## Homework 2

```
Homework 2
Question 1
Result.java
Rename.java (portion)
Splitter.java (portion)
Visual

Question 2
Question 3
Graph
```

## Question 1

#### Result.java

```
public class Result {
   public int m_down;
   public int m_right;
   public int m_id;

public Result(int down, int right, int id) {
      m_down = down;
      m_right = right;
      m_id = id;
   }
}
```

Rename.java (portion)

```
public Result rename() {
  boolean search = true;
  int down = 0, right = 0, id = -1;
  while (search) {
    Splitter splitter = null;
    try {
      splitter = m_splitters[down][right];
    } catch (Exception e) {
      return null;
    }
    if (splitter == null) {
      return null;
    }
    Direction direction = splitter.getDirection(Thread.currentThread()
         .getId());
    switch (direction) {
    case DOWN: {
      down++;
      break;
    }
    case RIGHT: {
      right++;
      break;
    }
    case STOP: {
      id = getId(down, right, m_range); // Use the function below to
calculate id from the returned Result
      search = false;
      break;
    }
    }
  return new Result(down, right, id);
```

```
public static int getId(int down, int right, int range) {
   int y = down * (down + 1) / 2 + 1;
   int x = 0;

// Arithmetic series
   if (right > 0) {
      int dx = down + 1;
      int ex = dx + right;
      int sx = dx * (dx + 1) / 2;
      x = ex * (ex + 1) / 2 - sx;
   }

// Based on grid position not the previous id
   return x + y;
}
```

### Splitter.java (portion)

```
public Direction getDirection(long pid) {
    m_pid.set(pid);

    if (m_stopped.get()) {
        return Direction.RIGHT;
    } else {
        m_stopped.set(true);
        if (m_pid.get() == pid) {
            return Direction.STOP;
        } else {
            return Direction.DOWN;
        }
    }
}

public void release() {
    m_stopped.set(false);
}
```

#### Visual



# Question 2

If you have an atomic read and non-atomic write, then the values returned during any duration of the write can yield an unexpected value. This behavior can cause the Bakery algorithm to fail. Bakery will fail in this case because it won't follow fairness anymore.

Thread A	Thread B	Labels
Enter CS	Waiting in while- loop	label[A] = 1; label[B] = 2;
Exits CS	OS halts	label[A] = 0; label[B] = 2;
Lock called		label[A] = 0; label[B] = 2;
Write label[A] = 3	Reads wrong label.	<pre>label[A] = -1; (Thread B reads Thread A wrong)</pre>

Enters CS

Thread A is able enter twice, while thread B has to wait; hence, fairness was not upheld and is not garunteed.

## Question 3

### Graph

