**Definition 0.1.** Let  $\vec{v} = (v_1, ..., v_n)$  be an n-dimentional vector. We will denote the (n-1)-dimentional vector in which the i-th coordinate is removed by  $\vec{v_{-i}} = (v_1, ..., v_{i-1}, v_{i+1}, ..., v_n)$ . Thus we have three equivalent notations:  $\vec{v} = (v_1, ..., v_n) = (v_i, \vec{v_{-i}})$ .

**Definition 0.2.** A (direct revelation) mechanism is a choice function f and a vector of utility functions  $u_1^f, ..., u_n^f$ 

$$f: V_1 \times \dots \times V_n \to A \tag{0.1}$$

$$u_i^f: V_1 \times \dots \times V_n \to \mathbb{R}$$
 (0.2)

**Definition 0.3.** A mechanism  $(f, u_1^f, ..., u_n^f)$  is called incentive compatible if  $\forall i, \exists v_i \in V_i, \forall v_i' \in V_i$ ,

$$u_i^f(v_i, v_{-i}) \ge u_i^f(v_i', v_{-i})$$
 (0.3)

**Definition 0.4.** A game with (independent private values and) strict incomplete information for a set of n players is given by the following ingredients:

- (i) For every player i, a set of actions  $X_i$ .
- (ii) For every player i , a set of  $types\ T_i$  . A value  $t_i\in T_i$  is the private information that i has.
- (iii)For every player i, a utility function  $u_i: T_i \times X_1 \times ... \times X_n \to \mathbb{R}$ , where  $u_i(t_i, x_1, ..., x_n)$  is the utility achieved by player i, if his type is  $t_i$ , and the profile of actions taken by all players is  $x_1, ..., x_n$ .

**Definition 0.5.** (i) A strategy of a player i is a function  $s_i: T_i \to X_i$ .

(ii) A strategy  $s_i$  is a (weakly) dominant strategy if for every  $t_i$  we have that the action  $s_i(t_i)$  is a dominant strategy in the full information game defined by  $t_i$ . Formally: For all  $t_i$ , all  $x_{-i}$  and all  $x_i'$  we have that

$$u_i(t_i, s_i(t_i), x_{-i}) \ge u_i(t_i, x_i', x_{-i})$$
 (0.4)

A profile  $s_1, ..., s_n$  is called a dominant strategy equilibrium if each  $s_i$  is a dominant strategy.

**Definition 0.6.** (i)A mechanism for n players is given by

- (a)players' type spaces  $T_1, ..., T_n$ ,
- (b) players' action spaces  $X_1, ..., X_n$ ,
- (c)an alternative set A,
- (d)an outcome function  $a: X_1 \times ... \times X_n \to A$  and,

The game with strict incomplete information induced by the mechanism is given by using the types spaces  $T_i$ , the action spaces  $X_i$ , and the utilities  $u_i(t_i, x_1, ..., x_n)$ .

(ii) The mechanism implements a choice function  $f: T_1 \times ... \times T_n \to A$  in dominant strategies if for some dominant strategy equilibrium  $s_1, ..., s_n$  of the induced game, where  $s_i: T_i \to X_i$ , we have that for all  $t_1, ..., t_n$ ,  $f(t_1, ..., t_n) = a(s_1(t_1), ..., s_n(t_n))$ .

**Proposition 0.1** (Revelation Principle). If there exists an arbitrary mechanism that implements f in dominant strategies, then there exists an incentive compatible mechanism that implements f. The payments of the players in the incentive compatible mechanism are identical to those, obtained at equilibrium, of the original mechanism.

*Proof.* The new mechanism will simply simulate the equilibrium strategies of the players . That is , let  $s_1, ..., s_n$  be a dominant strategy equilibrium of the original mechanism , we define a new direct revelation mechanism :

$$f(t_1, ..., t_n) := a(s_1(t_1), ..., s_n(t_n))$$
(0.5)

$$u_i^f(t_1, ..., t_n) := u_i(t(i), s_1(t_1), ..., s_n(t_n)), \text{ where } t(i) := t_i$$
 (0.6)

. Now since each  $s_i$  is a dominant strategy for player i , then for every  $t_i, x_{-i}, x_i^\prime$  we have that

$$u_i(t(i), s_i(t_i), x_{-i}) \ge u_i(t(i), x_i', x_{-i})$$
 (0.7)

$$u_{i}(t(i), s_{i}(t_{i}), s_{-i}(t_{-i})) \ge u_{i}(t(i), s_{i}(t'_{i}), s_{-i}(t_{-i}))$$

$$(0.8)$$

$$u_i^f(t_i, t_{-i}) \ge u_i^f(t_i', t_{-i})$$
 (0.9)

, which gives the definition of incentive compatibility of the mechanism  $(f,u_1^f,...,u_n^f)$  .  $\Box$