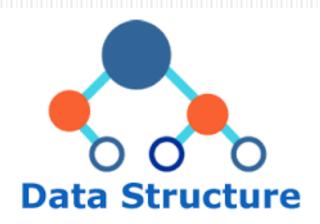
### Unit-01 Introduction to Data Structure



### Topic to be covered

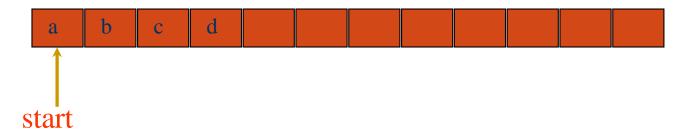
- Array:
  - Representation of arrays
    - One dimensional array
    - Two dimensional array
  - · Applications of arrays
    - Symbol Manipulation (matrix representation of polynomial equation)
    - Sparse matrix
  - Sparse matrix and its representation

## 1- D Array

- Simplest data structure that makes use of computer address to locate its elements is the 1-dimensional array or vector
- Number of memory locations is sequentially allocated to the vector
- A vector size is fixed and therefore requires a fixed number of memory locations

### 1-D Array Representation

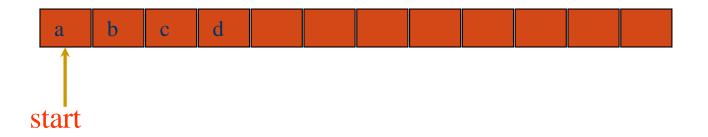
### Memory



- 1-dimensional array x = [a, b, c, d]
- map into contiguous memory locations
- location(x[i]) = start + i

### Space Overhead

### Memory



space overhead = 4 bytes for start

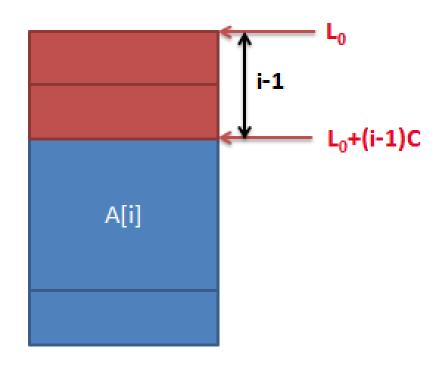
### 1-D Array Size

- A array will store five integer values, its name is num.
- it can be visualized as show.

index	data	
num[0]	10	
num[1]	20	
num[2]	30	
num[3]	40	
num[4]	50	
size of array=(upper bo	ound – lower bound) + 1	1
=4+0+1		
=5		

### 1-D Array Location

• Array A with subscript lower bound of "one" is represented as below.



- L<sub>0</sub> is the address of the first word allocated to the first element of vector A
- C words is size of each element or node
- The address of element Ai is Loc(Ai)=L<sub>0</sub>+(C\*(i-1))
- Let's consider the more general case of a vector A with lower bound for it's subscript is given by some variable b.
- The address of element Ai is

Loc (Ai) = 
$$L_0 + (C^*(i-b))$$

7

## 2-D Array

- 2-D arrays are also called table or matrix
- 2-D arrays have two subscripts- row & column

### 2-D Arrays

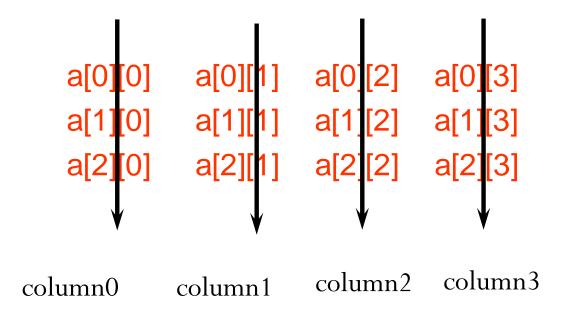
The elements of a 2-dimensional array a declared as: int a[3][4];

```
may be shown as a table
a[0][0] a[0][1] a[0][2] a[0][3]
a[1][0] a[1][1] a[1][2] a[1][3]
a[2][0] a[2][1] a[2][2] a[2][3]
```

# Rows of a 2-D Array

	2[U][1]	101101c	2[0][3]	row 0	
مراهارها	ا الاماله	مران الح	مراحارحا	1044-0	
(۱۱۱۱هـ	[1][1]د	2[1][2]	2[1][2]	row 1	
المال، الم	∞[ , ][ , ]	م[۱][۲]	الم][ ،	1011	
101101	101[4]	101101	101101	KOW 9	
		مركاركا	a[۷][۷]	TOW Z	

### **Columns of a 2-D Array**



### 2-D Array

Initialization of Two dimensional arrays: int arr[2][3]={

• Bellow show the data in 2-D grid.

	0	1	2
0	10	20	30
1	40	50	60

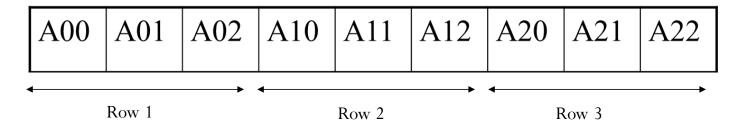
### Implementation of 2-D Array in memory

- A two-dimensional array can be implemented in a programming language in two ways:
  - Row-major implementation/ Row-major order matrix
  - Column-major implementation/ Column-major order matrix

## Row-major order matrix

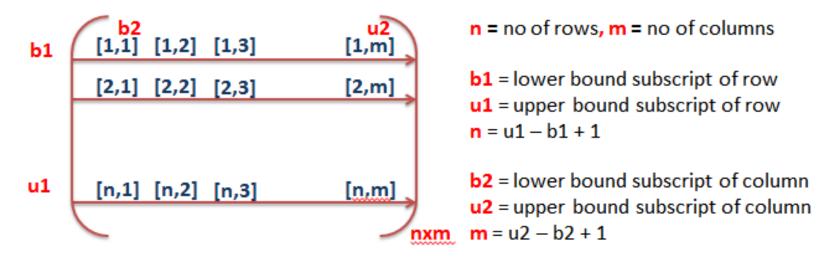
• 2-D array in which elements are stored row by row

For example an array [3][3] is stored in the memory as show bellow:



## 2-D array row-major order matrix

The storage can be clearly understood by arranging array as matrix:



The address element A [ i , j ] is given by

Loc (A [ i , j ]) = 
$$L_0$$
 + (i-b1)\*(u2-b2+1) + (j - b2)

## With Storage space

• storage space of each element is represented by W

Then formula will be-

$$Loc(A[i, j]) = Lo+W(m*(i-b_1)+(j-b_2))$$

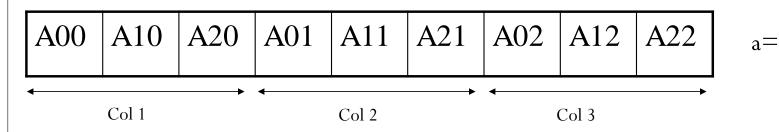
## **Example-**

A 2-D array defined as a [4:7,-1:3] requires 2 bytes of storage space for each element. If the array is stored in row-major form, then calculate the address of element at location a [6,2] given base address is 100.

- **Sol:** L<sub>0</sub>=100,  $b_1$ =4,  $b_2$ =-1,  $u_1$ =7,  $u_2$ =3, W=2(int size)
- i=6, j=2 and  $m(no. of col)= u_2-b_2+1=3-(-1)+1=5$
- Address of  $a[i][j] = L_0 + W(m(i-b_1) + (j-b_2))$
- Address of a[6][2]=100+2(5(6-4) + (2-(-1)) =100+(2(5\*2+3)) =100+26 =126

## 2-D array column-major order matrix

- 2-D array in which elements are stored column by column
- For example an array [3][3] is stored in the memory as show bellow:



A00	A01	A02
A10	A11	A12
A20	A21	A22

Col 1 Col 2 Col 3

## 2-D array column-major order matrix

The address of element A [ i , j ] can be obtained by expression

Loc (A [ i , j ]) = 
$$L_0$$
 + (j-1)\*2 + (i-1)  
Loc (A [2, 3]) =  $L_0$  + (3-1)\*2 + (2-1) =  $L_0$  + 5

## 2-D array column-major order matrix

In general for two dimensional array consisting of n rows and m columns the address element A [i,j] is given by

$$Loc(A[i, j]) = Lo+W(n*(j-b_2) + (i-b_1))$$

W=size of each element array element

n=the number of row  $(u_1-b_1+1)$ 

b<sub>1</sub> the lower bound of row

u<sub>1</sub> upper bound of row

b<sub>2</sub> is lower bound of column

u<sub>2</sub> upper bound of column

## 2-D Arrays: Column-major Implementation

### Example-

A two-dimensional array defined as a[-20:20,10:35] requires one bytes of storage space for each element. If the array is stored in column-major form, then calculate the address of element at location a[0,30] given base address is 500.

### **Solution-**

$$B=500, b_1=-20, b_2=10, u_1=20, u_2=35$$
W=1 byte
$$i=0,j=30 \text{ and } n \text{ (no. of rows)}= u_1-b_1+1$$

$$=20-(-20)+1=41$$
Address of a[i][j]= B+W(n(j-b\_2)+(i-b\_1))
Address of a[0][30]=500+1(41(30-10)+(0-(-20))
$$=500+1(41*20+20)$$

$$=500+840$$

$$=1340$$

### Q-1

Let A be a two-dimensional array declared as follows:

Assuming that each integer takes one memory location, the array is stored in **row-major** order and the first element of the array is stored at location 100, what is the address of the element A[7][12]?

Loc A[i][j] = 
$$100 + [(i-1)(15-1+1)+(j-1)] \times 1$$
  
=  $100 + (i-1)15 + j-1$   
=  $15i + j + 84$   
=  $15*7 + 12 + 84$   
=  $201$ 

## **Multi-Dimensional Array**

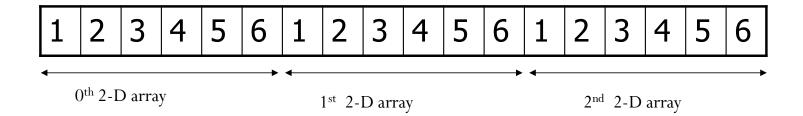
- The general form of a multi-dimensional array is: type array\_name[s1][s2][s3]....[sm];
- Suppose 3-D array can be group of an array of arrays.
- For example : int a[3][3][2];
- Here 3-D array which is collection of three 2-D arrays each contain 3 rows and 2 column.

### **Multi-Two-Dimensional Array**

• For example:

```
int a[3][3][2]={ { {1,2},{3,4},{5,6}},
 { {1,2},{3,4},{5,6}},
 { {1,2},{3,4},{5,6}},
 };
```

• Memory representation of above 3-D array is bellow:



### Different Implementations

Row Major-

$$Loc(A[k, i, j]) = Lo+[n*m(k-K)+m(i-L1)+(j-L2)]*W$$

• Column Major-

$$Loc(A[k, i, j]) = Lo+[n*m(k-K)+n(j-L2)+(i-L1)]*W$$

N=no of rows, m=no of column, K= Lowerbound, L1= lower bound of row, L2= lower bound of column

## Example

Let A be a 3-dimensional array declared as follows:

A: array [1 .... 8] [1 ..... 5] [1......7] of integer;

Given  $L_0 = 900$ 

Find location of A[5,3,6] element using row major & column major?

## Applications of Array

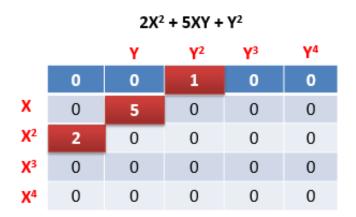
- 1. Symbol Manipulation (matrix representation of polynomial equation)
- 2. Sparse Matrix

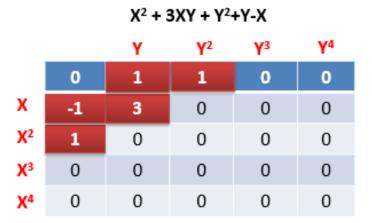
## Matrix representation of Polynomial Equation

- We can use array for different kind of operations in polynomial equation such as addition, subtraction etc.
- Array can be used to represent polynomial equation

## Example

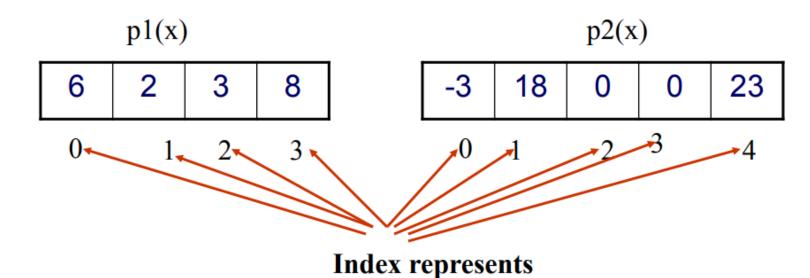
	Y	Y <sup>2</sup>	γз	Y <sup>4</sup>
Χ	XY	XY <sup>2</sup>	XY3	XY <sup>4</sup>
$X^2$	X₃A	$X^2Y^2$	$\chi^2 \gamma^3$	$X^2Y^4$
X3	X₃A	X <sup>3</sup> Y <sup>2</sup>	X <sub>3</sub> Y <sub>3</sub>	X <sup>3</sup> Y <sup>4</sup>
$X^4$	X <sup>4</sup> Y	$X^4Y^2$	$X^4Y^3$	$X^4Y^4$





### Polynomial ADT

- •Array Implementation:
- $p1(x) = 8x^3 + 3x^2 + 2x + 6$
- $p2(x) = 23x^4 + 18x 3$



exponents

Why arrays aren't good to represent polynomials?

• 
$$p3(x) = 16x^{21} - 3x^5 + 2x + 6$$

6	2	0	0	-3	0	•••••	0	16
---	---	---	---	----	---	-------	---	----

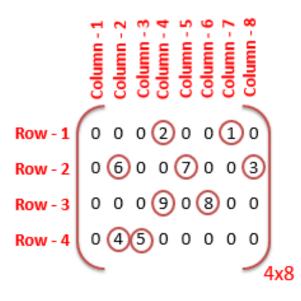
**WASTE OF SPACE!** 

### Exercise-

- 1. Write an algorithm to add two polynomials using linked list.
- 2. Write an algorithm to add two polynomials using Array.

### Sparse Matrix

- An mXn matrix is said to be *sparse* if "many" of its elements are zero.
- A matrix is not sparse is called a *dense matrix*.
- We can device a simple representation scheme whose space requirement equals the size of the non-zero elements.



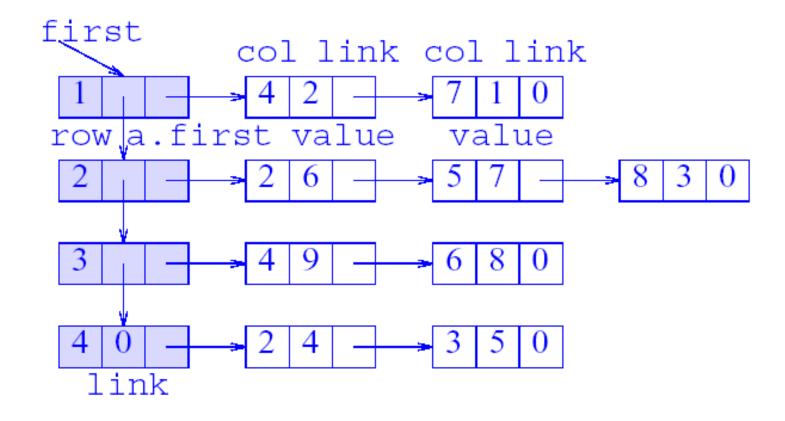
Terms	0	1	2	3	4	5	6	7	8
Row	1	1	2	2	2	3	3	4	4
Column	4	7	2	5	8	4	6	2	3
Value	2	1	6	7	3	9	8	4	5

Linear Representation of given matrix

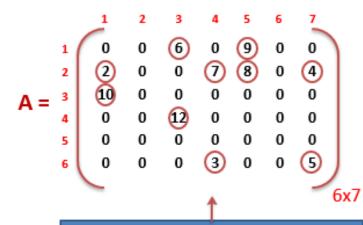
Any shortcoming of the 1-D array of a sparse matrix representation?

We need to know the number of nonzero terms in each of the sparse matrices when the array is created.

## Linked Representation of Sparse Matrix



### **Sparse Matrix**



Memory Space required to store 6x7 matrix

42 x 2 = 84 bytes

Memory Space required to store Linear Representation

 $30 \times 2 = 60 \text{ bytes}$ 

#### Linear representation of Matrix

Row	Column	A
1	3	6
1	5	9
2	1	2
2	4	7
2	5	8
2	7	4
3	1	10
4	3	12
6	4	3
6	7	5

Space Saved = 84 – 60 = 24 bytes

# Sparse Matrix

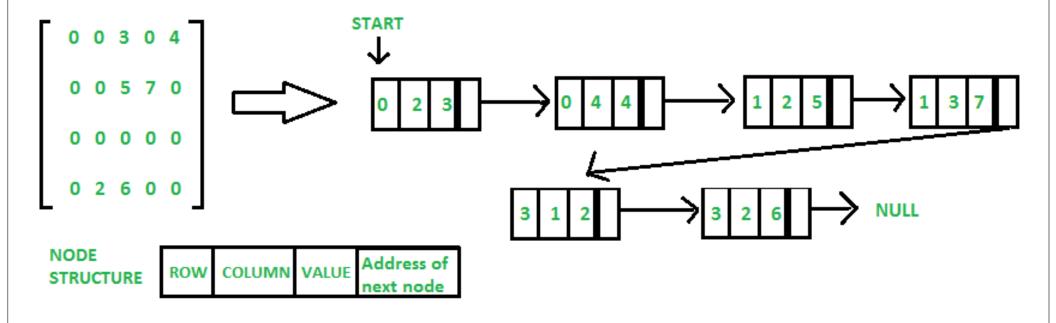
#### Linear representation of Matrix

Row	Column	Α
1	3	6
1	5	9
2	1	2
2	4	7
2	5	8
2	7	4
3	1	10
4	3	12
6	4	3
6	7	5

	Column	Α
<b>1</b>	3	6
2	5	9
3	1	2
4	4	7
5	5	8
6	7	4
7	1	10
8	3	12
9	4	3
10	7	5

Memory Space required to store Liner Representation = 26 x 2 = 42 bytes

## Linked List representation of sparse matrix



### Case study-

- A super market conducting a study of the mix items purchased by its customers.
- For this study data are gathered for the purchase made by 1000 customers.
- These data are organized into a matrix, purchases with purchases(i,j) being the quantity of item i purchased by customer j.
- Suppose that the super market has an inventory of 10,000 different items.
- The purchase matrix is therefore a 10,000 x 1,000 matrix
- If the average customer buys 20 different items only about 20,000 of 1,00,000,000 matrix entries are nonzero

