

## COURSE 11

### 5. Numerical methods for solving nonlinear equations in $\mathbb{R}$ (continuation)

Let  $f : \Omega \rightarrow \mathbb{R}$ ,  $\Omega \subset \mathbb{R}$ . Consider the equation

$$f(x) = 0, \quad x \in \Omega. \quad (1)$$

For  $m = 2$ .

$$F_2^T(x_i) = x_i - \frac{f(x_i)}{f'(x_i)}.$$

This method is called **Newton's method (the tangent method)**. Its order is 2.

When there exists a neighborhood of  $\alpha$  where the  $F$ -method is convergent, choosing  $x_0$  in such a neighborhood allows approximating  $\alpha$  by terms of the sequence

$$x_{i+1} = F_2^T(x_i) = x_i - \frac{f(x_i)}{f'(x_i)}, \quad i = 0, 1, \dots,$$

with a prescribed error  $\varepsilon$ .

**Lemma 1** *Let  $\alpha \in (a, b)$  be a solution of equation (1) and let  $x_n = F_2^T(x_{n-1})$ . Then*

$$|\alpha - x_n| \leq \frac{1}{m_1} |f(x_n)|, \quad \text{with } m_1 \leq m_1 f = \min_{a \leq x \leq b} |f'(x)|.$$

**Proof.** We use the mean formula

$$f(\alpha) - f(x_n) = f'(\xi)(\alpha - x_n),$$

with  $\xi \in$  to the interval determined by  $\alpha$  and  $x_n$ . From  $f(\alpha) = 0$  and  $|f'(x)| \geq m_1$  for  $x \in (a, b)$ , it follows  $|f(x_n)| \geq m_1 |\alpha - x_n|$ , that is

$$|\alpha - x_n| \leq \frac{1}{m_1} |f(x_n)|.$$

■

In practical applications the following evaluation is more useful:

**Lemma 2** *If  $f \in C^2[a, b]$  and  $F_2^T$  is convergent, then there exists  $n_0 \in \mathbb{N}$  such that*

$$|x_n - \alpha| \leq |x_n - x_{n-1}|, \quad n > n_0.$$

**Proof.** We start with Taylor formula

$$f(x_n) = f(x_{n-1}) + (x_n - x_{n-1}) f'(x_{n-1}) + \frac{1}{2} (x_n - x_{n-1})^2 f''(\xi),$$

where  $\xi$  belongs to the interval determined by  $x_{n-1}$  and  $x_n$ .

Since  $x_n = F_2^T(x_{n-1})$ , it follows that

$$x_n = x_{n-1} - \frac{f(x_{n-1})}{f'(x_{n-1})} \iff f(x_{n-1}) + (x_n - x_{n-1}) f'(x_{n-1}) = 0,$$

thus we obtain

$$f(x_n) = \frac{1}{2} (x_n - x_{n-1})^2 f''(\xi).$$

Consequently,

$$|f(x_n)| \leq \frac{1}{2} (x_n - x_{n-1})^2 M_2 f,$$

and Lemma 1 yields  $|\alpha - x_n| \leq \frac{1}{m_1} |f(x_n)|$  so

$$|\alpha - x_n| \leq \frac{1}{2m_1} (x_n - x_{n-1})^2 M_2 f.$$

Since  $F_2^T$  is convergent, there exists  $n_0 \in \mathbb{N}$  such that

$$\frac{1}{2m_1} |x_n - x_{n-1}| M_2 f < 1, \quad n > n_0.$$

Hence,

$$|\alpha - x_n| \leq |x_n - x_{n-1}|, \quad n > n_0.$$



**Remark 3** *The starting value is chosen randomly. If, after a fixed number of iterations the required precision is not achieved, i.e., condition  $|x_n - x_{n-1}| \leq \varepsilon$ , does not hold for a prescribed positive  $\varepsilon$ , the computation has to be started over with a new starting value.*

A modified form of Newton's method: - the same value during the computation of  $f'$ :

$$x_{k+1} = x_k - \frac{f(x_k)}{f'(x_0)}, \quad k = 0, 1, \dots$$

It is very useful because it doesn't request the computation of  $f'$  at  $x_j$ ,  $j = 1, 2, \dots$  but the order is no longer equal to 2.

### **Another way for obtaining Newton's method.**

We start with  $x_0$  as an initial guess, sufficiently close to the  $\alpha$ . Next approximation  $x_1$  is the point at which the tangent line to  $f$  at  $(x_0, f(x_0))$  crosses the  $Ox$ -axis. The value  $x_1$  is much closer to the root  $\alpha$  than  $x_0$ .

We write the equation of the tangent line at  $(x_0, f(x_0))$  :

$$y - f(x_0) = f'(x_0)(x - x_0).$$

If  $x = x_1$  is the point where this line intersects the  $Ox$ -axis, then  $y = 0$

$$-f(x_0) = f'(x_0)(x_1 - x_0),$$

and solving for  $x_1$  gives

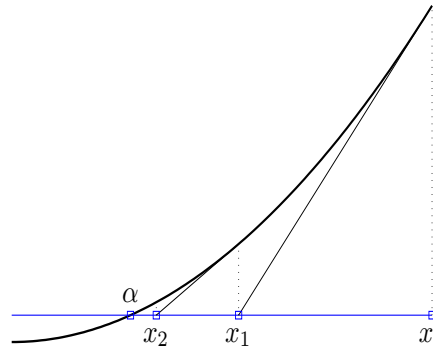
$$x_1 = x_0 - \frac{f(x_0)}{f'(x_0)}.$$

By repeating the process using the tangent line at  $(x_1, f(x_1))$ , we obtain for  $x_2$

$$x_2 = x_1 - \frac{f(x_1)}{f'(x_1)}$$

For the general case we have

$$x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}, \quad n \geq 0. \quad (2)$$



**The algorithm:**

Let  $x_0$  be the initial approximation.

**for**  $n = 0, 1, \dots, ITMAX$

$$x_{n+1} \leftarrow x_n - \frac{f(x_n)}{f'(x_n)}.$$

A stopping criterion is:

$$|f(x_n)| \leq \varepsilon \text{ or } |x_{n+1} - x_n| \leq \varepsilon \text{ or } \frac{|x_{n+1} - x_n|}{|x_{n+1}|} \leq \varepsilon,$$

where  $\varepsilon$  is a specified tolerance value.

**Example 4** Use Newton's method to compute a root of  $x^3 - x^2 - 1 = 0$ , to an accuracy of  $10^{-4}$ . Use  $x_0 = 1$ .

**Sol.** The derivative of  $f$  is  $f'(x) = 3x^2 - 2x$ . Using  $x_0 = 1$  gives  $f(1) = -1$  and  $f'(1) = 1$  and so the first Newton's iterate is

$$x_1 = 1 - \frac{-1}{1} = 2 \text{ and } f(2) = 3, f'(2) = 8.$$

The next iterate is

$$x_2 = 2 - \frac{3}{8} = 1.625.$$

Continuing in this manner we obtain the sequence of approximations which converges to 1.465571.



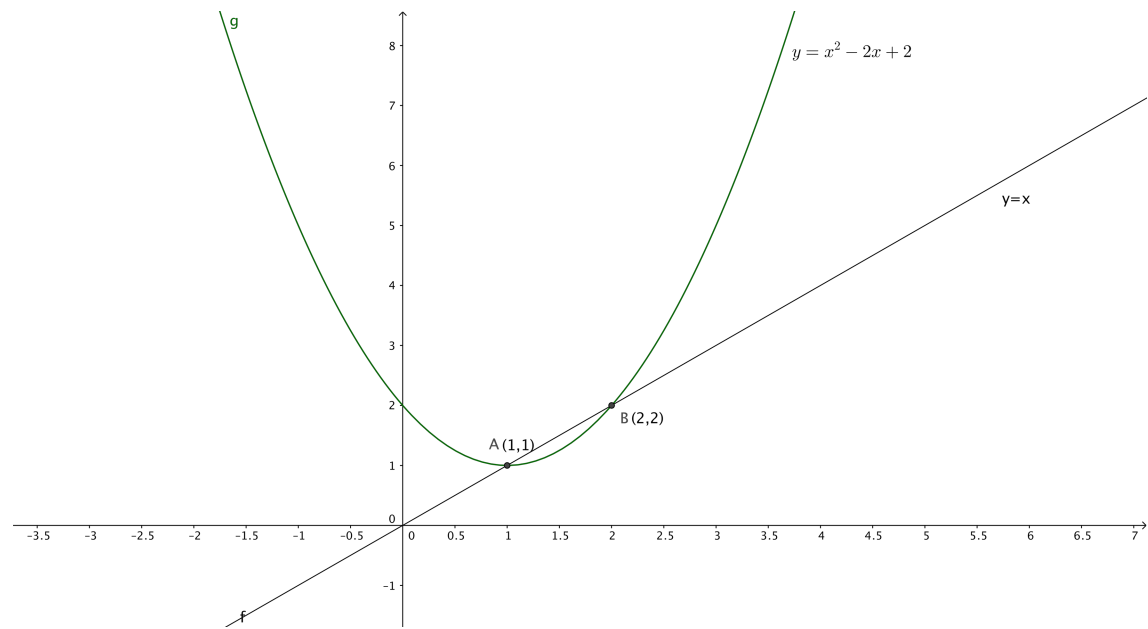
### 5.1.2. Fixed point iteration method (successive approximation method)

**Definition 5** *The number  $\alpha$  is called a **fixed point** of the function  $g : D \subseteq \mathbb{R} \rightarrow D$  if  $g(\alpha) = \alpha$ .*

**Example 6** *Find the fixed points of the function  $g(x) = x^2 - 2x + 2$ .*

*Sol.* A fixed point  $\alpha$  of  $g$  has the property  $\alpha = g(\alpha) = \alpha^2 - 2\alpha + 2$ , so  $0 = \alpha^2 - 3\alpha + 2 = (\alpha - 1)(\alpha - 2)$ . Whence, the fixed points of  $g$  are  $\alpha_1 = 1$  and  $\alpha_2 = 2$ .

Geometrically, the fixed points are the intersection points of the graph of the function  $g$  and the first bisection line ( $y = x$ ). (See the following figure.)



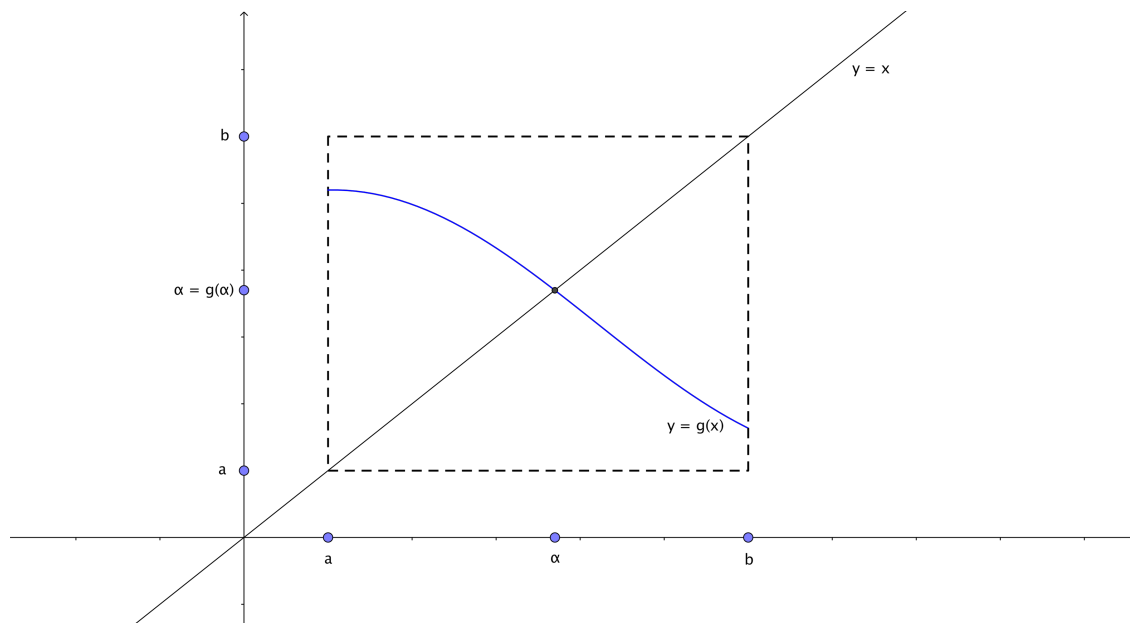
Sufficient condition for the existence and uniqueness of a fixed point:

**Theorem 7** 1. If  $g \in C[a, b]$  and  $g(x) \in [a, b]$  for any  $x \in [a, b]$ , then  $g$  has at least one fixed point in  $[a, b]$ . In fewer words, if  $g : [a, b] \rightarrow [a, b]$  and  $g \in C[a, b]$  then  $\exists \alpha \in [a, b]$  fixed point.

2. Moreover, if there exists  $g'(x)$  in  $(a, b)$  and

$$|g'(x)| < 1, \quad \forall x \in (a, b),$$

then the fixed point is unique in  $[a, b]$ .



**Example 8** Prove that  $g(x) = (x^2 - 4)/5$  has a unique fixed point in  $[-2, 2]$ .

*Sol.* The minimum and maximum of  $g(x)$  for  $x \in [-2, 2]$  are the limits of the interval, or at the points where  $g'(x) = 0$ . We have  $g'(x) = 2x/5$ ,  $g$  is continuous and there exists  $g'(x)$  in  $[-2, 2]$ . So, the minimum and maximum of  $g(x)$  on  $[-2, 2]$  are at  $x = -2$ ,  $x = 0$  or  $x = 2$ . We have  $g(-2) = 0$ ,  $g(2) = 0$ ,  $g(0) = -4/5$ , so  $x = -2$  and  $x = 2$  are points of absolute maximum and  $x = 0$  is a point of absolute minimum in  $[-2, 2]$ . Moreover,

$$|g'(x)| = \left| \frac{2x}{5} \right| \leq \left| \frac{4}{5} \right| < 1, \quad \forall x \in (-2, 2).$$

So,  $g$  satisfies the conditions of Theorem 7, so it follows that  $g$  has a unique fixed point in  $[-2, 2]$ .

Consider the equation

$$f(x) = 0, \tag{3}$$

where  $f : [a, b] \rightarrow \mathbb{R}$ . Assume that  $\alpha \in [a, b]$  is a zero of  $f(x)$ .

In order to compute  $\alpha$ , we transform (3) algebraically into *fixed point form*,

$$x = g(x), \tag{4}$$

where  $g$  is chosen so that  $g(x) = x \Leftrightarrow f(x) = 0$ .

A simple way to do this is, for example,  $x = x + f(x) =: g(x)$ .

Finding a zero of  $f(x)$  in  $[a, b]$  is then equivalent to finding a fixed point  $x = g(x)$  in  $[a, b]$ .

The fixed point form suggests *the fixed point iteration*

$$x_0 - \text{initial guess}, x_{k+1} = g(x_k), \quad k = 0, 1, 2, \dots$$

The hope is that iteration will produce a convergent sequence  $(x_n) \rightarrow \alpha$ .

For example, consider

$$f(x) = xe^x - 1 = 0. \quad (5)$$

A first fixed point iteration is obtained rearranging and dividing (5) by  $e^x$ :  $xe^x = 1 \Rightarrow x = e^{-x}$ , so  $x = g(x) = e^{-x}$  and

$$x_{k+1} = e^{-x_k}.$$

With the initial guess  $x_0 = 0.5$  we obtain the iterates  $x_1 = 0.6065306597$ ,  $x_2 = 0.5452392119, \dots, x_8 = 0.5664094527, x_9 = 0.5675596343, \dots, x_{28} = 0.56714328, x_{29} = 0.56714329$

So  $x_k$  seems to converge to  $\alpha = 0.5671432\dots$

A second fixed point form is obtained from  $xe^x = 1$  by adding  $x$  on both sides:  $xe^x + x = 1 + x \Rightarrow x(e^x + 1) = 1 + x \Rightarrow x = \frac{1+x}{e^x + 1}$ , we get

$$x = g(x) = \frac{1+x}{e^x + 1}.$$

This time the convergence is much faster (we need only three iterations to obtain a 10-digit approximation of  $\alpha$ ) :  $x_0 = 0.5$ ,  $x_1 = 0.5663110032$ ,  $x_2 = 0.5671431650$ ,  $x_3 = 0.5671432904$ .

Another possibility for a fixed point iteration is  $x = x + 1 - xe^x$ . But this iteration function does not generate a convergent sequence.

Finally we could also consider the fixed point form  $x = x + xe^x - 1$ . Also this iteration function does not generate a convergent sequence.

The question is: when does the iteration sequence converge?

Answer: when conditions of Theorem 7 are fulfilled.

For this example, we have two cases when  $|g'(x)| < 1$  and the algorithm converges and two cases when  $|g'(x)| > 1$  and the algorithm is not convergent.

A more general statement for the convergence is the theorem of Banach.

**Definition 9** A Banach space  $\mathcal{B}$  is a complete normed vector space over some number field  $K$  such as  $\mathbb{R}$  or  $\mathbb{C}$ . (Complete means that every Cauchy sequence converges in  $\mathcal{B}$ .)

**Definition 10** Let  $A \subset \mathcal{B}$  be a closed subset and  $G : A \rightarrow A$ .  $G$  is called **Lipschitz continuous** on  $A$  if there exists a constant  $L \geq 0$  such that  $\|G(x) - G(y)\| \leq L \|x - y\|$ ,  $\forall x, y \in A$ . Furthermore,  $G$  is called a **contraction** if  $L$  can be chosen such that  $L < 1$ .

**Theorem 11** (Banach Fixed Point Theorem) Let  $A$  be a closed subset of a Banach space  $\mathcal{B}$ , and let  $G$  be a **contraction**  $G : A \rightarrow A$ . Then:

a)  $G$  has a unique fixed point  $\alpha$ , which is the unique solution of the equation  $x = G(x)$ .



b) The sequence  $x_{n+1} = G(x_n)$  converges to  $\alpha$  for every initial guess  $x_0 \in A$ .

c) We have the estimate:  $\|\alpha - x_n\| \leq \frac{L^{n-l}}{1-L} \|x_{l+1} - x_l\|$ , for  $0 \leq l \leq n$  (or  $\|\alpha - x_n\| \leq \frac{L^n}{1-L} \|x_1 - x_0\|$ )

For practical applications is also useful the following estimation.

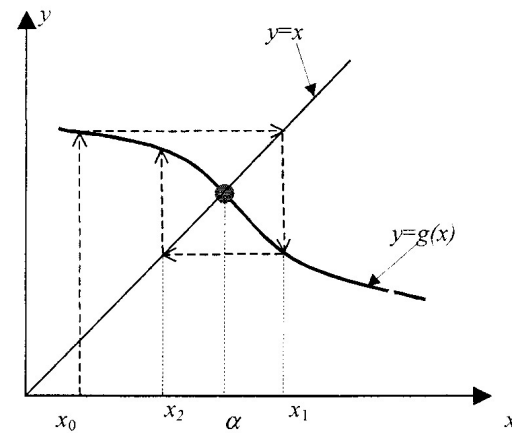
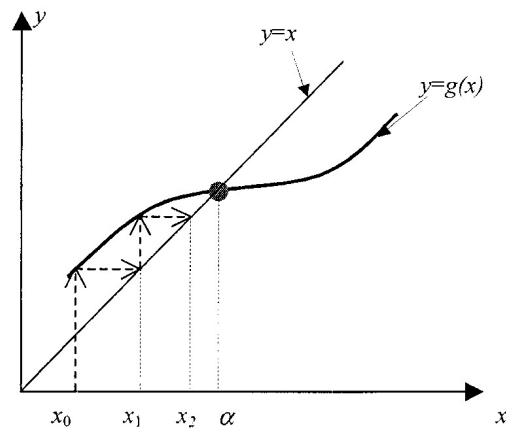
**Lemma 12** If  $\|G'(x)\| < L < 1$ ,  $x \in V(\alpha)$  then

$$\|\alpha - x_n\| \leq \frac{L}{1-L} \|x_n - x_{n-1}\|.$$

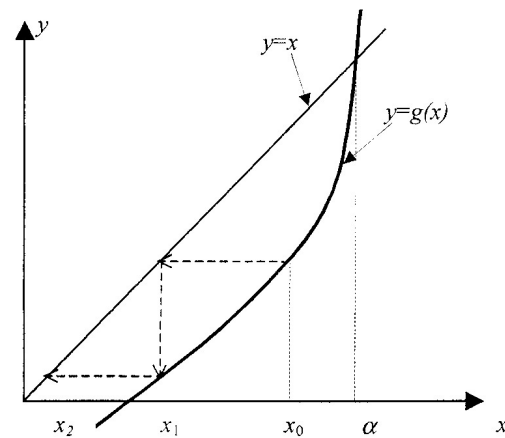
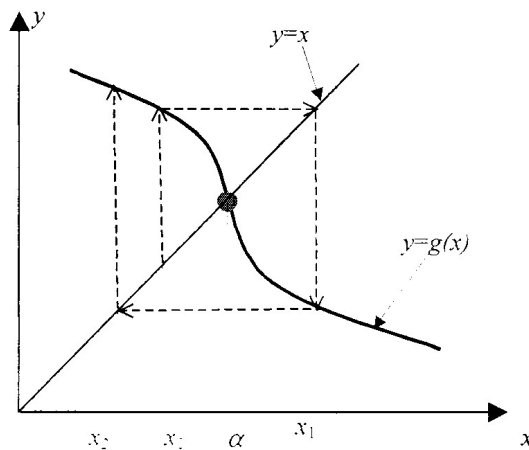
Geometric interpretation of the method: we plot  $y = G(x)$  and  $y = x$ . The intersection points of the two functions are the solutions of  $x = G(x)$ . The computation of the sequence  $\{x_k\}$  with  $x_0$  chosen initial value,  $x_{k+1} = G(x_k)$ ,  $k = 0, 1, 2, \dots$  can be interpreted geometrically via sequences of lines parallel to the coordinate axes:

$x_0$	start with $x_0$ on the $x$ -axis
$G(x_0)$	go parallel to the $y$ -axis to the graph of $G$
$x_1 = G(x_0)$	move parallel to the $x$ -axis to the graph $y = x$
$G(x_1)$	go parallel to the $y$ -axis to the graph of $G$
<i>etc.</i>	

Case of convergence  $|G'(x)| < 1$ .



Case of divergence  $|G'(x)| > 1$ .



When  $|g'(\alpha)| = 1$ , depending on the mapping  $g$ , the iterates may converge or not (e.g.,  $g(x) = x - x^3$  or  $g(x) = x + x^3$ ).

## 5.2. Multistep methods

### Lagrange inverse interpolation

Let  $y_k = f(x_k)$ ,  $k = 0, \dots, n$ , hence  $x_k = g(y_k)$ . We attach the Lagrange interpolation formula to  $y_k$  and  $g(y_k)$ ,  $k = 0, \dots, n$ :

$$g = L_n g + R_n g, \quad (6)$$

where

$$(L_n g)(y) = \sum_{k=0}^n \frac{(y-y_0)\dots(y-y_{k-1})(y-y_{k+1})\dots(y-y_n)}{(y_k-y_0)\dots(y_k-y_{k-1})(y_k-y_{k+1})\dots(y_k-y_n)} g(y_k). \quad (7)$$

Taking

$$F_n^L(x_0, \dots, x_n) = (L_n g)(0),$$

$F_n^L$  is a  $(n+1)$  – steps method defined by

$$\begin{aligned} F_n^L(x_0, \dots, x_n) &= \sum_{k=0}^n \frac{y_0 \dots y_{k-1} y_{k+1} \dots y_n}{(y_k - y_0) \dots (y_k - y_{k-1})(y_k - y_{k+1}) \dots (y_k - y_n)} (-1)^n g(y_k) \\ &= \sum_{k=0}^n \frac{y_0 \dots y_{k-1} y_{k+1} \dots y_n}{(y_k - y_0) \dots (y_k - y_{k-1})(y_k - y_{k+1}) \dots (y_k - y_n)} (-1)^n x_k. \end{aligned}$$

Concerning the convergence of this method we state:

**Theorem 13** *If  $\alpha \in (a, b)$  is solution of equation  $f(x) = 0$ ,  $f'$  is bounded on  $(a, b)$ , and the starting values satisfy  $|\alpha - x_k| < 1/c$ ,  $k = 0, \dots, n$ , with  $c = \text{constant}$ , then the sequence*

$$x_{i+1} = F_n^L(x_{n-i}, \dots, x_i), \quad i = n, n+1, \dots$$

*converges to  $\alpha$ .*

**Remark 14** *The order  $\text{ord}(F_n^L)$  is the positive solution of the equation*

$$t^{n+1} - t^n - \dots - t - 1 = 0.$$

**Particular cases.**

1) For  $n = 1$ , the nodes  $x_0, x_1$ , we get **the secant method**

$$F_1^L(x_0, x_1) = x_1 - \frac{(x_1 - x_0) f(x_1)}{f(x_1) - f(x_0)},$$

Thus,

$$x_{k+1} := F_1^L(x_{k-1}, x_k) = x_k - \frac{(x_k - x_{k-1}) f(x_k)}{f(x_k) - f(x_{k-1})}, \quad k = 1, 2, \dots$$

is the new approximation obtained using the previous approximations  $x_{k-1}, x_k$ .

The *order* of this method is the positive solution of equation:

$$t^2 - t - 1 = 0,$$

so  $\text{ord}(F_1^L) = \frac{(1+\sqrt{5})}{2} \approx 1.618$ . (the golden ratio).

A modified form of the secant method: if we keep  $x_1$  fixed and we change every time the same interpolation node, i.e.,

$$x_{k+1} = x_k - \frac{(x_k - x_1) f(x_k)}{f(x_k) - f(x_1)}, \quad k = 2, 3, \dots$$

2) For  $n = 2$ , the nodes  $x_0, x_1, x_2$  and we get

$$\begin{aligned} F_2^L(x_0, x_1, x_2) = & \frac{x_0 f(x_1) f(x_2)}{[f(x_0) - f(x_1)][f(x_0) - f(x_2)]} + \frac{x_1 f(x_0) f(x_2)}{[f(x_1) - f(x_0)][f(x_1) - f(x_2)]} \\ & + \frac{x_2 f(x_0) f(x_1)}{[f(x_2) - f(x_0)][f(x_2) - f(x_1)]}. \end{aligned}$$

The *order* of this method is the positive solution of equation:

$$t^3 - t^2 - t - 1 = 0,$$

so  $\text{ord}(F_2^L) = 1.8394$ .

*Comparing the Newton's method and secant method* with respect to the time needed for finding a root with some given precision, we have:

-Newton's method has more computation at one step: it is necessary to evaluate  $f(x)$  and  $f'(x)$ . Secant method evaluates just  $f(x)$  (supposing that  $f(x_{previous})$  is stored.)

-The number of iterations for Newton's method is smaller (its order is  $p_N = 2$ ). Secant method has order  $p_S = 1.618$  and we have that three steps of this method are equivalent with two steps of Newton's method.

- It is proved that if the time for computing  $f'(x)$  is greater than  $(0.44 \times \text{the time for computing } f(x))$ , then the secant method is faster.



**Remark 15** *The computation time is not the unique criterion in choosing the method! Newton's method is easier to apply. If  $f(x)$  is not explicitly known (for example, it is the solution of the numerical integration of a differential equation), then its derivative is computed numerically. If we consider the following expression for the numerical computation of derivative:*

$$f'(x) \approx \frac{f(x_n) - f(x_{n-1})}{x_n - x_{n-1}} \quad (8)$$

*then the Newton's method becomes the secant method.*

### **Another way of obtaining secant method.**

Based on approx. the function by a straight line connecting two points on the graph of  $f$  (not required  $f$  to have opposite signs at the initial points).

The first point,  $x_2$ , of the iteration is taken to be the point of intersection of the  $Ox$ -axis and the secant line connecting two starting points

$(x_0, f(x_0))$  and  $(x_1, f(x_1))$ . The next point,  $x_3$ , is generated by the intersection of the new secant line, joining  $(x_1, f(x_1))$  and  $(x_2, f(x_2))$  with the  $Ox$ -axis. The new point,  $x_3$ , together with  $x_2$ , is used to generate the next point,  $x_4$ , and so on.

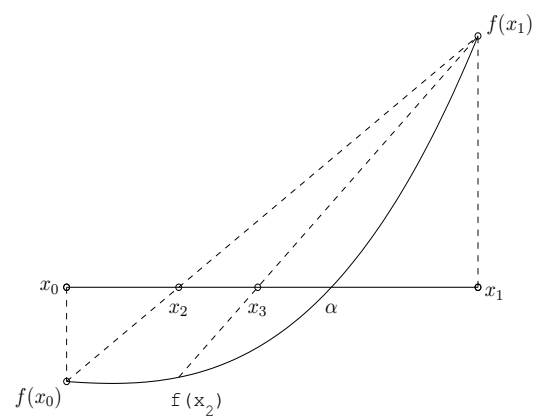
The formula for  $x_{n+1}$  is obtained by setting  $x = x_{n+1}$  and  $y = 0$  in the equation of the secant line from  $(x_{n-1}, f(x_{n-1}))$  to  $(x_n, f(x_n))$ :

$$\frac{x - x_n}{x_{n-1} - x_n} = \frac{y - f(x_n)}{f(x_{n-1}) - f(x_n)} \Leftrightarrow x = x_n + \frac{(x_{n-1} - x_n)(y - f(x_n))}{f(x_{n-1}) - f(x_n)},$$

we get

$$x_{n+1} = x_n - f(x_n) \left[ \frac{x_n - x_{n-1}}{f(x_n) - f(x_{n-1})} \right]. \quad (9)$$

Note that  $x_{n+1}$  depends on the two previous elements of the sequence  $\Rightarrow$  two initial guesses,  $x_0$  and  $x_1$ , for generating  $x_2, x_3, \dots$ .



## The algorithm:

Let  $x_0$  and  $x_1$  be two initial approximations.

**for**  $n = 1, 2, \dots, ITMAX$

$$x_{n+1} \leftarrow x_n - f(x_n) \left[ \frac{x_n - x_{n-1}}{f(x_n) - f(x_{n-1})} \right].$$

A suitable stopping criterion is

$$|f(x_n)| \leq \varepsilon \text{ or } |x_{n+1} - x_n| \leq \varepsilon \text{ or } \frac{|x_{n+1} - x_n|}{|x_{n+1}|} \leq \varepsilon,$$

where  $\varepsilon$  is a specified tolerance value.

**Example 16** Use the secant method with  $x_0 = 1$  and  $x_1 = 2$  to solve  $x^3 - x^2 - 1 = 0$ , with  $\varepsilon = 10^{-4}$ .

**Sol.** With  $x_0 = 1$ ,  $f(x_0) = -1$  and  $x_1 = 2$ ,  $f(x_1) = 3$ , we have

$$x_2 = 2 - \frac{(2 - 1)(3)}{3 - (-1)} = 1.25$$

from which  $f(x_2) = f(1.25) = -0.609375$ . The next iterate is

$$x_3 = 1.25 - \frac{(1.25 - 2)(-0.609375)}{-0.609375 - 3} = 1.3766234.$$

Continuing in this manner the iterations lead to the approximation 1.4655713.