Category	Instruction	Example	Meaning	Comments
Arithmetic	add	add \$s1,\$s2,\$s3	\$s1 = \$s2 + \$s3	Three operands; overflow detected
	subtract	sub \$s1,\$s2,\$s3	\$s1 = \$s2 - \$s3	Three operands; overflow detected
	add immediate	addi \$s1,\$s2,100	\$s1 = \$s2 + 100	+ constant; overflow detected
	add unsigned	addu \$s1,\$s2,\$s3	\$s1 = \$s2 + \$s3	Three operands; overflow undetected
	subtract unsigned	subu \$s1,\$s2,\$s3	\$s1 = \$s2 - \$s3	Three operands; overflow undetected
	add immediate unsigned	addiu \$s1,\$s2,100	\$s1 = \$s2 + 100	+ constant; overflow undetected
	move from coprocessor register	mfc0 \$s1,\$epc	\$s1 = \$epc	Copy Exception PC + special regs
	multiply	mult \$s2,\$s3	Hi, Lo = $$s2 \times $s3$	64-bit signed product in Hi, Lo
	multiply unsigned	multu \$s2.\$s3	Hi, Lo = $$s2 \times $s3$	64-bit unsigned product in Hi, Lo
	divide	div \$s2,\$s3	Lo = \$s2 / \$s3, Hi = \$s2 mod \$s3	Lo = quotient, Hi = remainder
	divide unsigned	divu \$s2,\$s3	Lo = \$s2 / \$s3, Hi = \$s2 mod \$s3	Unsigned quotient and remainder
	move from Hi	mfhi \$s1	\$s1 = Hi	Used to get copy of Hi
	move from Lo	mflo \$s1	\$s1 = Lo	Used to get copy of Lo
Data transfer	load word	lw \$s1,100(\$s2)	\$s1 = Memory[\$s2 + 100]	Word from memory to register
	store word	sw \$s1,100(\$s2)	Memory[\$s2 + 100] = \$s1	Word from register to memory
	load half unsigned	1hu \$s1,100(\$s2)	\$s1 = Memory[\$s2 + 100]	Halfword memory to register
	store half	sh \$s1,100(\$s2)	Memory[ $$s2 + 100$ ] = $$s1$	Halfword register to memory
	load byte unsigned	lbu \$s1,100(\$s2)	\$s1 = Memory[\$s2 + 100]	Byte from memory to register
	store byte	sb \$s1,100(\$s2)	Memory[\$s2 + 100] = \$s1	Byte from register to memory
	load upper immediate	lui \$s1,100	\$s1 = 100 * 2 <sup>16</sup>	Loads constant in upper 16 bits
Logical	and	and \$s1,\$s2,\$s3	\$s1 = \$s2 & \$s3	Three reg. operands; bit-by-bit AND
	or	or \$s1,\$s2,\$s3	\$s1 = \$s2   \$s3	Three reg. operands; bit-by-bit OR
	nor	nor \$s1,\$s2,\$s3	\$s1 = ~ (\$s2  \$s3)	Three reg. operands; bit-by-bit NOR
	and immediate	andi \$s1,\$s2,100	\$s1 = \$s2 & 100	Bit-by-bit AND with constant
	or immediate	ori \$s1,\$s2,100	\$s1 = \$s2   100	Bit-by-bit OR with constant
	shift left logical	sll \$s1,\$s2,10	\$s1 = \$s2 << 10	Shift left by constant
	shift right logical	srl \$s1,\$s2,10	\$s1 = \$s2 >> 10	Shift right by constant
Condi- tional branch	branch on equal	beq \$s1,\$s2,25	if (\$s1 == \$s2) go to PC + 4 + 100	Equal test; PC-relative branch
	branch on not equal	bne \$s1,\$s2,25	if (\$s1 != \$s2) go to PC + 4 + 100	Not equal test; PC-relative
	set on less than	slt \$s1,\$s2,\$s3	if $(\$s2 < \$s3) \$s1 = 1$ ; else $\$s1 = 0$	Compare less than; two's complement
	set less than immediate	slti \$s1,\$s2,100	if (\$s2 < 100) \$s1 = 1; else \$s1=0	Compare < constant; two's complement
	set less than unsigned	sltu \$s1,\$s2,\$s3	if (\$s2 < \$s3) \$s1 = 1; else \$s1=0	Compare less than; natural numbers
	set less than immediate unsigned	sltiu \$s1,\$s2,100	if $(\$s2 < 100)$ $\$s1 = 1$ ; else $\$s1 = 0$	Compare < constant; natural numbers
Uncondi- tional	jump	j 2500	go to 10000	Jump to target address
	jump register	jr \$ra	go to \$ra	For switch, procedure return
jump	jump and link	jal 2500	\$ra = PC + 4; go to 10000	For procedure call

**FIGURE 3.14 MIPS architecture revealed thus far.** The memory and registers of the MIPS architecture are not included for space reasons, but this section added the hi and lo registers to support multiply and divide. Color indicates the portions revealed since Figure 3.4 on page 175. MIPS machine language is listed in the MIPS summary reference card at the front of this book.