
COMP232 – Data Structures & Problem Solving
Fall 2023

Homework #6 – Sorting
[50 points]

1. Implement the `insertionSort` method in the `CS232SortableLinkedList` class in the `hw06` code. This method should perform an insertion sort on the linked list and must run in $O(n^2)$ time. Note: You do not need to change any of the links, swap the elements in the nodes instead. The `No1Tests` class contains tests that you can use to check your implementation of this functionality.

2. Consider each of the following computations on an initially unordered list of integer values:

a. Find the minimum value.

```
minVal <- list[0]
for I in range 1 to n(list length) - 1:
    if list[I] < minVal:
        minVal = list[I]
return minVal
```

$O(n)$ because it goes through the list once

b. Find the median (i.e., the middle value).

```
sort(list)
val <- length(list)
if length(val) is odd:
    return list[val/2]
else:
    return (list[(val/2) + 1] + list[val/2]) / 2
```

First sort the list. Then find the middle index value by splitting list in half

c. Find the 10 largest values.

```
sort(list)
for I in range 1 to 10:
    list1.add[size-i]
return list1
```

The pre sort takes $O(n \log n)$ and it's $O(1)$ to add each element into the array. The overall run time is $O(n \log n)$

For each computation above:

- i. Briefly describe, in a few sentences or pseudo code, an efficient algorithm to perform the computation.
- ii. Give and briefly justify an asymptotic upper bound on your algorithm's worst case running time.

3. In our implementation of Insertion Sort we worked backward from the $i-1^{\text{st}}$ position while swapping to find the location at which to insert the i^{th} value. This is essentially a linear search. However, because we know that the first $i-1$ values are already sorted we could have used a binary search to find the proper location at which to insert the i^{th} value. Is this a useful idea? Why or why not?

Binary search is more effective when trying to access elements at arbitrary indices and takes constant time, but since we are using a linked list, to find the middle element we'd have to traverse through the linked lists which would take linear time.

4. Complete the `heapSort` method in the `HeapSort` class so that given an array of integers they are sorted into descending order in $O(n \lg n)$ time using `HeapSort`. **You can run the `main` method to see if your sort works.**

5. In an application where we need to sort lists that we know will already be nearly sorted indicate which sort would you expect to run faster and briefly justify your answer:

- a. An unoptimized merge sort or insertion sort?

Insertion sort because it works better on nearly sorted lists since it requires shifting of elements. It would operate in $O(n)$ time now since it's doing next to no shifting

- b. An unoptimized merge sort or a heap sort?

Merge sort is more efficient and has lower swaps counts. Heap sort doesn't benefit from the nearly sorted structure because it always involves building a heap and heapifying it. These takes does more swaps.