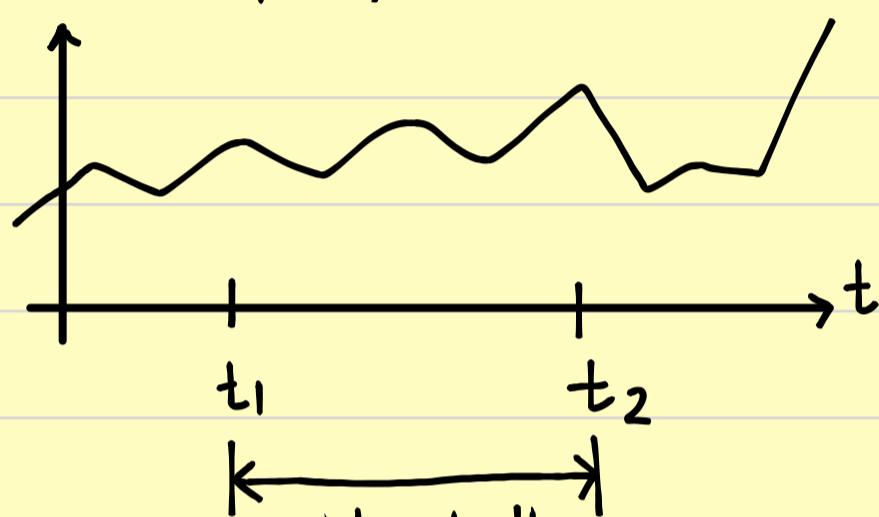


Datasets can be classified in 2 groups:

- EUCLIDEAN. if a metric can be defined to describe the distance.
- NON-EUCLIDEAN. otherwise.

Examples of EUCLIDEAN Datasets:

- Timeseries
(1D)

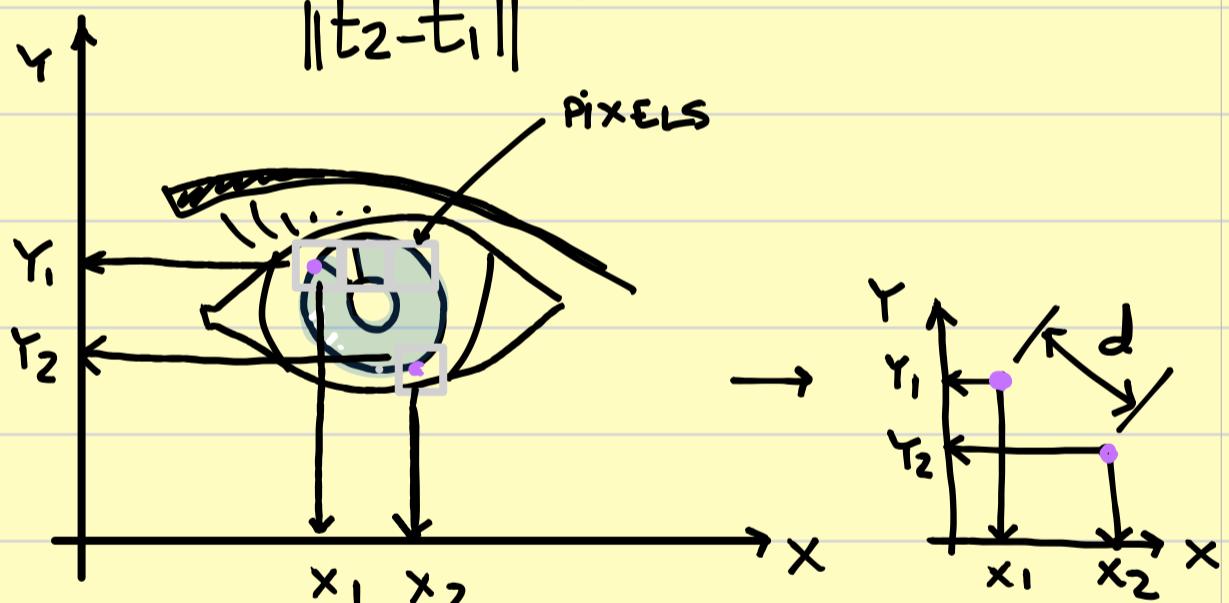


- Picture B&W
(2D)

$$\square : [0, 255]$$

8bits: 0 1 0 1 1 1 1 0 0

$$2^8 = 256$$



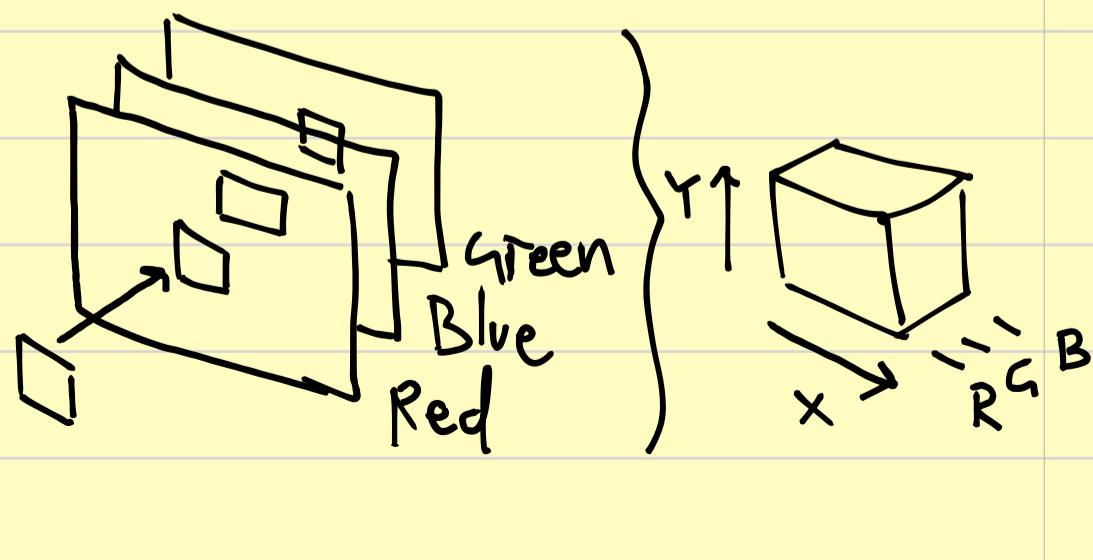
$$d = \sqrt{(x_2 - x_1)^2 + (y_2 - y_1)^2}$$

- Picture Color
(3D)

$$R : [0, 255]$$

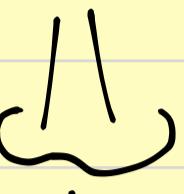
$$G : [0, 255]$$

$$B : [0, 255]$$



. Video Color

(3D) + time + sound \rightarrow (5D)

.  Olfactive
(12D)

. Haptic
(6D)

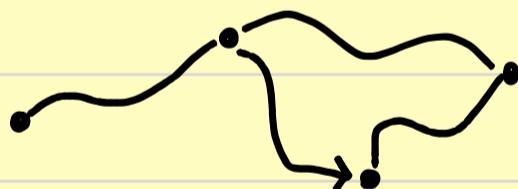
...

The hypothesis is that the underlying structure of the dataset allows for the calculation of a distance. This implies there is a way to ..measure.. between points in the space.

Example of non-euclidean datasets: NETWORKS .

What is a network ?

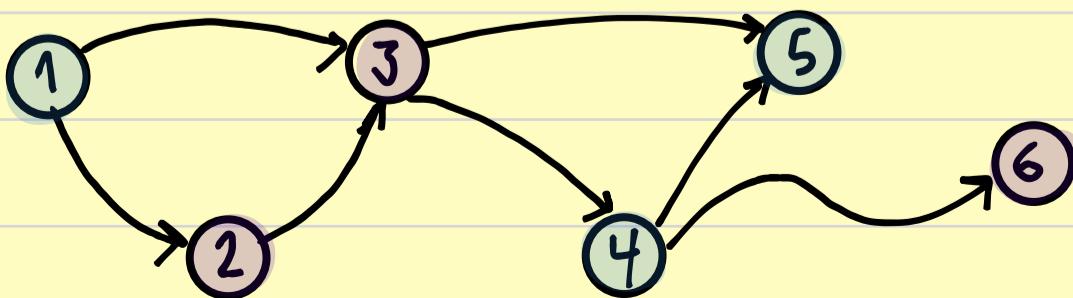
Nodes + Edges + Attributes



A network is defined (mathematically) by a set (group) of nodes and a set of edges. This group of sets is called .. GRAPH .. (G).

Network \equiv GRAPH $\equiv G(N \cup E)$ [N.Nodes, E.edges]

Example:



$$N: \{1, 2, 3, 4, 5, 6\} \quad E: \{(1 \rightarrow 2); (1 \rightarrow 3); (2 \rightarrow 3); \\ (3 \rightarrow 4); (3 \rightarrow 5); (4 \rightarrow 5); \\ (4 \rightarrow 6)\} \quad G = \{NUE\}$$

Attributes:



Networks can be directed or undirected:

- Directed if the edges .. have an arrow
 - Undirected otherwise
-

HOW CAN WE QUANTIFY & COMPARE NETWORKS?

1. AVERAGE PATH LENGTH (APL)
2. CLUSTERING COEFFICIENT (CC)
3. LAPLACIAN OF NETWORK

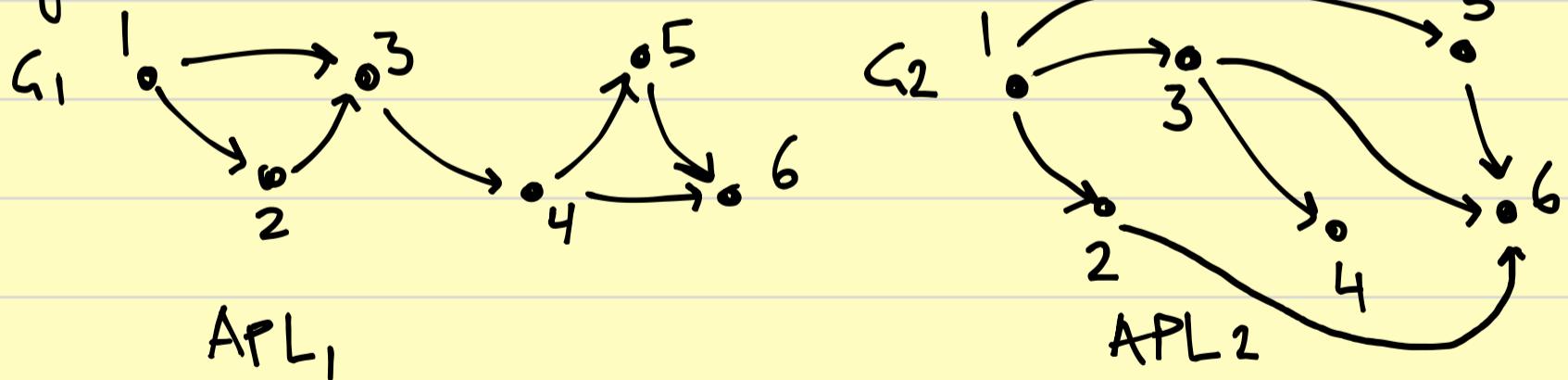
1. APL . Average distance (steps) between the nodes in the network.

$$APL = \frac{1}{N(N-1)} \cdot \sum_{i=1}^N \sum_{j=1}^N d_{ij}$$

- The maximum # of relationships in a network of N elements is $N \cdot (N-1)$ because each element relates to all others, hence $(N-1)$, N times.
- $\sum_{i=1}^N \sum_{j=1}^N d_{ij}$ represents the sum of all paths between nodes.

$$APL = \frac{1}{6.5} \cdot \left[\begin{array}{c} 1 \\ d_{12} \quad d_{13} \quad d_{14} \quad d_{15} \quad d_{16} \\ 1 + 1 + 2 + 2 + 3 \\ + \\ 2 \\ d_{21} \quad d_{23} \quad d_{24} \quad d_{25} \quad d_{26} \\ 1 + 1 + 2 + 2 + 3 \\ + \\ 3 \\ d_{31} \quad d_{32} \quad d_{34} \quad d_{35} \quad d_{36} \\ 1 + 1 + 1 + 1 + 2 \\ + \\ 4 \\ d_{41} \quad d_{42} \quad d_{43} \quad d_{45} \quad d_{46} \\ 2 + 2 + 1 + 1 + 1 \\ + \\ 5 \\ d_{51} \quad d_{52} \quad d_{53} \quad d_{54} \quad d_{56} \\ 2 + 2 + 1 + 1 + 2 \\ + \\ 6 \\ d_{61} \quad d_{62} \quad d_{63} \quad d_{64} \quad d_{65} \\ 3 + 3 + 2 + 1 + 2 \end{array} \right] = \dots$$

Example. Compare following two graphs in terms of performance related to APL.



When you compare two networks, the one with the shortest APL would be usually faster/more effective.

2. CC

Describes how good groups are created in the network.

$$CC = \frac{1}{N} \sum_{i=1}^N \frac{2L_i}{k_i(k_i-1)}$$

$$CC = \frac{1}{6} \left[\left[\frac{2 \cdot 1}{2(2-1)} \right] + L_1 = \text{node } \#1 \text{ has two neighbors with 1 rel.} \right]$$

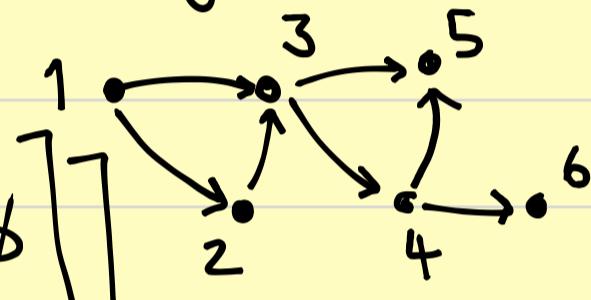
$$+ \left[\frac{2 \cdot 1}{2(2-1)} \right] + \left[\frac{2 \cdot 2}{4(4-1)} \right] +$$

$$+ \left[\frac{2 \cdot 1}{3(3-1)} \right] + \left[\frac{2 \cdot 1}{2 \cdot (2-1)} \right] + \left[\emptyset \right]$$

L_i = Number of relationships between the neighbors of node .. i

(how many of my friends are friends to each other?)

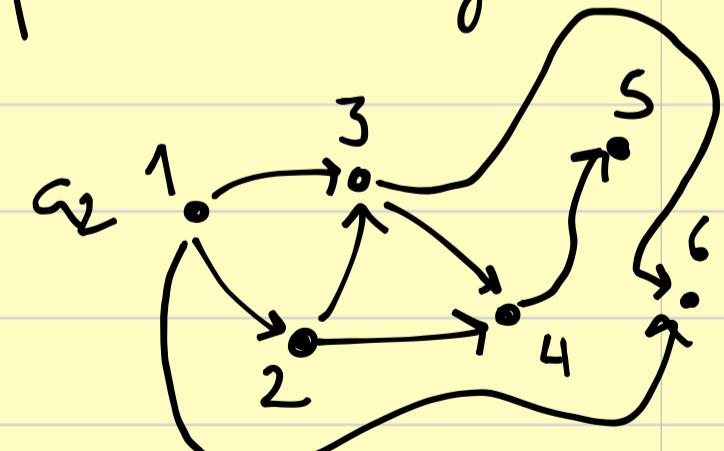
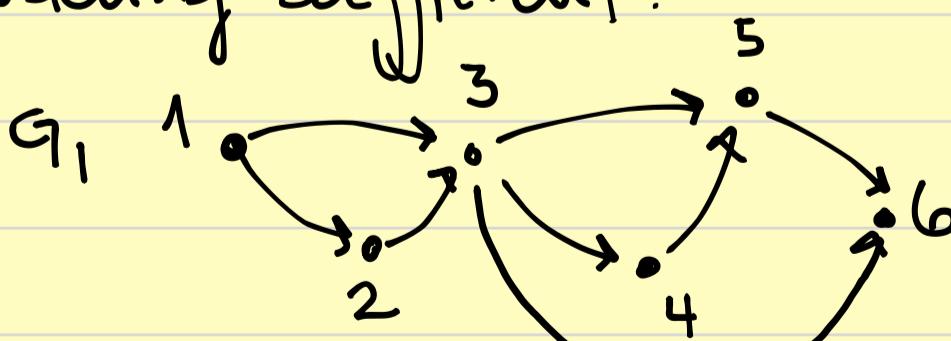
k_i = Number of neighbors of node .. i



= ...

The higher the clustering coefficient, the better the communication between the elements of the network.

Example. Compare these two graphs in terms of clustering coefficient.



3. LAPLACIAN NETWORK

Laplacian matrix of a network describes/contains all relevant information of a network.

Is defined as :

$$L = D - A$$

$L \equiv$ Laplacian Matrix

$D \equiv$ Degree Distribution Mat.

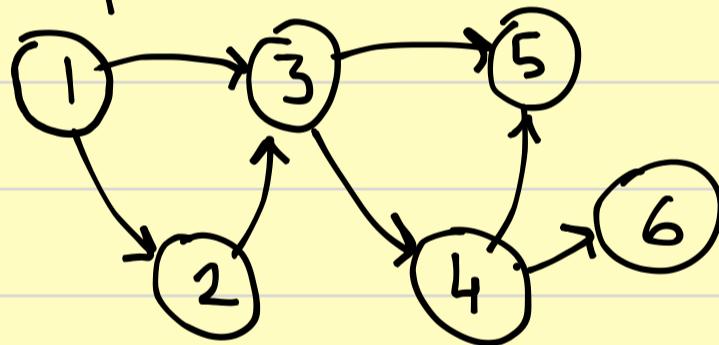
$A \equiv$ Adjacency Matrix

$K_i \equiv$ degree of node

$$D \equiv \text{degree matrix} \quad \left\{ \begin{array}{ll} k_i & i=j \\ 0 & i \neq j \end{array} \right.$$

$$A \equiv \text{adjacency matrix} \quad \left\{ \begin{array}{ll} 1 & \text{if } i \& j \text{ have a connection} \\ 0 & \text{otherwise.} \end{array} \right.$$

Example .



$$\begin{aligned}
 L &= \begin{bmatrix} 2 & & & & & \\ & 2 & & & & \\ & & 0 & & & \\ 3 & & & 4 & & \\ & & & & 3 & \\ & & & & & 2 \end{bmatrix} - \begin{bmatrix} 1 & 2 & 3 & 4 & 5 & 6 \\ 0 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix} = \\
 D &\quad A
 \end{aligned}$$

$$L = \begin{bmatrix} 1 & 2 & 3 & 4 & 5 & 6 \\ 2 & 0 & -1 & -1 & 0 & 0 & 0 \\ 3 & 0 & 0 & 2 & -1 & -1 & 0 \\ 4 & 0 & 0 & 0 & 3 & -1 & -1 \\ 5 & 0 & 0 & 0 & 0 & 2 & 0 \\ 6 & 0 & 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$

NETWORK

$$G = \{N, E\}$$

Non euclidean



LAPLACIAN MATRIX

$$L$$

Euclidean

- The eigenvectors of the laplacian matrix contain the spectral information of the network.

$$A\vec{v} = \lambda\vec{v} \rightarrow \det |A - \lambda I| = 0 \rightarrow \lambda \rightarrow \vec{v}$$

- The second eigenvector of the laplacian matrix of the graph is called the fiedler vector.
- The fiedler vector has an important property: it has a sign structure that shows the bottleneck of the graph.

$$G(N, E) \rightarrow L \rightarrow \vec{v}_2 = \begin{matrix} \text{Fiedler} \\ \text{vector} \end{matrix}$$

$$\begin{bmatrix} + \\ + \\ + \\ 0 \\ 0 \\ \vdots \\ 0 \end{bmatrix}^T_N$$

← Bottleneck of
the graph

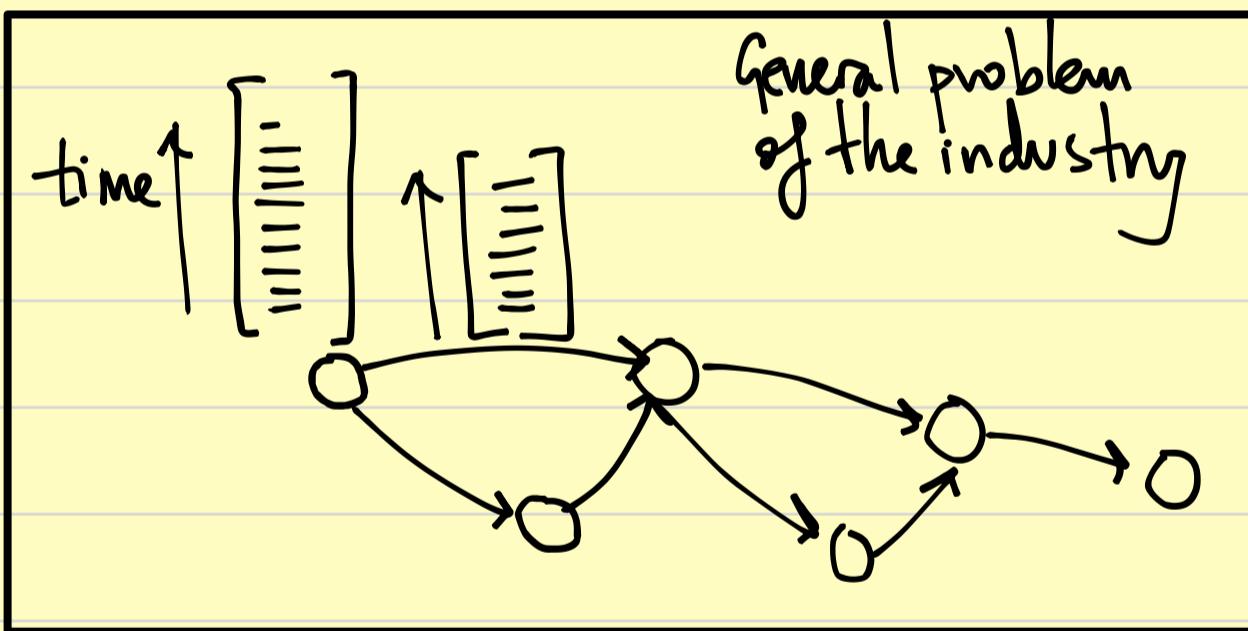
Exercise. please write a Python code 3.x that generates two graphs, each with 100 nodes, and each node attached to a vector of 300 elements that is normally distributed $N_1(k_i, \sqrt{k_i})$; and $N_2(\bar{n}_i^2, \sqrt[3]{k_i})$.

NETWORK 2

NETWORK 1

The code should then calculate the fiedler vector of each graph and you should interpret the results.

The networks should have a Barabasi structure with $\gamma = 2.4$ (exponent degree).



NETWORK TOPOLOGY

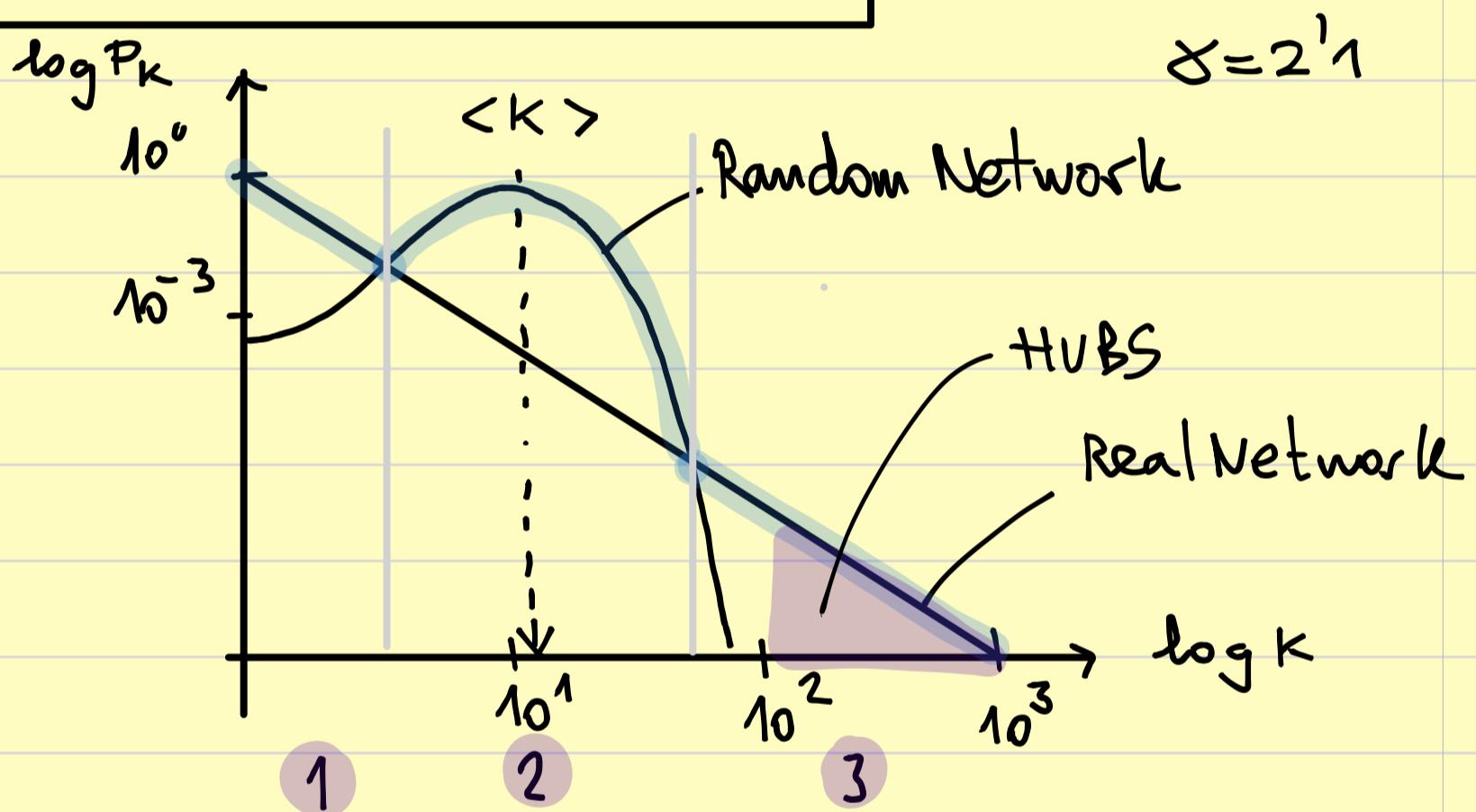
- The probability that a node with k neighbours attaches a new node in a random network is given by a Poisson distribution with parameter λ :

$$P_k = P_{\text{random}}(X=k) = \frac{\lambda^k e^{-\lambda}}{k!}$$

- The probability that a node is connected to another with .. k .. neighbours in **real networks** is given by a **POWER LAW** distribution:

$$P_k = P_{\text{Real}}(X=k) = k^{-\delta}; \quad \delta \equiv \text{Exponent degree}$$

NETWORK SCIENCE ; Barabasi 2016



- 1 RANDOM << REAL . For small k the power law of the real network is above the poisson(random), indicating that the real world has more nodes with few neighbours.
- 2 RANDOM > REAL . For nodes around the average degree $\langle k \rangle$, the random network has excess of nodes.
- 3 RANDOM << REAL . for large k the real network presents more nodes with high k .

The role of the exponent δ :

$\delta \leq 2$. Anomalous regime. Networks cannot exist long.

$2 < \delta < 3$. Scale free regime. " grow healthy"

$\delta > 3$. Random regime.

