$ .

$\lim_{n \rightarrow \infty} \frac{1}{n} \ln \Pr(Y_n \geq t) \geq \psi(\lambda_0) - (t + \delta)\lambda_0$ , and let  $\delta \rightarrow 0$ .

Large deviation theory vs CLT

Seemingly easy to get exponential decay in many cases, but hard to calculate.

Chernoff bounds concern large deviation

- Con: Generally weaker
- Pro: Always holds, not just asymptotically

Key assumption

**Independence!**



# One-order generalization of Chernoff Bound

<b>General</b> $f(X_1, \dots, X_n)$	<b>McDiarmid</b> <b>1989</b>	<b>Zhang, Liu et al.</b> <b>2019</b>	<b>Kontorovich et</b> <b>al. 2008</b>
	<b>Chernoff</b> <b>1948</b>	<b>Janson</b> <b>2004</b>	<b>Bosq 2012</b>
<b>Linear</b> $X_1 + \dots + X_n$	<b>Independent</b>	<b>Dependent</b> <b>(Qualitative)</b>	<b>Dependent</b> <b>(Quantitative)</b>

- ❶ <http://nowak.ece.wisc.edu/SLT07/lecture7.pdf>
- ❷ <https://www.math.illinois.edu/~psdey/414CourseNotes.pdf>
- ❸ When Do the Moments Uniquely Identify a Distribution
- ❹ <http://willperkins.org/6221/slides/largedeviations.pdf>