

## Median of Two Sorted Arrays

Solving the "Median of Two Sorted Arrays" problem - a very famous question with an optimal solution in  $O(\log(\min(m, n)))$  using binary search.

**Goal:** Find the median of two sorted arrays  $nums1$  and  $nums2$  in logarithmic time, not by merging.

**Core Idea - Binary Search Partitioning:**

We want to partition  $nums1$  and  $nums2$  such that:

- Left halves of both arrays together contain half of the total elements.
- All elements on the left side  $\leq$  all elements on the right side.
- We binary search only on the smaller array ( $nums1$ ), trying different cut points.

**Partition Logic:** Let:  $i$  = cut index in  $nums1$

$j = (m+n+1)/2 - i$  in  $nums2$  (so left half has  $(m+n+1)/2$  elements in total)

We define:  $maxLeftX = (i == 0) ? -\infty : nums1[i-1]$

$minRightX = (j == m) ? +\infty : nums1[i]$

$maxLeftY = (j == 0) ? -\infty : nums2[j-1]$

$minRightY = (j == n) ? +\infty : nums2[j]$

**Valid Partition:** If:  $maxLeftX \leq minRightY$  &  $maxLeftY \leq minRightX$   
Then: If  $(m+n)$  is even  $\rightarrow$  median =  $\max(maxLeftX, maxLeftY) + \min(minRightX, minRightY) // 2$   
If odd  $\rightarrow$  median =  $max(maxLeftX, maxLeftY)$

**Binary Search:** Adjust the binary search:

- If  $maxLeftX > minRightY \rightarrow$  move high  $= i - 1$
- Else  $\rightarrow$  move low  $= i + 1$

Example: For  $\text{nums1} = [1, 3]$ ,  $\text{nums2} = [2]$ :

Partition such that  $\text{left} = [1]$ ,  $\text{right} = [2]$  → median is 2

Time Complexity:  $O(\log(\min(m, n)))$  — binary search on shorter array.