LRU Cache

For this problem we need to design a date structure with O(1) get and part. The best approach combines;

"Hash Map (unordered-map in C++) - for O(1) access by key.

"Doubly Finled Fixt - to maintain the order of enage (most nearthy used at tail).

[Approach:]

1. Data structures:

· Doubly Linked List:

Each node stores (key, value).

· Most recently used node is near the head

· Zeast recently used nocle is near the tail.

· Hash Map:

· Maps key - node pointer in the linked list.

2. Sperations: get (key):

If key exests:

· Move the made to the head (most recently used).

· Retern the value.

· Elx seteery - 1

· pert (key, value):

If key exists:

· Update value and move made to head.

· Else:

· Breat new noole, envert at load.

'If size exceeds capacity, remove mode at tail (least secently eved).

Implomentation Detail:

Use a devining head and devining told for easier mode management.

· Frencteons:

add Node (mode): - insert mode right after head.
semone Node (node): - unlink node hom lest
remove Node (mode): - unlink node from lest. move To Flegol (mode): - remove and re-ensert at head.
· pgo Vail () - sem sue last node (before tail)
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vience 69mollary: O(x) 791 year and peur
fore tomplexity: D (copacity) 191 solling have map and lenked
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