

CIS507: Design & Analysis of Algorithms
Quiz 6, Spring 2012
(VERSION WITH ANSWERS)

Duration: 20 minutes
Total weight: 5%

Student Name: - - - - -

Student ID: - - - - -

Problem	Points Obtained	Points Possible
1		1
2		4
Total		5

Cheat list:

P: problems for which a polynomial time algorithm exists.

NP: problems which are polynomial time verifiable.

CoNP: problems which are complements of problems in *NP*.

NP – Hard: problems to which every problem in *NP* can be reduced in polynomial time.

CoNP – Hard: problems which are complements of *NP – Hard* problems.

NP – Complete = $NP \cap NP - Hard$

CoNP – Complete: problems which are complements of *NP – Complete* problems. Equivalently: problems in *CoNP* to which every problem in *CoNP* can be reduced in polynomial time.

1 True/False (1 point)

1. **(0.5 point)** Suppose there exists a *coNP – complete* problem *L* that is also in *NP*. Then $coNP \subseteq NP$.
2. **(0.5 point)** Suppose there exists a *coNP – complete* problem *L* that is also in *NP*. Then $NP \subseteq coNP$.

ANSWER:

1. True. Consider arbitrary problem $L' \in coNP$. Since L is $coNP - Complete$, then $L' \leq_p L$. And since L is in NP , then L' is also in NP .
2. True. Consider arbitrary problem $L'' \in NP$. Therefore problem “is $x \notin L''$?” is in $coNP$. Since L is $coNP - Complete$, then “is $x \notin L''$?” $\leq_p L$. Therefore, “is $x \in L''$?” \leq_p “is $x \notin L''$?”. But we assumed L is in NP , so $x \notin L''$ is in $coNP$. Therefore, “is $x \in L''$?” is also $coNP$.

2 Multiple Choice (4 points)

For each of the following, circle all (zero or more) correct answer(s). You will be graded 0.25 points on each choice.

1. **(1.5 point)** 3-SAT is:

(a) P	(b) NP	(c) CoNP
(d) NP-Hard	(e) coNP-Hard	(f) NP-Complete
2. **(1.5 point)** Consider the problem of determining, for a given boolean formula, whether every assignment to the variables satisfies it. The problem is:

(a) EXP	(b) NP	(c) CoNP
(d) NP-Hard	(e) coNP-Hard	(f) CoNP-Complete
3. **(1 point)** Showing a polynomial time reduction from 3-SAT to problem X proves that X is:

(a) P	(b) NP	(c) NP-Complete	(d) NP-Hard
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4. **(1 point)** If X is NP-Complete, this implies that X is:

(a) NP	(b) EXP	(c) P	(d) NP-Hard
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ANSWER:

1. (b), (d), (f): 3-SAT is NP-complete. So it is also NP, NP-Hard.
2. (a), (c), (e), (f): Tautology is the complimentary problem to Satisfiability, which is NP-complete, so Tautology is CoNP-complete, which also implies it is CoNP, CoNP-hard. EXP encompasses all of coNP.
3. (d) This only shows NP-hardness, since X is not NP .
4. (a), (b), (d): Definition of NP-complete implies NP and $NP - Hard$. And $NP \subseteq EXP$.