## How to Change the World with Donald Knuth

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Information Security Project Presentation

1 Discrete Logarithm

2 ElGamal Cryptosystem

# Discrete Logarithm in a Nutshell

The security of many cryptographic techniques depends on the intractability of discrete logarithm problem.

A partial list of these include:

- DiffieHellman key agreement and its derivatives.
- ElGamal encryption.
- ElGamal signature scheme and its variants.

General setting for algorithms in this section are:

- A (multiplicatively written) finite cyclic group G
- *n* is the order of group *G*
- ullet  $\alpha$  is a generator of group  $G^1$

<sup>&</sup>lt;sup>1</sup>For more math background, refer to [Ros12].

## Relevant Definitions

Cyclic group and its generator.

#### Definition

A group is *cyclic* if there is an element  $\alpha \in G$  such that for each  $b \in G$  there is an integer i with  $b = \alpha^i$ . Such an element  $\alpha$  is called a generator of G.

Discrete logarithm.

#### Definition

Let G be a finite cyclic group of order n. Let  $\alpha$  be a generator of G, and let  $\beta \in G$ . The discrete logarithm of  $\beta$  to the base  $\alpha$ , denoted  $\log_{\alpha} \beta$ , is the unique integer x,  $0 \le x \le n-1$ , such that  $\beta = \alpha^x[\mathsf{MVO96}]$ .

# A Discrete Logarithm Example

### Example

Let p=97. Then  $\mathbb{Z}_{97}^*$  is a cyclic group of order n=96. A generator of  $\mathbb{Z}_{97}^*$  is  $\alpha=5$ . Since  $5^{32}\equiv 35\mod 97$ ,  $\log_5 35=32$  in  $\mathbb{Z}_{97}^*$ .

## The DiffieHellman Problem

The DiffieHellman problem is closely related to the well-studied discrete logarithm problem.

#### **Definition**

The DiffieHellman problem is the following: given a prime p, a generator  $\alpha$  of  $\mathbb{Z}_p^*$ , and elements  $\alpha^a \mod p$  and  $\alpha^b \mod p$ , find  $\alpha^{ab} \mod p$ .

Wait! Could we just possibly do

$$\alpha^{a} \times \alpha^{b} \to \alpha^{ab} \tag{1}$$

Well, life is not as easy as it looks like...

$$\alpha^{a} \times \alpha^{b} = \alpha^{a+b} \tag{2}$$

# Links between Discrete Logarithm and DiffieHellman Problem

**Suppose** that the discrete logarithm problem in  $\mathbb{Z}_p^*$  could be efficiently solved<sup>2</sup>. Then given  $\alpha$ , p,  $\alpha^a \mod p$  and  $\alpha^b \mod p$ , one could first find a from  $\alpha$ , p and  $\alpha^a \mod p$  by what?!

Solving a discrete logarithm problem, and then compute  $(\alpha^b)^a = \alpha^{ab} \mod p$ .

<sup>&</sup>lt;sup>2</sup>In math, the assumption is as important as,

# ElGamal public-key encryption

The ElGamal public-key encryption scheme can be viewed as DiffieHellman key agreement<sup>3</sup> in key transfer mode.

Its security is based on the intractability of the discrete logarithm problem (Section 1) and the DiffieHellman problem (Section 2).

<sup>&</sup>lt;sup>3</sup>Yet another fancy nickname for key exchange

**Ensure:** A public key and its corresponding private key is created for every entity.

Steps to generate key pairs are described as follows:

- Generate a prime p that is large enough and cannot be predicted, i.e. it should be generated randomly. Find a generator  $\alpha$  of the multiplicative group  $\mathbb{Z}_p^*$  of integers modulo p.
- ② Randomly select an integer a satisfying  $1 \le a \le p-2$ . Then calculate  $\alpha^a \mod p$ .
- **3** The public key is returned as  $(p, \alpha, \alpha^a)$ ; The private key is returned as a.

Figure: Algorithm Key generation for ElGamal public-key encryption

## References I



