

## Dynamic Programming problem: regex matching

This is a classic dynamic programming problem: regex matching with support for  $.$  and  $*$ .

- Pattern Matching Rules:
- ' $a$ ' matches ' $a$ ' exactly.
  - ' $.$ ' matches any single character.
  - ' $a^*$ ' matches zero or more ' $a$ 's.
  - ' $.*$ ' matches any sequence of characters, including the empty string.
  - matching must consume the entire string  $s$ .

### Strategy: Dynamic Programming (DP):

Define a DP table:  $dp[i][j] = \text{true}$  if  $s[0..i-1]$  matches  $p[0..j-1]$

Use a 2D boolean table  $dp[s.length()+1][p.length()+1]$ , where:

- $dp[0][0] = \text{true}$  (empty string matches empty pattern)
- fill the table bottom up

### Transitions: For each $dp[i][j]$ , consider:

1. When  $p[j-1] \neq '*'$ : You can match  $s[i-1]$  and  $p[j-1]$  only if:  $s[i-1] == p[j-1]$  or  $p[j-1] == '.'$  and  $dp[i-1][j-1]$  is true
2. When  $p[j-1] == '*'$ : You need to look at the preceding character  $p[j-2]$  and consider:
  - $dp[i][j-2]$ : match 0 of the preceding element
  - $dp[i-1][j]$  & match ( $s[i-1], p[j-2]$ ): use  $*$  to match more.

Key Helper: You often need a  $\text{match}(i, j)$  function:  
 $s[i-1] == p[j-1] \parallel p[j-1] == '.'$

Edge Cases:

- pattern like "a\*b\*c" matches any empty string
- don't access  $s[i-1]$  &  $p[j-1]$  when  $i=0$  &  $j=0$