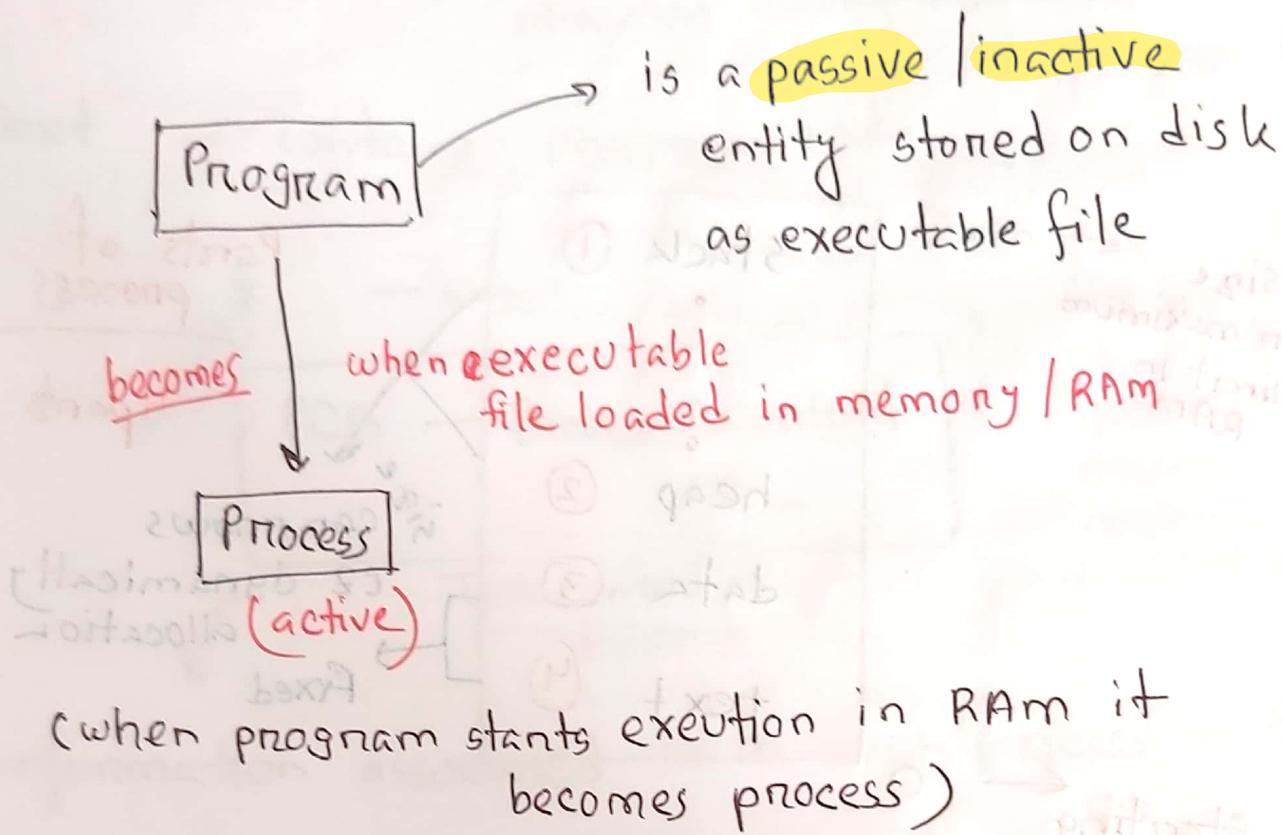


## Chapter-3

### Process



(one program can have several process)

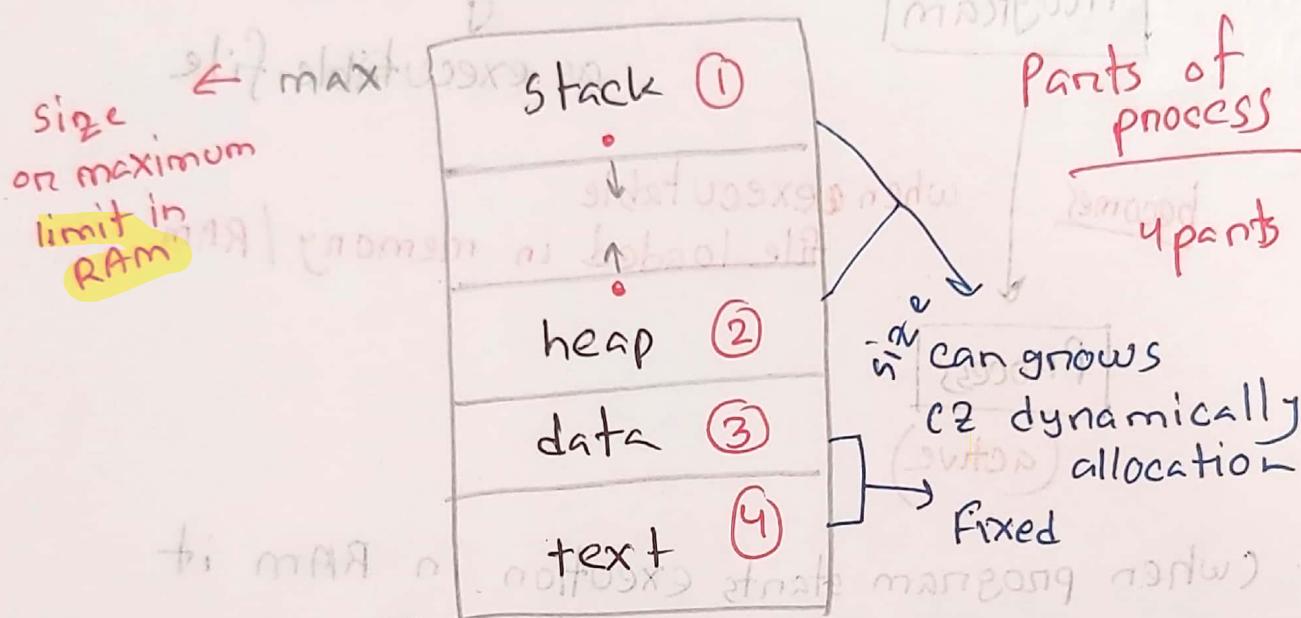
\* To start a program you must bring it to ram

\* program consists some instruction

## Process

## Components

→ Process is a program in execution



starting address of program

(How does 1 process look like in memory / RAM)

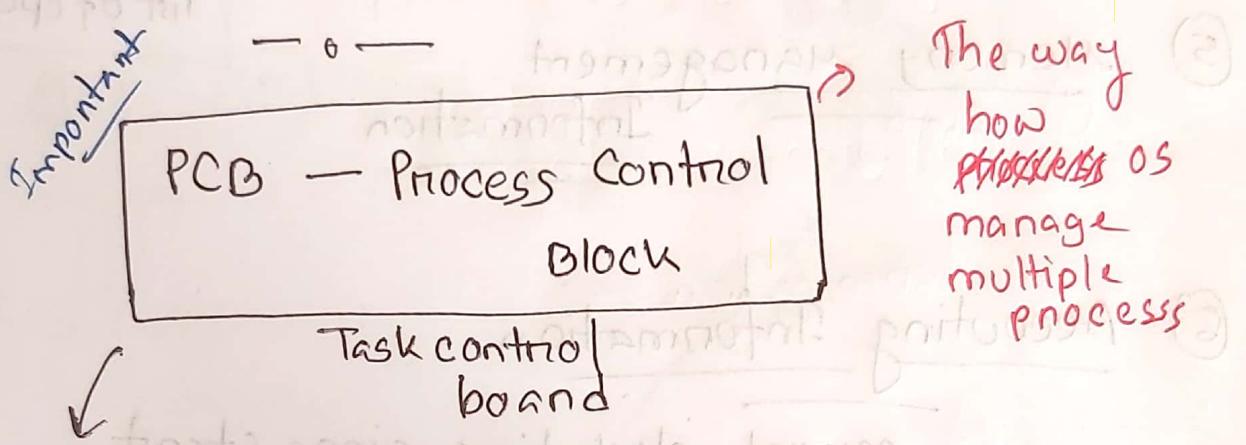
① stack → contains program code

→ contains temporary data

→ function parameters,  
return addresses  
Local variables

↓  
inside in function

- ② Data Section → contain global variables
- ③ Heap → contain memory dynamically allocated during run time
- ④ Text → contains program code



information associated with each process in the

### ① Process state:

- i - (new)
- ii - (running)
- iii - (waiting)
- iv - (ready)
- v - (terminated)

### ② Program counter : Location of instruction to next execute the (where starts of execution)

③

CPU Registers :- content of all process centric registers

Process state
Process number
Program counter
Registers
memory limits
list of open file

④

CPU scheduling Information

Priorities, scheduling pointers

⑤

Memory management

Information

⑥

Accounting Information

cpu used, clock time since start

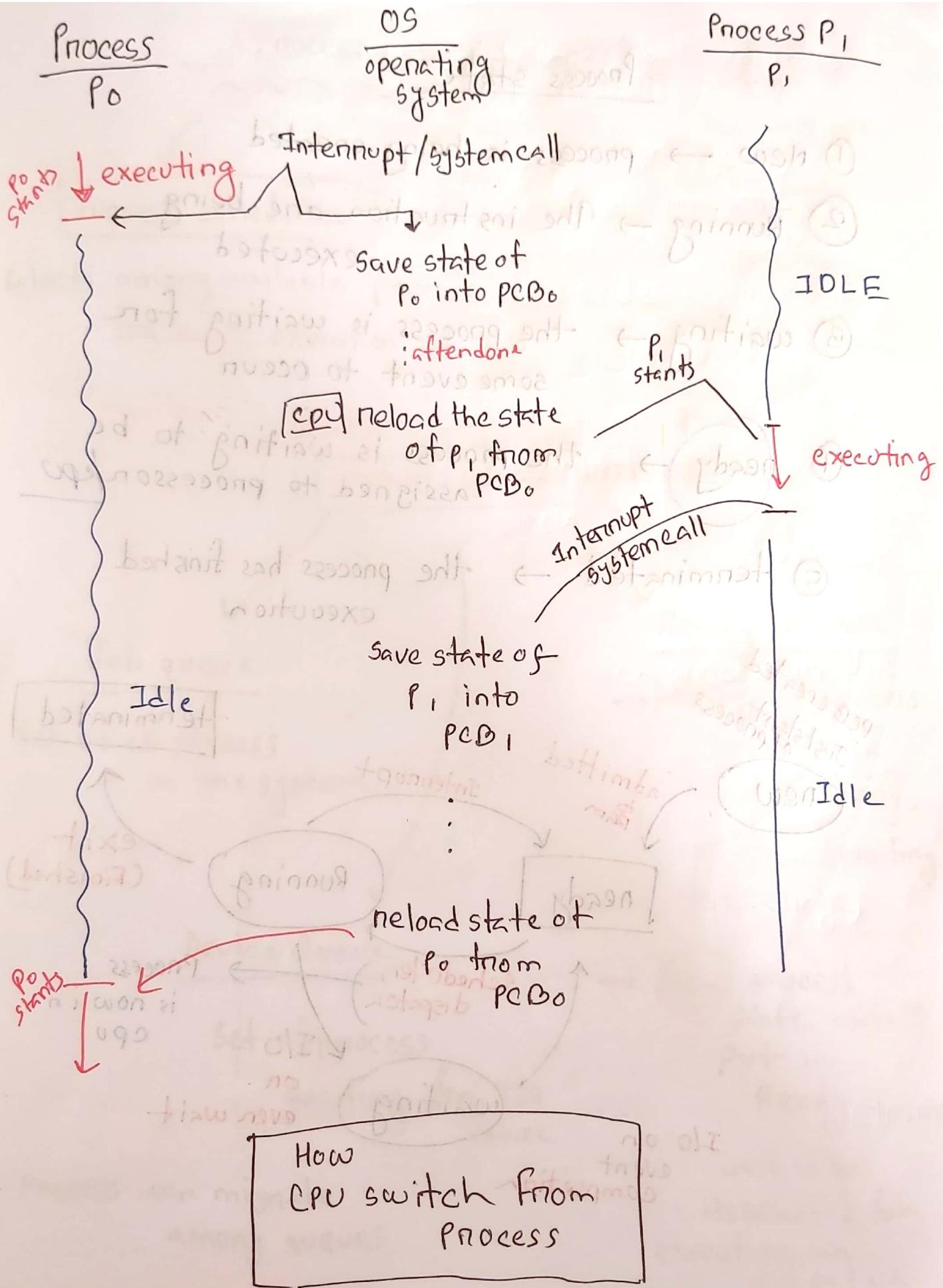
⑦

I/O status Information

I/O device allocated to process,  
list of open files

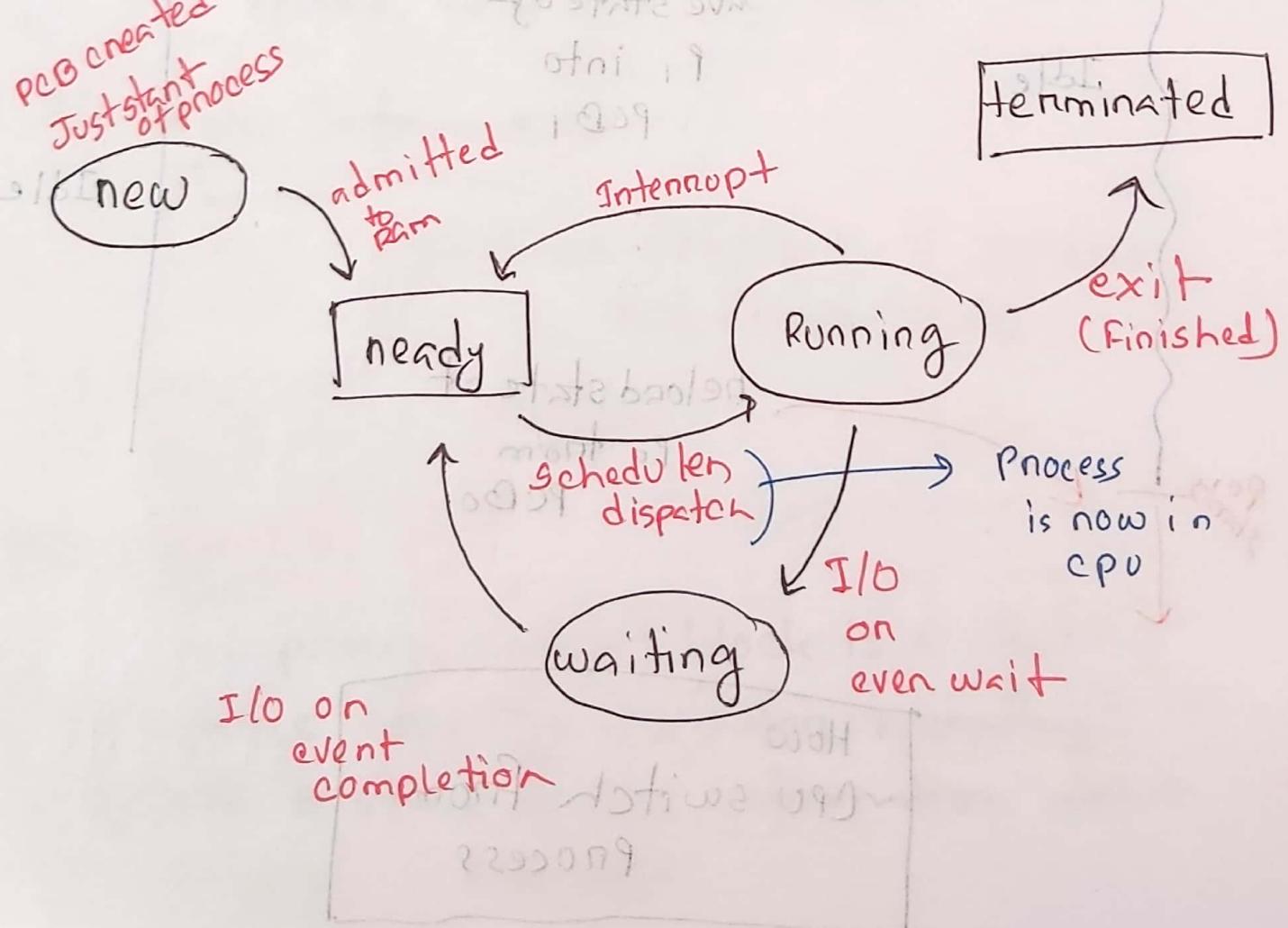
PCB definition

a process control block is a data structure used by computer operating system to store all the information about a process



## Process state

- ① New → process is being created
- ② Running → The instruction are being executed
- ③ waiting → the process is waiting for some event to occur
- ④ ready → the process is waiting to be assigned to processor/cpu
- ⑤ terminated → the process has finished execution



# Process scheduling

## Process scheduler

selects among available processes  
for next execution on CPU

maintains

scheduling

queues of processes

## Job queue

Set of all processes  
in the system

## Device Queue

Set of processes

waiting for I/O  
device

maximize CPU use  
quickly switch  
CPU for time sharing

## Ready Queue

→ Set of all processes  
residing in main memory,  
ready and waiting to execute

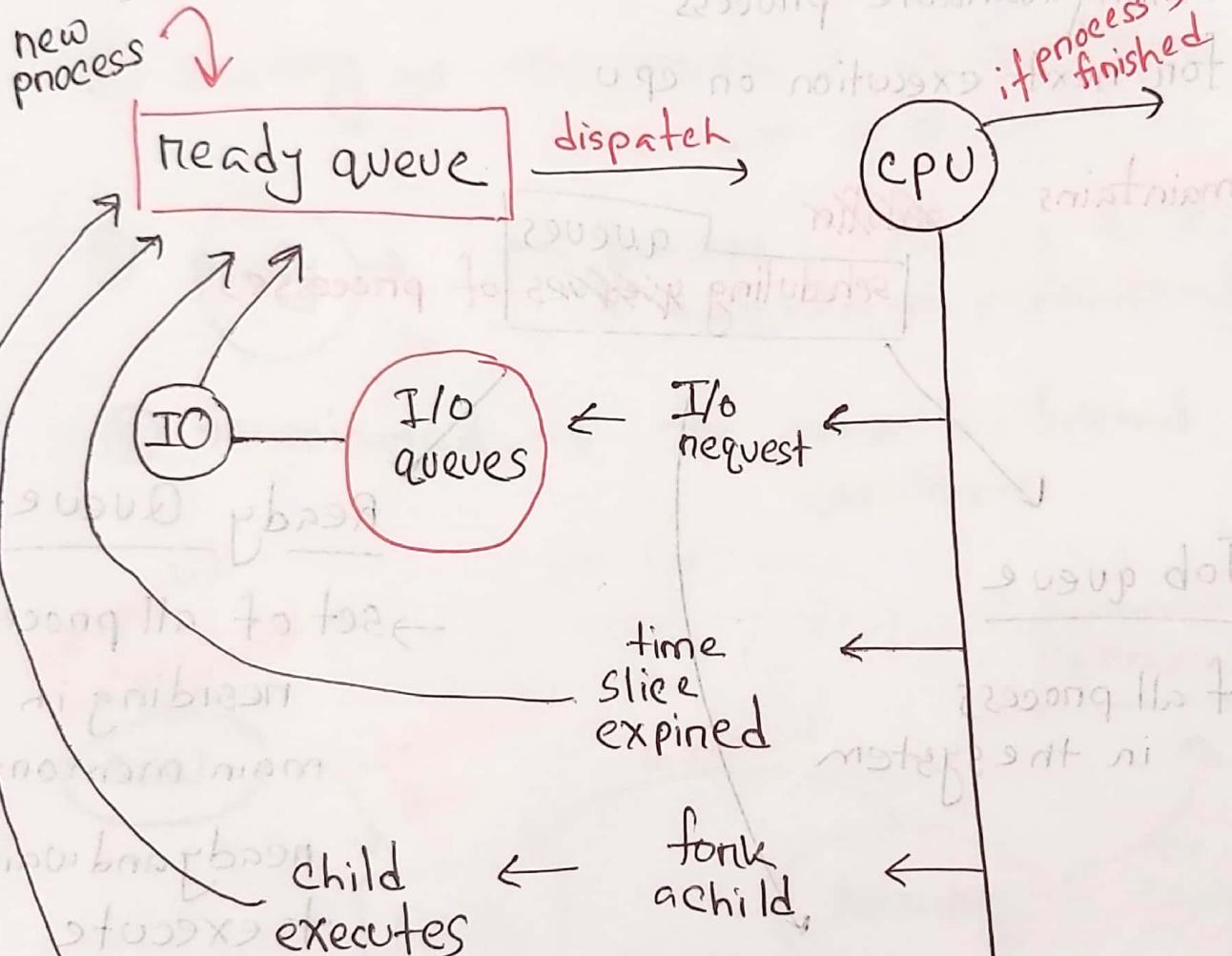
→ New process state initially put in Ready Queue

wait until its selected for execution on dispatch

\* Process can migrate  
among queues

## Representation of Process scheduling

### Queueing diagram



# Schedulers

## Short term scheduler

- CPU scheduler
- which sends a process to CPU
- select which process should be executed next and allocated to CPU
- must be very fast (is invoked / called frequently) (milliseconds)

## Long term scheduler

- Job scheduler
- selects which processes should be brought into the ready queue in main memory
- strives for good process mix
- may be slow (is invoked / called infrequently (seconds, minutes))

→ controls the degree of multiprogramming

definition ↓

how many programs should allow at same time on admitted to main memory

## Types of process

### ① I/O bound process :-

- spend more time doing I/O than computations
- many short CPU bursts

### ② CPU-bound process

- spend more time doing computations
- few very long CPU bursts

## Medium-term scheduler :

- remove process from memory
- store <sup>in</sup> disk
- bring back in RAM to continue execution

## Context Switch

*(definition important)*

① the system must save the state of old process

and load the saved state for new process via context switch

when a cpu switches

② to another process

it is called context switch

(context of process represented in PCB)

while context switch system does not do useful work

definition :- ① + ②

## Process Creation

Parent process create child process which in turn create other process , forming a tree of processes

① Chrome Browser is a multiprocess with 3 different type of process

④ Browser process : manage user interface, disk, network I/O

② Render process : render web pages, deals with HTML and javascript

③ Plug in process : for each type of plugin

⑤ + ① - ; nothing

— o —

Based no Interprocess communication

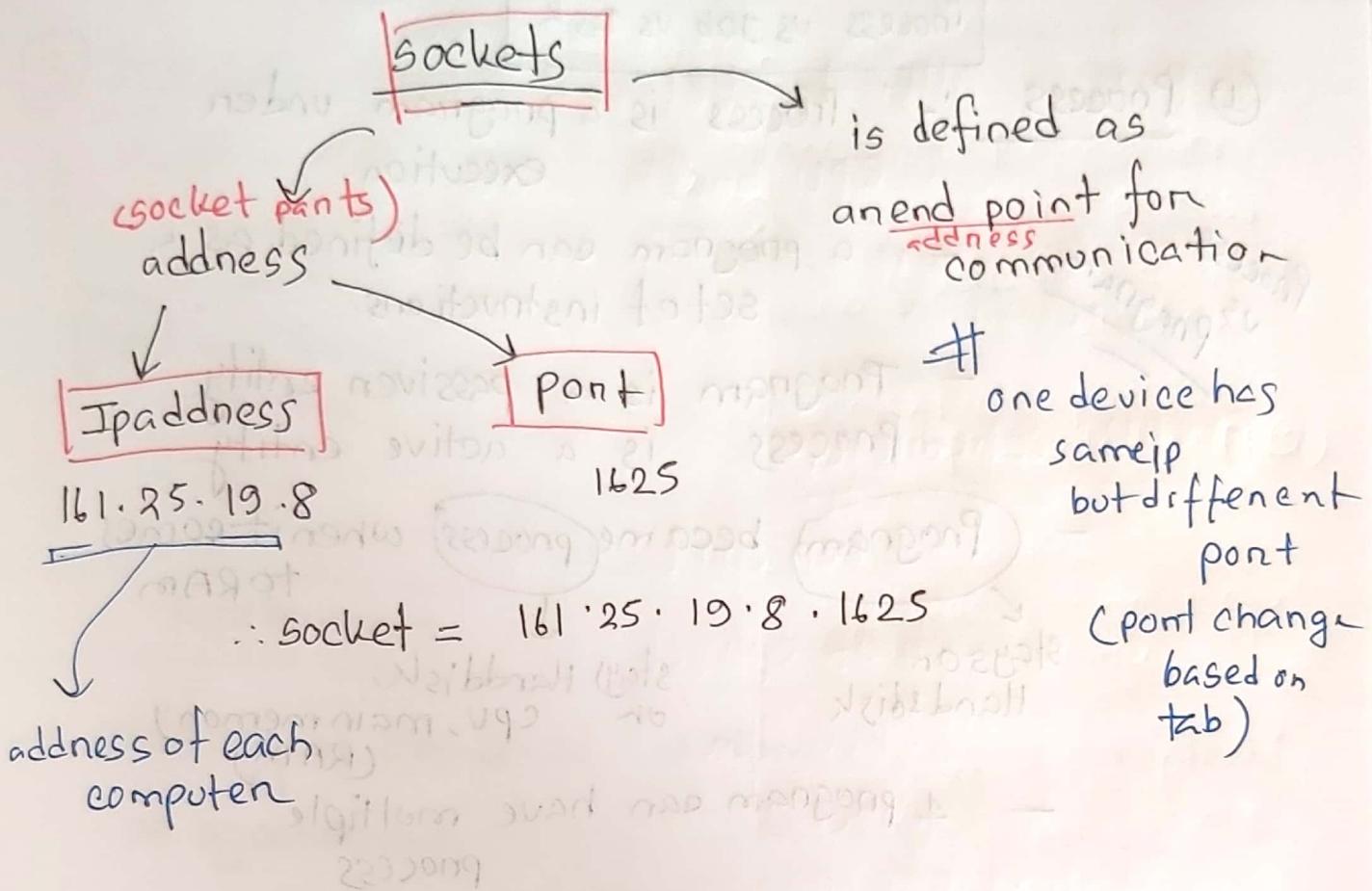
Process type

① Independent process

(not affected by other process)

② Cooperating process

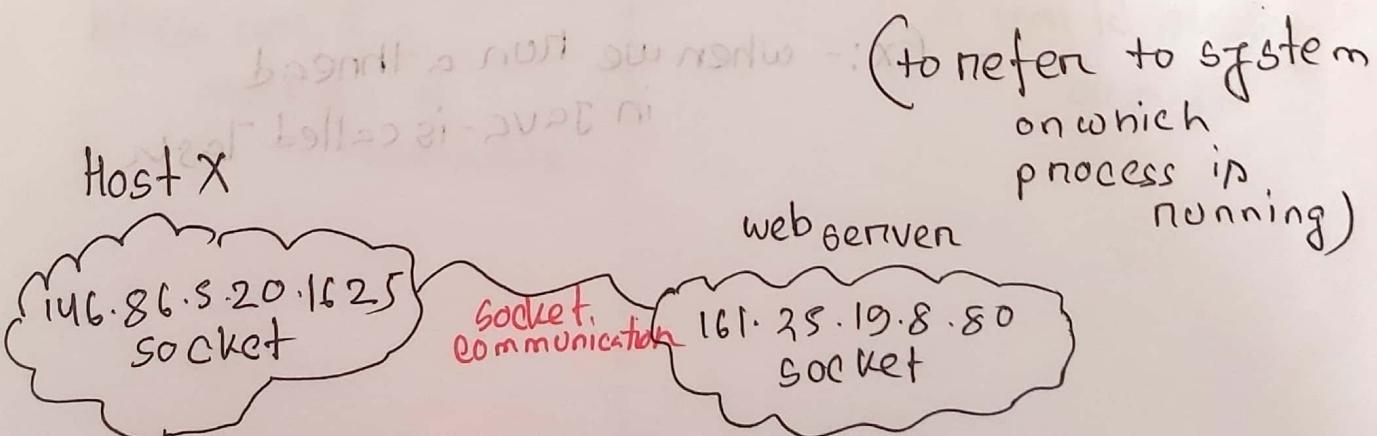
(affected by another process)



# communication  
consists between a pair of sockets

# port < 1024 → well known ports  
(used for standard  
Internet service)

# special IP address: 127.0.0.1



## Process vs Job vs Task

### ① Process

Process is a program under execution.

— a program can be defined as a set of instructions.

— Program is a passive entity  
Process is an active entity

- Program became process when it comes to RAM
- 1 program can have multiple process

stays on  
Harddisk

stays Harddisk  
or CPU, main memory  
(RAM)

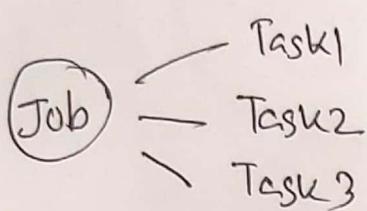
② Task :- Task is a unit of work being executed

Task in OS may be synonymous with processes

→ Task is a sub part of job

Ex:- when we run a thread in JAVA - is called Task

③ Job :- → a job is a complete unit of work under execution



- a job consists many task which in turn consists of many process

Synchronous

Asynchronous

→ Blocking is considered as synchronous

→ Blocking send has the sender block until the message is received

→ Blocking receive has the receiver block until a message is available

→ Nonblocking is considered as asynchronous

→ Non blocking send has the sender send the message and continue

→ Non blocking receive has the receiver receive a valid message or null

④ Process Altenate between  
CPU bursts and I/O bursts

⑤ CPU scheduler:-

when CPU becomes idle, it is the job of CPU scheduler / short term

Scheduler to select another process from the ready queue to run next

⑥ Scheduler dispatcher :- is a module that gives the control of CPU to the process selected by the scheduler

## Types of sockets

actually  
types

① Datagram socket  
(connectionless)

② Stream socket  
(connection-oriented)

③ Raw socket

④ Sequenced packet socket

access → Prot col

UDP  
(User Datagram Protocol)

TCP

Transfer control protocol  
(ICMP)

Internet control message protocol  
(ICMP)

(IDP)

Internet datagram protocol

## Degree of multiprogramming

Degree of multiprogramming describes the maximum number of processes that a single processor system

can accommodate efficiently.

## Process

## Thread

- |   |                                       |
|---|---------------------------------------|
| ① Process means any program is in execution | ① Thread means a segment of a process |
| ② Process takes more time to terminate      | ② Thread takes less time to terminate |
| ③ Takes more time for creation              | ③ Less time for creation              |
| ④ More time in context switching            | ④ Less time for context switching     |
| ⑤ Less efficient in communication           | ⑤ More efficient in communication     |
| ⑥ A system call is involved in it           | ⑥ No system call is involved          |
| ⑦ Process has Process control block         | ⑦ Thread has Thread control block     |

## Process scheduling

is the activity of process manager that handle the removal of the running processes from CPU and the selection of another process on the basis of a particular strategy.

## Loop back address

A loop back address is distinct reserved IP address 127.0.0.1 to refer to system on which process is running.

what is API?

## Application program Interface (API)

is code that allows two software programs to communicate with each other.

## The principle of least privilege

the principle of least privilege states that a subject should be given only those privileges needed for it to complete its task.

- In UNIX, user is domain. True or False?

- true

## what is kernel

→ is a central component of OS that manages operation of computer and hardware  
→ acts as a bridge between applications and data processing performed at hardware level using IPC and system calls

## Objectives

- To control disk, memory, task management
- To decide state of incoming process
- To establish communication between user level application and hardware

① what is system call ? explain different type of system call.

→ **System call** is a way for programs to interact with operating system.

5 types system call

② Process call :-

- deals with process such as creation & termination

Example:-

(Window)

CreateProcess()

Exit Process()

(Linux)

fork()

exit()

wait()

( )

② File management :-

- responsible for file manipulation  
creating, reading, writing into file

Example:-

(x00)

( )

(window)

• CreateFile()

ReadFile()

WriteFile()

CloseHandle()

(zwlobn)

• open()

read()

write()

close()

(Linux)

• open()

read()

write()

close()

### ③ Device management :-

- responsible for device manipulation
- reading / writing device buffer

Example :-

(Windows) ~~int main()~~ (Linux)

setconsolemode() ioctl()

ReadConsole() read()

### ④ Information Maintenance

- handle information and its transfer

(Xenix) between OS and user programs

Example :- (Windows) (Linux)  
(Driver) (Driver)  
(Handle) (Handle)  
SetTimer() alarm()  
Sleep() sleep()

### ⑤ Communication

-; transmission of?

→ useful for inter processes communication  
↓ it is achieved by using pipes

Example :- (Xenix)

(msgq) &  
(User) &  
(Semaphore)  
(Shared)

(Windows)

CreatePipe()

MapViewOfFile()

(Linux)

pipe()

mmap()

⑧ How does communication take place in client server system? Explain

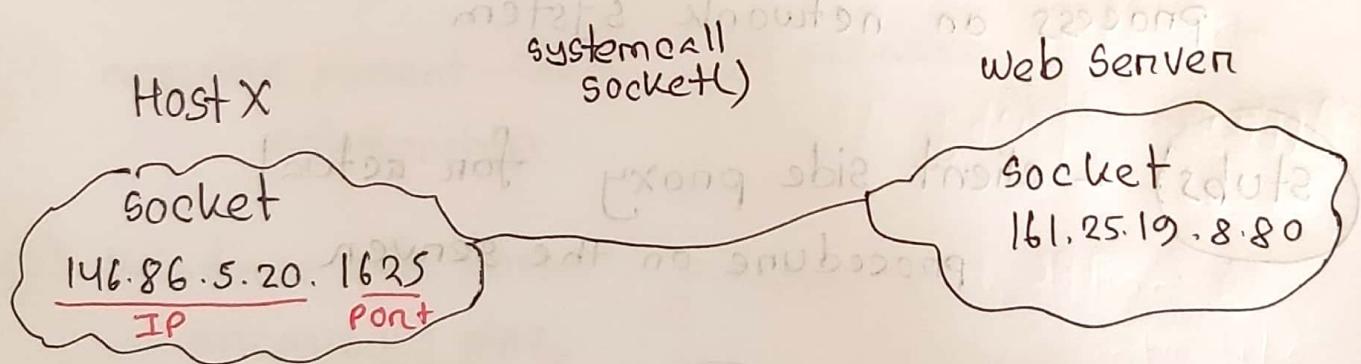
→ Client-server communication has two components client and server. The client sends request to the server and the server responds to the client request.

Three methods of client-server communication

① **socket** → low level communication channel

→ a socket is an endpoint for communication

→ a socket is identified by an IP address and a port number



Two major form

① Connection Oriented / TCP (Transmission Control Protocol)

→ guaranteed that all packets will arrive in good condition at other end and in order received

② Connectionless (UDP - user data protocol)  
→ no guarantee that packets will delivered  
with undamaged and will receive in  
particular order  
→ UDP ~~for~~ Transmission > TCP

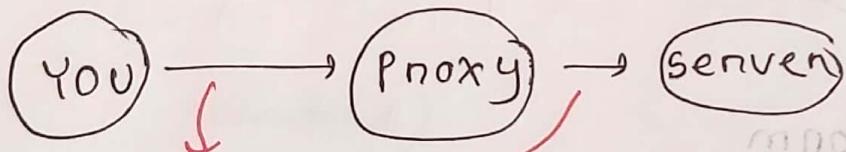
②

## RPC - Remote Procedure Calls

- \* in socket we communicate with process
  - \* In RPC "functions / procedure"
- RPC abstracts procedure calls between processes on network system

Stubs

client side proxy for actual procedure on the server



### ③ Pipe

- acts as a conduit between two process  
to communicate

#### Ordinary Pipe

- ① ordinary pipe allows communication in standard producer consumer style

- ② unidirectional

- ③ produces write to one end and consumer needs at other end

- ④ requires parent-child relationship

- ⑤ windows calls these anonymous pipe

#### Named pipe

- ① more powerful than ordinary pipe

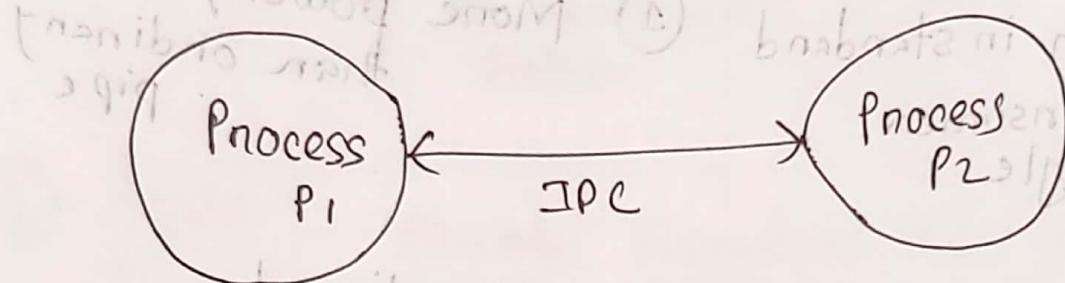
- ② bidirectional

- ④ no need parent child relationship

# IPC - Interprocess Communication

— is the mechanism provided by

OS that allows processes to communicate with each other



two types of process

- ① Independent process
- ② Co-operating process

Applic.

why IPC?

IPC helps achieving

- ① Computational speed up
- ② Modularity
- ③ Information & data sharing
- ④ Privilege separation
- ⑤ Process can communicate with each other and synchronize their action

# Approaches for Interprocess Communication

