

Profinite groups associated to symbolic dynamical systems

Benjamin Steinberg

City College of New York

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- It is compact and totally disconnected (in fact homeomorphic to the Cantor set).
- There is a natural action of \mathbb{Z} on $A^{\mathbb{Z}}$ via the **shift** map $\sigma: A^{\mathbb{Z}} \rightarrow A^{\mathbb{Z}}$ given by

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- A **symbolic dynamical system**, or **subshift**, or simply **shift**, is a closed, non-empty, shift-invariant subspace of $A^{\mathbb{Z}}$.

Conjugacy

- Subshifts $\mathcal{X} \subseteq A^{\mathbb{Z}}$ and $\mathcal{Y} \subseteq B^{\mathbb{Z}}$ are **conjugate** if there is a \mathbb{Z} -equivariant homeomorphism $\psi: \mathcal{X} \rightarrow \mathcal{Y}$, i.e.,

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- Usually one wants to classify shifts up to conjugacy, although sometimes weaker equivalence relations are considered.

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- Symbolic encodings of dynamical systems on manifolds.

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- A word $w \in A^{\mathbb{Z}}$ generates a minimal subshift iff it is **uniformly recurrent**.
- This means that if v is a finite factor of w , then there exists $N > 0$ so that each factor of w of length N contains v as a factor: the “**bounded gaps property**.”

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- The general case is obtained by replacing ab with an arbitrary finite word.

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- We will only be interested in irreducible subshifts.

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- So shifts are boundaries of factorial prolongable languages.
- \mathcal{X} is **irreducible** iff, for all $u, v \in L(\mathcal{X})$, there exists $w \in A^*$ so that $uwv \in L(\mathcal{X})$.

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- Up to conjugacy one can assume f is proper.

Examples of primitive endomorphisms

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$$f(a) = ab, \quad f(b) = ba$$

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- Any endomorphism of A^* extends to an endomorphism of the free group F_A .
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- An endomorphism of A^* that extends to an automorphism of F_A is called a **positive automorphism**.

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- Much attention has been devoted to classifying shifts of finite type up to conjugacy.

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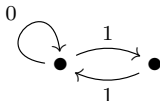
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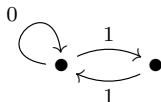
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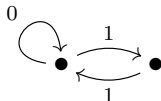
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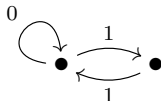
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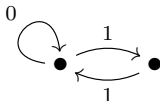
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- In the above example all vertices are initial and terminal.

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- A minimal sofic shift must be periodic (follow a cycle in the automaton).

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- In other words, L is regular iff it is saturated by a finite index congruence.

The algebraic characterization of regular languages

- There is an alternative definition of regular languages that is more algebraic.

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- The regular languages form a boolean algebra.

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- J. Almeida had the idea of relating symbolic dynamical systems to the universal compact totally disconnected boundary of a factorial prolongable language.
- Namely, he considered the boundaries of these languages inside the profinite completion of the free monoid.

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- $\widehat{A^*}$ is as the **Stone dual** of the boolean algebra of regular languages over A^* .

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$$\partial C = C \cap \partial\widehat{A}^*.$$

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- In particular, the ideal $I(\mathcal{X})$ associated to an irreducible subshift \mathcal{X} is prime.

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- So there is a profinite group invariant associated to an irreducible subshift via the free profinite monoid!

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Projective profinite groups

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- They are also the profinite groups of cohomological dimension one.

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- If M is a finitely generated profinite monoid, then $\text{End}(M)$ is a profinite monoid in the compact-open topology.
- So, if $f \in \text{End}(M)$, then $f^{n!}$ converges to an idempotent f^ω .

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Theorem (Almeida, Costa)

Let $f: A^ \rightarrow A^*$ be a proper aperiodic primitive endomorphism. Then*

$$G(\mathcal{X}_f) = \langle A \mid f^\omega(a) = a, a \in A \rangle$$

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- This answers negatively a part of Margolis's question.

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- I can now prove it using cohomology of profinite monoids.

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- Since $\text{cd } \widehat{A^*} = 1$, the theorem implies projectivity of maximal subgroups of $\widehat{A^*}$.

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- Standard profinite semigroup theory then implies $s \in C$.

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- $C = \{s^n, \dots, s^{n+m-1}\}$ is a finite cyclic subgroup with identity s^k where $n \leq k \leq n+m-1$ is divisible by m .
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Theorem (Rhodes, BS)

Every element of finite order in \widehat{A}^ is an idempotent.*

The free profinite group on a set converging to 1

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- Every metrizable profinite group has a countable set of generators converging to 1.

The maximal subgroup of the minimal ideal is free

Theorem (BS)

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- The proof relies on **Iwasawa's criterion**: a countably based profinite group G is free of countable rank iff given a diagram

$$\begin{array}{ccc} & G & \\ & \downarrow \varphi & \\ A & \xrightarrow{\alpha} & B \end{array}$$

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of epimorphisms (A and B are finite), there exists an **epimorphism** $\lambda: G \twoheadrightarrow A$ so that the diagram commutes.

Profinite groups associated to irreducible sofic subshifts

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- Minimal sofic shifts are periodic and hence have free procyclic associated groups.
- So the interesting case is the non-minimal case.

A topological characterization of sofic shifts

- The fact $\widehat{A^*}$ is the Stone dual of the boolean algebra of regular languages has a topological consequence.

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- This is almost enough to generalize my proof from the full shift to an arbitrary non-periodic sofic shift.
- But technically my proof scheme only works for irreducible sofic shifts containing a periodic subshift defined over a strictly smaller alphabet.

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Theorem (Costa, BS)

The profinite group associated to a non-periodic irreducible sofic shift is a free profinite group of countable rank.

Idempotent generators of the ideal of a sofic shift

- It is natural to ask where the idempotents corresponding to irreducible sofic shifts 'sit' in $\widehat{A^*}$.

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- So there is a dense set of idempotents whose associated maximal subgroups are free profinite.
- Given a strongly connected automaton accepting an irreducible sofic shift \mathcal{X} , we can effectively construct an idempotent generator of $I(\mathcal{X})$.

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- Is the profinite group associated to a minimal shift finitely generated?
- What are the possible finite subsemigroups of a free profinite monoid?

The end

THANK YOU FOR YOUR
ATTENTION!