

# The Černý Conjecture and Group Representations

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SATA September 4, 2008

# Synchronizing Automata

- An automaton  $\mathcal{A} = (Q, \Sigma)$  is *synchronizing* if there exists  $w \in \Sigma^*$  such that  $|Qw| = 1$ .
- Such a word  $w$  is called a *reset word*.

## Conjecture (Černý '64)

A synchronizing automaton with  $n$  states admits a reset word of length at most  $(n - 1)^2$ .

- The best known upper bound on the length of reset words is  $\frac{n^3 - n}{6}$ , due to Pin based on a non-trivial combinatorial result of Frankl.
- One can obtain a bound of  $\frac{n^3 - n}{3}$  with straightforward methods.
- Improving a bound by a factor of 2 can be hard work!

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# Pin's Theorem

- The literature on Černý's conjecture consists of a vast array of partial results, the first of which was Pin's Theorem.
- If  $(Q, \Sigma)$  is an automaton, we view  $\Sigma \subseteq T_Q$  (the semigroup of self-maps of  $Q$ ).

## Theorem (Pin '78)

*Let  $\mathcal{A} = (Q, \Sigma)$  be a synchronizing automaton so that  $|Q|$  is a prime  $p$  and some element of  $\Sigma$  cyclically permutes  $Q$ . Then:*

- 1  *$\mathcal{A}$  is synchronizing if and only if  $\Sigma$  contains a non-permutation;*
- 2 *In this case,  $\mathcal{A}$  has a reset word of length at most  $(p - 1)^2$ .*

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# Synchronizing Groups

- Motivated by the first part of Pin's Theorem, I defined in 2004 the notion of a synchronizing group.
- A group  $G \subseteq S_n$  is called *synchronizing* if, for all non-permutations  $t \in T_n$ , the monoid  $\langle G \cup t \rangle$  contains a constant map.
- Pin's Theorem implies that cyclic groups of prime order are synchronizing.
- It is easy to see that 2-transitive groups are synchronizing.
- With Arnold, I proved synchronizing groups are primitive and gave a sufficient condition for a group to be synchronizing in terms of representation theory that covers the above results.
- João Araújo independently came up with the notion in 2006 and found a beautiful group theoretic reformulation.
- Synchronizing groups have recently received quite a bit of attention from prominent group theorists including Peter Neumann, Jan Saxl, Peter Cameron and Csaba Schneider.



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# Dubuc's Theorem

- Dubuc extended the second part of Pin's Theorem to arbitrary automata containing a cyclic permutation via an ingenious linear algebraic argument.

## Theorem (Dubuc '98)

*Let  $\mathcal{A} = (Q, \Sigma)$  be a synchronizing automaton on  $n$  states such that  $\Sigma$  contains a cyclic permutation of the states. Then  $\mathcal{A}$  has a reset word of length at most  $(n - 1)^2$ .*

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# Černý Cayley Graphs

- Let  $G$  be a group of order  $n$  and  $\Delta$  a generating set of  $G$ .
- The automaton  $(G, \Delta)$  is called the *Cayley graph* of  $G$  with respect to  $\Delta$ . A typical transition is of the form  $g \xrightarrow{a} ga$  with  $g \in G$ ,  $a \in \Sigma$ .
- Let us say that an automaton  $\mathcal{A}$  *contains* the Cayley graph  $(G, \Delta)$  if  $\mathcal{A} = (G, \Sigma)$  where  $\Delta \subseteq \Sigma$ .
- So  $\mathcal{A}$  is obtained from the Cayley graph by adding new transitions but no new states.
- Call  $(G, \Delta)$  a *Černý Cayley graph* if every synchronizing automaton containing it has a reset word of length at most  $(n - 1)^2$ .
- Let's say  $G$  is a *Černý group* if all its Cayley graphs are Černý Cayley graphs.
- Dubuc's theorem says that  $(\mathbb{Z}, \{1\})$  is a Černý Cayley graph.
- Cyclic groups of prime power order are Černý groups.

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# Rystsov's Theorem

- The above notion was implicitly considered by Rystsov.
- In 1995, he proved a synchronizing automaton containing the Cayley graph of a group of order  $n$  admits a reset word of length  $\leq 2(n-1)^2$ .
- He proved in fact a slightly better result.
- Let  $(G, \Delta)$  be a Cayley graph with  $|G| = n > 1$ .
- Define  $\text{diam}_\Delta(G)$  to be the least  $m$  so that any two states of  $(G, \Delta)$  can be connected by a word of length at most  $m$ .
- $1 \leq \text{diam}_\Delta(G) \leq n-1$ .

## Theorem (Rystsov '95)

*A synchronizing automaton containing the Cayley graph  $(G, \Delta)$  has a reset word of length at most  $1 + (n-1 + \text{diam}_\Delta(G))(n-2)$ .*

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# Our Goal

- Recall Rystsov's bound is  $1 + (n - 1 + \text{diam}_\Delta(G))(n - 2)$  and  $(n - 1)^2 = 1 + n(n - 2)$ .
- So Rystsov's bound only achieves the Černý bound when the diameter is 1, i.e., all non-trivial elements of  $G$  belong to the generating set.
- We aim to improve his bound so that in many cases we achieve the Černý bound.
- Even when we do not achieve the Černý bound with our main result, our techniques often suffice to establish a family of Cayley graphs is Černý.
- Our results lead to several new families of Černý Groups.
- Our main tool is representation theory.

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- Our main tool is representation theory.

# Representation Theory

- Fix a field  $K$  of characteristic 0 (usually  $\mathbb{Q}$ ).
- All vector spaces  $V$  will be over  $K$  and finite dimensional.
- Fix a group  $G$  of order  $n$ .
- A *representation* of  $G$  is a homomorphism  $\varphi: G \rightarrow GL(V)$  to the group of invertible linear transformations on  $V$ .
- $\dim V$  is called the *degree* of  $\varphi$ .
- $K^G = \{f: G \rightarrow K\}$  is a vector space of dimension  $n$ .
- $\lambda: G \rightarrow GL(K^G)$  given by

$$\lambda_g(f)(x) = f(xg)$$

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# Invariant Subspaces

- Let  $\varphi: G \rightarrow GL(V)$  be a representation.
- $W \leq V$  is called a *G-invariant subspace* if  $\varphi_g W \subseteq W$  all  $g \in G$ .
- $0$  and  $V$  are the trivial  $G$ -invariant subspaces.
- For the regular representation  $\lambda: G \rightarrow GL(K^G)$  there are two important invariant subspaces.
- The subspace  $V_1$  of constant functions is  $G$ -invariant.
- If  $f(x) = c$  all  $x \in G$ , then

$$\lambda_g f(x) = f(xg) = c$$

and so  $\lambda_g f = f$ . Thus  $V_1$  is  $G$ -invariant.

- The space  $V_0$  of all functions  $f$  so that  $\sum_{x \in G} f(x) = 0$  is also  $G$ -invariant of dimension  $n - 1$ .

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# Irreducible Representations

- A representation  $\varphi: G \rightarrow GL(V)$  is *irreducible* if it admits no non-trivial  $G$ -invariant subspaces.
- A degree 1 representation is obviously irreducible.
- Any representation can be uniquely expressed as a direct sum of irreducible representations (*irreps*).
- Every irrep appears as summand in the decomposition of the regular representation.
- The degree of any irrep is between 1 and  $n - 1$ .

## Definition

Define  $m(G)$  to be the maximal degree of an irrep of  $G$  over  $\mathbb{Q}$ .

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# The Main Result

## Theorem (BS)

*Let  $(G, \Delta)$  be a Cayley graph of a group of order  $n$ . Then any synchronizing automaton containing  $(G, \Delta)$  admits a reset word of length at most*

$$1 + (n - m(G) + \text{diam}_\Delta(G))(n - 2).$$

*In particular, if  $\text{diam}_\Delta(G) \leq m(G)$ , then  $(G, \Delta)$  is a Černý Cayley graph.*

- The last statement follows since  $(n - 1)^2 = 1 + n(n - 2)$ .
- $m(G) = 1$  iff  $G \cong \mathbb{Z}_2^k$  for some  $k$ , as irreps separate points.
- So we beat Rystsov's bound of  $1 + (n - 1 + \text{diam}_\Delta(G))(n - 2)$  in essentially all cases.



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# Cyclic groups

- The regular representation of  $\mathbb{Z}_n$  is isomorphic to  $\mathbb{Q}[x]/(x^n - 1)$  where the generator acts by multiplication by  $x$ .
- $\mathbb{Q}[x]/(x^n - 1) \cong \bigoplus_{d|n} \mathbb{Q}(\zeta_d)$  ( $\zeta_d$  is a primitive  $d^{\text{th}}$ -root of unity).
- So  $m(\mathbb{Z}_n) = \phi(n)$ .
- As  $\text{diam}_{\{1\}}(\mathbb{Z}) = n - 1$ , we achieve the Černý bound if and only if  $\phi(n) = n - 1$ .
- In particular, we recover Pin's Theorem, but not Dubuc's Theorem (although we are very close).
- Suppose  $p < q$  are odd primes and  $n = pq$ .
- Then  $\text{diam}_{\{p,q\}}(\mathbb{Z}_n) = q - 1 + p - 1$  ( $\mathbb{Z}_n \cong \mathbb{Z}_q \times \mathbb{Z}_p$ ).
- $\phi(n) = (p - 1)(q - 1) \geq q - 1 + p - 1$ .
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# Direct Products of Cyclic Groups of Prime Order

- Let  $p$  be a prime.
- To show that  $\mathbb{Z}_p^k$  is a Černý group, it suffices to consider the Cayley graph with respect to a basis  $\Delta$ .
- $\text{diam}_\Delta(\mathbb{Z}_p^k) = k(p-1)$ .
- One can prove  $m(\mathbb{Z}_p^k) = p-1$ : each irrep factors through a map to  $\mathbb{Z}_p$ .
- Our bound therefore is not strong enough when  $k > 1$ . Nonetheless we can prove:

## Theorem (BS)

*The group  $\mathbb{Z}_p^k$  is a Černý group for  $p$  prime, all  $k \geq 1$ .*

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# Dihedral Groups

- Let  $D_n$  be the dihedral group of order  $2n$  (symmetry group of a regular  $n$ -gon).
- Let  $\Delta$  consist of a reflection and a rotation by  $2\pi/n$ .
- Then  $\text{diam}_\Delta(D_n) \leq \lceil \frac{n+1}{2} \rceil$ .
- One can prove  $m(D_n) = \phi(n)$ . (Act on  $\mathbb{Q}(\zeta_n)$  by having the reflection act as complex conjugation and the rotation act as multiplication by  $\zeta_n$ .)
- If  $n = p^a q^b$  where  $p \leq q$  are odd primes, then one verifies that  $\lceil \frac{n+1}{2} \rceil \leq \phi(n)$  and so we obtain a Černý Cayley graph.

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# Symmetric Groups

- It is known that the symmetric group  $S_n$  has  $p_n$  irreducible representations over  $\mathbb{Q}$  where  $p_n$  is the number of partitions of  $n$ .
- The sum of the squares of the degrees of the irreps of  $S_n$  is  $n!$ .
- Thus  $m(S_n)^2 p_n \geq n!$ , i.e.,  $m(S_n) \geq \sqrt{n!/p_n}$ .
- $p_n \sim \frac{\exp(\pi\sqrt{2n/3})}{4n\sqrt{3}}$  and  $n! \sim \sqrt{2\pi n} \left(\frac{n}{e}\right)^n$ .
- Therefore,  $m(S_n)$  grows extremely quickly as a function of  $n$ .
- With Coxeter-Moore generators  $(1\ 2), (2\ 3), \dots, (n-1\ n)$ , the diameter is  $\binom{n}{2}$  [think “Bubble Sort”] and so we obtain a Černý Cayley graph for  $n$  large enough.
- With the generating set  $(1\ 2), (1\ 2\ \dots\ n)$ , the diameter of  $S_n$  is at most  $\binom{n}{2}(n+1)$  and so we again get a Černý Cayley graph for  $n$  large enough.

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# Affine Groups

- Let  $p$  be a prime.
- The affine group  $AG(1, p)$  is the group of all functions  $f: \mathbb{Z}_p \rightarrow \mathbb{Z}_p$  of the form  $f(x) = ax + b$  with  $a \in \mathbb{Z}_p^*$  and  $b \in \mathbb{Z}_p$ .
- $AG(1, p) \cong \mathbb{Z}_p \rtimes \mathbb{Z}_p^*$ .
- It turns out  $m(AG(1, p)) = p - 1$  (act on  $\mathbb{Q}[\zeta_p]$ ), which is not in general larger than the diameter.

## Theorem (BS)

*Let  $\Delta$  be a generating set for  $AG(1, p)$  so that each translation can be represented by a word in  $\Delta^*$  of length at most  $p - 1$ . Then  $(AG(1, p), \Delta)$  is a Černý Cayley graph. This applies in particular if  $\Delta$  contains a translation.*

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# Special Linear Groups

- Let  $p$  be a prime.
- $SL(2, p) = \left\{ \begin{bmatrix} a & b \\ c & d \end{bmatrix} \mid a, b, c, d \in \mathbb{Z}_p, ad - bc = 1 \right\}$ .
- A standard generating set  $\Delta$  for  $SL(2, p)$  consists of the matrices

$$\begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix}, \begin{bmatrix} 1 & 0 \\ 1 & 1 \end{bmatrix}.$$

- The diameter with this generating set is no more than  $3p - 2$ .
- Estimating  $m(SL(2, p))$  is a bit more complicated.

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# Characters

- To estimate  $m(SL(2, p))$  we make use of character theory.
- Let  $\varphi: G \rightarrow GL(V)$  be a representation of  $G$  over  $K \leq \mathbb{C}$ .
- The *character*  $\chi_\varphi$  of  $\varphi$  is defined by  $\chi_\varphi(g) = \text{Trace}(\varphi_g)$ .
- $\chi_\varphi(1)$  is the degree of  $\varphi$  (it is the trace of the identity map).
- A representation is determined up to isomorphism by its character.
- If  $|G| = n$ , then  $\chi_\varphi(g)$  is a sum of  $n^{\text{th}}$ -roots of unity.
- The *character field*  $\mathbb{Q}(\chi_\varphi)$  is the field extension of  $\mathbb{Q}$  generated by the  $\chi_\varphi(g)$  with  $g \in G$ .
- It is a finite Galois extension of  $\mathbb{Q}$ .

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- Let  $\chi$  be an irreducible complex character of  $G$  and let  $H = \text{Gal}(\mathbb{Q}(\chi) : \mathbb{Q})$ .
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## $m(SL(2, p))$

- Computation of the irreducible complex characters of  $SL(2, p)$  go back to Frobenius and Schur.
- $SL(2, p)$  has irreducible complex characters  $\chi_1$  and  $\chi_2$  of degrees  $p + 1$  and  $p - 1$  (respectively) with respective character fields  $\mathbb{Q}(\cos \frac{2\pi}{p-1})$  and  $\mathbb{Q}(\cos \frac{2\pi}{p+1})$ .
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- The diameter of the Cayley graph of  $SL(2, p)$  with our generators was at most  $3p - 2$ .

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*Let  $p \geq 17$  be a prime. Then the Cayley graph of  $SL(2, p)$  with respect to the generators  $\begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix}, \begin{bmatrix} 1 & 0 \\ 1 & 1 \end{bmatrix}$  is a Černý Cayley graph.*

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# The Main Idea

- Let  $(G, \Sigma)$  be a synchronizing automaton containing the Cayley graph of  $(G, \Delta)$ .
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Thanks for your attention!