Minimum Size Subarray Sum

For this problem we need the shortest subarray where sum ≥ target.

Approach 1: Sliding Window (O(n)):

Key Idea:
- Use two pointers (start, end) to maintain a sliding window.
- Expand end to increase sum.
- Shrink start when sum ≥ target to minimize length.

Algorithm:

1. Initialize:
   - sum = 0
   - minLen = INT_MAX
   - start = 0
2. Loop end from 0 to n-1:
   - Add nums[end] to sum
   - While sum >= target:
     - Update minLen = min(minLen, end - start + 1).
     - Subtract nums[start] and increment start.
3. Return minLen if found, else 0.

Complexity:

- Time: O(n) (each element visited at most twice)
- Space: O(1)

Follow-up: O(n log n) Solution:

Idea:
- Use prefix sums + binary search:
  - Build prefix sum array.
  - For each i, binary search for the smallest j such that prefix[j] - prefix[i] >= target.