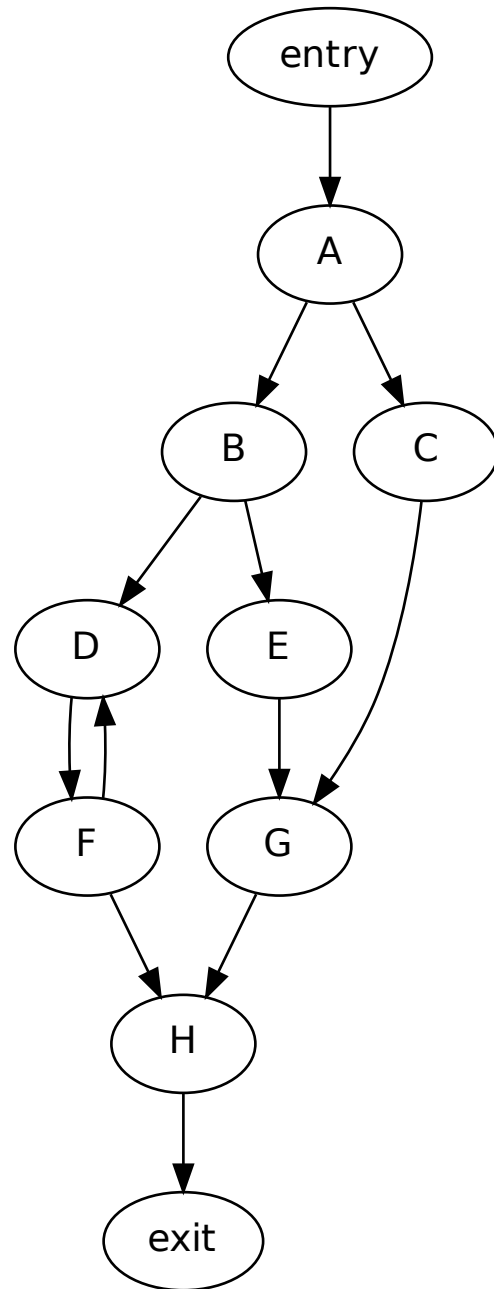


Homework 8 Solutions

1 Control Flow

1.1 Graph



1.2 Dominators Starting With A

	A	B	C	D	E	F	G	H
	A	N	N	N	N	N	N	N
1	A	A,B	A,C	A,B,D	A,B,E	A,B,D,F	A,G	A,H
2	A	A,B	A,C	A,B,D	A,B,E	A,B,D,F	A,G	A,H

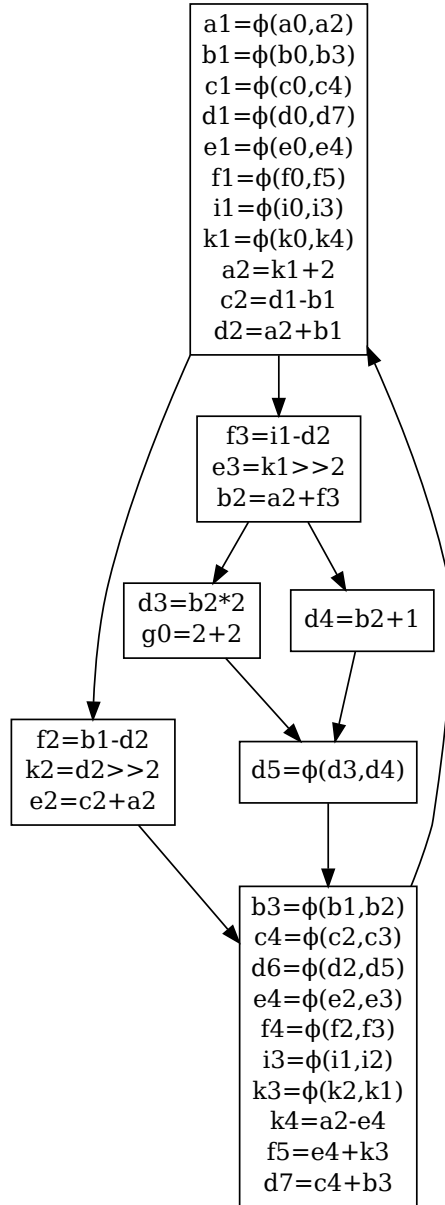
1.3 Dominators Starting With H

	H	G	F	E	D	C	B	A
	H	N	N	N	N	N	N	N
1	H	G	F	E	D	C	B	A
2	H	G	F,D	E,B	D,B	C,A	B,A	A
3	H	G	F,D,B	E,B,A	D,B,A	C,A	B,A	A
4	H	G,A	F,D,B,A	E,B,A	D,B,A	C,A	B,A	A
5	H,A	G,A	F,D,B,A	E,B,A	D,B,A	C,A	B,A	A
6	H,A	G,A	F,D,B,A	E,B,A	D,B,A	C,A	B,A	A

1.4 Dominance Frontier

- $DF(A) = \emptyset$
- $DF(B) = G,H$
- $DF(C) = G$
- $DF(D) = D,H$
- $DF(E) = G$
- $DF(F) = D,H$
- $DF(G) = H$
- $DF(H) = \emptyset$

2 SSA



3 SSA II

3.1 Critical Instructions

writeInt, cbr, jump, exit

3.2 Mark

```
writeInt x1
writeInt y0
jumpI -> B2
loadI 0          => t0
cmp_GTE x0, t0 => t1
cbr t1 -> B3, B4
phi(x0, x2) => x1
loadI 1          => t2
add x1, t2       => x2
loadI 10         => x0
loadI 2          => y0
```

3.3 Rewrite

```
B1:
loadI 10 => x0
loadI 2  => y0
```

```
B2:
phi(x0, x2) => x1
loadI 0      => t0
cmp_GTE x0, t0 => t1
cbr t1 -> B3, B4
```

```
B3:
loadI 1      => t2
add x1, t2   => x2
```

```
B4:
writeInt x1
writeInt y0
exit
```

4 Optimization Passes

4.1 Consequence of a compilation pass that creates code that sometimes behaves differently.

This is a failure of not honoring correctness. Could cause bugs and security vulnerabilities.

4.2 Consequence of pass that creates code that's slower than the original.

Will create a slower program. Could have very undesirable effects on real-time and high-performance systems.

4.3 Acceptable to intentionally have pass that changes the validity of some programs if it preserves correctness of most programs?

No.

4.4 Acceptable to intentionally have pass that produces worse performing code on some programs, but speeds up most programs?

Yes.

5 Machine-Dependant Optimization

Machine-independent optimizations that can be applied to the internal representations of compilers, and should speed up the majority of programs. Machine-dependant optimizations can be applied to machine-specific representations, and the resulting code will only work on machines that share the target's architecture.

6 Analysis/Optimization Scope

6.1 Local Scope

A local analysis scope means the analysis/optimization only takes into account instructions inside a basic block.

6.2 Global Scope

A global scope means the optimizer takes into account the interactions of all instructions.

6.3 Intraprocedure Scope

Intraprocedure scope means the analysis takes into account the interactions of instructions in the same procedure.