Lecture materials on system software development

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- System software
- Software interface
- Code structure
- Code compilation
- Error code
 - A role of the OS
 - OS booting process
 - Syscalls



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- File descriptor
- I/O streams of a process
 - Standard I/O streams
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 - Common access modes
 - Iseek(2)

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- close(2)
- dup(2) и dup2(2)
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- Useful functions
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Course structure VII

pmap(1) utility

Process

Process states

Useful links

```
http://src.illumos.org/
https://github.com/mit-pdos/xv6-public
```

```
https://se.ifmo.ru/~korg/
```

https://vk.com/korglings

Books:

- 1. Uresh Vahalia. UNIX Internals
- 2. A. S. Tanenbaum, A. S. Woodhull. Operating Systems: Design and Implementation



System software

- System software languages
- Syscalls

► I/O

► Threads and processes

УНИВЕРСИТЕТ ИТМО

Software interface

Each program receives arguments and environment variables

Error code is an integer describing the correctness of program termination



Code structure

```
int main(
    int argc,
    char *argv[],
    char *envp[]
) {
    /* ... */
    return 0;
}
```

Code structure

```
#include <stdlib.h>
int main(
   int argc,
   char *argv[],
   char *envp[]
   /* ... */
   return EXIT_SUCCES;
```

Code compilation

```
# gcc -c program.c
# gcc -o program program.c
# gcc -o program program.c
# cc -o program program.c
```



Error code

```
# rm -f /etc/passwd 2<&-
# echo $?
1
# echo Hello, world!
Hello, world!
# echo $?
0</pre>
```



A popular mistake

Using void main() is inappropriate!

```
# cat void.c
void main(void) {}
# ./void
# echo $?
16
```



A role of the OS

- Multitasking;
- Memory virtualization;
- Device management;
- Interrupt handling;
- Extending the available set of application-level operations.

OS booting process

- ► Reset vector: UEFI, BIOS, ...
- ► I/O, *PIC(IRQ), VGA
- ► POST + PCI BIOS
- Boot device detection
- Bootloader stage0 (boot sector)
- Bootloader stage1
- OS kernel

Syscalls

- Kernel functions calling
- Using hardware via the common API
- Have libc interfaces
- Have kernel privileges



Calling the OS

```
_exit(2);
.globl _start
_start:
pushq $2
movq $1, %rax
int $0x80
```

/* program termination with errcode 2 */

```
# cc -m64 -Wall -Wextra -Wno-comment \
-nostdlib -o main main.S
```

https://pastebin.com/knTdpZRe



File

What is a file?

Everything is a file!

Apart from threads and the kernel



File descriptor

```
http:
//src.illumos.org/source/xref/illumos-gate/
usr/src/uts/common/syscall/open.c#54
```

```
http://src.illumos.org/source/xref/
illumos-gate/usr/src/uts/common/fs/vnode.c#940
```

A file descriptor number is a positive integer that abstracts processes from files they are using.



I/O streams of a process

Number	File	Flags
0	/dev/tty	O_RDWR
		O_LARGEFILE
1	/dev/tty	O_RDWR
		O_LARGEFILE
2	/dev/tty	O_RDWR
		O_LARGEFILE
3	/etc/passwd	O_RDONLY
4	/dev/mtdblock3	O_RDWR
•••	•••	•••
255	•••	•••



Standard I/O streams

```
# grep FILENO /usr/include/unistd.h
#define STDIN_FILENO 0
#define STDOUT_FILENO 1
#define STDERR_FILENO 2
```



open(2)

```
int open(
   const char *path, /* file path */
   int oflag, /* access mode */
   /* mode_t mode */ /* access rights */
);
```

Returns a file descriptor number or an error code



Common access modes

O_RDONLY - Read-only

O_WRONLY – Write-only

O_RDWR – Read-write

O_CREAT – Create if not exists

O_APPEND – Append to the end of the file

O_TRUNC - Write from the beginning of the file

O_LARGEFILE - Long file position

O_EXCL - Long file position

Iseek(2)

```
off_t lseek(
   int fildes, /* номер открытого файла */
   off_t offset, /* смещение позиции */
   int whence /* действие */
);
```

Returns an updated offset in bytes or an error code



read(2)

```
ssize_t read(
   int fildes, /* номер открытого файла */
   void *buf, /* буфер чтения */
   size_t nbyte /* количество байт */
);
```

Returns the amount of bytes read successfully or an error code



write(2)

```
ssize_t write(
   int fildes, /* номер дескриптора */
   const void *buf, /* буфер записи */
   size_t nbyte /* количество байт */
);
```

Returns the amount of bytes written successfully or an error code



close(2)

```
int close(
int fildes, /* номер дескриптора */
);
```

Returns zero or an error code

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dup(2) и dup2(2)

```
int dup(
   int fildes /* номер открытого файла */
);
int dup2(int fildes, int fildes2);
```

Returns a number of a new file descriptor or an error code



stat(2)

```
int stat(
    const char *restrict path,
    /* путь к файлу*/
    struct stat *restrict buf
    /* результат */
);
```

Returns zero or an error code



Errors in syscalls

A return code of the syscall:

- below zero an error occured while syscall processing
- equals to zero successful execution of the syscall
- above zero the result of the successful exection



Error code standardization

- ► Error code unification
- errno variable
- perror(3) function
- strerror(3) function



Error example

```
if (read(7, buf, 1) < 0) {
   fprintf(stderr, "%d_", errno);
   perror("read");
   _exit(1);
}
/* 9 read: Bad file number */</pre>
```

Headers

- unistd.h UNIX declarations
- stdio.h standard input/output
- fcntl.h file operations
- sys/types.h system types
- sys/stat.h system statuses



Read/write example

```
#include <unistd.h>
int main(int argc, char *argv[]) {
   int bytes;
  char buf[256];
  while((bytes = read(STDIN_FILENO, buf,
  \rightarrow sizeof(buf))) > 0) {
      if (write(STDERR FILENO, buf, bytes)
  return 1;
   return bytes;
```

Makefile

```
PROJS=main
CC=gcc
CFLAGS=-m64

all: $(PROJS)
    @echo Done!

$(PROJS):
    $(CC) $(CFLAGS) -o $@ $(@:=.c)
```



make utility

```
# make
gcc -m64 -o main main.c
Done!
# ./main
Hello, world!
```



Useful functions

- ▶ isatty(3C)
- gethostbyname(3NSL) gethostbyaddr(3NSL)
- htons(3SOCKET) htonl(3SOCKET) ntohs(3SOCKET) ntohl(3SOCKET)
- usleep(3C)



I/O with offsets

Returns the amount of bytes or an error code



Useful functions

Returns the amount of bytes or an error code



iovec structure (I/O vector)

```
#include <sys/uio.h>
typedef struct iovec {
void *iov_base;
/* start address */
size_t iov_len;
/* segment length */
} iovec_t;
```



Cache flushing

```
void sync(
void /* не принимает аргументов */
);
```

Return code is meaningless



Working with file descriptors

Return code meaning depends on a particular command



fcntl(2) commands

F_DUPFD / F_DUP2FD F_FREESP F_GETFD / F_SETFD F_GETFL / F_SETFL similar to dup/dup2 free up some space close on exec flag access flags

F_GETLK / F_SETLK F_GETLKW / F_SETLKW file locking

F_RDLCK / F_WRLCK / F_UNLCK



Access rights check

```
int access(const char *path, int amode);
```

R_OK – read

W_OK - write

X_OK - execute

F_OK - existence

Returns zero or an error code



Access rights modification

```
int chmod(const char *path, mode_t mode);
int fchmod(int fildes, mode_t mode);
S_ISUID 04000
S IRWXU 00700
```

Returns zero or an error code

(S ISUID | S URWXU) 04700



Changing file owner

```
int chown(
   const char *path,
   uid_t owner,
   gid_t group
int fchown(
   int fildes,
   uid t owner,
   gid_t group
```

Returns zero or an error code

File creation mask

```
mode_t umask(
 mode_t cmask /* значение маски */
);
```

Returns the previous mask value

How should you get the current value?



Working with links

```
int link(
   const char *existing, /* путь к файлу*/
   const char *new /* путь к ссылке */
);
int unlink(const char *path);
```

Returns zero or an error code



Symbolic links

```
int symlink(
  const char *name1,
  const char *name2
ssize t readlink(
  const char *restrict path, /* ссылка*/
  char *restrict buf, /* буфер */
  size t bufsiz /* размер буфера */
```



Working with directories

```
int mkdir(
   const char *path, /* путь к каталогу */
   mode_t mode /* режим доступа */
);
int rmdir(const char *path);
```

Returns zero or an error code



Current working directory

```
int chdir(const char *path);
int fchdir(int fildes);

getcwd(3) returns either a pointer to a buffer
or a -1; has the following prototype
char *getcwd(char *buf, size_t size);
```

Reading directory

```
DIR *opendir(const char *dirname);
struct dirent *readdir(DIR *dirp);
void rewinddir(DIR *dirp);
int closedir(DIR *dirp);
```



dirent structure



Device files

```
int mknod(
   const char *path, /* путь к файлу */
   mode_t mode, /* режим доступа и тип */
   dev_t dev /* устройство */
);
```

Returns zero or an error code



Memory model of a process

	. OxFFFFFFF
Kernel	0xC0000000
Stack	0.0000000
Heap	
Data	
Code	0x00000000
	57.5555555



Memory allocation

Data segment extension:

```
int brk(void *endds);
void *sbrk(intptr_t incr);
```

New segment allocation from an Anonymous Memory:

```
void *mmap(
   void *addr,
   size_t len,
   int prot,
   int flags,
   int fildes,
   off_t off
);
```

pmap(1) utility

```
helios$ pmap $$
                           [ stack ]
08043000
           20K
                 rw---
           552K
08050000
                           /usr/hin/hash
                 r-x--
080E9000
           76K
                           /usr/bin/bash
                 rwx--
080FC000
           300K
                           [ heap ]
                 rwx--
FEB20000
           64K
                 rwx--
                           [ anon ]
                           /lib/module.so
FEB40000
           56K
                 r-x--
```



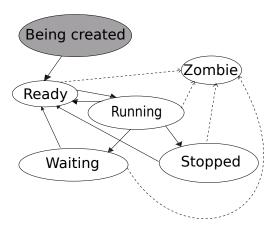
Process

A process is an aggregation of a program and a metadata describing the program's runtime ©KorG

Run in parallel; technically independent from each other



Process states





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