# **Linear Modules**

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$$MAN = \begin{bmatrix} 23 \\ 23 \\ 23 \\ 46 \end{bmatrix}$$

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## 1 Basic Categorical Module theory

### 1.1 Category of Modules

```
\textbf{LeftModule} \, :: \, \prod R \in \mathsf{RING} \, . \, \, \sum M \in \mathsf{Set} \, . \, (M \to M \to M) \times (R \to M \to M)
M, +, \cdot : \texttt{LeftModule} \iff (M, +) \in \mathsf{ABEL} \&
    & \forall a, b \in M : \forall \alpha \in M : \alpha(a+b) = \alpha a + \alpha b \&
    & \forall a \in M : \forall \alpha, \beta \in M : (\alpha + \beta)a = \alpha a + \beta b \&
    & \forall a \in M : \forall \alpha, \beta \in M : \beta(\alpha a) = (\beta \alpha)a \&
    & \forall a \in M . 1a = a
\texttt{RightModule} \, :: \, \prod R \in \mathsf{RING} \, . \, \sum M \in \mathsf{SET} \, . \, (M \to M \to M) \times (M \to R \to M)
(M,+,\cdot): \texttt{RightModule} \iff \left(M,+, \texttt{swap}(\cdot)\right): \texttt{LeftModule}(R^{\mathrm{op}})
Assume R: Ring,
moduleAsGroup :: LeftModule(R) \rightarrow ABEL
moduleAsSet((M, +, \cdot)) = implicit := (M, +)
\texttt{scalarMult} \ :: \ \prod M : \texttt{LeftModule}(R) \ . \ M \to M \to M
scalarMult(M, +, \cdot) = \cdot_M := \cdot
ZeroMultInModule :: \forall M : LeftModule(R) . \forall a \in M . 0a = 0
Proof =
(1) := \eth LeftModule(R)(M)\eth Neutral(R, +)(0)(1)\eth LeftModule(R)(M) :
    a = 1a = (1+0)a = 1a + 0a = a + 0a
(2) := \eth Neutaral(M, +)(0)(a) : a + 0 = 0,
(*) := TotalGroupMult(1)(2) : 0a = 0;
NegativeMultInModule :: \forall M : LeftModule(R) . \forall a \in M . (-1)a = -a
Proof =
(2) := \eth inverse(a) : 0 = a + (-a),
(*) := TotalGroupMult(1)(2) : (-1)a = -a;
\texttt{LeftLinear} \, :: \, \prod A, B : \texttt{LeftModule}(R) \, . \, ?(A \xrightarrow{\texttt{ABEL}} B)
T: \texttt{LeftLinear} \iff \forall a \in A . \forall \theta \in R . T(\theta a) = \theta T(a)
\texttt{RightLinear} \, :: \, \prod A, B : \texttt{LeftModule}(R) \, . \, ?(A \xrightarrow{\texttt{ABEL}} B)
T: \texttt{RightLinear} \iff \forall a \in A \ . \ \forall \theta \in R \ . \ T(a\theta) = T(a\theta) = T(a)\theta
```

```
IdIsiLeftLinear :: \forall A : LeftModule(R) . id<sub>A</sub> : LeftLinear(A, A)
Proof =
. . .
IdlsRightLinear :: \forall A : RightModule(R) . id_A : RightLinear(A, A)
Proof =
. . .
LeftLinearComp :: \forall A, B, C : LeftModule(R) . \forall T : LeftLinear(A, B) . \forall S : LeftLinear(B, C) .
   . S \circ T : LeftLinear(A, C)
Proof =
Assume a:A.
Assume \theta: R,
a := \eth \texttt{compose}(S, T) \eth \texttt{LeftLinearComp}(T) \eth \texttt{LeftLinear}(S) \eth^{-1} \texttt{compose}(S, T) :
   : S \circ T(\theta a) = S(T(\theta a)) = S(\theta T(a)) = \theta S(T(a)) = \theta S \circ T(a);
\rightsquigarrow (*) := \eth^{-1}RightLinear : S \circ T : RightLinear(A, C);
. S \circ T : \mathtt{RightLinear}(A, C)
Proof =
. . .
LeftModuleCategory :: RING \rightarrow CAT
\texttt{LeftModuleCategory}\left(R\right) = R\text{-}\mathsf{MOD} := \Big(\texttt{LeftModule}(R), \texttt{LeftLinear}, \mathsf{id}, \circ\Big)
RightModuleCategory :: RING \rightarrow CAT
\texttt{RightModuleCategory}\left(R\right) = \mathsf{MOD}\text{-}R := \Big(\mathtt{RightModule}(R), \mathtt{RightLinear}, \mathrm{id}, \circ \Big)
ABELIsZMOD :: ABEL \cong_{CAT} MOD-\mathbb{Z}
Proof =
. . .
{\tt CommutativeLeftModuleIsRight} \ :: \ \forall R \in {\sf ANN} \ . \ R{\textrm{-}MOD} \cong_{{\sf CAT}} {\sf MOD}{\textrm{-}}R
Proof =
. . .
```

```
Proof =
Assume T, S: A \xrightarrow{R-\mathsf{MOD}} B,
Assume a, a': A,
():=\eth \mathtt{mapOp}(B,+)(T,S)\eth(\mathtt{-MOD}R)(T,S)\eth^{-1}\mathtt{mapOp}(B,+)(T,S):
    : (T+S)(a+a') = T(a+a') + S(a+a') = T(a) + T(a') + S(a) + S(a') = (T+S)(a) + (T+S)(a');
\leadsto (1) := \eth \mathsf{ABEL} : \left\lceil T + S : A \xrightarrow{\mathsf{ABEL}} B \right\rceil,
Assume a:A.
Assume \theta:R,
():=\eth mapOp(B,+)(T,S):
    : (T+S)(\theta a) = T(\theta a) + S(\theta a) = \theta T(a) + \theta S(a) = \theta \left(T(a) + S(a)\right) = \theta (T+S)(a);
\leadsto \big(\big) := \eth R\text{-MOD}(1) : \Big\lceil T + S : A \xrightarrow{R\text{-MOD}} B \Big\rceil;
\rightsquigarrow (1) := I(\forall) : \forall T, S \in (-\mathsf{MOD}R) . T + S \in (-\mathsf{MOD}R),
Assume T: (-\mathsf{MOD}R),
Assume \theta:R,
Assume a, a' : A,
() := \eth \mathsf{mapOp}(B, \cdot)(\theta, T) \eth (-\mathsf{MOD}R)(T) \eth^{-1} \mathsf{mapOp}(B, \cdot)(\theta, T) :
    : (\theta T)(a+a') = \theta \Big( T(a+a') \Big) = \theta \Big( T(a) + T(a') = \theta T(a) + \theta T(a');
\leadsto (2) := \eth \mathsf{ABEL} : \left\lceil \theta T : A \xrightarrow{\mathsf{ABEL}} B \right\rceil;
Assume a:A,
Assume \rho: R,
() := \eth \mathsf{mapOp}(B, \cdot) \eth (-\mathsf{MOD}T) \eth \mathsf{ANN}(R) \eth^{-1} \mathsf{mapOp}(B, \cdot) (\theta, T) :
    : (\theta T)(\rho a) = \theta (T(\rho a)) = \theta \rho T(a) = \rho \theta T(a);
\rightsquigarrow () := \eth R-MOD(2) : \left[\theta T : A \xrightarrow{R\text{-MOD}} B\right];
\rightsquigarrow (2) := I(\forall) : \forall T : A \xrightarrow{R\text{-MOD}} B : \forall \theta \in R : \theta T : A \xrightarrow{R\text{-MOD}} B,
(*):=\eth(\operatorname{-MOD}{R})(1)(2):\operatorname{This};
Submodule :: \prod M \in R-MOD . ??M
S: \mathtt{Submodule} \iff S \subset_{R-\mathsf{MOD}} M \iff (S, +_M, \cdot_M) \in R-\mathsf{MOD}
Proof =
(1) := GroupImage : Im T \subset_{GRP} B,
Assume y : \operatorname{Im} T,
(x,(2)) := \eth image(T) : \sum x \in A . Tx = y,
Assume \rho:R,
(3) := (2)(\rho y) \delta R - MOD(A, B)(T)(x, \rho) : \rho y = \rho T x = T \rho x,
() := \Im image(T)(3) : \rho y \in \operatorname{Im} T;
\rightsquigarrow (*) := \eth^{-1}Suvset(B)(1) : Im T \subset_{R\text{-MOD}} B;
```

```
Proof =
. . .
{\tt SubspacrLinearImageIsSubspace} \ :: \ \forall A,B \in R{\textrm{-MOD}} \ . \ \forall S \subset_{R{\textrm{-MOD}}} A \ . \ \forall T : A \xrightarrow{R{\textrm{-MOD}}} B \ .
    T(S) \subset_{R\text{-MOD}} B
Proof =
. . .
T^{-1}(S) \subset_{R\text{-MOD}} A
Proof =
(1) := \texttt{GroupPreimage}(T, S) : T^{-1}(S) \subset_{\mathsf{GRP}} A,
Assume x:T^{-1}S,
(2) := \eth T^{-1}S(x) : T(x) \in S,
Assume \rho: R,
(3) := \eth R\text{-MOD}(A, B)T(x, \rho) \eth Submodule(B)(S)(2)(\rho, T(x)) : T(\rho x) = \rho T(x) \in S,
() := \eth T^{-1}S(3) : \rho x \in T^{-1}S;
\rightsquigarrow (*) := \eth^{-1}Submodule(A)(1) : T^{-1}S \subset_{R\text{-MOD}} A;
{\tt ZeroModule} \, :: \, \Big( \{\star\}, (\star, \star) \mapsto \star, (\rho, \star) \mapsto \star \Big) \in R{\text{-}MOD}
Proof =
. . .
zeroModule :: R-MOD
\texttt{zeroModule}\left(\right) = \star := \left(\{\star\}, (\star, \star) \mapsto \star, (\rho, \star) \mapsto \star\right)
ZeroElementStable :: \forall M \in R-MOD . \forall \alpha \in R . \alpha 0 = 0
Proof =
Assume (1): M \cong_{R\text{-MOD}} \star,
(2) := \eth \star : \alpha 0 = 0;
\sim (1) := I(\rightarrow) : M \cong_{LMODR} \star \Rightarrow \alpha 0 = 0,
Assume (2): M \ncong_{R\text{-MOD}} \star,
(m,(3)) := \eth \star (2) : \sum m \in M \cdot m \neq 0,
(4) := \eth R\text{-MOD}\eth \text{Neutral}(M, +)(0)(m) : \alpha m + \alpha 0 = \alpha (m + 0) = \alpha m,
(5) := \eth Neutral(M, +)(0)(m) : \alpha m + 0 = \alpha m,
() := GroupTotalMult(4)(5) : \alpha 0 = 0;
```

```
ZeroSubmodule :: \forall M \in R\text{-MOD} . \{0\} \subset_{R\text{-MOD}} M
Proof =
. . .
\texttt{kerIsSubmodule} \, :: \, \forall A,B \in R\text{-MOD} \, . \, \forall T:A \xrightarrow{R\text{-MOD}} B \, . \, \ker T \subset_{R\text{-MOD}} A
Proof =
(1) := \eth \ker T : \ker T = T^{-1}\{0\},\
(*) := LinearPreimageIsSubmodule(T)ZeroSubmodule(B)(1) : ker T \subset_{R-MOD} A;
{\tt quotScalarMult} \, :: \, \prod M \in R \text{-MOD} \, . \, \prod S \in S \text{-MOD} \, . \, R \to \frac{M}{\varsigma} \to \frac{M}{\varsigma}
quotScalarMult(\rho, [m]) = \rho[m] := [\rho m]
Assume s:S,
(*) := \eth quotScalarMult(m+s) \eth R-MOD(M) \eth Submodule(S)(\rho,s) \eth quotientGroup(M,S) :
    : \rho[m+s] = [\rho(m+s)] = [\rho m + \rho s] = [\rho m] = \rho[m];
QuotientModule :: \forall M \in R\text{-MOD} . \forall S \subset_{R\text{-MOD}} M . \left(\frac{S}{M}, + \frac{S}{M}, \text{quotScalarMult}\right)
Proof =
Proof =
. . .
. \forall (0): S \subset \ker T . \exists ! T': \frac{S}{M} \xrightarrow{R\text{-MOD}} N . \pi_S T' = T
Proof =
\left(T',(1)\right):= {\tt SubgroupProjUP}: \sum T': \frac{M}{S} \xrightarrow{{\tt GRP}} N \;.\; \pi_S T' = T \;\&\;
   & \forall T'': \frac{M}{S} \xrightarrow{\mathsf{GRP}} N : \pi_S T'' = T' \Rightarrow T' = T'',
Assume [m]: \frac{M}{S},
Assume \alpha:R,
(2) := \eth^{-1}\pi_S \eth \mathsf{quotScalarMult} \eth (1) \eth R - \mathsf{MOD}(M,N)(T)(m,\alpha) : T'\alpha[m] = \pi_S T'(\alpha m) = T(\alpha m) = \alpha T(m),
(3) := \eth^{-1}\pi_S(1) : T'[m] = \pi_S T'(m) = T(m),
() := (2)(3) : T'\alpha[m] = \alpha T'[m];
\leadsto (2) := \eth R\text{-MOD} : [T': \frac{M}{S} \xrightarrow{R\text{-MOD}} N],
(*) := (1)(2) : This;
```

```
\textbf{IntersectionOfSubmodule} \ :: \ \forall N \in R\text{-}\mathsf{MOD} \ . \ \prod I \in \mathsf{SET} \ . \ \forall S : I \to \mathtt{Submodule}(M) \ . \ \bigcap S_i \subset_{R\text{-}\mathsf{MOD}} M
Proof =
. . .
Proof =
. . .
{\tt UnionOfSubmodules} \ :: \ \forall M \in R{\textrm{-MOD}} \ . \ \prod I : {\tt Toset} \ . \ \forall S : {\tt Nondecreasing}(M)(I, {\tt Submodule}(M)) \ .
  \bigcup S_i \subset_{R\text{-MOD}} M
Proof =
. . .
Simple ::?R-MOD
M: \mathtt{Simple} \iff \Big\{M, \{0\}\Big\} = \mathtt{Submodule}(M)
Proof =
Assume y:B,
Assume \alpha:R,
x := T^{-1}(y) : A,
(1) := \partial R-MOD(A, B)\partial x : T(\alpha x) = \alpha T(x) = \alpha y
():=\eth \mathtt{Inverse}(T)(1) \eth x: T^{-1}(\alpha y)=\alpha x=\alpha T^{-1}(y);
\rightsquigarrow (*) := \eth R\text{-MOD} : [T^{-1} : A \xrightarrow{R\text{-MOD}} B];
(1) := LinearImageIsSubmodule : T(A) \subset_{R-MOD} B,
(2) := \partial Simple(A, B)(1)(0) : T(A) = B,
(3) := \mathtt{KerIsSubmodule}(T) : \ker T \subset_{R\text{-MOD}} A,
(4) := \partial Simple \partial \ker T(0)(3) : \ker T = \{0\},\
(*) := \eth^{-1} \mathbf{Iso}(R\text{-MOD}) \mathbf{InjectiveByKernel}(M, N, T)(4)(2) : [T : A \xleftarrow{R\text{-MOD}} B];
```

```
Proof =
\varphi := \Lambda T : A \xrightarrow{A\text{-MOD}} M \cdot T(1) : \mathcal{M}_{A\text{-MOD}}(A, M) \xrightarrow{A\text{-MOD}} M,
Assume T: A \xrightarrow{A\text{-MOD}} M,
Assume (1): \varphi(T) = 0,
(2) := \eth \varphi(0) : T(1) = 0,
Assume a:A,
() := \eth A-MOD(A, M)(T)() : T(a) = aT(1) = a0 = 0;
\rightsquigarrow () := E(\rightarrow, =) : T = 0;
\rightsquigarrow (1) := InjevtiveByKernel : [T : A \hookrightarrow M],
Assume m:M,
T := \Lambda a \in A \cdot am : A \xrightarrow{A - \mathsf{MOD}} M
() := \eth \varphi(T) : \varphi(T) = m;
\sim (2 := \eth \mathsf{Iso}(R\text{-MOD})(1) : [\varphi : A \overset{A\text{-MOD}}{\longleftrightarrow} M],
(*) := \eth Isomotphic(2) : \mathcal{M}_{A-MOD}(A, M) \cong_{A-MOD} M;
Proof =
\odot:=\Lambda T:\mathrm{End}_{A\text{-MOD}}(M)\ .\ \Lambda f\in A\Big[\mathbb{Z}_+\Big]\ .\ \Lambda m\in M\ .\ \sum_{i=0}f_iT^i(m):\mathrm{End}_{A\text{-MOD}}M\to \Big\{(M,\cdot):A\big[\mathbb{Z}_+\big]\text{-MOD}\Big\},
```

LinearMapsFromTheRing ::  $\forall A \in \mathsf{ANN} : \forall M \in A\text{-MOD} : \mathcal{M}_{A\text{-MOD}}(A,M) \cong_{A\text{-MOD}} M$ 

#### 1.2 Limits of Modules

```
{\tt DirectProductOfModulesIsAModule} \ :: \ \forall I \in {\sf SET} \ . \ \forall M: I \to R{\textrm{-MOD}} \ . \ \prod M_i \in R{\textrm{-MOD}}
Proof =
. . .
Proof =
. . .
{\tt DirectProductIsProduct} \, :: \, \Big( {\tt directProduct}, \pi \Big) : {\tt Product}(M)
Proof =
Assume I: SET,
Assume M: I \to R-MOD,
Assume N: R-MOD,
\texttt{Assume} \ T: \prod i \in I \ . \ N \xrightarrow{R\text{-MOD}} M_i,
\left(T',(1)\right) := \eth \mathsf{Product}(\mathsf{ABEL})(I,M,N,T) : \sum T' : N \xrightarrow{\mathsf{ABEL}} \prod_{i \in I} M_i \; . \; \forall i \in I \; . \; T'\pi_i = T_i,
Assume n:N,
Assume \alpha:R,
Assume i:I,
(2) := \eth^{-1}\pi_i(1)\eth R\text{-MOD}(N, M_i)(T_i) : \left(T'(\alpha n)\right)_i = T'\pi_i(\alpha n) = T_i(\alpha n) = \alpha T_i(n),
(3) := \eth^{-1}\pi_i : \left(T'(n)\right)_i = T'\pi_i(n) = T_i(n),
() := (2)(3) : (T'(\alpha n)) = \alpha (T'(n));
\rightsquigarrow () := \eth product : T'(\alpha n) = \alpha T'(n);
\rightsquigarrow (*) := \eth^{-1}R\text{-MOD} : \left| T' : N \xrightarrow{R\text{-MOD}} \prod_{i} M_i \right| ;
{\tt DirectSumOfModulesIsAModule} \, :: \, \forall I \in {\sf SET} \, . \, \forall M : I \to R \text{-}{\sf MOD} \, . \, \bigoplus M_i \in R \text{-}{\sf MOD}
Proof =
. . .
```

```
\textbf{InclusionIsLinear} \, :: \, \forall I \in \mathsf{SET} \, . \, \forall M: I \to R\text{-}\mathsf{MOD} \, . \, \forall i \in I \, . \, \iota_i : M_i \xrightarrow{R\text{-}\mathsf{MOD}} \bigoplus_{i \in I} M_j
Proof =
 . . .
 {\tt DirectSumIsCoproduct} \, :: \, \Big( {\tt directSum}, \iota \Big) : {\tt Coproduct}(M)
Proof =
Assume I: SET,
Assume M: I \to R\text{-MOD},
Assume N: R-MOD,
\text{Assume } T:\prod i\in I \text{ . } M_i\xrightarrow{R\text{-MOD}}N,
\Big(T',(1)\Big) := \eth \texttt{Product}(\mathsf{ABEL})(I,M,N,T) : \sum T' : \bigoplus i \in IM_i \to N \; . \; \forall i \in I \; . \; \iota_i T' = T_i,
Assume m: \bigoplus M_i,
Assume \alpha:R,
Assume i:I,
(2) := \eth^{-1}\iota_i(1)\eth R - \mathsf{MOD}(M_i, N)(T_i) : T'(\alpha m_i) = \iota_i T'(\alpha m) = T_i(\alpha m) = \alpha T_i(m),
(3) := \eth^{-1}\pi_i : T'(m_i) = \iota_i T'(m) = T_i(m),
():=(2)(3):T'(\alpha m_i)=\alpha T(m_i);
 \rightsquigarrow () := \eth directSum \eth R-MOD\eth R-MOD(\iota) :
     : T'(\alpha m) = \sum_{i \in I} T'(\alpha \iota_i(m_i)) = \sum_{i \in I} \alpha T'(\iota_i(m_i)) = \alpha \sum_{i \in I} T'(\iota_i(m_i)) = \alpha T'(m);
\rightsquigarrow (*) := \eth^{-1}R\text{-MOD} : \left| T' : \bigoplus_i M_i \xrightarrow{R\text{-MOD}} N \right| ;
 ZeroModuleIsZero :: ★: Zero (R-MOD)
Proof =
Assume M: M\text{-MOD},
0_1 := \star \mapsto 0 : \star \xrightarrow{R\text{-MOD}} M,
0_2 := m \mapsto \star : M \xrightarrow{R \text{-MOD}} \star.
\texttt{Assume}\ T: \star \xrightarrow{R\text{-MOD}} M.
(1) := NeutralImage(T) : T(\star) = 0,
() := \eth T(1) : T = 0_1;
\rightsquigarrow (1) := I(\forall)(T) : \forall T : \star \xrightarrow{R-\mathsf{MOD}} M : T = 0_1,
\operatorname{Assume} T: M \xrightarrow{R\operatorname{-MOD}} M
() := \eth \star (T) : T = 0_2;
\rightsquigarrow (2) := I(\forall) : \forall T : M \xrightarrow{R-MOD} \star . T = 0_2,
(*) := \eth^{-1} \operatorname{Zero}(1)(2) : [\star : \operatorname{Zero}(R-\mathsf{MOD})];
```

```
\textbf{fibredModule} \ :: \ \prod I \in \mathsf{SET} \ . \ \prod M : I \to R\text{-}\mathsf{MOD} \ . \ \prod N \in R\text{-}\mathsf{MOD} \ . \ \left(\prod_i M_i \xrightarrow{R\text{-}\mathsf{MOD}} N\right) \to R\text{-}\mathsf{MOD}
\mathtt{fibredModule}\left(\nu\right) = \prod_{i \ \subset \ I} \quad M_i := \bigcap_{i,j \in I} \ker\left(\pi_i \nu_i - \pi_j \nu_j\right)
                                    N \square \nu
FibredModuleIsPullback :: fibredModule : Pullback (R-MOD)
Proof =
Assume I: SET,
Assume M: I \to R\text{-MOD},
Assume N: R-MOD,
Assume \nu:\prod i\in I . M_i\xrightarrow{R	ext{-MOD}}N,
Assume P: R\text{-MOD},
Assume T: \prod i \in I . P \xrightarrow{R-MOD} M_i
Assume (1): \forall i, j \in I . T_i \nu_i = T_j \nu_j,
Assume (i, j) : I,
Assume p:P,
(2) := (1)(\pi_i \nu_i - \pi_j \nu_i) : \forall p \in P : (\pi_i \nu_i - \pi_j \nu_j) T(p) = T_i \nu_i(p) - T_j \nu_j(p) = 0,
() := \eth \ker : T(p) \in \ker T_i \nu_i - T_j \nu_j;

ightsquigarrow (2) := \eth^{-1} \mathtt{fibredModule} : \operatorname{Im} T \subset \prod M_i,
                                                                  i \in I
                                                                 N \square \nu
T' := T^{|\mathbf{fibredProduct}(I,M,N,\nu)} : P \xrightarrow{R-\mathsf{MOD}}
                                                                              M_i
                                                                  i \in I
                                                                 N \square \nu
N \square \nu
Assume (3): \forall i \in I . T''\pi_i = T_i,
() := \eth T'(3) : T'' = T';
\rightsquigarrow (*) := \eth^{-1}Pushout : |fibredModule : Pullback(R-MOD)|;
\mathtt{span} :: \prod M \in R\text{-}\mathsf{MOD} . ?M \to \mathtt{Submodule}(M)
\operatorname{span}(X) = \operatorname{span}(X) := 
                                        X{\subset}S{\subset}_{R{	ext{-}MOD}}M
\mathbf{fibredSum} \, :: \, \prod I \in \mathsf{SET} \, . \, \prod M : I \to R\text{-}\mathsf{MOD} \, . \, \prod N \in R\text{-}\mathsf{MOD} \, . \, \left(\prod M_i \xrightarrow{R\text{-}\mathsf{MOD}} N\right) \to R\text{-}\mathsf{MOD}
\mathsf{fibredSum}\,(\nu) = \bigoplus \quad M_i := \frac{\bigoplus_{i \in I} M_i}{\mathrm{span}\{\nu_i \iota_i(n) - \nu_i \iota_i(n) | i, j \in I, n \in N\}}
                               i \in I
```

 $N \square \nu$ 

```
FibredSumIsPushout :: fibredModule : Pushout (R-MOD)
 Proof =
 Assume I: SET,
 Assume M: I \to R\text{-MOD},
 Assume N: R-MOD,
Assume \nu:\prod i\in I . N\xrightarrow{R	ext{-MOD}} M_i,
 Assume P: R-MOD,
 Assume (1): \forall i, j \in I . \nu_i T_i = \nu_j T_j,
 Assume m : \operatorname{span} \{ \nu_i \iota_i(n) - \nu_j \iota_j(n) | i, j \in I, n \in \mathbb{N} \},
 L, i, j, n, 2) := \Im \operatorname{span}(m) : \sum_{l} L \in \mathbb{Z}_{+} . \sum_{l} i, j : L \to I . n : L \to N . m = \sum_{l} \nu_{i_{l}} \iota_{i_{l}}(n_{l}) - \nu_{j_{l}} \iota_{j_{l}}(n_{l}),
(3) := (2)(1) : \sum_{i=1}^{L} T_i(m_i) = \sum_{i=1}^{L} \nu_{i_l} T_{i_l}(n_l) - \nu_{j_l} T_{j_l}(n_l) = 0,
() := \eth \ker(3) : m \in \bigoplus_{I \in i} T_i(m);
\leadsto (T', ()) := \texttt{SubmoduleUP} : \sum T' : \bigoplus_{i \in I} M_i \to P . \ \forall m \in \bigoplus_{i \in I} M_i . \ T'[m] = \bigoplus_{i \in I} T_i(m) \ \& 
    & \forall T'': \bigoplus_{i \in I} M_i \to P. \left( \forall m \in \bigoplus_{i \in I} M_i \cdot T''[m] = \bigoplus T'(m) \right) \Rightarrow T'' = T';
                   N \square \nu
 \rightsquigarrow (*) := I(=, \rightarrow) : \left( \texttt{fibedModule} : \texttt{Pushout} \left( R - \mathsf{MOD} \right) \right);
 LeftModulesAreBicomplete :: R-MOD : Complete & Cocomplete
 Proof =
 . . .
```

$$\begin{aligned} &\operatorname{InnerDirectSum} :: \prod A \in \operatorname{ANN} . \ \prod M \in A\operatorname{-MOD} . \ ?\operatorname{Submodule}^2(M) \\ &(X,Y) : \operatorname{InnerDirectSum} \iff M = X \oplus Y \iff X + Y = M \ \& \ X \cap Y = \{0\} \end{aligned}$$
 
$$&\operatorname{InnerDirectSumIsDirectSum} :: \forall A \in \operatorname{ANN} . \ \forall M \in A\operatorname{-MOD} . \ \forall X,Y \subset_{A\operatorname{-MOD}} M \ . \\ &. M = X \oplus Y \Rightarrow M \cong_{A\operatorname{-MOD}} X \oplus Y \end{aligned}$$
 
$$&\operatorname{Proof} = \\ \varphi := \Lambda(x,y) : X \oplus Y . x + y : X \oplus Y \xrightarrow{A\operatorname{-MOD}} M, \\ [1] := \eth \operatorname{InnerDirectSum}(A,M,X,Y) : X + Y = M \ \& \ X \cap Y = \{0\}, \\ [*] := \eth^{-1}\operatorname{Iso}[1] : M \cong_{A\operatorname{-MOD}} X \oplus Y; \\ \square & \\ &\square & \\ &\operatorname{MultiInnerDirectSum} :: \prod A \in \operatorname{ANN} . \ \prod M \in .\operatorname{-MOD?} \sum I \in \operatorname{Set} . \ I \to \operatorname{Submodule}(M) \\ &(I,X) : \operatorname{MultiInnerDirectSum} \iff M = \bigoplus_{i \in I} X_i \iff M = \sum_{i \in I} X_i \ \& \ \forall i \in I . \ X_i \cap \sum_{j \in I, j \neq i} X_j = \{0\} \end{aligned}$$
 
$$&\operatorname{MultiInnerDirectSumIsDirectSum} :: \forall A \in \operatorname{ANN} . \ \forall M \in A\operatorname{-MOD} . \ \forall I \in \operatorname{SET} . \ \forall I : X \to \operatorname{Submodule}(M) \ . \\ &M = \bigoplus_{i \in I} X_i \Rightarrow M \cong_{A\operatorname{-MOD}} \bigoplus_{i \in I} X_i \end{aligned}$$
 
$$&\operatorname{Proof} = \\ \varphi := \Lambda x : \bigoplus_{i \in I} X_i . \sum_{i \in I} x_i : \bigoplus_{i \in I} X_i \xrightarrow{A\operatorname{-MOD}} M,$$
 
$$&[1] := \eth \operatorname{InnerDirectSum}(A,M,X,Y) : M = \bigoplus_{i \in I} X_i \Rightarrow M \cong_{A\operatorname{-MOD}} \bigoplus_{i \in I} X_i,$$
 
$$&[*] := \eth^{-1}\operatorname{Iso}[1] : \bigoplus_{i \in I} X_i \cong_{A\operatorname{-MOD}} \bigoplus_{i \in I} X_i; \end{aligned}$$

### 1.3 Free Modules and Generation of Submodules

```
freeModule :: Covariant(SET, R-MOD)
\mathtt{freeModule}\left(X\right) = R^{\oplus X} := \bigoplus_{} R
\texttt{freeModule}\left(X,Y,f\right) = R_{X,Y}^{\oplus f} := \Lambda v \in \mathcal{R}^{\oplus f} \; . \; \sum_{x \in X} v_x \iota_{f(x)}(1)
FreeModuleIsAdjoint :: freeModule ⊢ forgetful(R-MOD, SET)
Proof =
Assume X : \mathsf{SET},
Assume M: R-MOD,
\phi:=\Lambda f:X\to M \text{ . } \Lambda \sum_{x\in X} v_x \iota_x(1)\in R^{\oplus X} \text{ . } \sum_{x\in Y} v_x f(x):(X\to M)\to (R^{\oplus X}\xrightarrow{R\text{-MOD}})M,
\psi:=\Lambda T:R^{\oplus X}\xrightarrow{R\text{-MOD}}M\;.\;\Lambda x\in X\;.\;T\big(\iota_x(1)\big):(R^{\oplus X}\xrightarrow{R\text{-MOD}}M)\to (X\to M),
Assume f: X \to M,
() := \eth \phi \eth \psi : \psi \circ \phi(f) = \psi \left( \Lambda \sum_{x \in Y} v_x \iota_x \in R^{\oplus X} \cdot \sum_{x \in Y} v_x f(x) \right) = \Lambda x \in X \cdot f(x) = f;
\sim (1) := \eth^{-1} \text{RightInverse} : |\psi : \text{RightInverse}(\phi)|,
\operatorname{Assume} T: R^{\oplus X} \xrightarrow{R\operatorname{-MOD}} M
():=\eth\psi\eth\phi:\phi\circ\psi(T)=\phi\Big(\Lambda x\in X\ .\ T\big(\iota_x(1)\big)\Big)=\Lambda\sum_{x\in Y}v_x\iota_x(1)\ .\ \sum_{x\in Y}v_xxT\big(\iota_x)=T;
\rightsquigarrow () := \eth^{-1}inverse : \psi = \phi^{-1};
\rightsquigarrow (*) := \eth^{-1}LeftAdjoint : This;
\texttt{basisVector} :: \prod X \in \mathsf{SET} : X \to R^{\oplus X}
basisVector(x) = e_x := \iota_x(1)
linearCombination :: \prod X \in \mathsf{SET} . \prod M \in R\text{-MOD} . R^{\otimes X} \times M^X \to M
\texttt{linearCombination}\left(a,m\right) = am := \sum_{x \in X} a_x m_x
\texttt{spanWithFamily} :: \prod X \in \mathsf{SET} \;. \; \prod M \in R\text{-MOD} \;. \; (X \to M) \to \mathsf{Subspace}(M)
\operatorname{spanWithFamily}(v) = \operatorname{span}(v_x)_{x \in X} := \left\{ av | a \in R^{\oplus X} \right\}
SpanIsSpan :: \forall X \in \mathsf{SET} . \forall M \in R\text{-MOD} . \forall m : X \to M . \mathrm{span}(m_x)_{x \in X} = \mathrm{span}\,\mathrm{Im}\,m
Proof =
. . .
```

```
\label{eq:finitelyGeneratedModule} FinitelyGeneratedModule :: ?R-MOD \\ M: FinitelyGeneratedModule \iff \exists F: \texttt{Finite}(M): M = \text{span}(F) \\ \\ \texttt{Noetherian} :: ?R-\mathsf{MOD} \\ M: \texttt{Noetherian} \iff \forall A \subset_{R-\mathsf{MOD}} M \ . \ A: \texttt{FinitelyGeneratedModule}(R) \\ \\ \end{cases}
```

FGMBySubspace ::  $\forall M \in R\text{-MOD}$  .  $\forall N \subset_{R\text{-MOD}} M$  .

. 
$$N, \frac{M}{N}$$
 : FinitelyGeneratedModule  $\Rightarrow M$  : FinitelyGeneratedModule

Proof =

$$\begin{split} &\left(n,a,(1)\right) := \eth \texttt{FinitelyGeneratedModule}(R)(N) : \sum n \in \mathbb{N} \;.\; \sum a : n \to N \;.\; N = \mathrm{span}(a_i)_{i=1}^n, \\ &\left(m,[b],(2)\right) := \eth \texttt{FinitelyGeneratedModule}(R) \; \frac{M}{N} : \sum m \in \mathbb{N} \;.\; \sum [b] : m \to \frac{M}{N} \;.\; \frac{M}{N} = \mathrm{span}\left([b_i]\right)_{i=1}^m, \end{split}$$

Assume x:M,

$$\left(\alpha,(3)\right) := \eth \frac{M}{N}[x] : \sum \alpha : m \to R . \sum_{i=1}^{m} \alpha_i[b_i] = \left[\sum_{i=1}^{m} \alpha_i b_i\right],$$

$$(y,4) := \eth \frac{M}{N}(3) : \sum y \in N : x = y + \sum_{i=1}^{m} \alpha_i b_i,$$

$$(\beta, 5) := (2)(y) : \sum \beta : n \to R : y = \sum_{i=1}^{n} \beta_i a_i,$$

() := (5)(4) : 
$$x = \sum_{i=1}^{n} \beta_i a_i + \sum_{i=1}^{m} \alpha_i b_i$$
;

$$\leadsto (*) := \eth^{-1} \texttt{FinitelyGeneratedModule}(R) : \Big[ M : \texttt{FinitelyGeneratedModule} \Big];$$

 ${\tt NoetherianBySubspace} :: \forall M : {\tt FinitelyGeneratedModule}(R) : \forall N \subset_{R{\textrm{-}MOD}} M : {\tt Model}(R) : \forall N \in R{\textrm{-}MOD}(R) : {\tt Model}(R) : {\tt Model}(R)$ 

. 
$$N, \frac{M}{N}: \operatorname{Noetherian}(R) \iff M: \operatorname{Noetherian}(R)$$

Proof =

Assume B: Submodule(A),

$$(3) := {\tt SecondIsomorphism}(M,N,B) : \frac{B}{B \cap N} \cong_{R{\textrm{-}MOD}} \frac{B+N}{N},$$

$$(4) := \eth \mathtt{Noetherian} \; \frac{M}{N}(3) : \left[ \frac{B}{B \cap N} : \mathtt{FinitelyGeneratedModule}(R) \right],$$

$$(5) := \eth Noetherian(N)(N \cap B) : [N \cap B : FinitelyGeneratedModule(R)],$$

$$() := {\tt FGMBySubspace}(4)(6) : \Big[B : {\tt FinitelyGeneratedModule}(R)\Big];$$

$$\leadsto (*) := \eth^{-1} \mathtt{Noetherian} : [M : \mathtt{Noetherian}(R)],$$

```
Noetherian ByNoetherian :: \forall A: Noetherian . \forall M: FinitelyGeneratedModule(A) . M: Noetherian(A)
Proof =
\left(n,N,(1)\right):=\operatorname{\widetilde{o}FGM}(A)(M):\sum n\in\mathbb{N}\;.\;\sum N\subset_{A\text{-MOD}}A^{\oplus n}\;.\;M=\frac{A^{\oplus n}}{N},
\label{eq:definition} \mbox{$\begin{picture}(10,0) \put(0,0){$\begin{picture}(0,0) \put(0,0){$\beg
(2) := \eth Noetherian(A) : \Diamond (1),
Assume m:\mathbb{N},
Assume (3): \Diamond(n),
(4) := \eth \mathtt{quotientModule}(A^{\oplus n+1}, A^{\oplus n}) : \frac{A^{\oplus n+1}}{{}^{A \oplus}} \cong_{A\text{-MOD}} A,
(5) := NoetherianBySubspace(2)(3)(4) : \forall n + 1,
 \rightsquigarrow (3) := \ethInductiveSet(\mathbb{N}) : \forall n \in \mathbb{N} . \not \subseteq (n),
(4) := (3) : \xi(n),
(*) := NoetherianBySubspace(1)(4) : [M : Noetherian];
Generating :: \prod M \in R\text{-MOD} . ? \sum X \in \mathsf{SET} . X \to M
m: Generating \iff \operatorname{span}(m_i)_{i \in X} = M
 \texttt{LinearlyIndependent} :: \prod M \in R\text{-MOD} . ? \sum X \in \mathsf{SET} . X \to M 
m: \texttt{LinearlyIndependent} \iff \forall \alpha \in R^{\oplus X} \; . \; \alpha m = 0 \iff \alpha = 0
Basis := \Lambda M \in R\text{-MOD}. Generating & LinearlyIndependent(M): R\text{-MOD} \rightarrow :
FreeModule ::?R-MOD
M: \mathtt{FreeModule} \iff \exists X \in \mathsf{SET} \ . \ M \cong_{R\mathsf{-MOD}} R^{\oplus X}
\texttt{BasisIffFree} :: \prod M \in R\text{-}\mathsf{MOD} \;.\; \exists (X,m) : \mathtt{Basis}(M) \iff M : \mathtt{FreeModule}
Proof =
Assume (X, m): Basis(M),
\varphi := \Lambda \alpha \in A^{\oplus X} \ . \ \sum_{x \in A} \alpha_x m_x : A^{\oplus X} \xrightarrow{R \text{-MOD}} M,
(1) := \eth Generating(X, m) \eth^{-1} Surjective : [\varphi : A^{\oplus X} \rightarrow M],
(2) := \eth \texttt{LinearlyIndependent}(X, n) \eth^{-1} \texttt{Injective} : [\varphi : A^{\oplus X} \twoheadrightarrow M],
(3) := {\tt LinearInversion\eth Bijective}(2)(1) : [\varphi : A^{\oplus X} \xleftarrow{R{\text{-}MOD}} M],
(4) := \eth Isomorphic : [A^{\oplus X} \cong_{R-MOD} M],
() := \eth^{-1}FreeModule : [M : FreeModule];
 \rightsquigarrow (1) := I(\Rightarrow) : \exists (X, m) : Basis(M) \Rightarrow M : FreeModule,
Assume (2): [M: FreeModule],
(X,(3)) := \eth \mathsf{FreeModule}(X) : [M = R^{\oplus X}],
() := \eth^{-1} \mathtt{Basis} \eth \mathtt{basis} : [e(R, X) : \mathtt{Basis}(X)];
```

```
{\tt LinearlyIndependentSet} \ :: \ \prod M : R{\textrm{-}{\sf MOD}} \ . \ ?M
S: \texttt{LinearlyIndependentSet} \iff \exists m: \texttt{LinearlyIndependent}(M) . S = \operatorname{Im} m
\texttt{MaximalLinearIndependentExists} :: \forall M : R \text{-}\mathsf{MOD} . \forall S : \texttt{LinearlyIndependentSet}(M).
     \exists S' \in \max \texttt{LinearlyIndependentSet}(M) . S \subset S'
Proof =
 Use Zorn Lemma.
 {\tt LinearIndependenceOverFaF} \, :: \, \forall A : {\tt IntegralDomain} \, . \, \forall X,Y \in {\sf SET} \, . \, \forall m : Y \to A^{\oplus X} \, .
     (Y,m): LinearlyIndependent(A^{\oplus X}) \iff (Y,m): LinearlyIndependent(\operatorname{Frac} A^{\oplus X})
Proof =
Assume (1): [(Y, m): LinearlyIndependent(A^{\oplus X})],
Assume \frac{\alpha}{\beta}: Y \to \operatorname{Frac}(A)^{\oplus X},
Assume (2): \frac{\alpha}{\beta}m = 0,
(3) := \eth A\text{-MOD}(2) : 0 = \frac{\alpha}{\beta} m = \sum_{y \in Y} \frac{\alpha_y}{\beta_y} m_y = \frac{\sum_{y \in Y} \alpha_y \left(\prod_{y' \neq y} \beta_y\right) m_y}{\prod_{y \in Y} \beta_y},
(4) := \eth \texttt{VectorSpace}(\operatorname{Frac}(A))(3) : \sum_{y} \alpha_y \left( \prod_{x} \beta_{y'} \right) m_y = 0,
(5) := \eth \texttt{LinearlyIndependent}(m)(a)(4) : \forall y \in Y : \alpha_y \prod \beta_{y'} = 0,
(6) := \eth \operatorname{Frac} A(5) : \forall y \in Y . \alpha_y = 0,
() := \eth \operatorname{Frac} A\left(\frac{\alpha}{\beta}\right) : \frac{\alpha}{\beta} = 0;
\leadsto (*) := \eth^{-1} \texttt{LinearlyIndependent}(\operatorname{Frac}(A)^{\oplus X}) : [m : \texttt{LinearlyIndependent}(\operatorname{Frac}(A)^{\oplus X})];
 П
\texttt{BasisIso} \, :: \, \forall M, N : R\text{-}\mathsf{MOD} \, . \, \forall (X,f) : \texttt{Basis}(M) \, . \, \forall T : M \xleftarrow{R\text{-}\mathsf{MOD}} N \, . \, \Big(X,T(f)\Big) : \texttt{Basis}(M)
Proof =
Assume \alpha:A^{\oplus X}
Assume (1): \alpha T(f) = 0,
(2) := \eth R-MOD(T)(1) : T(\alpha f) = 0,
(3) := \eth \mathsf{Iso} \Big( R \mathsf{-MOD} \Big) (T)(2) : \alpha f = 0,
() := \eth LinearlyIndependent()(3) : \alpha = 0;
 \rightsquigarrow (1) := \eth^{-1}LinearlyIndependent(a) : [T(f) : LinearlyIndependent(N)],
Assume n:N,
(m,(2)):=\operatorname{\mathtt{\widetilde{d}Iso}}(T)(n):\sum m\in M . n=T(m),
(\alpha, (3)) := \eth Generating(M)(f)(m) : m = \alpha f,
(4) := \eth R\text{-MOD}(M)(T) : T(m) = T(\alpha f) = \alpha T(f);
 \sim (5) := \eth^{-1} \mathtt{Basis} : [T(f) : \mathtt{Basis}(M)],
```

```
MaximalLinearIndDominates :: \forall A: IntegralDomain . \forall M: FreeModule A.
         \forall E \in \max \texttt{LinearlyIndependentSet}(M) : \forall S \in \texttt
         |E| \le |S|
Proof =
\left(X,e\right):=\eth \texttt{LinearlyIndependentSet}(M)(E):\sum X\in \mathsf{SET}\;.\;e:X\overset{\mathsf{SET}}{\longleftrightarrow}E\;\&\;\mathtt{LinearlyIndependent}(M),
\left(Y,s\right):=\eth \texttt{LinearlyIndependentSet}(M)(S):\sum X\in \mathsf{SET}\;.\;s:Y\overset{\mathsf{SET}}{\longleftrightarrow}S\;\&\;\mathtt{LinearlyIndependent}(S),
Y' := Y \sqcup \{0\} : \mathsf{SET},
(1) := WellOrderingTheorem(X', 0) : [X' : WellOrdered & 0 = min X'],
I_0 := \emptyset : \emptyset \to X,
m_0 := e : X \to E
Assume y:Y,
(\alpha, \beta, (2)) := \eth \max 	exttt{LinearlyIndependentSet}(\operatorname{Im} m_{y--})(s_y) : \sum \alpha : A^{\oplus X} \cdot \sum \beta \in A^* .
       \sum_{x} \alpha_x m_{y--,x} = \beta s_y,
\Big(a,b,(3)\Big) := \eth m_{y--} : \sum a : I_{y--}\Big(\mathrm{Less}(y)\Big) \to S \; .
        . \sum b; X \setminus I_{y--}(\operatorname{Less}(y)) \to E \cdot m = a \oplus b,
(\xi,\zeta,(4)) := \operatorname{decomp}(\alpha,\operatorname{Im}I_{y--}) : \sum \xi : \operatorname{Im}I_{y--} \to A.
       \sum X \setminus \in I_{y--} \to ... \alpha = \xi \oplus \zeta,
(5) := (3)(4) : \beta s_y = \sum_{x} \xi_x a_x + \sum_{x} \zeta_x b_x,
(x,(6)) := \eth \texttt{LinearlyIndependentSet}(M)(S)(a,b)(5) : \sum x \in X . \zeta_x \neq 0,
I_y := I_{y--} \oplus (y \mapsto x) : \mathsf{Less}(y++) \hookrightarrow X,
m_y:=\Lambda x\in X . if x\in {
m Im}\, I then s\Bigl(I^{-1}(x)\Bigr) else e(x):X	o M,
Assume \gamma: A^{\oplus X},
Assume (7): \gamma m_u = 0,
Assume (8): \gamma \neq 0,
(9) := \eth LinearlyIndependent(m_{y--})(8) : \gamma_x \neq 0,
(10) := (2)(9) : 0 = \alpha_x^{-1} \beta_x \gamma_x e_x + \sum_{x' \in X : x' \neq x} \alpha_x^{-1} \beta \gamma_x m_{x'} - \gamma_{x'} m_{x'},
(11) := \eth LinearlyIndependent(m_{y--})(10) : \alpha_x^{-1}\beta\gamma_x = 0,
() := (8)(6)(11) : \bot;
 \rightarrow (12) := \eth \text{LinearlyIndependent}(M) : [m_u : \text{LinearlyIndependent}(M)],
(13) := \eth m_y(6)(5) : e_x \in \text{span}(m_{y,x'})_{x' \in X},
(14) := \eth \max \texttt{LinearlyIndependentSet}(M)(\operatorname{Im} m_{y--})(13) : [\operatorname{Im} m_y \in \max \texttt{LinearlyIndependentSet}(M)];
 \leadsto \Big(I,m,(2)\Big) := I\left(\sum\right) : \sum I : \prod y \in Y' \text{ . Less}(y++) \hookrightarrow X \text{ . } \sum m : \prod y \in Y' \text{ . } X \to M \text{ . }
        . Im m \in \max \texttt{LinearlyIndependentSet}(M) \& \forall x \in \text{Im } I : m_x = s_{I^{-1}(x)} \& \forall x \notin \text{Im } I : m_x = e_x,
(I') := TransfiniteInduction(Y')(2) : I' : X \hookrightarrow Y,
(*) := CardByInclusion(I')CardIso(e, s) : |S| \le |E|;
```

```
\texttt{FreeReflectsIso} \, :: \, \forall X,Y \in \mathsf{SET} \, . \, \forall A : \mathtt{IntegralDomain} \, . \, \forall (0) : A^{\oplus X} \cong A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} Y = A^{\oplus Y} \, . \, X \cong_{\mathsf{SET}} 
Proof =
T := \eth \mathsf{Isomorphic}(0) :: A^{\oplus X} \overset{A-\mathsf{MOD}}{\longleftrightarrow} A^{\oplus Y}.
 (1) := {\tt Isobasis}(T, e(A, X)) : \Big[T\big(e(A, X)\big) : {\tt Basis}(A^{\oplus Y})\Big],
 (2) := \texttt{MaximalLinearIndDomianates}(1)(T(e(A, X))) : |X| \le |Y|,
 (3) := \texttt{MaximalLinearIndDominates}(1)(e(A, Y), T(e(A, X))) : |Y| \le |X|,
 () := \text{EqCard}((2), (3)) : |Y| = |X|;
  \mathtt{idealModule} \, :: \, \prod A \in \mathsf{ANN} \, . \, \prod M \in A\mathtt{-MOD} \, . \, \mathtt{Ideal}(A) \to \mathtt{Subspace}(M)
 idealModule(I) = IM := span\{am | a \in I, m \in M\}
Proof =
\odot := \Lambda[a] \in \frac{A}{I} \cdot \Lambda m \in M \cdot am : \Lambda \frac{A}{I} \times M \to M,
Assume a:A,
 Assume i:I,
 Assume m:M,
(1):=\eth\odot\eth A\text{-}\mathsf{MOD}(M)(0)\eth^{-1}\odot:[a+i]\odot m=(a+i)m=am+im=am=[a]\odot;
 \rightsquigarrow () := \eth A\text{-MOD}(M)\eth \odot : \left| (M, \odot) \in \left( \frac{A}{I} \right) \text{-MOD} \right| ;
  Proof =
  . . .
  InvariantBasisProperty ::?RING
 R: Invariant Basis Property \iff free Module(R): Conservative (SET, R-MOD)
```

```
CRingIsIBP :: \forall A \in \mathsf{ANN} \ . \ A : InvariantBasisProperty
Proof =
I := MaximalIdealExists(A) : MaximalIdeal(A),
Assume X, Y : \mathsf{SET},
\mathrm{Assume}\ (1): A^{\oplus X} \cong_{A\text{-MOD}} A^{\oplus Y},
T:=\eth \mathtt{Isomorphic}:A^{\oplus X}\cong_{A\text{-MOD}}A^{\oplus Y},
f:=T\Big(e(A,X)\Big): {\tt Basis}\Big(A^{\oplus Y}\Big),
T' := \Lambda \sum_{x \in X} [a_i][e_x] \; . \; \sum_{x \in X} [a_x][f_x] : \left(\frac{A}{I}\right)^{\oplus X} \xleftarrow{\left(\frac{A}{I}\right)\text{-MOD}} \left(\frac{A}{I}\right)^{\oplus Y},
(1) := FreeModCreate \eth Isomorphic(A) : |X| = |Y|;
\mathtt{rank} \, :: \, \prod A \in \mathsf{ANN} \, . \, \mathtt{FinitelyGeneratedModule}(A) \to \mathbb{N}
\operatorname{rank}(M) = \operatorname{rank} M := |X| \quad \text{where} \quad A^{\oplus X} \cong_{A\operatorname{-MOD}} M
{\tt NoetherianACC} :: \forall M : {\tt Noetherian}(R) . \forall N : {\tt Nondescending}\Big(\mathbb{N}, {\tt Submodule}(M)\Big) \; .
     \exists n \in \mathbb{N} : \forall k \in \mathbb{N} : k \ge n \Rightarrow N_k = N_n
Proof =
V := \bigcup_{n=1}^{\infty} N_n : \mathtt{Submodule}(R),
(m,v,1) := \eth \mathtt{Noetherian}(A)(M)(V) : \sum m \in \mathbb{N} \; . \; \sum v : m \to V \; . \; V = \mathrm{span}(v_i)_{i=1}^m,
(n,2) := \eth \mathtt{Nondescending}(N) \eth \mathtt{union}(N)(v) : \exists n \in \mathbb{N} . N_n = V,
(3) := \eth Nondescending(N)UnionContains(N, V) : THIS;
```

## 1.4 Chain Complexes and Exact Sequences

```
{\tt MorphismChain} := \Lambda R \in {\sf RING} \;.\; \sum V : \mathbb{Z} \to R \text{-MOD} \;.\; \sum \prod^{\infty} \; \phi : V_i \xrightarrow{R \text{-MOD}} V_{i-1} : {\sf RING} \to {\sf Type};
\texttt{ChainComplex} :: \prod R \in \mathsf{RING} \:. \: ? \sum V : \mathbb{Z} \to R \text{-MOD} \:. \: \sum \phi : V_i \xrightarrow{R \text{-MOD}} V_{i-1}
(V,\phi): \mathtt{ChainComplex} \iff \forall i \in \mathbb{Z} \ . \ \mathrm{Im} \ \phi_{i+1} \subset \ker \phi_i
Exact :: \prod R \in \mathsf{RING} . ?ChainComplex(R)
(V,\phi): \mathtt{Exact} \iff \forall i \in \mathbb{Z} \;.\; \mathrm{Im}\, \phi_{i+1} = \ker \phi_i
Finite :: \prod R \in \mathsf{RING} . ?ChainComplex(R)
V, \phi : \texttt{Finite} \iff \exists n \in \mathbb{N} : \exists m \in \mathbb{N} : \forall i \in \mathbb{Z} : i < -n \ \& \ i > m \Rightarrow V_i = 0
\texttt{finiteChain} \, :: \, \prod R \in \mathsf{RING} \, . \, \left( \sum n \in \mathbb{N} \, . \, \sum V : n \to R \text{-}\mathsf{MOD} \, . \, \sum \phi : \prod i \in (n-1) \, . \, V_{i+1} \xrightarrow{R \text{-}\mathsf{MOD}} V_i \right) \to 0
     \rightarrow MorphismChain(R)
\texttt{finiteChain}\left(V,\phi\right) = V_n \xrightarrow{\phi_{n-1}} V_{n-1} \xrightarrow{\phi_{n-2}} \dots \xrightarrow{\phi_1} V_1 :=
     := \Big(\Lambda i \in \mathbb{Z} 	ext{ . if } i \in n 	ext{ then } V_i 	ext{ else } 0, \Lambda i \in \mathbb{Z}. 	ext{if } 2 \leq i \leq n 	ext{ then } \phi_{i-1} 	ext{ else } 0\Big)
RightChain :: \prod R \in \mathsf{RING} . ?MorphismChain(R)
V, \phi : \mathtt{RightChain} \iff \forall i \in \mathbb{Z} : i > 0 \Rightarrow V_i = 0
LeftChain :: \prod R \in \mathsf{RING} . ?MorphismChain(R)
V, \phi : \texttt{LeftChain} \iff \forall i \in \mathbb{Z} : i < 0 \Rightarrow V_i = 0
{\tt InjectionByRightChain} :: \forall (V,\varphi); {\tt RightChain} \ \& \ {\tt Exact}(R) \ . \ \varphi_1 : V_1 \hookrightarrow V_0
Proof =
[1] := RightChain(V, \varphi)\delta R-MOD(0) : \varphi_1 = 0,
[2] := \eth^{-1} \operatorname{Im} \varphi_0[1] : \operatorname{Im} \varphi_1 = 0,
[3] := \eth \texttt{Exact}[2] : \ker \varphi_0 = 0,
[*] := ZeroKernelTHM[4] : (\varphi_1 : V_1 \hookrightarrow V_0);
 \texttt{SurjectionByLeftChain} :: \forall (V, \varphi); \texttt{LeftChain} \& \texttt{Exact}(R) . \varphi_0 : V_0 \twoheadrightarrow V_{-1}
Proof =
[1] := \mathbf{LeftChain}(V, \varphi) \eth R - \mathsf{MOD}(0) : \varphi_1 = 0,
[2] := \eth^{-1} \ker \varphi_1[1] : \ker \varphi_1 = V_0,
[3] := \eth \mathsf{Exact}[2] : \mathrm{Im}\,\varphi_0 = V_0,
[*] := \eth^{-1}Surjective : (\varphi_0 : V_0 \rightarrow V_- 1);
```

```
\texttt{MorphismOfChains} \, :: \, \prod(V,\varphi), (W,\psi) : \texttt{ChainComplex}(R) \, . \, ? \, \stackrel{\infty}{\prod} \, V_i \xrightarrow{R\texttt{-MOD}} W_i
f: \texttt{MorphismOfChains} \iff \forall i \in \mathbb{Z} : f_{i-1}\psi_i = \phi_i f_i
{\tt MorphismOfChainsComposition} \ :: \ \prod(V,\varphi), (W,\psi), (U,\eta) : {\tt ChainComplex}(R) \ .
     . \ \forall f : \texttt{MorphismOfChains}\Big((V,\varphi),(W,\psi)\Big) \ . \ \forall g : \texttt{MorphismOfChains}\Big((W,\psi),(U,\eta)\Big) \ . \\
     . g \circ f : MorphismOfChains \Big((V, \varphi), (U, \eta)\Big)
Proof =
Assume i: \mathbb{Z},
[i.*] := \eth^2 \texttt{MorphismOfChains}\Big((V,\varphi),(W,\psi)\Big)\Big((W,\psi),(U,\eta)\Big)(f)(g) : f_{i-1}g_{i-1}\eta_i = f_{i-1}\psi_ig_i = \varphi_if_ig_i;
\sim [*] := \eth^{-1} \texttt{MorphismOfChains} : \bigg(g \circ f : \texttt{MorphismOfChains}\Big((V,\varphi), (U,\eta)\Big)\bigg);
CategoryOfChains :: RING \rightarrow CAT
	ext{CategoryOfChains}(R) = R	ext{-CH} := \Big(	ext{ChainComplex}, 	ext{MorphismOfChains}, (\circ)_*, \Lambda n \in \mathbb{Z} \text{ . id}\Big)
IsoByExact :: \forall A \xrightarrow{\varphi} B : \text{Exact}(R) : \varphi : A \xleftarrow{R} B
Proof =
Combine InjectionByRightChain and SurjectionByLeftChain.
ShortExact :: \prod R \in \mathsf{RING} . ?ExactR
(V,\varphi): \mathtt{ShortExact} \iff \exists A,B,C \in R\mathtt{-MOD}: \exists \alpha: A \xrightarrow{R\mathtt{-MOD}} B: \exists \beta: B \xrightarrow{R\mathtt{-MOD}} C: (V,\varphi) = A \xrightarrow{\alpha} B \xrightarrow{\beta} C
ShortExactProperty :: \forall A \xrightarrow{\alpha} B \xrightarrow{\beta} C : ShortExact(R) . C \cong_{R\text{-MOD}} \frac{B}{\operatorname{Im} \alpha}
Proof =
Combine exactens with the first isomorphism theorem.
\texttt{ShortExactOfHomo} \, :: \, \forall R \in \mathsf{RING} \, . \, \forall A, B \in R\text{-}\mathsf{MOD} \, . \, \forall T : A \xrightarrow{R\text{-}\mathsf{MOD}} B \, .
     . \ker T \xrightarrow{\iota} A \xrightarrow{T^{|\operatorname{Im} T}} \operatorname{Im} T : \operatorname{ShortExact}(R)
Proof =
. . .
```

```
\textbf{ExactIsNaturalForChains1} \ :: \ \forall (V,\varphi), (W,\psi) \in R\text{-CH} \ . \ \forall f: (V,\varphi) \xleftarrow{R\text{-CH}} (W,\psi) \ .
    (V, \varphi) : \mathsf{Exact}(R) \Rightarrow (W, \varphi) : \mathsf{Exact}(R)
Proof =
\mathtt{Assume}\ [1]: \Big((V,\varphi): \mathtt{Exact}(R)\Big),
Assume i: \mathbb{Z},
[2] := \eth \mathsf{Iso}(f) \eth R\text{-}\mathsf{CH} : f_i : V_i \xleftarrow{R\text{-}\mathsf{MOD}} W_i \ \& \ f_{i-1} : V_{i-1} \xleftarrow{R\text{-}\mathsf{MOD}} W_{i-1},
[3] := \eth \mathsf{Exact}(V, \varphi)(i-1) : \ker \varphi_{i-1} = \operatorname{Im} \varphi_i,
[4] := \eth(-\mathsf{CH}R)(f)(i-1) : f_{i-1}\psi_{i-1} = \varphi_{i-1}f_{i-2},
[5] := \Im \mathsf{Iso}(f)[2][4] : \ker \psi_{i-1} = f_{i-1}(\ker \varphi_{i-1}),
[6] := \eth(-\mathsf{CH}R)(f)(i) : f_i \psi_i = \varphi_i f_{i-1},
[7] := \eth Iso(f)[2][6] : Im \psi_i = f_{i-1}(Im \varphi_i),
[1.*] := [3][5][7] : \operatorname{Im} \psi_i = \ker \psi_{i-1};
\sim [*] := I(\Rightarrow)\eth^{-1}(-\mathsf{CH}R)(V,\varphi)(W,\psi):(V,\varphi):\mathsf{Exact}(R)\Rightarrow(W,\psi):\mathsf{Exact}(R);
\textbf{ExactIsNaturalForChains2} \ :: \ \forall (V,\varphi), (W,\psi) \in R\text{-CH} \ . \ \forall f: (V,\varphi) \xleftarrow{R\text{-CH}} (W,\psi) \ .
    (V, \varphi) : \mathsf{Exact}(R) \iff (W, \varphi) : \mathsf{Exact}(R)
Proof =
[1] := \text{ExactIsNaturalForChains}(f) : (V, \varphi) : \text{Exact}(R) \Rightarrow (W, \psi) : \text{Exact}(R),
[2] := \texttt{ExactIsNaturalForChains}(f^{-1}) : (W, \psi) : \texttt{Exact}(R) \Rightarrow (V, \varphi) : \texttt{Exact}(R),
[*] := I(\iff)[1][2] : This;
```

#### 1.5 Split Exact Sequences

```
\texttt{DirectSumShortExact} :: \forall R \in \mathsf{RING} : \forall A, B \in R\text{-}\mathsf{MOD} : A \xrightarrow{\iota_1} A \oplus B \xrightarrow{\pi_2} B : \texttt{ShortExact}(R)
 . . .
 Split :: \prod R \in \mathsf{RING} . ?ShortExact(R)
(V,\varphi): \mathtt{Split} \iff \exists A,B \in R\mathtt{-MOD} : A \xrightarrow{\iota_1} A \oplus B \xrightarrow{\pi_2} B \cong_{R\mathtt{-CH}} (V,\varphi)
 \texttt{LeftInverseBySplit} :: \ \forall R \in \mathsf{RING} \ . \ \forall A,B \in R\text{-}\mathsf{MOD} \ . \ \forall \varphi : A \xrightarrow{R\text{-}\mathsf{MOD}} B \ .
     \varphi: \texttt{LeftInvertible}(A, B) \iff A \xrightarrow{\varphi} B \xrightarrow{\pi} \operatorname{coker} \varphi: \texttt{Split}
Proof =
\Big(\psi,[1]\Big):=\eth \mathsf{LeftInvertible}(\varphi):\sum \psi:B\xrightarrow{R\mathsf{-MOD}}A . \varphi\psi=\mathrm{id},
[2] := [1]ZeroKernelTHM : (\varphi : A \hookrightarrow B),
[3] := \eth^{-1} \operatorname{Iso}[2] : (\varphi^{|\operatorname{Im}\varphi} : A \stackrel{R\operatorname{-MOD}}{\longleftrightarrow} \operatorname{Im}\varphi),
T:=\Lambda(a,t)\in\operatorname{Im}\varphi\oplus\ker\psi\ .\ a+t:\operatorname{Im}\varphi\oplus\ker\psi\xrightarrow{R\text{-MOD}}B,
S:=\Lambda b\in B \ . \ \Big(\psi\varphi(b),b-\psi\varphi(b)\Big): B\xrightarrow{R\text{-MOD}}\operatorname{Im}\varphi\oplus\ker\psi,
[5] := \forall (b,t) \in \operatorname{Im} \varphi \oplus \ker \psi \ . \ \eth T[1] \eth S : \forall (b,t) \in \operatorname{Im} \varphi \oplus \ker \psi \ . \ ST(a,t) = S\big(a+t\big) = (a,t),
[6] := \forall b \in B \ . \ \eth S[1] \eth T : \forall b \in B \ . \ TS(b) = S\big(\psi \varphi(b), b - \psi \varphi(b)\big) = b,
[7] := [5][6] \eth \mathtt{Iso} : \left(T : \ker \varphi \oplus \operatorname{Im} \psi \overset{R\text{-MOD}}{\longleftrightarrow} A\right),
[8] := \mathbf{IsomorphismTHM3}(\operatorname{Im}\varphi)[1] : \left(\pi_{|\ker\psi}^{-1} : \operatorname{coker}\varphi \overset{R\text{-MOD}}{\longleftrightarrow} \ker\psi\right),
[9] := \forall a \in A \cdot [1](a) : \forall a \in A \cdot \varphi^{|\operatorname{Im}\varphi}\iota(a) = (\varphi(a), 0) = (\varphi\psi\varphi(a), \varphi(a) - \varphi\psi\varphi(a)) = \varphi S(a),
[10] := I(=, \rightarrow)[9] : \varphi^{|\operatorname{Im}\varphi}\iota = \varphi S,
[11] := \forall b \in B . \ldots : \pi_2 S(b) = b - \psi \varphi(b) = \pi_{\operatorname{Im} \varphi} \pi_{\ker \psi}^{-1}(b),
[12] := I(=, \to)[11] : \pi_2 S = \pi_{\operatorname{Im}\varphi} \pi_{\ker \psi}^{-1},
[0.*] := \eth^{-1} \operatorname{Split} \eth^{-1} R \operatorname{-CH}[10][12] : A \xrightarrow{\varphi} B \xrightarrow{\pi} \operatorname{coker} \varphi : \operatorname{Split};
 \sim [0] := I(\Rightarrow) : Left \Rightarrow Right,
Assume [1]: \left(A \xrightarrow{\varphi} B \xrightarrow{\pi} \operatorname{coker} \varphi : \operatorname{Split}\right),
\left(X,Y,[2]\right):=\eth \mathtt{Split}[1]: \sum X,Y \in R\text{-MOD} \ . \ X \xrightarrow{\iota_1} X \oplus Y \xrightarrow{\pi_2} Y \cong_{R\text{-CH}} A \xrightarrow{\varphi} B \xrightarrow{\pi} \operatorname{coker} \varphi,
[3] := \eth \mathtt{Isomorphic} \big( (\mathsf{-CH} R) \big) : \sum f : X \xrightarrow{\iota_1} X \oplus Y \xrightarrow{\pi_2} Y \xleftarrow{R\mathtt{-CH}} A \xrightarrow{\varphi} B \xrightarrow{\pi} \operatorname{coker} \varphi,
[4] := \eth \iota \eth \pi : \iota_1 \pi_1 = \mathrm{id},
[5] := \eth R-CH(f)[4]\ethInverse(f_0): \varphi f_{-1}\pi_1 f_0^{-1} = f_0 \iota_1 \pi_1 f_0^{-1} = \mathrm{id},
[1.*] := \eth LeftInvertible : (\varphi : LeftInvertible(A, B));
 \sim [1] := I(\Leftarrow)[0]I(\iff) : \text{This};
```

```
\textbf{RightInverseBySplit} \, :: \, \forall R \in \mathsf{RING} \, . \, \forall A, B \in R\text{-MOD} \, . \, \forall \varphi : A \xrightarrow{R\text{-MOD}} B \, .
    \varphi: \mathtt{RightInvertible}(A, B) \iff \ker \varphi \xrightarrow{\iota} A \xrightarrow{\varphi} B: \mathtt{Split}
Proof =
Assume [0]: (\varphi: RightInvertible(A, B)),
\left(\psi,[1]\right):=\eth \mathtt{RightInvertible}(\varphi): \sum \psi: B \xrightarrow{R\text{-MOD}} A \ . \ \psi\varphi=\mathrm{id},
[2] := [1]\eth^{-1}Surjective : (\varphi : A \rightarrow B),
[3] := \mathbf{IsomorphismTHM}[2] : \left(\hat{\varphi} : \frac{A}{\ker(\varphi)} \overset{R\text{-MOD}}{\longleftrightarrow} B\right),
f := pi_{\ker \varphi \mid \operatorname{Im} \psi} \hat{\varphi} : \operatorname{Im} \psi \stackrel{R-\mathsf{MOD}}{\longleftrightarrow} B,
T:=\Lambda(b,t)\in\ker\varphi\oplus\operatorname{Im}\psi\;.\;b+t:\ker\varphi\oplus\operatorname{Im}\psi\xrightarrow{R\text{-MOD}}A,
S:=\Lambda a\in A \ . \ \left(a-\varphi\psi(a),\varphi\psi(a)\right): A\xrightarrow{R\text{-MOD}}\ker\varphi\oplus\operatorname{Im}\psi,
[5] := \forall (a,t) \in \ker \varphi \oplus \operatorname{Im} \psi \cdot \eth T[1] \eth S : \forall (a,t) \in \ker \varphi \oplus \operatorname{Im} \psi \cdot ST(a,t) = S(a+t) = (a,t),
[6] := \forall b \in B \ . \ \partial S[1] \partial T : \forall a \in A \ . \ TS(a) = S(\varphi \psi(a), a - \varphi \psi(a)) = a,
[7] := [5][6]\eth \mathsf{Iso} : \left(S : A \overset{R\text{-MOD}}{\longleftrightarrow} \ker \varphi \oplus \operatorname{Im} \psi\right),
[8] := \forall a \in ker \varphi . [1](a) : \forall a \in ker \varphi . \iota_1(a) = (a,0) = (a-\varphi\psi(a), \varphi\psi(a)) = \iota S(a),
[9] := I(=, \to)[8] : \iota_1 = \iota S,
[10] := \forall a \in A . ... : \varphi f^{-1}(a) = \varphi \psi(a) = S\pi_2(b),
[11] := I(=, \to)[10] : \varphi f^{-1} = S\pi_2,
[12] := \eth^{-1} \operatorname{Split} \eth^{-1} R \operatorname{-CH}[10][12] : \ker \varphi \xrightarrow{\iota} A \xrightarrow{\varphi} B : \operatorname{Split};
\leadsto [0] := I(\Rightarrow) : \texttt{Left} \Rightarrow \texttt{Right},
Assume [1]: \left(\ker\varphi\xrightarrow{\iota} A\xrightarrow{\varphi} B\xrightarrow{\pi}: \operatorname{Split}\right),
\left(X,Y,[2]\right):=\eth \mathtt{Split}[1]: \sum X,Y \in R\text{-MOD} \ . \ X \xrightarrow{\iota_1} X \oplus Y \xrightarrow{\pi_2} Y \cong_{R\text{-CH}} \ker \varphi \xrightarrow{\iota} A \xrightarrow{\varphi} B,
[3] := \eth \mathtt{Isomorphic} \big( (\mathsf{-CH} R) \big) : \sum f : X \xrightarrow{\iota_1} X \oplus Y \xrightarrow{\pi_2} Y \xleftarrow{R \mathsf{-CH}} \ker \varphi \xrightarrow{\iota} A \xrightarrow{\varphi} B,
[4] := \eth \iota \eth \pi : \iota_2 \pi_2 = \mathrm{id},
[5] := \eth R\text{-CH}(f)[4] \eth Inverse(f_0) : f_{-2}\iota_2 f_{-1}^{-1} \varphi = f_{-2}\iota_2 \pi_2 f_{-2}^{-1} = \mathrm{id},
[1.*] := \eth RightInvertible : (\varphi : RightInvertible(A, B));
 \sim [1] := I(\Leftarrow)[0]I(\iff) : \text{This};
```

## 1.6 Homology And Snake Lemma

```
homology :: ChainComplex(R) \to \mathbb{Z} \to R\text{-MOD}
\operatorname{homology}((V,\varphi),i) = H_i(V,\varphi) := \frac{\ker \varphi_i}{\operatorname{Im} \varphi_{i+1}}
. \ \forall (\lambda, \mu, \nu) : L_1 \xrightarrow{\alpha_1} M_1 \xrightarrow{\beta_1} N_1 \xrightarrow{R-\mathsf{CH}} L_0 \xrightarrow{\alpha_0} M_0 \xrightarrow{\beta_0} N_0 \ .
    \exists \delta : \ker \nu \xrightarrow{R\text{-MOD}} \operatorname{coker} \lambda : \ker \lambda \xrightarrow{\alpha_{1|\ker \lambda}} \ker \mu \xrightarrow{\beta_{1|\ker \mu}} \ker \nu \xrightarrow{\delta} \operatorname{coker} \lambda \xrightarrow{\hat{\alpha}_0} \operatorname{coker} \mu \xrightarrow{\hat{\beta}_0} \operatorname{coker} \nu : \operatorname{Exact}(R)
Proof =
Assume a : \ker \nu,
[1] := \eth ShortExact(\beta_1) : (\beta_1 : M_1 \rightarrow N_1)
[2] := \eth Surjective(\beta_1) : \{b \in M_1 : \beta_1(b) = a\} \neq \emptyset,
Assume b: M_1,
Assume [4]: \beta_1(b) = a,
c := \mu(b) : M_0,
[5] := \eth(-\mathsf{CH}R)(\mu)\eth\ker\nu: \beta_0(c) = \mu\beta_0(b) = \beta_1\nu(b) = 0,
[6] := \eth^{-1} \ker \beta_0 : c \in \ker \beta_0,
[7] := \eth \texttt{Exact}[6] : c \in \operatorname{Im} \alpha,
Assume d, d': L_0,
Assume [8]: \alpha_0(d') = c \& \alpha_0(d) = c,
\delta(a) := [d] : \operatorname{coker} \lambda,
[9] := \eth ShortExact(\alpha_0) : (\alpha_0 : \hookrightarrow (L_0, M_0)),
[d.*] := \eth Injective(\alpha_0)[8] : a = a';
 \sim (\delta(a), [8]) := \eth \text{WellDefined} : \sum \delta(a) \in \operatorname{coker} \lambda : \exists d \in L_0 : [d] = \delta(a)\alpha_0(d) = c,
Assume b': M_1,
Assume [9]: \beta_1(b') = a,
[10] := \eth R\text{-MOD}\beta_1[9] - [4] : 0 = \beta_1(b') - \beta_1(b) = \beta_1(b'-b),
[11] := \eth^{-1} \ker \beta_1[10] : b' - b \in \ker \beta_1,
[12] := \eth \mathsf{Exact} : b' - b \in \mathrm{Im} \, \alpha_1,
(z, [13]) := \eth \operatorname{Im}[12] : \sum z \in L_1 \cdot \alpha_1(z) = b' - b,
[14] := \eth R-CH[13] : \lambda \alpha_0(z) = \alpha_1 \mu(z) = \mu(b'-b),
[a.*] := \eth^{-1} \operatorname{Im} \lambda[14] : \mu(b'-b) \in \operatorname{Im} \lambda;
 \sim (\delta, [1]) := \eth^{-1} \mathtt{WellDefined} : \sum \delta : \ker \nu \xrightarrow{R\text{-MOD}} \operatorname{coker} \lambda .
     . \forall a \in \ker \nu : \exists d \in L_0 : \delta(a) = [d] \& \exists b \in M_1 : \beta_1(b) = a \& \mu(b) = \alpha_0(d),
```

```
Assume a : \ker \lambda,
[2] := \eth R\text{-CH}(\lambda, \mu, \nu) : \lambda \alpha_0 = \alpha_1 \mu,
[3] := [2]\eth \ker \lambda(a) : 0 = \lambda \alpha_0(a) = \alpha_1 \mu(a),
[a.*] := \eth^{-1} \ker \alpha_1 : \alpha_1(a) \in \ker \mu;
\sim [2] := \eth^{-1}image\eth^{-1}Subset : \alpha_1(\ker \lambda) \subset \ker \mu,
Assume a : \ker \mu,
[a.2] := \eth R \text{-CH}(\mu, \mu, \nu) : \mu \beta_0 = \beta_1 \nu,
[a.3] := [a.2] \eth \ker \mu(a) : 0 = \mu \beta_0(a) = \beta_1 \nu(a),
[a.*] := \eth^{-1} \ker \beta_1 : \beta_1(a) \in \ker \nu;
 \sim [3] := \eth^{-1} image \eth^{-1} Subset : \beta_1(\ker \mu) \subset \ker \nu,
Assume a: \operatorname{Im} \lambda,
[a.1] := \eth R \text{-CH}(\lambda, \mu, \nu) : \lambda \alpha_0 = \alpha_1 \mu,
(b, [a.2]) := \eth \operatorname{Im} \lambda a : \sum b \in L_1 \cdot a = \lambda(b),
[a.3]:=[a.2][a.1]\eth \texttt{quotientModule}:\alpha_0\pi_{\mathrm{Im}\,\mu}(a)=\lambda\alpha_0\pi_{\mathrm{Im}\,\mu}(b)=\alpha_1\mu\pi_{\mathrm{Im}\,\mu}(b)=0,
[a.*] := \eth^{-1} \ker[a.3] : a \in \ker \alpha_0 \pi_{\operatorname{Im} \mu};
\sim (\hat{\alpha}_0, [4]) := \texttt{RestrictMorphism} : \sum \hat{\alpha}_0 : \operatorname{coker} \lambda \xrightarrow{R\text{-MOD}} \operatorname{coker} \mu : \hat{\alpha}_0[a] = [\alpha_0(a)],
Assume a: \operatorname{Im} \mu,
[a.1] := \eth R\text{-CH}(\lambda, \mu, \nu) : \mu\beta_0 = \beta_1\nu,
(b, [a.2]) := \eth \operatorname{Im} \mu a : \sum b \in L_1 . a = \mu(b),
[a.3]:=[a.2][a.1]\eth \texttt{quotientModule}:\beta_0\pi_{\mathrm{Im}\,\nu}(a)=\mu\beta_0\pi_{\mathrm{Im}\,\nu}(b)=\beta_1\nu\pi_{\mathrm{Im}\,\nu}(b)=0,
[a.*] := \eth^{-1} \ker[a.3] : a \in \ker \beta_0 \pi_{\operatorname{Im} \nu};
\sim (\hat{\beta}_0, [5]) := \texttt{RestrictMorphism} : \sum \hat{\beta}_0 : \operatorname{coker} \mu \xrightarrow{R - \mathsf{MOD}} \operatorname{coker} \nu : \hat{\beta}_0[a] = [\beta_0(a)],
Assume a : \ker \nu \cap \beta_1(\ker \mu),
(d.b,[6]) := [1](a) : \sum d \in L_0 . \sum b \in M_1 . \delta(a) = [d] \& \beta_1(b) = a \& \mu(b) = \alpha_0(d),
[7] := [6] \eth a : b \in \ker \mu,
[8] := [7]\eth \ker \mu[6] : \alpha_0(d) = 0,
[9] := \eth Injective[8] : d = 0,
[a.*] := \eth^{-1} \ker \delta[9][6] : a \in \ker \delta;
\sim [6] := \eth^{-1} Subset : Im \beta_{1|\ker \mu} \subset \ker \delta,
Assume c: \operatorname{Im} \delta,
Assume (a,[c.1]):\sum a\in\ker\nu , a=\delta(c),
\left(d,b,[c.2]\right):=[1](a):\sum d\in L_0\;.\;\sum b\in M_1\;.\;\delta(a)=[d]\;\&\;\beta_1(b)=a\;\&\;\mu(b)=\alpha_0(d),
[c.3] := [c.1][c.2][4] dquotient Module : \hat{\alpha}_0(c) = \hat{\alpha}_0\delta(a) = \hat{\alpha}_0[d] = [\alpha_0(d)] = [\mu(b)] = 0,
[c.*] := \eth^{-1} \ker \hat{\alpha}_0[c.3] : c \in \ker \hat{\alpha}_0;
 \sim [7] := \eth Subset : Im \alpha_0 \subset \ker \hat{\alpha}_0,
```

```
Assume a : \ker \delta,  \begin{pmatrix} d, b, [a.1] \end{pmatrix}) := [1](a) : \sum d \in L_0 \ . \ \sum b \in M_1 \ . \ \delta(a) = [d] \ \& \ \beta_1(b) = a \ \& \ \mu(b) = \alpha_0(d),   [a.2] := \eth \ker \delta(a)[a.1] : d \in \operatorname{Im} \lambda,   (c, [a.3]) := \eth \ker \delta : \sum c \in L_1 \ . \ d = \lambda(c),   [a.4] := [a.3]\eth(-\operatorname{CH}R)(\lambda, \mu, \nu) : \mu(b) = \lambda \alpha_0(c) = \alpha_1 \mu(c),   [a.5] := \ldots : 0 = \mu(b - \alpha_1(c)),   [a.6] := R-\operatorname{MOD}(\beta_1)\eth\operatorname{ChainComplex}(\ldots)[a.3] : \beta_1(b - \alpha_1(c)) = \beta_1(b) = a,   [a.*] := [a.6][a.5] : a \in \beta_1(\ker \mu);   \leadsto [8] := \eth^{-1}\operatorname{Subset}\eth^{-1}\operatorname{SetEq}[6] : \beta_1(\ker \mu) = \ker \delta,   \operatorname{Assume} [d] : \ker \hat{\alpha}_0,   [d.1] := \eth \ker \hat{\alpha}_0 \eth \operatorname{coker} \mu : \alpha_0(d) \in \operatorname{Im} \mu,   (b, [d.2]) := \eth \operatorname{Im} \mu[d.1] : \sum b \in M_1 \ . \ \alpha_0(d) = \mu(b),   [d.*] := \eth\delta[d.2] : \delta\beta_1(b) = [d];   \leadsto [8] := \eth^{-1}\operatorname{Subset}\eth^{-1}\operatorname{SetEq}[6] : \beta_1(\ker \mu) = \ker \delta,   [*] := \eth^{-1}\operatorname{Exact}[2][3][4][5][6][7][8] : \operatorname{This};
```

## 1.7 Torsion, Presentation and Resolution

```
Torsion :: \prod M \in R-MOD . ?M
m: {\tt Torsion} \iff \{m\} \; ! \; {\tt LinearlyIndependent}(M)
Assume A: IntegralDomain,
torsion :: \prod M \in A\operatorname{-MOD} . \operatorname{Submod}(M)
torsion(M) = tor M := \{ m \in M : m : Torsion(M) \}
TorsionFree :: ?A-MOD
M: {\tt TorsionFree} \iff {\tt tor}\, M = \{0\}
Torsion :: ?A-MOD
M: \mathtt{Toraion} \iff \mathtt{tor}\, M = M
{\tt TorsionFreeSubmoduleIsTorsionFree} :: \forall M : {\tt TorsionFree}(A) . \forall S \subset_{A\texttt{-MOD}} M . S : {\tt TorsionFree}(A)
Proof =
  . . .
  {\tt TorsionFreeSumIsTorsionFree} :: \forall I \in {\tt SET} \ . \ \forall M: I \to {\tt TorsionFree}(A) \ . \ \bigoplus M_i : {\tt TorsionFree}(R)
Proof =
  . . .
  FreeModuleIsTorsionFree :: \forall X \in \mathsf{SET} \ . \ A^{\oplus X} : \mathsf{TorsionFree}(A)
Proof =
  Cyclic ::?R-MOD
M: \mathtt{Cyclic} \iff \exists m \in M : M = \langle m \rangle
\texttt{FieldByCyclic} :: \left( \forall M : \texttt{Cyclic}(A) \;.\; M : \texttt{TorsionFree} \right) \Rightarrow A : \texttt{FieldByCyclic} :: \left( \forall M : \texttt{Cyclic}(A) \;.\; M : \texttt{TorsionFree} \right) \Rightarrow A : \texttt{FieldByCyclic} :: \left( \forall M : \texttt{Cyclic}(A) \;.\; M : \texttt{TorsionFree} \right) \Rightarrow A : \texttt{FieldByCyclic} :: \left( \forall M : \texttt{Cyclic}(A) \;.\; M : \texttt{TorsionFree} \right) \Rightarrow A : \texttt{FieldByCyclic} :: \left( \forall M : \texttt{Cyclic}(A) \;.\; M : \texttt{TorsionFree} \right) \Rightarrow A : \texttt{FieldByCyclic} :: \left( \forall M : \texttt{Cyclic}(A) \;.\; M : \texttt{TorsionFree} \right) \Rightarrow A : \texttt{FieldByCyclic} :: \left( \forall M : \texttt{Cyclic}(A) \;.\; M : \texttt{TorsionFree} \right) \Rightarrow A : \texttt{FieldByCyclic} :: \left( \forall M : \texttt{Cyclic}(A) \;.\; M : \texttt{TorsionFree} \right) \Rightarrow A : \texttt{FieldByCyclic} :: \left( \forall M : \texttt{Cyclic}(A) \;.\; M : \texttt{TorsionFree} \right) \Rightarrow A : \texttt{FieldByCyclic} :: \left( \forall M : \texttt{Cyclic}(A) \;.\; M : \texttt{TorsionFree} \right) \Rightarrow A : \texttt{FieldByCyclic} :: \left( \forall M : \texttt{Cyclic}(A) \;.\; M : \texttt{TorsionFree} \right) \Rightarrow A : \texttt{FieldByCyclic} :: \left( \forall M : \texttt{Cyclic}(A) \;.\; M : \texttt{TorsionFree} \right) \Rightarrow A : \texttt{FieldByCyclic} :: \left( \forall M : \texttt{Cyclic}(A) \;.\; M : \texttt{TorsionFree} \right) \Rightarrow A : \texttt{FieldByCyclic} :: \left( \forall M : \texttt{Cyclic}(A) \;.\; M : \texttt{C
Proof =
  . . .
```

```
annihilator :: \prod M : R\text{-MOD} : M \to \text{LeftIdeal}(R)
annihilator (m) = Ann(m) := \{r \in R : rm = 0\}
annihilator :: R-MOD \rightarrow LeftIdeal(R)
annihilator(M) = Ann M := \bigcap Ann(m)
TorsionHasNontrivalAnnihilator :: \forall M: Torsion & FinitelyGeneratedModule(A). Ann M \neq \{0\}
Proof =
(n,m,1) := \eth \texttt{FinitelyGeneratedModule}(A)(M) : \sum n \in \mathbb{N} \; . \; m : n \to M \; . \; \mathrm{span}(m_i)_{i=1}^n,
(2):=\forall i\in n \ . \ \eth \texttt{Torsion}(A)(M)(a_i)\eth^{-1} \operatorname{Ann}(a_i): \forall i\in n \ . \ \operatorname{Ann}(a_i)\neq \{0\},
(3) := \eth IntegralDomain(A)(2) : \prod_{i=1}^{n} Ann(a_i) \neq \{0\},
(4) := \eth \operatorname{Ann} M(3) : \{0\} \neq \prod_{i=1}^{n} \operatorname{Ann}(a_i) \subset \operatorname{Ann} M;
 Presentation :: R-MOD \rightarrow?Exact(R-MOD)
ig([A,B,C,D],[\phi,\pi,z]ig) : Presentation \iff \Lambda M \in R	ext{-MOD} .
    . C=M\ \&\ D=0\ \&\ A,B: {\tt FreeModule}(R)
FinitelyPresented ::?R-MOD
M:FinitelyPresented \iff \existsPresentation(M)
OverNoetherianRingFGMIsFI :: \forall A : Noetherian . \forall M : FinitelyGeneratedModule(A) .
    M: FinitelyPresented(A)
Proof =
\Big(n,m,(1)\Big) := \eth \texttt{FinitelyGeneratedModule}(A)(M) : \sum n \in \mathbb{N} \;.\; m : n \to M \;.\; \mathrm{span}(m_i)_{i=1}^n,
T := \Lambda \alpha \in A^n \cdot \alpha m : A^n \xrightarrow{A - \mathsf{MOD}} M,
(2) := \eth T(2) : [T : A^n \to M],
(3) := \eth \texttt{Noetherian}(A^n, \ker T) \eth \texttt{NoetherianByNoetherian}(A, A^n) : \Big\lceil \ker T : \texttt{FinitelyGeneratedModule}(A) \Big\rceil,
\left(k,v,(4)\right) := \eth \texttt{FinitelyGeneratedModule}(A)(\ker T) : \sum k \in \mathbb{N} \;.\; v : k \to \ker T \;.\; \ker T = \mathrm{span}(v_i)_{i=1}^k, \; k \in \mathbb{N} \;.
T' := \Lambda \alpha \in A^k \cdot \alpha v : A^k \to A^n
(*) := \eth^{-1} \texttt{Presentation} : \left\lceil \left( [A^k, A^n, M, 0], [T', T, 0] \right) : \texttt{Presentation}(M) \right\rceil;
```

```
FreeResolution :: R-MOD \rightarrow?Exact(R-MOD)
(P,\varphi): FreeResolution \iff \Lambda M \in R\text{-MOD}. P_{-1}=0 \& P_0=M \& P_0=M
          & \forall n \in \mathbb{N} . P_n : \mathtt{FreeModule}(R)
FreeResolutionExtends1 ::?RING
A: \texttt{FreeResolutionExtends1} \iff \forall M: R\texttt{-MOD} \ . \ \forall n \in \mathbb{Z}_+ \ . \ \forall T: R^n \xrightarrow{R\texttt{-MOD}} M \ . \ \forall (0): \mathrm{Im} \ T = M \ .
          . \exists m \in \mathbb{N} . \exists S: R^m \xrightarrow{R\text{-MOD}} R^n: \operatorname{Im} S \subset \ker T \ \& \ \ker S = 0
A: PrincipleIdealDomain
Proof =
Assume I: Ideal(A),
\left(m,T,(1)\right):= \eth \texttt{FreeResolutionExtends1}\left(\frac{A}{I},1,\pi_I\right): \sum m \in \mathbb{N} \;.\; \sum T: A^n \xrightarrow{A} A \;.
          . Im T = \ker \pi_I \& \ker T = \{0\},\
(2) := \eth \pi_I(1) : \operatorname{Im} T = I,
Assume (3): n > 1,
(4) := \eth A \text{-MOD}(A^n, A)(T) : T\Big(\big(T(e_2), -T(e_1)\big) \oplus 0\Big) = T(e_1)T(e_2) - T(e_1)T(e_2) = 0,
(4) := \eth e(1)_2 : T(e_1) \neq 0 \& T(e_2) \neq 0,
() := (1)_2(5)(4) : \bot;
 \rightsquigarrow (3) := E(\perp) : n = 1,
() := (3)(1)_1 : I = \langle T(1) \rangle;
 \sim (*) := \eth^{-1}PrincipleIdealDomain : [A : PrincipleIdealDomain];
{\tt MapsToTorsionFreeIsTorsionFree} :: \forall M : {\tt TorsionFree}(A) . \forall N : A{\tt -MOD} .
          \forall \mathcal{M}_{A\text{-MOD}}(N,M) : \mathsf{TorsionFree}(A)
Proof =
  . . .
  Increasing Annihilator :: \forall A \in \mathsf{ANN} : \forall M \in A \text{-}\mathsf{MOD} : \forall m \in M : \forall a \in A : \mathsf{Ann}(m) \subset \mathsf{Ann}(am)
Proof =
 . . .
```

#### 1.8 Associated Primes

```
{\tt SubmoduleByIdeal} \ :: \ \forall A \in {\sf ANN} \ . \ \forall M : A \text{-}{\sf MOD} \ . \ \forall I : {\tt Ideal}(A) \ . \ \forall m \in M \ . \ \forall (0) : {\sf Ann}(m) = I \ .
    \exists S \subset_{A\text{-MOD}} M : S \cong_{A\text{-MOD}} \frac{M}{I}
Proof =
(*) := \dots : \operatorname{span}\{m\} \cong_{A\operatorname{-MOD}} \frac{M}{I};
associated :: \prod A \in \mathsf{ANN} . \prod M \in A\text{-MOD} . ?Ideal(A)
\texttt{associated}\,() = \operatorname{Ass}(M) := \{\operatorname{Ann}(m) | m \in M \ \& \ \operatorname{Ann}(m) : \operatorname{\texttt{Prime}}(M)\}
MaximalAnnihilatorsAreAssociated :: \forall a \in A . \forall M : A \text{-MOD}.
    \max\{\operatorname{Ann}(m)|m\in M\setminus\{0\}\}\subset\operatorname{Ass}(M)
Proof =
Assume Ann(m) : max\{Ann(m) | m \in M \setminus \{0\}\},\
Assume a, b: A,
Assume (1): ab \in ANN(m),
(2) := \eth \operatorname{Ann}(m) \eth^{-1} \operatorname{Subset}(a, b) : \operatorname{Ann}(m) \subset \operatorname{Ann}(am) \& \operatorname{Ann}(m) \subset \operatorname{Ann}(bm),
Assume (3): am \neq 0,
(4) := (1)\eth \operatorname{Ann}(bm) : b \in \operatorname{Ann}(am),
() := \eth \max(4)(2)I(|) : b \in Ann(m)|a \in Ann(m);
\rightsquigarrow (3) := I(\Rightarrow) : am \neq 0 \Rightarrow b \in \text{Ann}(m) | a \in \text{Ann}(m),
Assume (4): am = 0,
() := \eth \operatorname{Ann}(m)(4)I(|) : a \in \operatorname{Ann}(m)|b \in \operatorname{Ann}(m);
\rightsquigarrow () := E(|)(2)(3)I(\Rightarrow): b \in \text{Ann}(m)|a \in \text{Ann}(m);
\sim (1) := \eth^{-1} \text{Prime} : [\text{Ann}(m) : \text{Prime}(M)],
() := \eth^{-1} \operatorname{Ass}(M) : \operatorname{Ann}(m) \subset \operatorname{Ass}(M);
\rightsquigarrow (*) := \eth^{-1}Subset : \max\{\text{Ann}(m)|m \in M \setminus \{0\}\} \subset \text{Ass}(M);
AssociatedPrimesExistInNoetherian :: \forall A : Noetherian(M).
    \forall M : A \text{-MOD} : \exists I \in \mathrm{Ass}(M)
Proof =
. . .
```

```
\exists (N,P): n \to \mathtt{Submodule}(M): \exists P: (n-1) \to \mathtt{Prime}(n) . N_n = M \& N_1 = \{0\} \& n \in \mathbb{N}
     & \forall i \in (n-1). N_i \subset N_{i+1} & \frac{N_{i+1}}{N_i} \cong \frac{A}{P_i}
Proof =
N_1' := M : A\text{-MOD},
f_1 := \mathrm{id}_M : M \xrightarrow{A\text{-MOD}} N_1'
Assume n:\mathbb{N},
(1) := NoetherianFactor(N'_n) : [N'_n : Noetherian],
\Big(m,P_n,(2)\Big):= {	t AssociatedPrimesExistInNoetherian}(A)(N_n'):
     P_n: \sum m \in M . \sum P_n: \mathtt{Ideal}(A) . P_n = \mathrm{Ann}(m) \ \& \ N_n' \neq 0 \Rightarrow P_n: \mathtt{Prime}(A)
\left(S,(3)\right) := {\tt SubmoduleByIdeal}\left(N_n',P_n,(2)\right) : \sum S \subset_{A{\textrm{-MOD}}} N_N' \; . \; S \cong_{A{\textrm{-MOD}}} \frac{A}{P_n},
N'_{n+1} := \frac{N'_n}{S} : A\operatorname{-MOD},
f_{n+1} := f_n \pi_S : M \xrightarrow{A-\mathsf{MOD}} N'_{n+1};
\rightsquigarrow \Big(N',P,f,(1)\Big) := \mathtt{IductiveConstr} : \sum N' : \mathbb{N} \rightarrow A\mathtt{-MOD} \; . \; \sum P : \mathbb{N} \rightarrow \mathtt{Ideal}(A) \; .
    \sum f:\prod n\in\mathbb{N}\;.\;M\xrightarrow{A\text{-MOD}}N_n'\;.\;\forall n\in\mathbb{N}\;.\;\Im f_n=N_n'\;\&\;N_n'\neq 0\Rightarrow P_n:\operatorname{Prime}(A)
     \& \exists g: N'_n \xrightarrow{A\text{-MOD}} N'_{n+1}: \ker g \cong_{A\text{-MOD}} \frac{A}{P_n} \& f_{n+1} = f_n g,
N := \ker f : \mathbb{N} \to \mathtt{Submodule}(M).
(2) := \eth^{-1} \mathtt{Nodecreasing}(1) : \left[ N : \mathtt{Nondecreasing} ig( \mathbb{N}, \mathtt{Submodule}(M) ig) \right]
\Big(n,(3)):= \texttt{NoetherianACC}(N)(2): \sum n \in \mathbb{N} \; . \; \forall k \in \mathbb{N} \; . \; k \geq N \Rightarrow N_k=N_n,
(4) := (3)(n+1) : N_n = N_{n+1},
(5) := (1)(n) : \operatorname{Im} f_{n+1} = N'_{n+1},
(g,(6)) := (1)(n) : \sum g : N'_n \xrightarrow{A-MOD} N'_{n+1} : f_n g = f_{n+1} \& \ker g \cong \frac{A}{P},
(7) := (6)_1(1)(3) : \ker g = \{0\},\
(8) := (6)_2(7) : \frac{A}{P} \cong \{0\},\
(9) := \eth \operatorname{quotientRing}(A, P_n)(8) : P_n = A,
(10) := (1) \eth \mathtt{Prime}(A)(9) : N_n' = 0,
(11) := \eth N \eth \ker(10) : N_n = M,
Assume k:(n-1),
(k.1) := (3)(k+1) : \operatorname{Im} f_{k+1} = N'_{n+1},
(g,(k.2)) := (3)(k) : \sum g : N'_k \xrightarrow{A-MOD} N'_{k+1} : f_{k+1} = f_k g \& \ker g \cong \frac{A}{P_k},
(k.*) := \eth N(k.2)_2CompositionKernel(f_k, g)(k.1)\ethqutientModule(k.2)_2:
     : \frac{N_{k+1}}{N_k} = \frac{\ker f_{k+1}}{\ker f_k} = \frac{\ker f_k g}{\ker f_k} = \frac{\ker f_k \oplus \ker g}{\ker f_k} \cong \ker f \cong \frac{A}{P_k};
 \rightsquigarrow (*) := I(\forall) : This(n, N_{\mid n})
```

PrimeFactorizationSeria ::  $\forall A$  : Noetherian(M) .  $\forall M$  : FinitelyGeneratedModule(A) .  $\exists n \in \mathbb{N}$  :

```
\texttt{listOfAssPrimes} :: \prod A : \texttt{Noetherian} \;. \; \prod M : \texttt{FinitelyGeneratedModule}(A) \;. \; \sum n \in \mathbb{N} \;. \; n \to \mathsf{Ass}(M)
lisOfAssPrimes() = (n(M), p(M)) := PrimeFactorizationSeria(M)
AllTheAssesAreListed :: \forall A: Noetherian . \forall M: FinitelyGeneratedModule(A) . Im \mathfrak{p}(M) = \mathrm{Ass}(M)
Proof =
Assume M: FinitelyGeneratedModule(A),
Assume P : Ass(M),
Q := \mathfrak{p}_1(M) : \mathrm{Ass}(M),
\left(m,n,(M.P.1)\right):=\left(\eth\operatorname{Ass}(M)\right)^2(P)(Q):\sum m,n\in M\;.\;P=\operatorname{Ann}(m)\;\&\;Q=\operatorname{Ann}(n),
Assume (M.P.2): P \neq Q,
Assume \alpha, \beta : A,
Assume (M.P.2.1): \alpha m = \beta n,
a := \alpha m : M,
(M.P.2.1.1) := IncreasingAnnihilator\eth a(M.P.2.1) : Ann(m) \subset Ann(a) \& Ann(n) \subset Ann(a),
(M.P.2.1.2) := \eth Subset(M.P.2)(M.P.2.1.1) : Ann(m) \subseteq Ann(a) \& Ann(n) \subseteq Ann(a),
(M.P.2.1.3) := \texttt{MaximalAnnihilatorsAreAssociated}(M.P.2.1.2) : Ann(a) = Ann(0) = A,
(P.2.*) := \eth \operatorname{Ann}(a) \eth A - \operatorname{\mathsf{MOD}}(M)(M.P.2.1.3) : a = 0;
\leadsto (M.*) := \mathtt{InverseImplication} I(\Rightarrow) I(\forall) I(\Rightarrow) \eth^{-1} \mathtt{span} : P = Q | \operatorname{span}(n) \cap \operatorname{span}(m) = \{0\};
\rightsquigarrow (*) := \eth(n(M), \mathfrak{p}(M)) : This;
FGMHasFiniteAss :: \forall A: Noetherian . \forall M: FinitelyGeneratedModule(A) . |\operatorname{Ass}(M)| < \infty
Proof =
. . .
\bigcup \quad \operatorname{Ann}(m) = \bigcup \operatorname{Ass}(M)
UnionOfAnnihilators :: \forall A \in \mathsf{ANN} \ . \ \forall M \in A\mathsf{-MOD} \ .
                                                                           m \in M: m \neq 0
Proof =
. . .
```

#### 1.9 Free Modules over PID

```
IdealIsLinearlyARingInPID :: \forall A: PrincipleIdealDomain . \forall I: Proper(A) . I \cong_{A\text{-MOD}} A
Proof =
\Big(a,[1]\Big):=\eth 	exttt{PrincipleIdealDomain}(A)(I):\sum a\in A \ .\ I=\langle a
angle,
[2] := \eth Proper(I) \eth genIdeal[1] : a \neq 0,
\varphi := \Lambda x \in A \cdot xa : A \xrightarrow{A - \mathsf{MOD}} I,
[3] := [2] \eth Integral Domain(A) : \ker \varphi = \{0\},\
[4] := ZeroKernelTHM(2) : (\varphi : A \hookrightarrow I),
[5] := \eth genIdeal[1] : (\varphi : A \rightarrow I),
[6] := \eth \mathsf{ANN}(A)[4][5] : (\varphi : A \overset{A\text{-MOD}}{\longleftrightarrow} I),
[*] := \eth^{-1} \mathbf{Isomorphic}(\mathsf{ANN})[6] : A \cong_{A\text{-MOD}} I;
NonZeroFunctionalExist :: \forall A: PrincipleIdealDomain . \forall M: FreeModule(A) .
    . \forall N \subset_{A\text{-MOD}} . \forall [0] : N \neq \{0\} . \exists \varphi : M \xrightarrow{A} A . \varphi(N) \neq 0
Proof =
\Big(n,[1]\Big):=\eth \mathtt{Singleton}[0]:\sum n\in N \ . \ n\neq 0,
(k,x) := \eth \mathtt{FreeModule}(A)(M) : \sum k \in \mathsf{SET} \ . \ M \xleftarrow{A\mathtt{-MOD}} A^k,
(i,[2]:=\eth \mathrm{Iso}(x)[1]:\sum i\in k \ . \ x_i(n)\neq 0,
[*] := \eth image(x_i, N)[2] : x_i(N) \neq \{0\},\
OneDimensionalSubspaceDecomposition :: \forall A: PrincipleIdealDomain.
     \forall M : \mathtt{FreeModule}(A) : \forall N \subset_{A\mathtt{-MOD}} M : \forall (0) : N \neq \{0\}.
     \exists a \in A :: \exists m \in M : \exists n \in N : \exists M' \subset_{A\text{-MOD}} M : \exists N' \subset_{A\text{-MOD}} N :
     : n = am \& N' = M' \cap N \& M = \operatorname{span}(m) \oplus M' \& N = \operatorname{span}(n) \oplus N'
Proof =
Assume \varphi: M \xrightarrow{A\text{-MOD}} A,
I_{\varphi} := \varphi(N) : \mathtt{Submodule}(A),
[\varphi.*] := \eth \mathsf{ANN}(A)(I_{\varphi}) : (I_{\varphi} : \mathsf{Ideal}(A));
\rightsquigarrow I := I(\rightarrow) : M \xrightarrow{A \text{-MOD}} A \rightarrow \text{Ideal}(A),
[1] := \eth(M \xrightarrow{A \text{-MOD}} A)(0) : 0 \in \operatorname{Im} I,
[2] := \eth NonEmpty([1]) : Im I \neq \emptyset,
J := \max I : \mathbf{Ideal}(A),
\left(\alpha,[3]\right):=\eth J:\sum\alpha:M\xrightarrow{A\text{-MOD}}A\to\operatorname{Ideal}(A):\alpha(N)=J,
[4] := NonZeroFunctionalExists(A, M, N, [0])(\eth \alpha) : \alpha \neq 0,
\Big(a,[5]\Big):=\eth 	exttt{PrincipleIdealDomain}(A)(J):\sum a\in A:\langle a
angle=J,
[6] := [3][4][5] : a \neq 0,
\Big(n,[7]\Big):=\eth \mathtt{genIdeal}[3][4]:\sum n\in N\;.\;a=\alpha(n),
```

```
\mathtt{Assume}\ \varphi: M \xrightarrow{A\text{-MOD}} A,
\Big(b,[\varphi.1]\Big) := \eth \mathtt{PrincipleIdealDomain}(A) \langle a,\varphi(n)\rangle : \sum b \in A : \langle a,\varphi(n)\rangle = \langle b\rangle,
\Big(s,r,[\varphi.2]\Big]):=\eth \mathtt{Ideal}[\varphi.1]:\sum s,r\in A\;.\;b=sa+r\varphi(n),
\psi := s\alpha + r\varphi : M \xrightarrow{A\text{-MOD}} A.
[\varphi.3] := \eth \psi[\varphi.2] : b \in \psi(N),
[\varphi.4] := [5][\varphi.2][\varphi.3] : J \subset \psi(N),
[\varphi.5] := \eth J[\varphi.4] : J = \psi(N),
[\varphi.*] := \eth Principle(J)[\varphi.5] : a|\varphi(n),
\sim [8] := I(\forall) : \forall \varphi : M \xrightarrow{A-\mathsf{MOD}} A \cdot a | \varphi(n),
(k,x) := \eth \mathtt{FreeModule}(A)(M) : \sum k \in \mathsf{SET} \ . \ M \xleftarrow{A\mathtt{-MOD}} A^k,
[9] := \forall i \in k . [8](x_i) : \forall i \in k . a | x_i(n),
(r, [10]) := \eth \mathtt{Divides}[9] : \sum r : A^k \cdot x(n) = ar,
m := x^{-1}(r) : M,
[11] := \eth m[10] : n = am,
[12] := [7][11] \eth A - \mathsf{MOD}(M, A)(\alpha) : a = \alpha(n) = \alpha(am) = a\alpha(m),
[13] := \eth IntegralDomain(A)[12] : \alpha(m) = 1,
M' := \ker \alpha : \mathtt{Submodule}(M),
N' := M' \cap N : \mathtt{Submodule}(N),
Assume v:M,
[v.1] := \eth Functor(inverse, () \alpha(v)m) + v : v = \alpha(v)m + (v - \alpha(v)m),
[v.2] := [13] \delta A - \mathsf{MOD}(M, A) \alpha \delta \mathsf{Inverse} : \alpha(v - \alpha(v)m) = \alpha(v) - \alpha(v) = 0,
[v.*] := [v.1][v.2]\eth^{-1}M' : (v - \alpha(v)m) \in M';
\sim [14] := \mathtt{DirectDecmposition} : M = \langle m \rangle \oplus M',
Assume w:N.
\Big(c.[w.1]\Big) := \eth \texttt{genIdeal}[5](w) : \sum c \in A \ . \ \alpha(w) = ca.
[w.2] := [w.1][11] : \alpha(w)m = cam = cn \in \text{span}(n),
[w.3] := \eth inverse(\alpha(w)m) + w : w = \alpha(w)m + (w - \alpha(w)m),
[w.4] := [13] \partial A - \mathsf{MOD}(M, A)(\alpha) \partial \mathsf{Inverse} : \alpha(w - \alpha(w)m) = \alpha(w) - \alpha(w) = 0,
[w.*] := [w.3][w.4] \eth Sumodule(M)(N) \eth M' \eth N' : (w - \alpha(w)m) \in N',
 \sim [*] := DirectDecomposition : N = \text{span}(n) \oplus N';
```

```
{\tt SubmodOfFreeIsFree} :: \forall A : {\tt PrincipleIdealDomain} . \forall F : {\tt FreeModule \& FinitelyGeneratedModule}(A) .
                 \forall M \subset_{A\text{-MOD}} F \cdot M : FreeModule(A)
 Proof =
 V_0 := M : \mathtt{Submodule}(F),
 [1.0] := \eth V_0 : V_0 = M,
 Assume n:\mathbb{N},
 Assume [n.1]: V_{n-1} = \{0\},\
 v_n := 0 : V_{n-1},
 V_n := V_{n-1} : \mathtt{Submodule}(F),
 [1.n] := [1.(n-1)] \eth V_n \eth v_n : M = \bigoplus_{i=1}^n \operatorname{span}(v_i) \oplus V_n;
   \rightsquigarrow [n.1] := I(\Rightarrow) : V_{n-1} = \{0\} \Rightarrow (\ldots),
 Assume [n.2]: V_{n-1} \neq \{0\},\
  \Big(\ldots,V_n,v_n,[n.3]\Big):={	t OneDimensionalSubspaceDecompasition}(F,V_{n-1},[n.2]):
              : \dots : \sum V_n : \mathtt{Submodule}(F) : \sum v_n \in V_{n-1} : V_{n-1} = \mathrm{span}(v_i) \oplus V_n,
[1.n] := [1.(n-1)][n.3] : M = \bigoplus_{i=1}^{n} \operatorname{span}(v_i) \oplus V_n;
   \sim [n.*] := I(\Rightarrow)E(|)[n.1] : (\ldots)
   \sim \left(V, v, [2]\right) := I\left(\sum\right) : \sum V : \mathbb{N} \to \mathtt{Submodule}(F) \; . \; \sum v : \mathbb{N} \to M \; . \; \forall n \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n, \text{ and } v_i \in \mathbb{N} \; . \; M = \bigoplus_{i=1}^n \{v_i\} \oplus V_n \oplus 
[3] := \emptysetFinitelyGeneratedModule(F)(v) : \exists k \in (\operatorname{rank} F + 1) . (v_i)_{i=1}^k ! LinearlyIndependent<math>(F),
[4] := \eth \mathsf{InnerDirectSum} \eth(v, V, [2])[3] : M = \bigoplus_{i=1}^{k-1} \mathrm{span}(v_i),
[*] := \eth^{-1} \texttt{FreeModule}(A) \texttt{IdealIsLinearlyARingInPID}[4] : (M : \texttt{FreeModule}(A));
    Snowball :: \prod A: IntegralDomain . ?A^n
 a: \mathtt{Snowball} \iff \forall i \in (n-1) . a_{i-1} | a_i
 FreeSubmodBasisTHM :: \forall A: PrincipleIdealDomain . \forall F: FreeModule & FinitelyGeneratedModule(A) .
                 . \ \forall M \subset_{A\text{-MOD}} F \ . \ \forall [0]: M \neq \{0\} \ . \ \exists (e_i)_{i=1}^n : \mathtt{Basis}(F): (f_i)_{i=1}^m : \mathtt{Basis}(M) \ . \ \exists a : \mathtt{Snowball}(m,A): ae_{|m} = f(a_i)_{i=1}^m : \mathtt{Basis}(M) = f(a_
                           where n = \operatorname{rank} F, m = \operatorname{rank} M
  V_0 := F: FreeModule & FinitelyGeneratedModule(A),
  W_0 := N : Submodule(V_0),
  [1.0] := \eth V_0 : V_0 = F,
 [2.0] := \eth W_0 : W_0 = M,
 [3.0] := [0] \eth W_0 : W_0 \neq \{0\},
 Assume k:m,
  \Big(V_k,W_k,v_k,w_k,a_k,[k.1]\Big):=\mathtt{OneDimensionalSubspaceDecomposition}\Big(V_{k-1},W_{k-1},[3.(k-1)]\Big):=\mathtt{OneDimensionalSubspaceDecomposition}\Big(V_{k-1},W_{k-1},[3.(k-1)]\Big):=\mathtt{OneDimensionalSubspaceDecomposition}\Big(V_{k-1},W_{k-1},[3.(k-1)]\Big):=\mathtt{OneDimensionalSubspaceDecomposition}\Big(V_{k-1},W_{k-1},[3.(k-1)]\Big):=\mathtt{OneDimensionalSubspaceDecomposition}\Big(V_{k-1},W_{k-1},[3.(k-1)]\Big):=\mathtt{OneDimensionalSubspaceDecomposition}\Big(V_{k-1},W_{k-1},[3.(k-1)]\Big):=\mathtt{OneDimensionalSubspaceDecomposition}\Big(V_{k-1},W_{k-1},[3.(k-1)]\Big):=\mathtt{OneDimensionalSubspaceDecomposition}\Big(V_{k-1},W_{k-1},[3.(k-1)]\Big):=\mathtt{OneDimensionalSubspaceDecomposition}\Big(V_{k-1},W_{k-1},[3.(k-1)]\Big):=\mathtt{OneDimensionalSubspaceDecomposition}\Big(V_{k-1},W_{k-1},[3.(k-1)]\Big):=\mathtt{OneDimensionalSubspaceDecomposition}\Big(V_{k-1},W_{k-1},[3.(k-1)]\Big):=\mathtt{OneDimensionalSubspaceDecomposition}\Big(V_{k-1},W_{k-1},[3.(k-1)]\Big):=\mathtt{OneDimensionalSubspaceDecomposition}\Big(V_{k-1},W_{k-1},[3.(k-1)]\Big):=\mathtt{OneDimensionalSubspaceDecomposition}\Big(V_{k-1},W_{k-1},[3.(k-1)]\Big):=\mathtt{OneDimensionalSubspaceDecomposition}\Big(V_{k-1},W_{k-1},[3.(k-1)]\Big):=\mathtt{OneDimensionalSubspaceDecomposition}\Big(V_{k-1},W_{k-1},[3.(k-1)]\Big):=\mathtt{OneDimensionalSubspaceDecomposition}\Big(V_{k-1},W_{k-1},[3.(k-1)]\Big):=\mathtt{OneDimensionalSubspaceDecomposition}\Big(V_{k-1},W_{k-1},[3.(k-1)]\Big):=\mathtt{OneDimensionalSubspaceDecomposition}\Big(V_{k-1},W_{k-1},[3.(k-1)]\Big):=\mathtt{OneDimensionalSubspaceDecomposition}\Big(V_{k-1},W_{k-1},[3.(k-1)]\Big):=\mathtt{OneDimensionalSubspaceDecomposition}\Big(V_{k-1},W_{k-1},[3.(k-1)]\Big):=\mathtt{OneDimensionalSubspaceDecomposition}\Big(V_{k-1},W_{k-1},[3.(k-1)]\Big):=\mathtt{OneDimensionalSubspaceDecomposition}\Big(V_{k-1},W_{k-1},[3.(k-1)]\Big):=\mathtt{OneDimensionalSubspaceDecomposition}\Big(V_{k-1},W_{k-1},V_{k-1},[3.(k-1)]\Big)
                : \sum V_k \subset_{A\text{-MOD}} V_{k-1} \ . \ \sum W_k \subset_{A\text{-MOD}} W_{k-1} \ . \ \sum v_k \in V_{k-1} \ . \ \sum w_k \in W_{k-1} \ . \ \sum a_n \in A \ .
                w_k = a_k v_k \& \operatorname{span}(w_k) \oplus W_k = W_{k-1} \& \oplus \operatorname{span}(v_k) \oplus V_k = V_{k-1} \& W_k = V_k \cap W_{k-1}
```

```
[k.2] := {\tt SubmodOfFreeIsFree}(V_{k-1}, V_k) : (V_k : {\tt FreeModule}(A)),
[1.k] := [k.1][1.(k-1)] : F = \bigoplus_{i=1}^{\kappa} \operatorname{span}(v_i) \oplus V_k,
[2.k] := [k.1][2.(k-1)] : M = \bigoplus^{k} \operatorname{span}(w_i) \oplus W_k,
[3.k] := \eth m \eth \operatorname{rank}[2.k] : k = m | W_k \neq \{0\};
 \rightsquigarrow \Big(V,W,v,w,a,[4]\Big) := I\left(\sum\right): \sum V: m \rightarrow \mathtt{Submodule}(F) \;.\; \sum W: m \rightarrow \mathtt{Submodule}(M) \;.
    . \sum v: \prod k \in m . V[k-1] . \sum w: \prod k \in m . W_{k-1} . \sum a: m \to A .
    . w = av \& \forall k \in m . F = \bigoplus_{i=1}^{k} \operatorname{span}(v_i) \oplus V_k \& M = \bigoplus_{i=1}^{k} \oplus W_k \& W_k = V_k \cap V_{[k-1]},
v' := \mathtt{FreeHasBasis}(V_m) : \mathtt{Basis}(V_m),
e := v \oplus v' : \mathtt{Basis}(F),
f := w : Basis(M),
[5] := \eth f \eth 4[4]_1 : ae_{|m} = f,
Assume k:(m-1),
\varphi := \mathbf{free}(\{k, n\}, A)(i \mapsto \mathbf{if} \ i \leq k+1 \ \mathbf{then} \ 1 \ \mathbf{else} \ 0) : V_{i-1} \xrightarrow{A-\mathsf{MOD}} A,
[k.1] := \eth free \eth \varphi : \varphi(e_k) = 1 \& \varphi(e_{k+1}) = 1,
[k.2] := [5] \eth A - \mathsf{MOD}(V_{k-1}, A)(\varphi)[k.1] : \varphi(f_k) = \varphi(a_k e_k) = a_k \varphi(e_k) = a_k,
[\varphi.3] := \eth a_k \eth \mathsf{OneDimensionalSubspaceDecomposition}[\varphi.2] : \varphi(W_{k-1}) = \langle a_k \rangle,
[\varphi.4] := [5] \eth A \text{-MOD}(V_{k-1}, A)(\varphi)[k.1] : \varphi(f_{k+1}) = \varphi(a_{k+1}e_{k+1}) = a_{k+1}\varphi(f_{k+1}) = a_{k+1},
[k.*] := [k.4][k.3] : a_k | a_{k+1};
\sim [*] := \eth^{-1} \text{Snowball} : (a : \text{Snowball}(A, m));
 PIDExtendsToFreeResolution :: \forall A: PrincipleIdealDomain . A: FreeResolutionExtends1
Proof =
Assume M: FinitelyGeneratedModule(A),
Assume n:\mathbb{N},
Assume \pi : \text{Epi}((-\mathsf{MOD}A), A^n, M),
[1] := SubmodOfFreeIsFree(A, ker \pi) : (ker \pi : FreeModule(A)),
(m,x):=\eth \mathtt{FreeModule}(A)(\ker \pi): \sum m \in \mathbb{N} \ . \ x:A^m \overset{\ker}{\longleftrightarrow} \pi,
[*] := \eth^{-1}FreeResoluton\eth x : 0 \xrightarrow{0} A^m \xrightarrow{x} A^n \xrightarrow{x} M \xrightarrow{0} 0 : FreeResolution(A, M);
 \rightarrow [*] := \eth^{-1}FreeResolutionExtends1 : (A : FreeResolutionExtend1);
```

## 1.10 Classification of Finitely Generated Modules over PID

```
PidClassificationTheorem1 :: \forall A: PrincipleIdealDomain . \forall M: FinitelyGeneratedModule(A) .
     . \exists n \in \mathbb{N} : \exists a : \mathtt{Snowball}(n,A) \ . \ M \cong A^{\mathrm{rank}\,M} \oplus \bigoplus_{i=1}^n \frac{A}{\langle a_i \rangle}
Proof =
(m,\pi):=\eth \mathtt{FinitelyGeneratedModule}(A)(M):\sum m\in \mathbb{N}\;.\;\pi:\mathtt{Epi}(A\mathtt{-MOD},A^m,M),
\left(n,x\right):= \texttt{PIDExtensToFreeResolution}(m,\pi,M): \sum n \in \mathbb{N} \;.\; \sum x: A^n \xleftarrow{A\text{-MOD}} \ker \pi,
\left(e,a,f,[1]\right) := \mathtt{FreeSubmodBasisTHM}(A,M,\ker\pi) : \sum e : \mathtt{Basis}(A^m) \; . \; \sum a : \mathtt{Snowball}(A,n) \; .
    . \sum f : Basis(\ker \pi) . f = ae_{|n},
[*] := [1] \eth \ker \pi \eth x \eth^{-1} \mathbf{InnerDirectSum} : M = A^{\operatorname{rank} M} \oplus \bigoplus^{n} \frac{A}{\langle a_i \rangle};
 PidClassificationTheorem2 :: \forall A: PrincipleIdealDomain . \forall M: FinitelyGeneratedModule(A) .
     . \ \exists n \in \mathbb{N} : \exists p : n \to \mathtt{PrimeElement}(A) : \exists k : n \to \mathbb{N} : \exists t : \prod i \in n \ . \ k_i \to \mathbb{N} : M \cong A^{\mathrm{rank}\,M} \oplus \bigoplus_{i=1}^n \bigoplus_{j=1}^{k_i} \frac{A}{\langle p_i^{t_j} \rangle}
Proof =
a_i = \prod_{i=1}^n p_{i,j}^{k_j}
\frac{A}{\langle a_i \rangle} = \bigoplus_{i=1}^n \frac{A}{\langle p_{i,i}^{t_j} \rangle}
{\tt snowballOfMod} \ :: \ \prod A : {\tt PrincipleIdealDomain} \ . \ {\tt FinitelyGeneratedModule}(A) \rightarrow

ightarrow \sum n \in \mathbb{Z}_+ . Snowball(A,n)
\mathtt{snowballOfMod}(M) := \mathtt{PidClassificationTheorem1}(A, M)
{\tt primesOfMod} \ :: \ \prod A : {\tt PrincipleIdealDomain} \ . \ {\tt FinitelyGeneratedModule}(A) \rightarrow \\
     \rightarrow \sum n \in \mathbb{N} : \sum p : n \hookrightarrow \mathtt{PrimeElement}(A) : \sum k : n \rightarrow \mathbb{N} : \sum t : \prod i \in n \; . \; k_i \rightarrow \mathbb{N}
primesOfMod(M) := PidClassificationTheorem2(A, M)
AnnihilatorBySnowball :: \forall A: PrincipleIdealDomain . \forall M: FinitelyGeneratedModule & Torsion(A).
     . Ann(M) = \langle a_n \rangle where (n, a) = \text{snowballOfMod}(M)
Proof =
```

#### 1.11 Graded Modules

```
{\tt LeftGradedModule} \, :: \, \prod \Delta : {\tt CommutativeMonoid} \, . \, \, \prod (R,H) \in {\sf GRING}(\Delta) \, .
             . ? \sum M \in R\text{-}\mathsf{MOD} : \Delta \to \mathtt{LeftSubmodule}(M)
(M,S): \texttt{LeftGradedModule} \iff (M,\Delta,S): \texttt{GradedAbelean} \ \& \ \forall a,b \in \Delta \ . \ H_aS_b \subset S_{a+b}
{\tt RightGradedModule} \, :: \, \prod \Delta : {\tt CommutativeMonoid} \, . \, \, \prod (R,H) \in {\sf GRING}(\Delta) \, .
             . ? \sum M \in \mathsf{MOD}\text{-}R : \Delta \to \mathtt{RightSubmodule}(M)
(M,S): RightGradedModule \iff (M,\Delta,S): GradedAbelean & \forall a,b\in\Delta. S_bH_a\subset S_{a+b}
{\tt GradedModule} \, :: \, \prod \Delta : {\tt CommutativeMonoid} \, . \, \, \prod (R,H) \in {\sf GRING}(\Delta) \, .
             . ? \sum M \in R	ext{-MOD} \ \& \ \mathsf{MOD}	ext{-}R . \Delta \to \mathtt{Submodule}(M)
(M,S): GradedModule \iff (M,\Delta,S): GradedAbelean & \forall a,b \in \Delta. H_aS_b \cup S_bH_a \subset S_{a+b}
{\tt HomogeneousSubmodule} :: \forall (M,S) : {\tt LeftGadedModule}(\Delta,R,H) . \ \forall [0] : (\Delta : {\tt Cancelable}) \ .
              \forall \delta \in \Delta : M_{\delta} : R_0-MOD
Proof =
  . . .
   \texttt{GradedLeftModuleHomo} \; :: \; \prod(M,S), (M',S') : \texttt{LeftGradedModule}(\Delta,R,H) \; . \; ?M \xrightarrow{R\texttt{-MOD}} M'
f: {\tt GradedLeftModuleHomo} \iff \exists \delta \in \Delta: \ . \ \forall \alpha \in \Delta \ . \ f\left(S_{\alpha}\right) \subset S_{\alpha+\delta}
\texttt{GradedRightModuleHomo} \; :: \; \prod(M,S), (M',S') : \texttt{RightGradedModule}(\Delta,R,H) \; . \; ?M \xrightarrow{\texttt{MOD-}R} M' = (M,S), (M',S') : \texttt{RightGradedModule}(\Delta,R,H) \; . \; ?M \xrightarrow{\texttt{MOD-}R} M' = (M,S), (M',S') : \texttt{RightGradedModule}(\Delta,R,H) \; . \; ?M \xrightarrow{\texttt{MOD-}R} M' = (M,S), (M',S') : \texttt{RightGradedModule}(\Delta,R,H) \; . \; ?M \xrightarrow{\texttt{MOD-}R} M' = (M,S), (M',S') : \texttt{RightGradedModule}(\Delta,R,H) \; . \; ?M \xrightarrow{\texttt{MOD-}R} M' = (M,S), (M',S') : \texttt{RightGradedModule}(\Delta,R,H) \; . \; ?M \xrightarrow{\texttt{MOD-}R} M' = (M,S), (M',S') : \texttt{RightGradedModule}(\Delta,R,H) \; . \; ?M \xrightarrow{\texttt{MOD-}R} M' = (M,S), (M',S') : \texttt{RightGradedModule}(\Delta,R,H) \; . \; ?M \xrightarrow{\texttt{MOD-}R} M' = (M,S), (M',S') : \texttt{RightGradedModule}(\Delta,R,H) \; . \; ?M \xrightarrow{\texttt{MOD-}R} M' = (M,S), (M',S') : \texttt{RightGradedModule}(\Delta,R,H) \; . \; ?M \xrightarrow{\texttt{MOD-}R} M' = (M,S), (M,S) : \texttt{RightGradedModule}(\Delta,R,H) \; . \; ?M \xrightarrow{\texttt{MOD-}R} M' = (M,S), (M,S) : \texttt{RightGradedModule}(\Delta,R,H) \; . \; ?M \xrightarrow{\texttt{MOD-}R} M' = (M,S) : \texttt{RightGradedModule}(\Delta,R,H) \; . \; ?M \xrightarrow{\texttt{MOD-}R} M' = (M,S) : \texttt{RightGradedModule}(\Delta,R,H) \; . \; ?M \xrightarrow{\texttt{MOD-}R} M' = (M,S) : \texttt{RightGradedModule}(\Delta,R,H) \; . \; ?M \xrightarrow{\texttt{MOD-}R} M' = (M,S) : \texttt{MOD-}R : \texttt
f: {\tt GradedRightModuleHomo} \iff \exists \delta \in \Delta: \ . \ \forall \alpha \in \Delta \ . \ f\Big(S_\alpha\Big) \subset S_{\alpha+\delta}
\texttt{degreeOfMorphism} :: \prod (M,S), (M',S') : \texttt{GradedLeftModule}(R,\Delta,H) \; . \; \texttt{GradedLeftModuleHomo}(R,\Delta,H) \to \texttt{GradedLeftModuleHomo}(R,\Delta,H) \; . \; \; \texttt{GradedLeftModuleHomo}(R,A,H) \; . \; \; \texttt{GradedLeftModule
\operatorname{degreeOfMorphism}(\varphi) = \operatorname{deg} \varphi := \operatorname{\eth GradedRightModuleHomo} RH\Delta(M,S)(M',S')(\varphi)
\texttt{degreeOfMorphism} :: \prod (M,S), (M',S') : \texttt{GradedRightModule}(R,\Delta,H) \; . \; \texttt{GradedRightModuleHomo}(R,\Delta,H) \; . \; \\
\operatorname{degreeOfMorphism}(\varphi) = \operatorname{deg} \varphi := \eth \operatorname{GradedRightModuleHomo} RH\Delta(M,S)(M',S')(\varphi)
degreeOfMorphismComp :: \forall (M, S), (M', S'), (M'', S'') \in GradedLeftModule.
              \forall \varphi : \mathtt{GradedLeftModuleHomo}(R, \Delta, H)(M, S)(M', S').
              . \forall \psi : \mathtt{GradedLeftModuleHomo}(R, \Delta, H)(M', S')(M'', S'') . \deg \varphi \psi = \deg \varphi + \deg \psi
Proof =
  . . .
```

```
categoryOfGradedLeftModules :: GRING \rightarrow CAT
 {\tt categoryOfGradedLeftModules}\ (R,\Delta,H) = (R,H) \text{-}{\sf GMOD}(\Delta) := {\tt categoryOfGradedLeftModules}\ (R,\Delta,H) = (R,H) \text{-}{\sf CategoryOfGradedLeftModules}\ (R,L) + (R,H) + (R,H) + (R,H) + (R,H) + (R,H) + (R,H) + (
                  := \Big( \mathtt{GradedLeftModules}(R, \Delta, H), \Big)
              \{f: \mathtt{GradedLeftModuleHomo}(R,\Delta,H): \deg f=0\}, \circ, \mathrm{id} \}
 categoryOfGradedRightModules :: GRING \rightarrow CAT
 categoryOfGradedRightModules(R, \Delta, H) = GMOD(\Delta) - (R, H) :=
                 := \Big( \mathtt{GradedRightModules}(R, \Delta, H), \Big)
              \{f: \mathtt{GradedRightModuleHomo}(R,\Delta,H): \deg f=0\}, \circ, \mathrm{id} \ 
{\tt GradedSubmodule} :: \prod (M,S) \in (R,H) - {\sf GMOD}(\Delta) . ? (R,H) - {\sf GMOD}(\Delta)
(N,Z): \texttt{GradedSubmodule} \iff (N,Z) \subset_{(R,H)\text{-}\mathsf{GMOD}(\Delta)} (M,S) \iff N \subset M \ \& \ \forall \delta \in \Delta \ . \ Z_{\delta} = S_{\delta} \cap N = S_{\delta} \cap S_{\delta
{\tt GradedSubmoduleByHomogeneousPart} :: \forall (M,S), (N,Z) : (R,H) - {\sf GMOD}(\Delta) .
                  (M,S) \subset_{(R,H)\text{-GMOD}(\Delta)} (N,Z) \iff \forall n \in N : \forall \delta \in \Delta : n_{\delta} \in N
Proof =
    . . .
    GeneratedGratedSubmodule :: \forall (M, S), (N, Z) : (R, H) \text{-GMOD}(\Delta).
                  A:(N,Z)\subset_{(R,H)\text{-}\mathsf{GMOD}(\Delta)}(M,S)\iff \exists A:?\mathtt{Homogeneous}(M).\ N=\mathrm{span}(A)
Proof =
   . . .
    HomogeneousGeneration :: \forall (M,S) : (R,H) \text{-}\mathsf{GMOD}(\Delta) . \forall (N,Z) \subset_{(R,H)\text{-}\mathsf{GMOD}(\Delta)} (M,S).
                  . \forall A : \texttt{Generating}(N) . \{a_{\delta} | a \in A, \delta \in \Delta\} : \texttt{Generating}(N)
Proof =
    . . .
    FiniteHomogeneousGeneration :: \forall (M, S) : (R, H) \text{-GMOD}(\Delta).
                  \forall [0] : (M : \texttt{FinitelyGeneratedModule}(R)) : \exists F : ?\texttt{Homogeneous} \& \texttt{Finite} \& \texttt{Generating}(M, \Delta, H)
Proof =
    {\tt GradedQuotient} :: \forall (M,S) : (R,H) - {\sf GMOD}(\Delta) \; . \; \forall (N,Z) \subset_{(R,H) - {\sf GMOD}(\Delta)} (M,S) \; .
                \left(\frac{M}{N}, \left(\frac{S_{\delta}+N}{N}\right)_{S\in \Delta}\right): (R,H)\operatorname{-\mathsf{GMOD}}(\Delta)
Proof =
```

```
{\tt GradedImage} \, :: \, \forall (M,S), (M',S') : (R,H) \, - \, \mathsf{GMOD}(\Delta) \, \, . \, \, \forall \phi : (M,S) \xrightarrow{(R,H) - \, \mathsf{GMOD}(\Delta)} (M',S') \, \, .
                     (\phi(M),\phi(S))\subset_{(R,H)\text{-}\mathsf{GMOD}(\Delta)}(M',S')
 Proof =
    . . .
    {\tt GradedKernel} \, :: \, \forall (M,S), (M',S') : (R,H) \, - \, {\sf GMOD}(\Delta) \, \, . \, \, \forall \phi : (M,S) \xrightarrow{(R,H) - \, {\sf GMOD}(\Delta)} (M',S') \, \, .
                     . \forall [0] : \deg \phi : \mathtt{Cancelable}(\Delta) . (\ker \varphi, \ker \varphi \cap S) \subset_{(R,H)\text{-}\mathsf{GMOD}(\Delta)} (M,S)
    . . .
    \textbf{GradedIsomorphismTheorem} \ :: \ \forall (M,S), (M',S') : (R,H) \text{-}\mathsf{GMOD}(\Delta) \ . \ \forall \phi : (M,S) \xrightarrow{(R,H)\text{-}\mathsf{GMOD}(\Delta)} (M',S') \ .
                    . \ \forall [0] : \deg \phi = 0 \ . \ \exists \psi : \left(\frac{M}{\ker \phi}, \frac{S + \ker \phi}{\ker \phi}\right) \xleftarrow{(R, H) \text{-}\mathsf{GMOD}(\Delta)} \left(\mathrm{Im}(\phi), \mathrm{Im}(\phi) \cap S'\right) : \forall x \in M \ . \ \psi[x] = \phi(x)
 Proof =
    . . .
    {\tt GradedSum} :: \forall (M,S) : (R,H) - {\sf GMOD}(\Delta) . \forall I \in {\sf SET} .
                     . \ \forall (N,Z): I \rightarrow \mathtt{GradedLeftSubmodule}(R,\Delta,H)(M,S) \ . \ \left(\sum_{i \in I} N_i, \sum_{i \in I} Z_i\right): (R,H) - \mathsf{GMOD}(\Delta)
 Proof =
    . . .
    {\tt GradedIntersect} :: \forall (M,S) : (R,H) - {\sf GMOD}(\Delta) . \forall I \in {\sf SET} .
                    . \ \forall (N,Z): I \rightarrow \mathtt{GradedLeftSubmodule}(R,\Delta,H)(M,S) \ . \ \left(\bigcap_{i \in I} N_i, \bigcap_{i \in I} Z_i\right) : (R,H) \text{-}\mathsf{GMOD}(\Delta)
 Proof =
    . . .
     {\tt GradedAnnihilator} :: \forall (M,S) : (R,H) - {\tt GMOD}(\Delta) \; . \; \forall \delta : {\tt Cancelable}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Gr
 Proof =
    . . .
     {\tt GradedAnnihilator} :: \forall (M,S) : (R,H) - {\tt GMOD}(\Delta) \; . \; \forall \delta : {\tt Cancelable}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Ann}(x) : {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt GradedLeftIde}(\Delta) \; . \; \forall x \in S_\delta \; . \; {\tt Gr
 Proof =
    . . .
```

```
\begin{aligned} & \texttt{GradedAnnihilator} :: \, \forall (M,S) : (R,H) \text{-}\mathsf{GMOD}(\Delta) \; . \; \forall [0] : \left(\Delta : \mathtt{Cancelable}(\Delta)\right) \; . \\ & . \; \; \mathsf{Ann}(M) : \mathtt{GradedLeftIdeal}(R,\Delta,H) \end{aligned} \mathsf{Proof} \; = \qquad \cdots
```

# 2 Linear Operators and Matrices

#### 2.1 Matrices as Tables

```
\texttt{transpose} \, :: \, \prod X, n, m \in \mathsf{SET} \, . \, X^{n \times m} \to X^{m \times n}
\mathtt{transpose}\left(A\right) = A^{\top} := \Lambda(j,i) \in m \times n \; . \; A_{i,j}
\texttt{Symmetric} :: \prod X, n \in \mathsf{SET} \:.\: ?X^{n \times n}
A: \mathtt{Symmetric} \iff A^{\top} = A
\texttt{column} \, :: \, \prod X, n, m \in \mathsf{SET} \, . \, X^{n \times m} \to m \to X^n
\mathtt{column}\left(A,j
ight) = \mathcal{C}_{j}(A) := \Lambda i \in n \;.\; A_{i,j}
\texttt{row} \, :: \, \prod X, n, m \in \mathsf{SET} \, . \, X^{n \times m} \to n \to X^m
row(A, i) = \mathcal{R}_i(A) := \Lambda j \in m \cdot A_{i,j}
fromColumns :: \prod X, n, m \in SET : m \to X^n \to X^{n \times m}
fromColumns(C) := \Lambda(i,j) \in n \times m . C_i(j)
from Rows :: \prod X, n, m \in \mathsf{SET} : n \to X^m \to X^{n \times m}
fromRows(R) := fromColumns(R)^{\top}
diagonal :: \prod X, n \in \mathsf{SET} : X^{n \times n} \to X^n
\operatorname{diagonal}(A) = \operatorname{diag}(A) := \Lambda i \in n . A_{i,i}
\texttt{Submatrix} \; :: \; \prod X, n, m, n', m' \in \mathsf{SET} \; . \; X^{n \times m} \to ?X^{n' \times m'}
A': \mathtt{Submatrix} \iff \Lambda A \in X^{n \times m} . \exists k: n' \hookrightarrow n: \exists l: m' \hookrightarrow m . A' = A_{k,l}
\texttt{fromBlocks} \, :: \, \prod X, n, m, k, l \in \mathsf{SET} \, . \, \prod p : \mathsf{Partition}(n,k) \, . \, \prod q : \mathsf{Partition}(m,l) \, .
     . \left(\prod(i,j) \in k \times l : X^{p_i \times q_j}\right) \to X^{n \times k}
fromBlocks (B) := \Lambda(i,j) \in n \times m. B_{i,j}(i',j')
        where i' = \eth Partition(n, k)(p)(i), j' = \eth Partition(m, l)(q)(j)
\texttt{TransposeIsConvolution} \ :: \ \forall X, n, m \in \mathsf{SET} \ . \ \forall A \in A^{n \times m} \ . \ A^{\top \top} = A
Proof =
 . . .
 \texttt{blockDiagonal} :: \prod R \in \mathsf{RING} \;. \; \prod n, k \in \mathsf{SET} \;. \; \prod p : \mathsf{Partition}(n,k) \;. \; \left(\prod i \in k \;. \; X^{p_i \times p_i}\right) \to X^{n, \times n}
\texttt{blockDiagonal}\,(A) := \texttt{fromBlocks}(X, n, m, k, l)(p, p)(\Lambda(i, j) \in n \times n \;.\; \texttt{if}\; i == j \;\texttt{then}\; A_{i, j} \;\texttt{else}\; 0)
```

## 2.2 Elementary Matrix algebra

```
\texttt{matrixMult} :: \prod R \in \mathsf{RING} \;. \; \prod n, m, k \in \mathbb{Z} \;. \; R^{n \times m} \times R^{m \times k} \to R^{n \times k}
\mathtt{matrixMult}(A, B) = AB := \Lambda(i, j) \in n \times k . \mathcal{R}_i(A)\mathcal{C}_j(B)
\texttt{matrixMultIsAssoc} :: \forall R \in \mathsf{RING} . \ \forall n: 4 \to \mathbb{Z} . \ \forall A \in R^{n_1 \times n_2} . \ \forall B \in R^{n_2 \times n_3} . \ \forall C \in R^{n_3 \times n_4} .
     A(AB)C = A(BC)
Proof =
Assume (i, j) : n_1 \times n_4,
[j,*] := \eth \mathsf{matixMult}((AB),C) \eth \mathcal{R}_i \eth \mathsf{RING}(R) \eth^{-1} C_i(BC) \eth^{-1} \mathsf{matrixMult}:
     : \left( (AB)C \right)_{i,j} = \mathcal{R}_i(AB)\mathcal{C}_j(C) = \sum_{i=1}^{n_3} \mathcal{R}_i(A)\mathcal{C}_t(B)C_{t,j} = \mathcal{R}_i(A)\sum_{i=1}^{n_3} \mathcal{C}_t(B)C_{t,j} =
     = \mathcal{R}_i(A)\mathcal{C}_j(BC) = \left(A(BC)\right)_{i,i};
 \rightsquigarrow [*] := I(=, \rightarrow) : (AB)C = A(BC);
	exttt{MatrixAlgebra} :: \forall R \in \mathsf{RING} \ . \ \forall n \in \mathbb{N} \ . \ \left( R^{n 	imes n}, \mathtt{matrixMult} \right) : \mathsf{AssociativeAlgebra}(R)
Proof =
 . . .
 from
Diag :: \prod R \in \mathsf{RING} \;.\; \prod n \in \mathbb{N} \;.\; R^n \to R^{n \times n}
\mathtt{fromDiag}\left(a\right):=\Lambda(i,j)\in n	imes n . if i==j then a_{i} else 0
identityMatrix :: \prod R \in \mathsf{RING} . \prod n \in \mathbb{N} . R^n
indentityMatrix() = I := fromDiag(i \mapsto 1)
IdentityMatrixIsIdentity :: \forall R \in \mathsf{RING} : \forall n \in \mathbb{N} : \forall A \in R^{n \times n} : IA = A = AI
Proof =
Assume i, j: n \times n,
[j,*] := \eth \mathsf{matrixMult}(i,j) \eth \mathsf{identityMatrix} : (IA)_{i,j} = \mathcal{R}_i(I)\mathcal{C}_j(A) = A_{i,j} = \mathcal{R}_i(A)\mathcal{C}_j(I) = (AI)_{i,j};
 \sim [*] := I(=, \rightarrow) : IA = A = AI,
 \texttt{MatrixAlgebra2} \ :: \ \forall R \in \mathsf{RING} \ . \ \forall n \in \mathbb{N} \ . \ \left(R^{n \times n}, \mathtt{matrixMult}, I\right) : \mathtt{UnitaryAlgebra}(R)
Proof =
 IdentityIsSymmetric :: \forall R \in \mathsf{RING} . \forall n \in \mathbb{N} . I_{R,n} : \mathsf{Symmetric}(R,n,n)
Proof =
 . . .
```

```
\texttt{TransposeIsLinear} :: \ \forall R \in \mathsf{RING} \ . \ \forall n,m \in \mathbb{N} \ . \ \texttt{transpose}(R,n,m) : R^{n \times m} \xrightarrow{R\text{-MOD}} R^{m \times n}
Proof =
. . .
 TransposeIsAlgebraAntihomomorphism :: \forall R \in \mathsf{ANN}.
     . \forall n, m, k \in \mathbb{N} . \forall A \in R^{n \times m} . \forall B \in R^{m \times k} . (AB)^{\top} = B^{\top} A^{\top}
Proof =
Assume (i,j): k \times n,
j.* := \eth \mathtt{transpose}(AB) \eth \mathtt{MatrixMult} \eth^{-2} \mathtt{transpose}(A)(B) \eth \mathsf{ANN}(R) \eth^{-1} \mathtt{matrixMult} :
     : (AB)_{i,j}^{\top} = (AB)_{j,i} = \mathcal{R}_j(A)\mathcal{C}_i(B) = \mathcal{C}_j(A^{\top})\mathcal{R}_i(B^{\top}) = \mathcal{R}_i(B^{\top})\mathcal{C}_j(A^{\top}) = \left(B^{\top}A^{\top}\right)_{i,j};
\sim [*] := I(=m, \to) : (AB)^{\top} = B^{\top}A^{\top};
rowExchange :: \prod R \in \mathsf{RING} . \prod n \in \mathbb{N} . n \times n \to R^{n \times n}
\texttt{rowExchange}\,(i,j) := \texttt{fromRows}(\Lambda k \in n \;.\; \texttt{if}\; k == i \;\texttt{then}\; \delta^j \;\texttt{else}\; \texttt{if}\; k == j \;\texttt{then}\; \delta^i \;\texttt{else}\; \delta^k)
rowScalarMult :: \prod R \in \mathsf{RING} . \prod n \in \mathbb{N} . n \to R^* \to R^{n \times n}
\operatorname{rowScalarMult}(i,\alpha) := \operatorname{fromRows}(\Lambda k \in n \;.\; \text{if}\; k == i \; \text{then}\; \alpha \delta^i \; \text{else}\; \delta^k)
{\tt rowAddScalarMult} \, :: \, \prod R \in {\sf RING} \, . \, \, \prod n \in \mathbb{N} \, . \, n \to n \to R \to R^{n \times n}
rowAddScalarMult (i, j, \alpha) := I + \text{FromRows}(\Lambda k \in n \text{ . if } k == i \text{ then } \alpha \delta^j \text{ else } 0)
ElementaryRowOperation :: \prod R \in \mathsf{RING} \ . \ \prod n \in \mathbb{N} \ . \ ?R^{n \times n}
E: \texttt{ElementaryRowOperation} \iff \exists i,j \in n: \exists \alpha \in R^*: \exists \beta \in R \ . \ E = \texttt{rowExchange}(i,j)
    |E = rowScalarMult(i, \alpha)|
    |E = rowAddScalarMult(i, j, \beta)
\texttt{RowEquivalent} :: \prod R \in \mathsf{RING} \;.\; \prod n,m \in \mathbb{N} \;.\; ?\Big(R^{n \times m} \times R^{n \times m}\Big)
A,B: \texttt{RowEquivalent} \iff \exists N \in \mathbb{Z}_+ \; . \; \exists E:N \to \texttt{ElementaryRowOperations}(R,n): \left(\prod^N E_i\right)A = B
ColumnEquivalent :: \prod R \in \mathsf{RING} . \prod n, m \in \mathbb{N} . ?(R^{n \times m} \times R^{n \times m})
A,B: \texttt{ColumnEquivalent} \iff \exists N \in \mathbb{Z}_+ \ . \ \exists E: N \to \texttt{ElementaryRowOperations}(R,m): A\left(\prod^N E_i\right) = B
\texttt{RowColumnEquivalent} :: \prod R \in \mathsf{RING} \;. \; \prod n,m \in \mathbb{N} \;. \; ? \Big( R^{n \times m} \times R^{n \times m} )
A, B : \texttt{RowColumnEquivalent} \iff \exists N, N' \in \mathbb{Z}_+ : \exists E : N \to \texttt{ElementaryOperations}(R, n) :
     : \exists E': N' \rightarrow \mathtt{ElementaryOperation}(R,m): \left(\prod_{i=1}^N E_i\right) A \left(\prod_{i=1}^N E_i'\right) = B
```

```
SmithNormalForm :: \prod R \in \text{IntegralDomain}. \prod n, m \in \mathbb{N}. ?R^{n \times m}
A: \mathtt{SmithNormalForm} \iff \exists r \in \min(n,m) \ . \ \exists a: \mathtt{Snowball}(R,r) \ .
   A = \texttt{fromColumns}(\Lambda i \in n : \Lambda j \in m : \texttt{if } i == j \& i \leq r \texttt{ then } a_i \texttt{ else } 0)
((X,A): {	t RowColumnEquivalent}(R,n,m))
Proof =
M^{0,0} := X : R^{n \times m}
\vec{\sigma} := \Lambda K \in \min(n, m) \cdot \forall k \in [1, K]_{\mathbb{N}} \cdot \forall j \in n \cdot \forall j' \in m .
    i, j \neq k \Rightarrow M_{i,k}^{K,0} = 0 \& j' \neq k \Rightarrow M_{k,j'}^{K,0} = 0 : \min(n,m) \to \mathsf{Type},
Q := \Lambda K \in \min(n,m) : \forall i \in k : \forall j \in i : M_{i,j}^{K,0} | M_{i,i}^{K,0} : \min(n,m) \to \mathsf{Type},
1.0 := \eth \emptyset \eth \sigma : \sigma(0),
2.0 := \eth \emptyset \eth \varsigma : \varsigma(0),
Assume K : \min(n, m),
Assume [1.(K-1)]: O(K-1),
Assume [2.(K-1)]: Q(K-1),
Assume [3.0]: \exists (i,j) \in [k,n] \times [k,m]: M^{(K-1),0} \neq 0,
v_0 := +\infty : \mathbb{Z}_+ \cup \{+\infty\},\
Assume k : \mathbb{N},
(I, J, [4]) := \texttt{ArgMinExists}([3.k-1])(\texttt{normOfEuclid}(R)(M)) :
    : \sum (I,J) \in [K,n] \times [K,m] \; . \; M_{I,J}^{K,(k-1)} \neq 0 \; \& \; \forall (i,j) \in [K,n] \times [K,m] \; . \; |M_{I,J}^{(K-1),(k-1)}| \leq |M_{i,j}^{(K-1)),(k-1)}|,
N := \texttt{RowExchange}(K, I) M^{(K-1),(k-1)} : R^{n \times m}.
v_k := |N_{K,K}| : \mathbb{Z}_+,
\mathcal{I} := \{i \in [K+1, n] : N_{i,K} \neq 0\} : ?[K+1, n],
\mathcal{J} := \{ j \in [K+1, m] : N_{K,j} \neq 0 \} : ?[K+1, n] \neq 0,
Assume i:\mathcal{I},
\left(a_i, r_i, [i.*]\right) := \eth \texttt{EuclideanRing}(R)(N_{i,K}, N_{K,K}) : \sum a_i \in R \;.\; \sum r_i \in R \;.\; N_{i,K} = a_i N_{K,K} + r_i |r_i| < N_{K,K};
\rightsquigarrow \left(a, r, [k.1]\right) := I\left(\sum\right) : \sum(a, r) : \mathcal{I} \rightarrow R^2 . \forall i \in I . N_{i,K} = a_i N_{K,K} + r_i |r_i| < |N_{K,K}|,
Assume j: \mathcal{J},
\left(b_j,r_j',[j.*]\right) := \eth \texttt{EuclideanRing}(R)(N_{K,j},N_{K,K}) : \sum b_j \in R \;.\; \sum r_j' \in R \;.\; N_{K,j} = b_j N_{K,K} + r_j' |r_j'| < |N_{K,K}|;
\rightsquigarrow (b, r', [k.2]) := I(\sum) : \sum (b, r') : \mathcal{J} \to R^2 . \forall j \in J . N_{K,j} = a_j N_{K,K} + r'_j |r'_j| < |N_{K,K}|,
Assume [5]: \mathcal{I} \neq \emptyset | \mathcal{J} \neq \emptyset,
M^{(K-1),k} := \Big(\prod_{i \in \mathcal{I}} \texttt{rowAddScalarMult}(i,K,-a_i)\Big) N\Big(\prod_{i \in \mathcal{I}} \texttt{rowAddScalarMult})(j,K,-b_k) : R^{n \times m},
[3.k] := \eth N \eth M^{(K-1),k} : M_{K,K}^{(K-1),k} \neq 0;
Assume [5]: \mathcal{I} = \emptyset = \mathcal{J},
```

Assume  $(i, j) : [K + 1, n] \times [K + 1, m],$ 

```
Assume [6]: N_{(K-1),K} / N_{i,j},
M^{(K-1),k} := \texttt{rowAddScalarMult}(K,i,1)N : R^{n \times m},
[3.k] := \eth N[5] \eth M^{(K-1),k} : M_{K,K}^{(K-1),k} \neq 0;
Assume [6]: \forall (i,j) \in [K+1,n] \;.\; N_{K,K}|N_{i,j},
M^{(K-1),k} := N : R^{n \times m}.
[3.k] := \eth N \eth M^{(K-1),k} : M_{K,K}^{(K-1),k} \neq 0;
 \rightsquigarrow \left( M^{(K-1),v,[K.1]} \right) := I\left( \sum \right) : \sum M^{(K-1)} : \mathbb{N} \rightarrow R^{n \times m} \; . \; \sum v : \texttt{Nonincreasing}(\mathbb{N},\mathbb{Z}_+) \; . \; \ldots,
\Big(k,[4]\Big):=\mathtt{\eth}\mathtt{WellOrdered}(\mathbb{Z}_+)(v):\sum k\in\mathbb{N}: \forall k':\mathtt{after}(k)\;.\;v_k=v_k',
[5] := [4][K.1] : \forall k' : \mathbf{after}(k) . M^{(K-1),k'} = M^{(K-1),k')},
M^{K,0} := M^{K-1,k} : R^{n \times m}
1.K := [5][K.1][1.(K-1)] : \mathcal{O}(K),
2.K := [5][K.1][2.(K-1)] : \mathcal{Q}(K);
Assume [3.0]: \forall (i,j) \in [K,n] \times [K,m]. M^{(K-1),(k-1)} = 0,
M^{K,0} := M^{K-1,0} : R^{n \times m}.
1.K := \eth M^{K,0}[3.0][1.K-1] : \sigma(K),
2.K := \eth M^{K,0}[2.0][2.K - 1] : \wp(K);
\sim (M,[1],[2]) := I(\left(\sum\right))I\left(\prod\right) : \sum M : \min(n,m) \to R^{n,m} . \prod k \in \min(n,m) . \sigma(k) \circ (k),
A := M^{\min(n,m)} : R^{n,m}
[3] := \eth^{-1}SmithNormalForm\eth A1.\min(n,m)2.\min(n,m):(A:SmithNormalForm),
[*] := \eth M \eth A : \Big( \left( X, A \right) : \texttt{RowColumnEquivalent}(R, n, m) \Big);
```

## 2.3 Matrices as Linear Operators, Change of Basis

```
\mathtt{matrixOfOperator} :: \prod R \in \mathsf{RING} . \prod A, B : \mathtt{FreeModule}(R) .
         (A \xrightarrow{R\text{-MOD}} B) \to \operatorname{Basis}(A) \to \operatorname{Basis}(B) \to R^{\dim B \times \dim A}
\texttt{matrixOfOperator}\left(T,e,f\right) = T^{e,f} := \Lambda x \in \dim A \;.\; \Lambda y \in \dim B \;.\; r_y \quad \text{where} \quad Ae_x = \sum_{x} r_y f_y = \sum_{x} r_x f_x 
{\tt operatorFromMatrix} \, :: \, \prod R \in {\sf RING} \, . \, \prod A, B : {\tt FreeModule}(R) \, \& \, {\tt FinitelyGeneratedModule}(R) \, .
         .\ R^{\dim B \times \dim A} \to \mathtt{Basis}(A) \to \mathtt{Basis}(B) \to (A \xrightarrow{R\text{-MOD}} B)
\texttt{operatorFromMatrix}\left(M,e,f\right) = M_{e,f} := \Lambda ae \in A \;.\; a_i M_{j,i} f_j
{\tt ChoiceOfBasisDefinesIso} :: \forall R \in {\sf RING} . \forall A, B : {\tt FreeModule}(R) \& {\tt FinitelyGeneratedModule}(R) \ .
         . \ \forall e: \mathtt{Basis}(A) \ . \ \forall f: \mathtt{Basis}(B) \ . \ (\cdot)^{e,f}: R\text{-}\mathsf{MOD}(A,B) \xleftarrow{R\text{-}\mathsf{MOD}} R^{\dim B \times \dim A}
Proof =
 . . .
 \mathtt{MatricesAreLinearMaps} :: \forall R \in \mathsf{RING} . \forall A, B : \mathsf{FreeModule}(R) \& \mathsf{FinitelyGeneratedModule}(R).
         . R	ext{-MOD}(A,B)\cong_{R	ext{-MOD}} R^{\dim B 	imes \dim A}
Proof =
 . . .
 ChoiceOfBasisDefinesAlgIso :: \forall R \in \mathsf{ANN} \ . \ \forall A \in \mathsf{FreeModule} \ \& \ \mathsf{FinitelyGeneratedModule}(R) \ .
       \forall e: \mathtt{Basis}(A) \ . \ (\cdot)^{e,e}: \mathrm{End}_{R\text{-}\mathsf{MOD}}(A) \xleftarrow{R\text{-}\mathsf{ALG}} R^{\dim A \times \dim A}
Proof =
n := \operatorname{rank} A : \mathbb{Z}_+,
Assume T, S : \operatorname{End}_{R\text{-MOD}}(A),
Assume i:n,
(a_{\cdot,i},[1.i]):=\eth \mathtt{Basis}(e)(Te_i):\sum a:n	o A . a_{\cdot,i}e=Te_i;
\rightsquigarrow (a,[1]) := I(\sim) : \sum a : n \to n \to A . \forall i \in n . C_i(a)e = Te_i,
Assume i:n,
(b_{\cdot,i},[2.i]) := \eth \mathtt{Basis}(e)(Se_i) : \sum b : n \to A \;.\; b_{\cdot,i}e = Se_i;
 \rightsquigarrow (b, [2]) := I(\sim) : \sum b : n \rightarrow n \rightarrow A . \forall i \in n . C_i(b)e = Se_i,
Assume i, j:n,
[j.1] := [1.j]EinsteinSummation\partial \mathcal{C}_i \forall k \in n. [2.i]EinstenSummation\partial \mathcal{C}_k(b):
         : STe_j = S\mathcal{C}_j(a)e = Sa_{k,j}e_k = a_{k,j}Se_k = a_{k,j}\mathcal{C}_k(b)e = a_{k,j}b_{t,k}e_t,
[j.*] := \eth \mathtt{matrixOfOperator}[j.1] \eth^{-1} \mathcal{R}_i(b) \mathcal{C}_j(a) :: (ST)_{i,j}^{e,e} = a_{k,j} b_{i,k} = \mathcal{R}_i(b) \mathcal{C}_j(a);
 \sim [3] := I(\forall) : \forall i, j \in n . (ST)_{i,j}^{e,e} = \mathcal{R}_i(b)\mathcal{C}_j(a),
[*] := [3] \eth a \eth b \eth^{-1} \mathtt{matrixOfOperator} : (ST)^{e,e} = S^{e,e} T^{e,e};
```

```
MatricesAreOperators :: \forall R \in \mathsf{RING} \ . \ \forall A : \mathsf{FreeModule}(R) \ \& \ \mathsf{FinitelyGeneratedModule}(R) \ .
    . \operatorname{End}_{R\text{-MOD}}(A) \cong_{R\text{-ALG}} R^{\dim A \times \dim A}
Proof =
. . .
 \verb|changeOfBasisMatrix| :: \prod R \in \mathsf{ANN} \; . \; \prod A : \verb|FreeModule| \; \& \; \verb|FinitelyGeneratedModule| (R) \; .
    . \operatorname{Basis}(A) \to \operatorname{Basis}(A) \to R^{\dim A \times \dim A}
changeOfBasisMatrix (e, f) = C^{e \to f} := \text{coordinate}(f, e)
\texttt{GeneralLinearGroup} = \mathrm{GL} := \Lambda R \in \mathsf{RING} \ . \ \Lambda n \in \mathbb{N} \ . \ \left(R^{n \times n}\right)^* : \mathsf{RING} \to \mathbb{N} \to \mathsf{GRP};
BasisMatrixInvertible :: \forall R \in \mathsf{ANN} : \forall A : \mathsf{FreeModule} \& \mathsf{FinitelyGeneratedModule}(R).
    \forall f, e : \mathtt{Basis}(A) . C^{e \to f} \in \mathrm{GL}(R, \mathrm{rank}\,R)
Proof =
n := \operatorname{rank} A : \mathbb{Z}_+,
E := C^{e \to f} : R^{\operatorname{rank} A \times \operatorname{rank} A}.
[1] := \eth E : \forall i \in n . E_{f,f} f_i = e_i,
Assume v:A,
(\alpha, [v.1]) := \eth \mathsf{Basis}(A)(e)(v) : \sum \alpha \in R^n . v = \alpha e,
[v.*] := \eth R\text{-MOD}(A, A)(E_{f,f})(\alpha f)[1][v.1] : E_{f,f}\alpha f = \alpha e = v;
\sim [2] := \eth^{-1} \text{Surjective} : \left( E_{f,f} : \text{Surjective}(A,A) \right),
Assume v : \ker E_{f,f},
(\alpha, [v.1]) := \eth Basis(A)(f)(v) : \sum \alpha \in \mathbb{R}^n . v = \alpha f,
[v.2] := \eth R\text{-MOD}(A, A)(E_{f,f})(\alpha f)[1][v.1] : 0 = E_{f,f}v = E_{f,f}\alpha f = \alpha e,
[v.*] := \eth \mathtt{Basis}(e)[v.2][v.1] : v = 0;
\sim [3] := \eth \mathtt{Iso}()\mathtt{-MOD}[2] \mathtt{ZeroKernelTHM} : \Big( E_{f,f} : A \overset{R\mathtt{-MOD}}{\longleftrightarrow} A \Big),
[*] := ChoiceOfBasisDefinesAlgIso[3] : E \in GL(R, n);
 . \forall e,f: \mathtt{Basis}(A) \ . \ C^{e \to f}C^{f \to e} = I = C^{f \to e}C^{e \to f}
Proof =
[1] := \eth C \eth^{-1} coordinates : C^{e \to f} C^{f \to e} = coordinate(f, e)coordinates(e, f) = coordinates(e, e) = I,
[2] := \eth C \eth^{-1} \texttt{coordinates} : C^{f \to e} C^{e \to f} = \texttt{coordinate}(e,f) \texttt{coordinates}(f,e) = \texttt{coordinates}(f,f) = I,
[*] := [1][2] : This,
```

```
LinearMapCategory :: RING \rightarrow CAT
 LinearMapCategory() = R-LMAP :=
                = \left(\sum A \in \mathsf{FreeModule} \ \& \ \mathsf{FinitelyGeneratedModule}(R) \ . \ \mathsf{Basis}(A), R\text{-}\mathsf{MOD}, \mathrm{id}, \circ \right)
{\tt MatrixCategory} :: {\sf RING} \to {\sf CAT}
MatrixCategory() = R-MAT:=
                = \Big(\sum A \in {\tt FreeModule} \ \& \ {\tt FinitelyGeneratedModule}(R) \ . \ {\tt Basis}(A),
            ,(A,e),(B,f)\mapsto R^{\operatorname{rank} B\times \operatorname{rank} A},I,\operatorname{\mathtt{matrixMult}})
\mathtt{inCoord} \, :: \, \prod R \in \mathsf{RING} \, . \, R\text{-}\mathsf{LMAP} \overset{\mathsf{CAT}}{\longleftrightarrow} R\text{-}\mathsf{MAT}
 inCoord(A, e) := (A, e)
 inCoord((A, e), (B, f), T) := T^{e,f}
ChangeOfBasis :: \forall R \in \mathsf{RING} \ . \ \prod A, B : \mathsf{FreeModule} \ \& \ \mathsf{FinitelyGeneratedModule}(R) \ .
               . \ \forall T \in A \xrightarrow{R\text{-MOD}} B \ . \ \forall e,e' : \texttt{Basis}(A) \ . \ \forall f,f' : \texttt{Basis}(B) \ . \ T^{e',f'} = C^{f \to f'} T^{e,f} C^{e' \to e'} T^{e,f} T^{e,f} T^{e,f} C^{e' \to e'} T^{e,f} T^{e,f} T^{e,f} T^{e,f} T^{e,f} T^{e,f} T^{e,f} T^{
Proof =
n := \operatorname{rank} A : \mathbb{Z}_+,
m := \operatorname{rank} B : \mathbb{Z}_+,
 Assume i:n,
[i.*] := \eth \texttt{CovariantinCoord}(R\text{-LMAP}, R\text{-MAT}) \eth C^{e' \to e}_{e',e} \eth C^{f \to f'} \eth \texttt{operatorFromMatrix}:
               : \left( C^{f \to f'} T^{e,f} C^{e' \to e} \right)_{{}_{e'} \ f'} e'_i = C^{f \to f'}_{f,f'} T^{e,f}_{e,f} C^{e' \to e}_{e',e} e'_i = C^{f \to f'}_{f,f'} T e'_i = T e'_i;
  \sim [1] := I(=, \rightarrow) : \left(C^{f \rightarrow f'} T^{e, f} C^{e' \rightarrow e}\right)_{e', f'} = T,
[*] := \texttt{ChoiceOfBasisDefinesIso}[1] : C^{f \to f'} T^{e,f} C^{e' \to e} = T^{e',f'};
  EquivalentMatrices :: \prod R \in \mathsf{RING} . \prod n, m \in \mathbb{N} . ?\Big(R^{n \times m} \times R^{n \times m}\Big)
A,B: \texttt{EquivalentMatices} \iff A pprox B \iff \exists T: R^m \xrightarrow{R-\mathsf{MOD}} R^n: \exists e,e': \mathtt{Basis}(R^m): \exists f,f': \mathtt{Basis}(R^n): \exists f,f': \mathtt{Basis}(R^n):
                : T^{e,f} = A \& T^{e',f'} = B
 {\tt InvertibleIsChangeOfBasis} :: \forall R \in {\sf ANN} \ . \ \forall A \in {\tt FreeModule} \ \& \ {\tt FinitelyGeneratedModule}(R) \ .
           \forall M \in \left(R^{\operatorname{rank} A \times \operatorname{rank} A}\right)^* . \exists e, f : \operatorname{Basis}(A) : A = C^{e \to f}
Proof =
 e := FreehasBasis(A) : Basis(A),
 [1] := ChoicOfBasisDefineAlgIso(e, M) : M_{e,e} \in Aut_{R-MOD}(A),
 f := M_{e,e}e : \operatorname{rank} A \to A,
[2] := \eth \operatorname{Aut}_{R\text{-MOD}}(A)(M_{e,e})[1]\eth^{-1}\operatorname{Basis} : (f : \operatorname{Basis}(A)),
 [*] := \eth^{-1}C^{f \to e}[2]\eth f : M = C^{f \to e};
```

```
AltMatrixEquivalence :: \forall R \in \mathsf{ANN} \ . \ \forall n,m \in \mathbb{N} \ . \ \forall A,B \in R^{n \times m} \ .
   A \approx B \iff \exists M \in \operatorname{GL}(R, n) : \exists N \in \operatorname{GL}(R, m) : B = MAN
Proof =
. . .
MatrixeEquivalence :: \forall n, m \in \mathbb{N} : \forall R \in \mathsf{ANN}.
    . EquivalentMatrices(R, n, m): Equivalence(R^{n \times m})
Proof =
Assume A, B, C : \mathbb{R}^{n \times m},
Assume [1]: A \approx B \& B \approx C,
(T, e, e', f, f', [1.1]) := \eth \mathsf{EquivalentMatrices}[1]_1 :
    : \sum T: R^m \xrightarrow{R\text{-MOD}} R^n \;.\; \sum e, e' \texttt{Basis}(R^m) \;.\; \sum f, f': \texttt{Basis}(R^n) \;.\; T^{e,f} = A \;\&\; T^{e',f'} = B,
(S, e'', e''', f'', f''', [1.2]) := \eth \texttt{EquivalentMatrices}[1]_2 :
    : \sum S: R^m \xrightarrow{R\text{-MOD}} R^n \;.\; \sum e'', e''' \texttt{Basis}(R^m) \;.\; \sum f'', f''' : \texttt{Basis}(R^n) \;.\; S^{e'',f''} = B \;\&\; S^{e''',f'''} = C,
[1.3] := \texttt{ChangeOfBasis}([1.1]) : C^{f \rightarrow f'} A C^{e' \rightarrow e} = B,
[1.4] := \texttt{ChangeOfBasis}([2.2]) : C^{f'' \rightarrow f'''}BC^{e''' \rightarrow e''} = C,
[1.5] := [1.3][1.4] : C^{f'' \to f'''} C^{f \to f'} A C^{e' \to e} C^{e''' \to e''} = C,
[*] := ChangeOfBasisInversion(...)AltMatixEquivalence^{-1}(...) : A \approx C;
```

#### 2.4 Determinant and Trace

```
\texttt{matrixDeterminant} \, :: \, \prod R \in \mathsf{ANN} \, . \, \prod n \in \mathbb{N} \, . \, R^{n \times n} \to R
\mathtt{matrixDeterminant}\left(A\right) = \det A := \sum_{\sigma \in S} (-1)^{\sigma} \prod_{i=1}^{n} A_{i,\sigma(i)}
DetTranspose :: \forall R \in \mathsf{ANN} : \forall n \in \mathbb{N} : \forall A \in R^{n \times n} : \det A^{\top} = \det A
Proof =
 . . .
 DetMultilinear :: \forall R \in \mathsf{ANN} : \forall n \in \mathbb{N} : \det \circ \mathsf{fromRows}(n) \in \mathcal{L}_R(n; i \mapsto R^n; R)
Proof =
Assume A: \mathbb{R}^{n \times n}.
Assume v: \mathbb{R}^n,
Assume \alpha:R,
Assume B: R^{n \times n}
Assume i:n,
Assume [1]: \mathcal{R}_i(A) + v = \mathcal{R}_i(B) \& \forall j \in n : i \neq j \Rightarrow \mathcal{R}_j(A) = \mathcal{R}_j(B),
C:=\mathtt{fromRows}\Big(\Lambda j\in n \ . \ \mathtt{if} \ j==i \ \mathtt{then} \ v \ \mathtt{else} \ \mathcal{R}_j(A)\Big): R^{n	imes n},
[1.*] := \eth \det B[1] \eth \mathsf{ANN}(R) \eth^{-1} C :
     : det B = \sum_{\sigma \in S_n} (-1)^{\sigma} \prod_{i=1}^n B_{i,\sigma(i)} = \sum_{\sigma \in S_n} (-1)^{\sigma} (A_{i,\sigma(i)} + v_{\sigma(i)}) \prod_{j=1: j \neq i}^n A_{j,\sigma(j)} =
     = \sum_{\sigma \in S} (-1)^{\sigma} \prod_{i=1}^{n} A_{i,\sigma(i)} + \sum_{\sigma \in S} (-1)^{\sigma} v_{\sigma(i)} \prod_{i=1, i \neq i}^{n} A_{i,\sigma(i)} = \det A + \det C;
 \sim [1] := I(\Rightarrow) : (\ldots),
Assume [2]: \mathcal{R}_i(B) = \alpha \mathcal{R}_i(B) \& \forall j \in n : i \neq j \Rightarrow \mathcal{R}_i(A) = \mathcal{R}_i(B),
[2.*] := \eth \det B[2] \eth \operatorname{Ann}(A) \eth^{-1} \det A :
     : \det B = \sum_{\sigma \in S_n} (-1)^{\sigma} \prod_{i=1}^n B_{i,\sigma(i)} = \sum_{\sigma \in S_n} (-1)^{\sigma} \alpha A_i \prod_{j=1: j \neq i}^n A_{j,\sigma(j)} \alpha \sum_{\sigma \in S_n} (-1)^{\sigma} \prod_{j=1}^n A_{i,\sigma(i)} = \alpha \det A;
 \rightsquigarrow [2] := I(\Rightarrow) : (\ldots)
[*] := \eth^{-1}Multilinear(R, n, i \mapsto R^n, R)[1][2] : \det \circ fromRows \in \mathcal{L}_R(n; i \mapsto R^n; R);
```

```
DetAntisymmetric :: \forall R \in \mathsf{ANN} : \forall n \in \mathbb{N} : \det \circ \mathsf{fromRows}(n) : \mathsf{Antisymmetric}(R, n, i \mapsto R^n, R)
Proof =
Assume A: \mathbb{R}^{n \times n},
Assume \tau: 2\text{-Cycle}(n),
B := fromRows \circ \tau(\mathcal{R}(A)) : R^{n \times n}
[\tau .*] := \eth \det B\eth B\eth sign \eth \mathsf{GRP}(S_n) \eth \mathsf{GRP}(S_n, sign) sign \eth \mathsf{GRP}(S_n) \eth^{-1} \det A :
     : \det B = \sum_{\sigma \in S} (-1)^{\sigma} \prod_{i=1}^{n} B_{i,\sigma(i)} = \sum_{\sigma \in S_n} (-1)^{\sigma} \prod_{i=1}^{n} A_{\tau(i),\sigma(i)} - \sum_{\sigma \in S_n} (-1)^{\sigma} (-1)^{\tau} \prod_{i=1}^{n} A_{i,\sigma\tau(i)} =
     = -\sum_{\sigma \in S_n} (-1)^{\sigma \tau} \prod_{i=1}^n A_{i,\sigma \tau(i)} = -\sum_{\sigma \in S_n} (-1)^{\sigma} \prod_{i=1}^n A_{i,\sigma(i)} = -\det A;
 \sim [*] := \eth^{-1}AntisymmetricByTranspositions : (\det \circ fromRows(n) : Antysymmetric(R, n, i \mapsto R^n, R));
DetHomo1 :: \forall R \in \mathsf{ANN} . \forall n \in \mathbb{N} . \forall A, B \in R^{n \times n} . \det AB = \det A \det B
Proof =
[*] := \eth \det AB \eth \mathcal{C} \eth \mathcal{R} \eth \mathsf{ANN}(R) \eth \mathsf{sign} \eth \mathsf{GRP}(S_n, \mathsf{Signs}) \mathsf{signMultIsBij}(S_n) \eth \mathsf{ANN}(R) \eth^{-1} \det :
     : det AB = \sum_{\sigma \in S} (-1)^{\sigma} \prod_{i=1}^{m} \mathcal{R}_{i}(A) \mathcal{C}_{\sigma(i)}(B) = \sum_{\sigma \in S} (-1)^{\sigma} \prod_{i=1}^{m} \sum_{j=1}^{m} A_{i,j} B_{j,\sigma(i)} =
     = \sum_{\sigma \in S_n} (-1)^{\sigma} \sum_{J: n \to n} \prod_{i=1}^n A_{i, J_i} B_{J_i, \sigma(i)} = \sum_{\sigma \in S_n} (-1)^{\sigma} \sum_{\tau \in S_n} \prod_{i=1}^n A_{i, \tau(i)} \prod_{i=1}^n B_{\tau(i), \sigma(i)} =
     =\sum_{\sigma \in S} \sum_{\tau \in S} (-1)^{\sigma \tau^{-1}} (-1)^{\tau} \prod_{i=1}^{n} A_{i,\tau(i)} \prod_{i=1}^{n} B_{i,\sigma \tau^{-1}(i)} = \sum_{\sigma \in S} \sum_{\tau \in S} (-1)^{\sigma} (-1)^{\tau} \prod_{i=1}^{n} A_{i,\tau(i)} \prod_{i=1}^{n} B_{i,\sigma(i)} =
     = \left(\sum_{i=1}^{n} (-1)^{\sigma} \prod_{i=1}^{n} A_{i,\sigma(i)}\right) \left(\sum_{i=1}^{n} (-1)^{\sigma} \prod_{i=1}^{n} B_{i,\sigma(i)}\right) = \det A \det B;
 Proof =
[1] := \eth I \eth \det : \det I = 1,
[*] := detHomo\eth^{-1}GRP(GL(R, n), R^*) : (det : GL(R, n) \xrightarrow{GRP} R^*);
 {\tt DetBasisInvariant} \ :: \ \forall R \in {\tt ANN} \ . \ \forall M \in {\tt FreeModuleFinitelyGeneratedModule}(R) \ .
     . \forall T \in \operatorname{End}_{R\text{-MOD}}(M) . \forall e, f : \operatorname{Basis}(M) . \det T^{e,e} = \det T^{f,f}
Proof =
[1] := \mathtt{ChangeOfBasis}(T, e, f) : T^{f,f} = C^{e \to f} T^{e,e} C^{f \to e}
[*] := det[1]DetHomo(...) \partial ANN(R)ChangeOfBasisInversion(e, f)DetHomo2(...) :
     : \det T^{f,f} = \det C^{e \to f} T^{e,e} C^{f \to e} = \det C^{e \to f} \det T^{e,e} \det C^{f \to e} = \det C^{e \to f} \det C^{f \to e} \det T^{e,e} = \det T^{e,e}
```

```
\texttt{determinantOfTheOperator} :: \prod R \in \mathsf{ANN} \;. \; \prod M : \texttt{FreeModule} \; \& \; \texttt{FinitelyGeneratedModule}(R) \;.
     . \operatorname{End}_{R\text{-MOD}}(M) \to R
\mathtt{determinantOfTheOperator}\left(T\right) = \det T := \det T^{e,e} \quad \mathtt{where} \quad e = \mathtt{FreeHasBasis}(M)
\forall A, B \in \text{End}_{R\text{-MOD}}(M). det AB = \det A \det B
Proof =
 . . .
 {\tt DetOperatorHomo1} :: \forall R \in {\tt Ann} \ . \ \forall M : {\tt FreeModule} \ \& \ {\tt FinitelyGeneratedModule}(R) \ .
    . det: \operatorname{Aut}_{R\text{-MOD}}(M) \xrightarrow{\mathsf{GRP}} R^*
Proof =
. . .
 antiindex :: n \to (n-1) \to n
\mathtt{antiindex}\,(i) = \hat{i} := \Lambda j \in (n-1) \;. \; \mathtt{if} \; j < i \; \mathtt{then} \; j \; \mathtt{else} \; j+1
\texttt{minor} :: \prod R \in \mathsf{ANN} \:. \: \prod n \in \mathbb{N} \:. \: R^{n \times n} \to (n \times n) \to R
minor(A, i, j) = \Delta_{i,j}(A) := (-1)^{i+j} \det A_{\hat{i},\hat{j}}
DeterminantComputation :: \forall R \in \text{Ann} : \forall n \in \mathbb{N} : \forall A \in \mathbb{R}^{n \times n} : \forall i \in n : \det A = \sum_{i=1}^n A_{i,j} \Delta_{i,j}(A)
Proof =
 . . .
 \texttt{adjointMatrix} :: \prod R \in \mathsf{ANN} \;. \; \prod n \in \mathbb{N} \;. \; R^{n \times n} \to R^{n \times n}
adjointMatrix(A) = adj A := \Delta^{\top}(A)
```

```
CramerMatrixInversion :: \forall R \in \text{Ann} : \forall n \in \mathbb{N} : \forall A \in \text{GL}(R,n) : A^{-1} = \frac{\text{adj } A}{\text{dot } A}
Proof =
Assume i:n,
[1] := \eth \mathsf{matrixMult} \eth \operatorname{adj} A \mathsf{DeterminantComputation}(A) \eth \mathsf{Inverse}(\det A) :
     : \left(\frac{A \operatorname{adj} A}{\det A}\right)_{i,i} = \frac{1}{\det A} \sum_{i=1}^{n} A_{i,j} \Delta_{i,j}(A) = \frac{\det A}{\det A} = 1,
Assume j:n,
Assume [2]: j \neq i,
C:= {\tt fromRows}(\Lambda k \in i \;.\; {\tt if}\; k == j \;{\tt then}\; \mathcal{R}_i(A) \;{\tt else}\; \mathcal{R}_k(A)): R^{n \times n},
[2.*] := \eth \mathsf{matrixMult} \eth \operatorname{adj} A \mathsf{DeterminantComputation}(C) \mathsf{AntisymmetricZero}(\det C) :
     : \left(\frac{A \operatorname{adj} A}{\det A}\right)_{i,i} = \frac{1}{\det A} \sum_{i=1}^{n} A_{i,k} \Delta_{i,k}(A) = \frac{\det C}{\det A} = 0;
\sim [*] := \eth^{-1}I : \frac{A \operatorname{adj} A}{\det A} = I;
\texttt{CramerMatrixInversionCorollary} :: \forall R \in \mathsf{ANN} . \forall n \in \mathbb{N} . \forall A \in R^{n \times n} . \det A \in R^* \Rightarrow A \in \mathrm{GL}(R,n)
Proof =
. . .
 \mathtt{specialLinearGroup} :: \mathsf{ANN} \to \mathbb{N} \to \mathsf{GRP}
specialLinearGroup (R, n) = SL(R, n) := \Lambda R \in Ann . \Lambda n \in \mathbb{N} . \det_{R, n}^{-1}(1)
trace :: \prod R \in \mathsf{RING} . \prod n \in \mathbb{N} . R^{n \times n}R\text{-MOD}R
trace(A) = tr A := A_{i,i}
TraceTransInvariant :: \forall R \in \mathsf{RING} : \forall n \in \mathbb{N} : \forall A \in R^{n \times n} : \operatorname{tr} A = \operatorname{tr} A^{\top}
Proof =
. . .
 TraceProduct :: \forall R \in \text{Ann} : \forall n \in \mathbb{N} : \forall A, B \in R^{n \times n} : \text{tr } AB^{\top} = A_{i,i}B_{i,i}
Proof =
 . . .
 ShiftInTrace :: \forall R \in \text{Ann} : \forall n \in \mathbb{N} : \forall A, B \in R^{n \times m} : \text{tr } AB = \text{tr } BA
Proof =
 . . .
```

```
TraceIsBasisInvariant :: \forall R \in \text{Ann} : \forall M : \text{FreeModule } \& \text{FinitelyGeneratedModule}(R).
    . \forall T \in \operatorname{End}_{R\text{-MOD}}(M) . \forall e,f: \mathtt{Basis}(M) . \operatorname{tr} T^{e,e} = \operatorname{tr} T^{f,f}
Proof =
[1] := \mathtt{ChangeOfBasis}(T, e, f) : T^{f,f} = C^{e \to f} T^{e,e} C^{f \to e},
[*] := \det[1] \texttt{ShiftInTraceChangeOfBasisInversion}(e,f) :
    :\operatorname{tr} T^{f,f}=\operatorname{tr} C^{e\to f}T^{e,e}C^{f\to e}=\operatorname{tr} C^{f\to e}C^{e\to f}\operatorname{det} T^{e,e}=\operatorname{tr} T^{e,e};
{\tt traceOfTheOperator} \, :: \, \prod R \in {\tt Ann} \, \, . \, \, \prod M : {\tt FreeModule\,\&\,FinitelyGeneratedModule}(R) \, .
    . \operatorname{End}_{R\operatorname{\mathsf{-MOD}}}(M) \xrightarrow{R\operatorname{\mathsf{-MOD}}} R
{\tt traceOfTheOperator}\,(T) = \operatorname{tr} T := \operatorname{tr} T^{e,e} \quad {\tt where} \quad e = {\tt FreeHasBasis}(M)
\forall A, B \in \operatorname{End}_{R\text{-MOD}}(M). \operatorname{tr} AB = \operatorname{tr} BA
Proof =
. . .
{\tt DetOfBlockDiagonalMatrix} :: \forall R \in {\tt Ann} \ . \ \forall n,k \in \mathbb{N} \ . \ \forall p : {\tt Partition}(n,k) \ . \ \forall A \in R^{n \times n} \ .
    . \forall B: \prod i \in k . R^{p_i \times p_i} . \forall [0]: A = \texttt{blockDiagonal}(n,k,p,B) . \det A = \prod^n \det B_i
Proof =
. . .
```

## 2.5 Upper and Lower Triangular Matrices

```
UpperTriangularMatrix :: \prod R \in \mathsf{RING} . \prod n \in \mathbb{N} . ?R^{n \times n}
 A: \texttt{UpperTriangularMatrix} \iff \forall i, j \in n : i > j \Rightarrow A_{i,j} = 0
LowerTriangularMatrix :: \prod R \in \mathsf{RING} . \prod n \in \mathbb{N} . ?R^{n \times n}
A: \texttt{LowerTriangularMatrix} \iff \forall i,j \in n : i < j \Rightarrow A_{i,j} = 0
Diagonal
Matrix :: \prod R \in \mathsf{RING} . \prod n \in \mathbb{N} . ?R^{n \times n}
 A: Diagonal Matrix \iff A: Upper Triangular Matrix (R, n) \& A: Lower Triangular Matrix (R, n)
\texttt{DetOfUpperTriangular} \, :: \, \forall R \in \mathsf{ANN} \, . \, \forall A : \texttt{UpperTriangularMatrix}(A) \, . \, \det A = \prod^{n} A_{i,i}
Proof =
  {\tt DetOfLowerTriangular} :: \ \forall R \in {\sf ANN} \ . \ \forall A : {\tt LowerTriangularMatrix}(A) \ . \ \det A = \prod A_{i,i}
Proof =
  . . .
  UpperTriangularizable :: \prod M \in \text{FinitelyGeneratedModule } \& \text{FreeModule}(R) . ? \text{End}_{R\text{-MOD}}(M)
T: 	extsf{UpperTriangularizable} \iff \exists e: 	extsf{Basis}(M): T^{e,e}: 	extsf{UpperTriangularMatix}
\texttt{LowerTriangularizable} :: \prod M \in \texttt{FinitelyGeneratedModule} \ \& \ \texttt{FreeModule}(R) \ . \ ? \\ \texttt{End}_{R\text{-MOD}}(M) \\ \texttt{Module}(R) : \\ \texttt{Modu
 T: \texttt{LowerTriangularizable} \iff \exists e: \texttt{Basis}(M): T^{e,e}: \texttt{LowerTriangularMatix}
Diagonalizable :: \prod R \in \mathsf{RING} . \prod n \in \mathbb{N} . ?R^{n \times n}
 A: Diagonalizable \iff \exists e: \mathtt{Basis}(M): T^{e,e}: \mathtt{DiagonalMatix}
```

## 2.6 Eigenelemments and Simmilarity

```
SimmilarMatrices :: \prod R \in \mathsf{RING} . \prod n \in \mathbb{N} . ?(R^{n \times n} \times R^{n \times n})
(A,B): \mathtt{SimmilarMatrices} \iff A \sim B \iff \exists M \in \mathtt{FreeModule} \ \& \ \mathtt{FinitelyGeneratedModule}(R) \ .
            . \exists T \in \operatorname{Aut}_{R\text{-MOD}}(M) . \exists e, f : \operatorname{Basis}(M) . A_{e,e} = T = B_{f,f}
characteristicIdeal :: \prod R \in \mathsf{ANN} : R^{n \times n} \to \mathsf{Ideal}(R[\mathbb{Z}_+])
\texttt{characteristicIdeal} \ (N) = \mathcal{A}(N) := \mathrm{Ann}_{R[\mathbb{Z}_+]}(N)
\texttt{characteristicIdealOfSimmilarMatricesAgree} \ :: \ \forall R \in \mathsf{ANN} \ . \ \forall n \in \mathbb{N} \ . \ \forall X,Y \in R^{n \times n} \ .
             X \sim Y \Rightarrow \mathcal{A}(X) = \mathcal{A}(Y)
Proof =
Assume f: \mathcal{A}(X),
[f.1] := \eth applyPolynimial(f,Y)[1] \forall i \in \deg f . \texttt{ConjugatePower}(A,X) \eth (-\mathsf{ALG}R)(R^{n \times n})
         \eth^{-1} \text{applyPolynomial}(F,X) \eth \mathcal{A}(X)(f): f(Y) = \sum_{i=0}^{\deg f} f_i Y^i = \sum_{i=0}^{\deg f} f_i \Big(AXA^{-1})^i = \sum_{i=0}^{\deg f} f_i AX^i A^{-1} = \sum_{i=0}^{\deg f} 
            = A\left(\sum_{i=0}^{\deg f} f_i X^i\right) A^{-1} = Af^{-1}(X)A^{-1} = 0,
[f.*] := \partial \mathcal{A}(Y)[f.1] : f \in \mathcal{A}(Y);
 \sim [2] := \eth^{-1}Subset : \mathcal{A}(Y) \subset \mathcal{A}(X),
Assume f: \mathcal{A}(Y),
[f.1] := \eth applyPolynomial(f, X)[1] \forall i \in \deg f. ConjugatePower(A, Y) \eth (-\mathsf{ALG} R)(R^{n \times n})
         \eth^{-1} \text{applyPolynomial}(F,X) \eth \mathcal{A}(Y)(f): f(X) = \sum_{i=0}^{\deg f} f_i X^i = \sum_{i=0}^{\deg f} f_i \Big(A^{-1}YA\big)^i = \sum_{i=0}^{\deg f} f_i A^{-1}Y^i A = \sum_{i=0}^{\deg f}
           = A^{-1} \left( \sum_{i=1}^{\deg f} f_i Y^i \right) A = A^{-1} f^{-1}(Y) A = 0,
[f.*] := \eth \mathcal{A}(Y)[f.1] : f \in \mathcal{A}(X);
 \rightsquigarrow [*] := \eth^{-1} \mathtt{SetEq}[2] : \mathcal{A}(X) = \mathcal{A}(Y),
Eigenvalue :: \prod R \in \mathsf{RING} . \prod M \in R\text{-MOD} . \operatorname{End}_{R\text{-MOD}}(M) \to ?R
\lambda : \mathtt{Eigenvalue} \iff \Lambda T \in \mathtt{End}_{R\text{-MOD}}(M) . \exists m \in M \setminus \{0\} . Tm = \lambda m
Eigenspace :: \prod R \in \mathsf{RING} . \prod M \in R\text{-MOD} . \operatorname{End}_{R\text{-MOD}}(M) \to R \to ?M
m: \mathtt{Eigenspace} \iff \Lambda T \in \mathrm{End}_{R\text{-MOD}}(M) . \Lambda \lambda \in R . \ker \lambda \operatorname{id} - T
\texttt{GeneralizedEigenelement} \ :: \ \prod R \in \mathsf{RING} \ . \ \prod M \in R\text{-}\mathsf{MOD} \ . \ \mathrm{End}_{R\text{-}\mathsf{MOD}}(M) \to R \to ?M
m: \texttt{GeneralizedEigenelement} \iff \Lambda T \in \operatorname{End}_{R\text{-MOD}}(M) \ . \ \Lambda \lambda \in R \ . \ \bigcup_{k=1}^{k} \ker(\lambda \operatorname{id} - T)^k
```

```
characteristic
Polynomial :: \prod R \in \mathsf{ANN} . \prod n \in \mathbb{N} . R^{n \times n} \to R\mathbb{Z}_+
characteristicPolynomial (A) = \chi_A(\lambda) := \det A - \lambda I
characteristicPolynomial2 :: \prod R \in \mathsf{ANN} \cdot M: FreeModule & FinitelyGeneratedModule(R).
   \operatorname{End}_{R\text{-MOD}}(M) \to R\mathbb{Z}_+
\texttt{characteristicPolynomial2}\left(A\right) = \chi_A(\lambda) := \det A - \lambda \operatorname{id}
characteristicPolynomialsOfSimmilarAgree :: \forall R \in \mathsf{ANN} : \forall n \in \mathbb{N} : \forall X, Y \in R^{n \times n}.
    X \sim Y \Rightarrow \chi_X(\lambda) = \chi_Y(\lambda)
Proof =
[*] := \eth \chi_X(\lambda) \eth \mathsf{GRP}(\mathsf{GL}(n,R),R^*)(\det)[1] \eth^{-1} \chi_Y(\lambda) :
    : \chi_X(\lambda) = \det(X - \lambda I) = \det A \det(X - \lambda I) \det A^{-1} = \det \left(AXA^{-1} - \lambda AA^{-1}\right) = \det(Y - \lambda I) = \chi_Y(\lambda);
 {\tt SimmilarOperators} \, :: \, \prod R \in {\sf RING} \, . \, \, \prod M \in R\text{-}{\sf MOD} \, . \, \, ? \big( {\tt End}_{R\text{-}{\sf MOD}}(M) \times {\tt End}_{R\text{-}{\sf MOD}}(M) \big)
X,Y: \texttt{SimmilarOperators} \iff X \sum Y \iff \exists A \in \operatorname{Aut}_{R\text{-MOD}}(M) \; . \; Y = AXA^{-1}
SimmilarOperatorsHaveSameEigenvalues :: \forall R \in RING . \forall M \in R-MOD . \forall (X,Y) : Simmilar(M) .
    . Eigenvalue(X) = Eigenvalue(Y)
Proof =
\left(A,[1]\right):=\eth \mathtt{Simmilar}(X,Y):\sum A\in \mathrm{Aut}_{R\text{-MOD}}(M)\;.\;Y=AXA^{-1},
Assume \lambda: Eigenvalue(X),
ig(m,[\lambda.1]ig):= \eth \mathtt{Eigenvalue}(X): \sum m \in M \ . \ m 
eq 0 \ \& \ Xm = \lambda m,
[\lambda.2] := \eth \operatorname{Aut}_{R\text{-MOD}}(M)(A) \operatorname{ZeroKerTHM}(A) : Am \neq 0,
[\lambda.3] := [1][\lambda.1] \eth R \text{-}\mathsf{ALG}(\mathrm{End}_{R\text{-}\mathsf{MOD}}(M)) : YAm = AXA^{-1}Am = AXm = \lambda Am,
[\lambda.*] := \eth^{-1}Eigenvalue[\lambda.1][\lambda.2] : (\lambda : Eigenvalue(Y));
\sim [2] := \eth^{-1}Subset : Eigenvalue(X) \subset Eigenvalue(Y),
Assume \lambda: Eigenvalue(X),
\big(m,[\lambda.1]\big) := \eth \mathtt{Eigenvalue}(Y) : \sum m \in M \;.\; m \neq 0 \;\&\; Ym = \lambda m,
[\lambda.2] := \eth \operatorname{Aut}_{R-\mathsf{MOD}}(M)(A^{-1}) \operatorname{ZeroKerTHM}(A^{-1}) : A^{-1}m \neq 0,
[\lambda.3] := [1][\lambda.1] \eth R \text{-} \mathsf{ALG}(\mathrm{End}_{R\text{-}\mathsf{MOD}}(M)) : XA^{-1}m = A^{-1}YA^{-1}Am = A^{-1}Ym = \lambda A^{-1}m,
[\lambda.*] := \eth^{-1} \text{Eigenvalue}[\lambda.1][\lambda.2] : (\lambda : \text{Eigenvalue}(X));
\sim [*] := \eth^{-1}SetEq[2] : EigenValue(X) = Eigenvalue(Y);
 PolynomialModuleStructure :: \prod R \in \mathsf{ANN} . \prod M \in R\text{-MOD} . \operatorname{End}_{R\text{-MOD}}(M) \to R[\mathbb{Z}_+]\text{-MOD}
polynomial Module Structure(T) = M_T := Polinimial Module Structure(M, R)
```

```
SimmilarityByPolynomialModuleStructure :: \forall R \in \mathsf{ANN} : \forall M \in R\text{-MOD} : \forall X, Y \in \operatorname{End}_{R\text{-MOD}}(M).
          X \sim Y \iff M_X \cong M_Y
Proof =
Assume [1]: X \sim Y,
\left(A,[1.1]\right):= \texttt{ChangeOfBasis}(\ldots) \eth \texttt{Simmilar}(X,Y): \sum A \in \mathrm{Aut}_{R\text{-MOD}}(M) \; . \; Y = AXA^{-1},
Assume f: R[\mathbb{Z}_+],
Assume m:M,
[f.*] := \eth M_X \eth \operatorname{Aut}_{R-\mathsf{MOD}}(M)(A) \forall i \in \deg f. ConjugatePower(X,A)[1.1] \eth^{-1} M_Y:
           : Af \cdot_{M_X} m = A \sum_{i=0}^{\deg f} f_i X^i m = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A m = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A m = \sum_{i=0}^{\deg f} f_i \Big( A X A^{-1} \Big)^i A m = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A m = \sum_{i=0}^{\deg f} f_i \Big( A X A^{-1} \Big)^i A m = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A m = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A m = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A m = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A m = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A m = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A m = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A m = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A m = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A m = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A m = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A m = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A m = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A m = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^{-1} A M = \sum_{i=0}^{\deg f} f_i A X^i A^
           =\sum_{i=0}^{\operatorname{deg} f} f_i Y^i A m = f \cdot_{M_Y} A m;
  \sim [1.2] := \eth^{-1} \mathsf{Iso} : (A : M_X \overset{R[\mathbb{Z}_+] \text{-MOD}}{\longleftrightarrow} M_Y),
[1.*] := \eth^{-1} \mathbf{Isomorphic}[1.2] : M_X \cong M_Y;
  \rightsquigarrow [1] := I(\Rightarrow) : X \sim Y \Rightarrow M_X \cong M_Y
Assume [2]: M_X \cong M_Y,
A := \eth \mathsf{Isomorphic}[2] : M_X \stackrel{R[\mathbb{Z}_+]-\mathsf{MOD}}{\longleftrightarrow} M_Y,
Assume m:M,
[m.*] := \eth A : AX(m) = Ax \cdot_{M_X} m = x \cdot_{M_Y} Am = YAm;
  \sim [2.1] := I(=, \rightarrow) : AX = YA,
[2.2] := [2.1]A^{-1} : AXA^{-1} = Y,
[2.*] := \eth^{-1} \text{Simmilar}[2.2] : X \sim Y;
  \sim [*] := I(\iff)[1] : X \sim Y \iff A_X \cong A_Y;
```

## 2.7 Elementary Duality

```
\mathtt{dual} \ :: \ \prod R \in \mathsf{RING} \ . \ \mathtt{Contravariant}(R\text{-}\mathsf{MOD}, R\text{-}\mathsf{MOD})
\operatorname{dual}(M) = M^* := \mathcal{M}_{R\text{-MOD}}(M, R)
\mathtt{dual}\,(A,B,T) = T^* := \Lambda f \in B^* \;.\; f \circ T
{\tt dualBasis} \, :: \, \prod R \in {\tt RING} \, . \, \, \prod M : {\tt FreeModule}(R) \, . \, {\tt Basis}(M) \to {\tt rank} \, M \to M^*
\mathtt{dualBasis}(e,i) = e_*^i := \mathtt{free}(\Lambda j \in \mathrm{rank}\,M \cdot \delta_j^i)
DualBasisTHM :: \forall R \in \mathsf{RING} : \forall M \in \mathsf{FreeModule} \& \mathsf{FinitelyGeneratedModule}(R).
    \forall e : \mathtt{Basis}(M) \cdot e_* : \mathtt{Basis}(M^*)
Proof =
n := \operatorname{rank} M : \mathbb{Z}_+,
Assume f: M^*,
Assume m:M,
(\alpha,[1]) := \eth \mathtt{Generating}(e)(m) : \sum \alpha \in R^n \; . \; m = \alpha e,
m.* := [1] \partial M^* \partial e_* \partial M^* [1] : f(m) = f(\alpha e) = \alpha_i f(e_i) = \alpha_i f(e_j) e_*^j (e_i) = f(e_i) e_*^i (m);
\sim [f.*] := I(=,to)) : f = f(e)e_*;
\sim [1] := \eth^{-1}Generatiting : (e_* : \text{Generating}(M^*)),
Assume \alpha: \mathbb{R}^n,
Assume [\alpha.1]: \alpha e_* = 0,
Assume i:n,
[i.*] := [\alpha.1] \eth e_* : 0 = \alpha e_*(e_i) = \alpha_i;
\sim [\alpha.*] := I(\rightarrow, =) : \alpha = 0;
\sim [2] := \eth^{-1}LinearlyIndependent(M^*): (e_*: LinearlyIndependent(M^*)),
[*] := \eth^{-1} Basis[1][2] : (e_* : Basis(M));
M^*: FiniteGroup & FinitelyGeneratedModule(R)
Proof =
. . .
RankOfFreeDual :: \forall R \in \text{Ring} . \forall M \in \text{FiniteGroup} \& \text{FinitelyGeneratedModule}(R) . \text{rank} M^* = \text{rank} M
Proof =
. . .
```

```
DualMapByTranspose :: \forall R \in \mathsf{RING} \ . \ \forall M, N \in \mathsf{FreeModule} \ \& \ \mathsf{FinitelyGeneratedModule}(R) \ .
           . \ \forall T: M \xrightarrow{R\text{-MOD}} N \ . \ \forall e: \mathtt{Basis}(M) \ . \ \forall f: \mathtt{Basis}(N) \ . \ \left(T^{e,f}\right)^\top = \left(T^*\right)^{f_*,e_*}
Proof =
n := \operatorname{rank} N : \mathbb{Z}_+,
m := \operatorname{rank} M : \mathbb{Z}_+,
Assume j:n,
Assume i:m,
[(i,j).*] := \eth T^* \eth^{-1} T^{e,f} \eth f_* \eth \mathcal{C}_i : T^* f_*^j(e_i) = f_*^j (Te_i) = f_*^j \Big( \mathcal{C}_i (T^{e,f}) f \Big) = \Big( \mathcal{C}_i (T^{e,f}) \Big)_i = T_{i,j}^{e,f};
 \sim [*] := I(=, \to) : (T^*)^{f_*, e_*} = (T^{e,f})^\top;
Reflexive :: \prod R \in \mathsf{RING} \cdot ?R\text{-MOD}
M: \mathtt{Reflexive} \iff M \cong_{R-\mathsf{MOD}} M^*
FinitelGenerateFreeModuleIsReflexive :: \forall R \in \mathsf{RING}.
            \forall M \in \mathsf{FreeModule} \ \& \ \mathsf{FinitelyGeneratedModule}(R) \ . \ M : \mathsf{Reflexive}(R)
Proof =
\texttt{naturalEmbedding} :: \prod R \in \mathsf{RING} . \prod M \in R\text{-MOD} . M \xrightarrow{R\text{-MOD}} M^{**}
naturalEmbedding(m) = \omega_M(m) := \Lambda f \in M^* . f(m)
NaturalEmbeddingIsNatural :: \forall R \in \mathsf{RING} \ . \ \omega : \mathtt{NaturalTransform} \big( \mathrm{id}_{R\text{-}\mathsf{MOD}}, (\cdot)^{**} \big)
Proof =
Assume M, N : R-MOD,
Assume T: M \xrightarrow{R\text{-MOD}} N.
Assume m:M,
Assume f:N^*,
\left[(m,f).*\right]:=\eth \mathtt{compose} \eth \omega_N \eth^{-1} T^* \eth^{-1} \omega_M(m) \eth^{-1} T^{**} \eth^{-1} \mathtt{compose}: \left(T\omega_N(m)\right)(f)=\omega_N \big(T(m)\big)(f)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(T(m)\big)=f\big(
  \sim [(M, N).*,] :=: T\omega_N = \omega_M T^**;
 \sim [*] := \eth^{-1} \mathtt{NaturalTransform} : \left(\omega : \mathtt{NaturalTransform}(\mathrm{id}_{R\text{-MOD}}, (\cdot)^* *)\right);
  DoubleDualNaturalIsomorphism :: \forall R \in \mathsf{RING} \ . \ \forall M : \mathsf{FreeModule} \ \& \ \mathsf{FinitelyGeneratedModule} \ .
        \omega_M: M \xrightarrow{R\text{-MOD}} M^{**}
Proof =
  . . .
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- 3 Advanced Categorical Module Theory
- 3.1 Duality