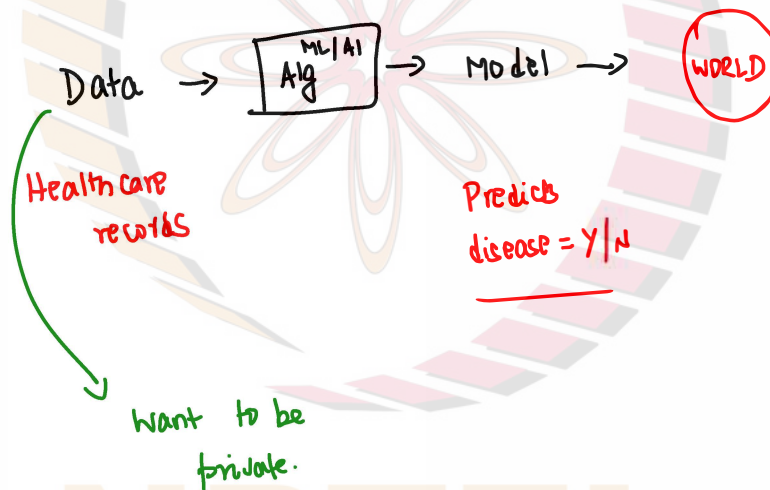


PRIVACY IN AI/ML

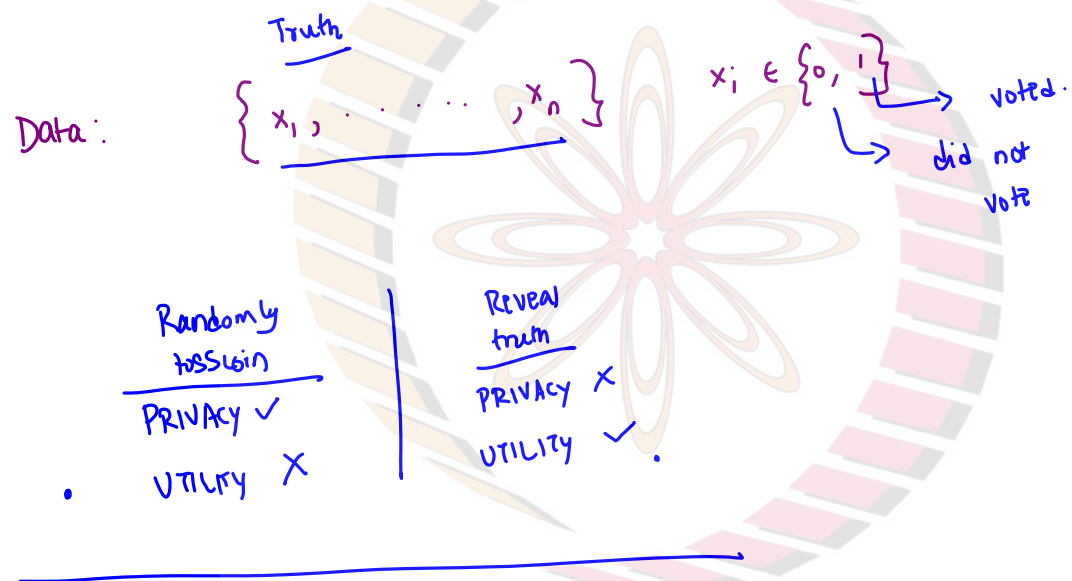


Anonymousisation

Netflix



- CRYPTOGRAPHIC SOLUTION X
- STATISTICAL SOLUTION TO PRIVACY ✓



RANDOMIZED RESPONSE

Truth $\{x_1, \dots, x_n\}$ $x_i \in \{0, 1\}$

Revealed

$$y_i = \begin{cases} x_i \\ 1 - x_i \end{cases}$$

with probability

$$\frac{e^\epsilon}{1 + e^\epsilon}$$

TRUTH

with probability

$$\frac{1}{1 + e^\epsilon}$$

FALSE HOOD

Consider two datasets

$$X = \{x_1, \dots, x_n\}$$

$$X' = \{x'_1, \dots, x'_n\}$$

Assume

$$x_i = x'_i \quad \forall i \in \{1, \dots, n-1\}$$

NPTEL

Revealed

$$RR(x) = y = [y_1, \dots, y_n]$$

$$RR(x') = y' = [y'_1, \dots, y'_n]$$

$$y_i \in \{0, 1\}$$

$$y'_i \in \{0, 1\}$$

$$\text{Prob} \left(RR(x) = b \right) = \text{Prob} \left([y_1, y_2, \dots, y_n] = [b_1, b_2, \dots, b_n] \right)$$

$$[b_1, b_2, \dots, b_n]$$

$$\text{eg: } [0, 1, 0, 1, 1]$$

$$= \prod_{i=1}^n P(y_i = b_i)$$

[Independence of
 y_1, \dots, y_n]

$$= \left[\prod_{i=1}^{n-1} P(y_i = b_i) \right] \underline{P(y_n = b_n)} \quad - \textcircled{1}$$

$$\begin{aligned}
 \text{Prob} \left(\text{RR}(x) = b \right) &= \prod_{i=1}^n P(y_i' = b_i) \\
 &= \left(\prod_{i=1}^{n-1} P(y_i' = b_i) \right) P(y_n' = b_n) \\
 &= \left(\prod_{i=1}^n P(y_i = b_i) \right) \underline{P(y_n' = b_n)} \quad \text{--- (2)}
 \end{aligned}$$

$$\frac{\checkmark P(y_n = b_n)}{\checkmark P(y_n' = b_n)} = \begin{cases} \frac{e^{\epsilon} / (1 + e^{\epsilon})}{1 / (1 + e^{\epsilon})} = e^{\epsilon} & \underline{b_n = x_n} \quad \checkmark \\ \frac{1 / (1 + e^{\epsilon})}{1 / (1 + e^{\epsilon}) / e^{\epsilon} (1 + e^{\epsilon})} = \underline{e^{-\epsilon}} & b_n \neq x_n \quad (\text{i.e., } b_n = 1 - x_n) \quad \checkmark \end{cases}$$

In both cases,

$$\frac{P(y_n = b_n)}{P(y_n' = b_n)} \leq e^\epsilon$$

\Rightarrow

$$P(y_n = b_n) \leq e^\epsilon \cdot P(y_n' = b_n)$$

— (3)

$$\frac{P_Y(RR(x) = b)}{P_Y(RR(x') = b)} = \frac{P(y_n = b_n)}{P(y_n' = b_n)} \leq e^\epsilon$$

\downarrow (1) & (2) \downarrow (3)

$$\Pr(RR(x)=b) \leq e^{\epsilon} \Pr(RR(x')=b)$$

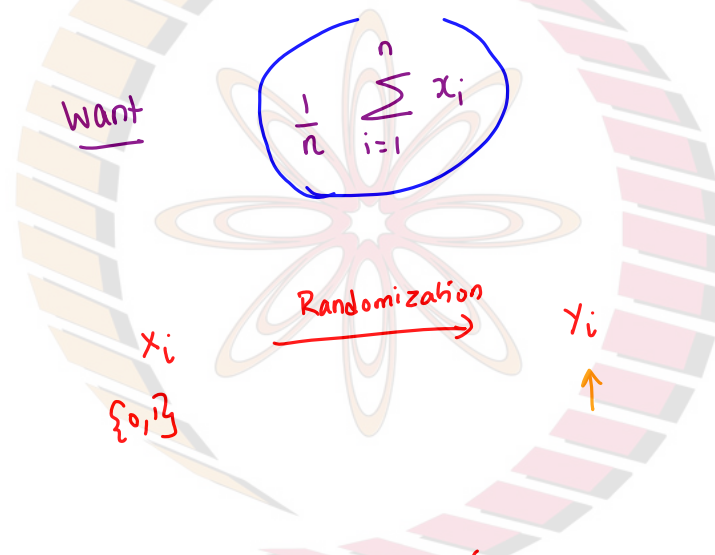
observe

\Rightarrow

$$\Pr(RR(x')=b) \cdot e^{-\epsilon} \leq \Pr(RR(x)=b) \leq e^{\epsilon} \Pr(RR(x')=b)$$

UTILITY of the mechanism

NPTEL



$$\underline{E[y_i]} = \left(\frac{e^\epsilon}{1+e^\epsilon} \right) x_i + \left(\frac{1}{1+e^\epsilon} \right) (1-x_i)$$

$$= \frac{e^\epsilon x_i + 1 - x_i}{1+e^\epsilon}$$

$$= x_i \left(\frac{e^\epsilon - 1}{e^\epsilon + 1} \right) + \frac{1}{1+e^\epsilon}$$

$$x_i \rightarrow y_i \xrightarrow{\text{unbiasing}} z_i$$

$$z_i = \left(y_i - \frac{1}{1+p^k} \right) \cdot \left(\frac{p^k + 1}{p^k - 1} \right)$$

$$E[z_i] = x_i \quad [\text{Exercise}]$$

$$\begin{array}{ccc} x_1, & \dots & x_n \\ \downarrow & & \downarrow \\ y_1 & \dots & y_n \\ \downarrow & & \downarrow \\ z_1 & \dots & z_n \end{array}$$

$$\bar{x} = \frac{1}{n} \sum_{i=1}^n x_i$$

Guess for $\bar{x} = \bar{z} = \frac{1}{n} \sum_{i=1}^n z_i$

$$E[\bar{z}] = ? = \bar{x}$$

Utility :

$$\left| \underbrace{\frac{1}{n} \sum_{i=1}^n x_i} - \underbrace{\frac{1}{n} \sum_{i=1}^n z_i} \right|$$

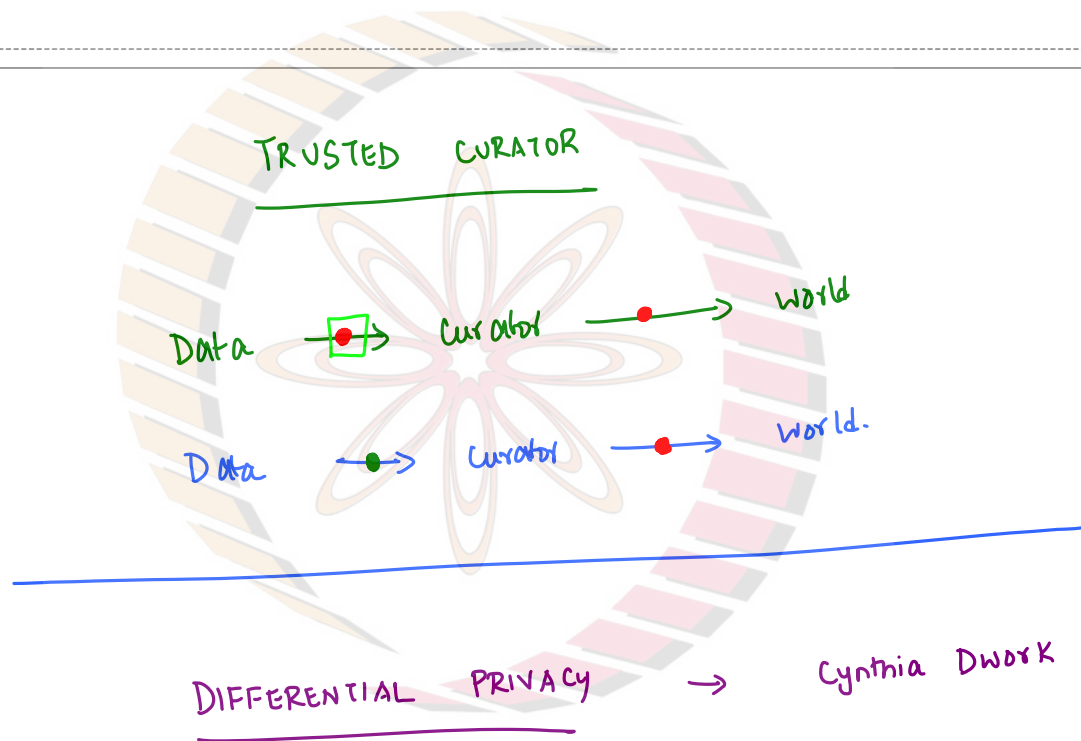
with high probability

$$\leq O\left(\frac{1}{\epsilon \sqrt{n}}\right)$$

← UTILITY

ϵ Controls Privacy-utility tradeoff

NPTEL



NPTEL

$$\{x_1, \dots, x_n\}$$

$$x_i \in \{0, 1\}$$

$$RR(x) = \{y_1, \dots, y_n\}$$

$$x, x'$$

Fix any $b \in \{0, 1\}^n$

$$\forall b \in \{0, 1\}^n$$

$$Pr(RR(x) = b)$$

$$Pr(RR(x') = b)$$

$$\leq e^{\epsilon}$$

Privacy

$$y_i = x_i \quad \text{w.p.} \quad \frac{e^\epsilon}{1+e^\epsilon}$$

$$= 1-x_i \quad \frac{1}{1+e^\epsilon}$$

Utility

w.h.p

$$\begin{array}{ccc} x_1, \dots, x_n & & \\ \downarrow & & \downarrow \\ y_1, \dots, y_n & & \\ \downarrow & & \downarrow \\ z_1, \dots, z_n & & \end{array}$$

$$\rightarrow \frac{1}{n} \sum_{i=1}^n z_i$$

$$\left| \frac{1}{n} \sum x_i - \frac{1}{n} \sum z_i \right| \leq O\left(\frac{1}{\epsilon \sqrt{n}}\right)$$

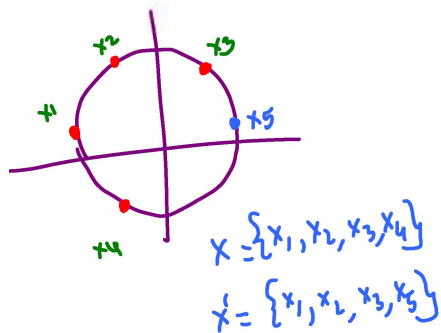
Trusted Curator Model

Differential Privacy [Cynthia Dwork]

Let $M: \mathcal{X}^n \rightarrow \mathcal{Y}$
 $\{0,1\}^n \rightarrow \{0,1\}^n$ Consider two "neighboring" datasets x and $x' \in \mathcal{X}^n$

M is ϵ -differentially private. if for all x, x' neighbouring.

and all $S \subseteq \mathcal{Y}$



NPTTEL

$$\frac{\Pr(M(x) \in \underline{S})}{\Pr(M(x') \in S)} \leq \frac{\epsilon}{\epsilon}$$

In the previous example

$$\mathcal{Y} = \{0, 1\}^n$$

$$S \subseteq \{0, 1\}^n$$

$$S = \{b_1, b_{25}, b_{100}\}$$

$$\Pr(M(x) \in S)$$

$$= \Pr(M(x) = b_1) + \Pr(M(x) = b_{25}) + \Pr(M(x) = b_{100})$$

$$\Pr(M(x') \in S)$$

$$\Pr(M(x') = b_1) + \dots$$

LAPLACE MECHANISM

Sensitivity

$$f: \mathcal{X}^n \rightarrow \mathbb{R} \quad (\text{average})$$

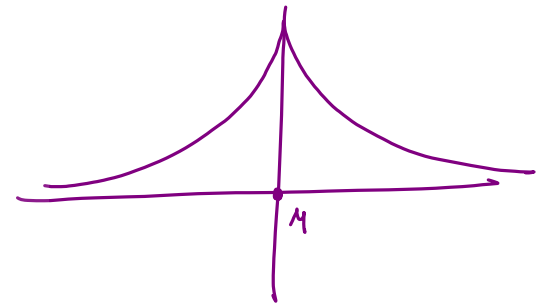
$$\Delta = \max_{\substack{x, x' \\ \text{neighbouring}}} |f(x) - f(x')|$$

NPTEL

$$\begin{aligned} & \left| \frac{1}{n} \sum_{i=1}^n x_i - \frac{1}{n} \sum_{i=1}^n x'_i \right| \\ &= \left| \frac{\cancel{\sum_{i=1}^{n-1} x_i} + x_n}{n} - \frac{\cancel{\sum_{i=1}^{n-1} x'_i} + x'_n}{n} \right| \\ &= \frac{1}{n} |x_n - x'_n| \\ \Delta &= \frac{1}{n} \end{aligned}$$

LAPLACE DISTRIBUTION

$$f_{\text{Lap}}(x) = \frac{1}{2b} e^{-\frac{|x-\mu|}{b}}$$



Laplace mechanism :

$$f(x) + \eta$$

noise

$$\text{Laplace}\left(0, \left\{\frac{\Delta}{\epsilon}\right\}\right)$$

Sensitivity of f

Privacy parameter

average + η

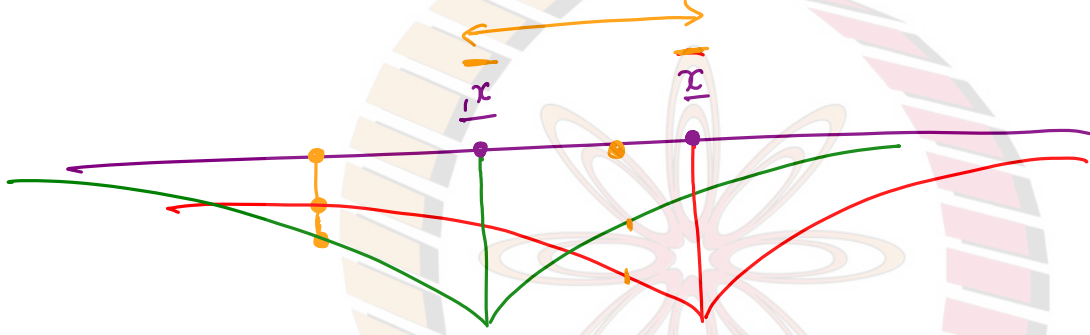
$$\eta \sim \text{Lap}\left(0, \frac{1}{n\epsilon}\right)$$

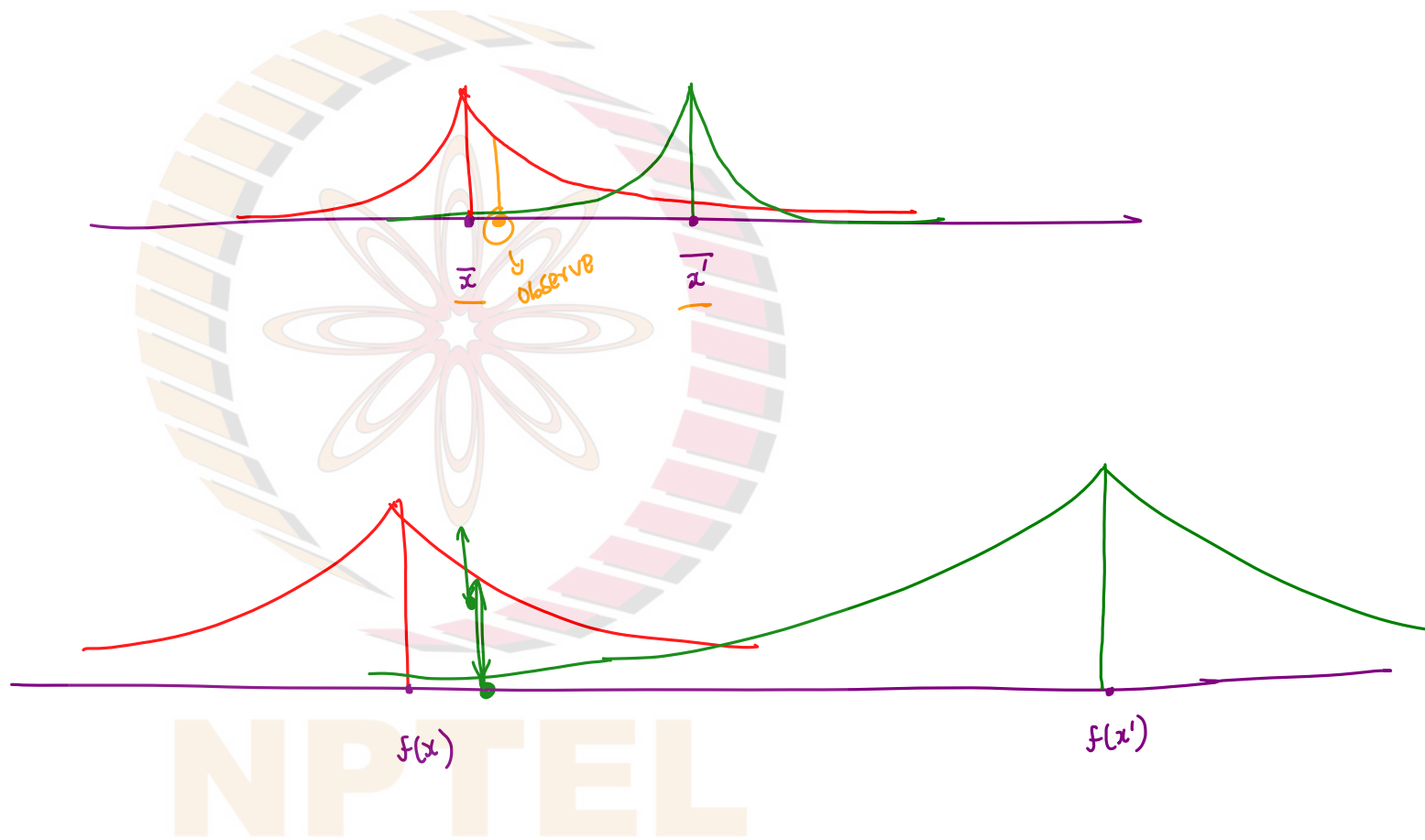
NPTEL

• Can argue that the Laplace mechanism is ε-DP

$$x_i = \frac{1}{n} \sum_{i=1}^n x_i$$

$$x_i = \frac{1}{n} \sum_{i=1}^n x_i$$





Utility

w.h.p

deviation from truth.

$$\left| \underbrace{\frac{1}{n} \sum_{i=1}^n x_i}_{\text{Truth}} - \underbrace{\left(\frac{1}{n} \sum_{i=1}^n x_i + \eta \right)}_{\text{actual released value } F(x) + \eta} \right|$$

Truth

actual
released
value
 $F(x) + \eta$

NPTEL

LAPLACE	R.R
$O\left(\frac{1}{\epsilon \sqrt{n}}\right)$	$O\left(\frac{1}{\epsilon \sqrt{n}}\right)$