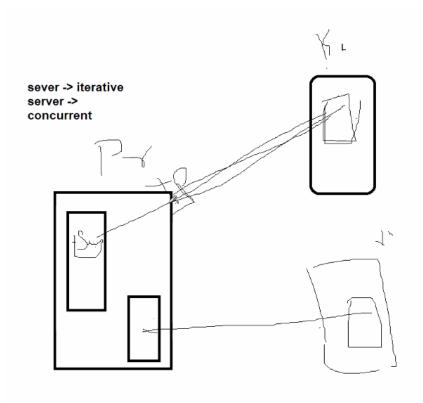
SOCKETS

5 mechanism(pipe/fifo/sm/mq/sem) => communication within same system But through socket=> Either in same or different machines Socket=>2 processes in same or diff machines can communicate Note: It is not system communication but process communication 2 types a)Connection oriented [TCP] (client server) (slow) (more reliable) b) Connectionless [UDP] (no client server) (takes address of dst and data content) and simply sends it) (fast) (less reliable) eg:live streaming has some data loss Note: Every system has 2 ip address a) local ip address b)ethernet/wlan ip address cmd: ip a Note: when both processes are in same system=>local/eth ip different system=>ethernet A server can handle only 1 client at a time=>ITERATIVE SERVER A server can handle many client at a time =>CONCURRENT SERVER 1 client comes ...creates a child using fork()

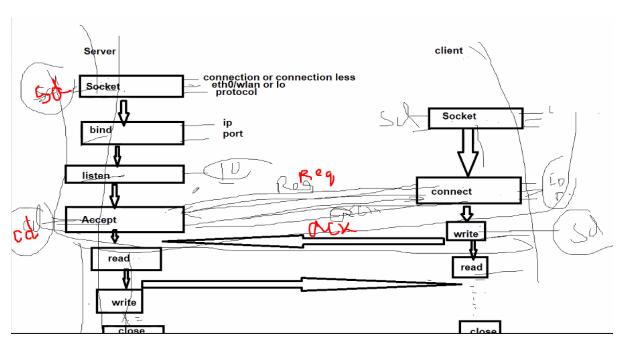


client should know the $\ensuremath{\text{ip}}$ address and port no of server server need not know

2 processes in server
They have same ip address but diff port no

CONNECTION ORIENTED COMMUNICATION

Socket steps



```
addr.sin_port=8900;
if(bind(st, (struct sockaddr *)&addr, sizeof(addr)) == -1)
{
    perror("bind");
    exit(0);
}
else
{
    isten(sd,10);
    printf("Server is started\n");
    printf("Waiting for clients\n");
    printf("Waiting for clients\n");
    read(cd,buf,1024);
    if(connect(sd, (struct sockaddr *)&addr, sizeof(addr)) == fortif("connect(sd, (struct sockaddr *)&addr, sizeof(addr)) == fortif("connect(sd, (struct sockaddr *)&addr, sizeof(addr)) == fortif("client request Accepted\n");
    read(cd,buf,1024);
    else
```

In order to communicate , we need a node=>socket Socket is a communication end point

Socket determines is it conn oriented or conn less Socket determines is it local or eth ip address Socket determines the protocol (tcp/udp)

Note: port no should be > 8080

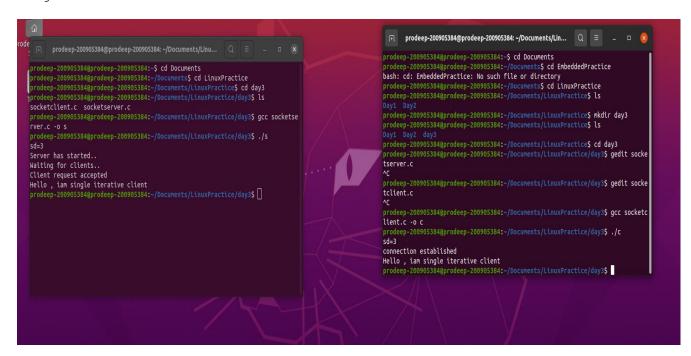
Only when server is in listen mode =>it is operational and can accept request from client

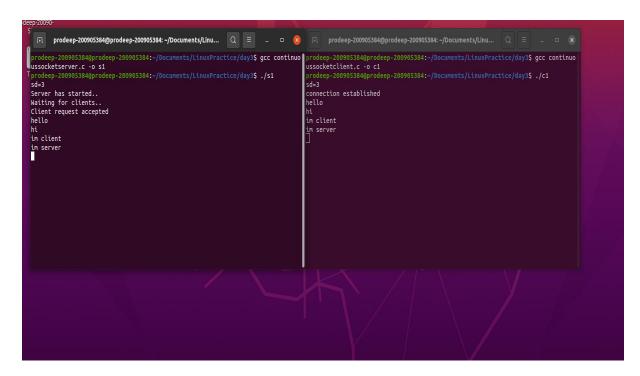
Listen =>decides max no of clients

Accept=>blocking system call & waits for client connections

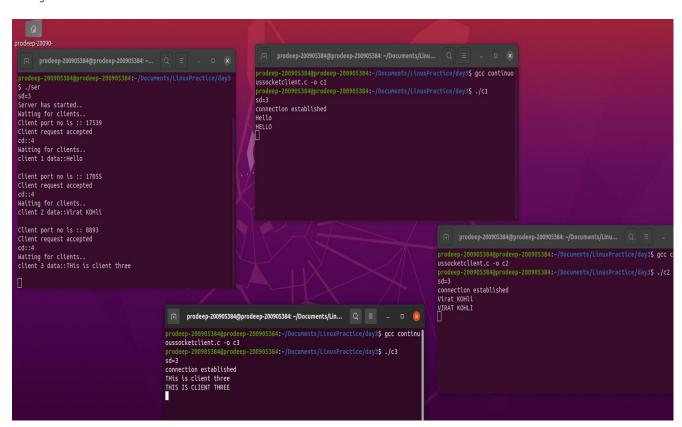
Server has diff client descriptor for every client

Program for iterative server





Program for concurrent server



//server parent process is just accepting connections
//server child is converting to uppercase

```
//cd changes every time new client added ...in accept()
//same client prg c1,c2,c3
```

MULTITHREADING

Processes: Heavy and bulky

System is loaded with a lot of resources
Lot of overhead needed =>pcb ,tables

We can achieve multitasking by thread which is better than process

Threads:Lightweight processes
Many Threads exist inside 1 process
Atleast 1 thread is running =>Main thread
Even if we dont create a thread , 1 thraed already exists =>
Main Thread

multithreading

thread LWP
Light Weight Process

Os doesnt know about threads OS only knows about process

eg:P1 has 3 threads and 1 main thread OS has alloted 10ms to P1

th1/th2/th3/main th=>each get 2.5 ms

MULTICORES

eg:quadcore=>4 cores

In the above example, in p1
1st core=>th1=>10ms and so on

Note: Multicore utilized better in multithreading

```
fns execute serially
f1();
f2();
f3();
//if f1 goes in infinite loop ..f2 & f3 wont execute
threads execute concurrently
_____
Thread->pthreads.h file
to compile:
gcc thread demo.c -o test -lpthread
./test
Note: main thread has to wait for other threads to complete
Threads alone cannot survive without main thread
We can have same func , executed by many threads but different arg
_____
limitations of thread
1) If one thread has issue out of 5 threads in a process,
=>entire process will be terminated
2) race condition
to achieve synchronization in threads =>mutex locks
makes threads to wait & run sequentially
application: print multiplication table
void myfun1()
     int count=0;
     while(1)
           printf("in myfun1\n");
           //printf("in myfun1 process id = %d thread id = lu\n",getpid(),pthread_self());
           printf("count = %d\n",count);
```

count++;
if(count == 60)

if(count == 100)

if(count == 140)

usleep(500000);

pthread_cancel(tid2);

pthread_cancel(tid3);

pthread_cancel(pthread_self());

All threads share the same memory,pcb,code/data/stack/heap segment

```
prodeep-200905384@prodeep-200905384:-/Documents/LinuxPractice/day3$ gcc threadsync2.c -o ts2 -lpthread prodeep-200905384@prodeep-200905384:-/Documents/LinuxPractice/day3$ ./ts2
Thread 1s ts trying to take lock
Thread 1s takes lock
Multiplication Table of 15
Thread 6 ts trying to take lock
Thread 12 ts trying to take lock
15 x 1=15
15 x 2=36
15 x 4=45
15 x 4=75
15 x 6=90
15 x 7=105
15 x 7=105
15 x 9=130
15 x 9=135
15 x 10=150

Thread 6 takes lock
Multiplication Table of 6
6 x 1=6
6 x 2=12
6 x 3=18
6 x 4=24
6 x 8=48
6 x 0=54
6 x 1=6
6 x 2=12
6 x 3=18
6 x 4=24
6 x 3=36
6 x 1=6
6 x 1=6
6 x 1=6
6 x 1=6
6 x 2=2
12 x 1=12
12 x 1=12
12 x 2=2
12 x 1=12
12 x 2=2
12 x 1=12
12 x 2=2
12 x 1=12
13 x 1=12
14 x 1=12
15 x 1=12
16 x 1=12
17 x 1=12
18 x 1=12
18
```