

Analysis of Algorithms
Assignment 5
Protik Dey

Ans: to the Ques: No: 1

find-median(A, B, sa, ea, sb, eb):

ma = $\lfloor (sa + ea) / 2 \rfloor$; // median of A

mb = $\lfloor (sb + eb) / 2 \rfloor$; // median of B

if (len(A) == 1 && len(B) == 1)

return (A[0] + B[0]) / 2;

else if (len(A) == 2 && len(B) == 2)

return (max(A[0], B[0]) +

min(A[1], B[1])) / 2;

else if (A[ma] > B[mb])

return find-median(A, B, sa, ma-1, mb+1, eb)

else if (A[ma] < B[mb])

return find-median(A, B, ma+1, ea, sb, mb-1)

else

return A[ma]

Time Complexity: At each step, the problem size

is reduced by half. So the run time is

$O(\log n)$.

Ans: to the Ques: No: 2(a)

A Red-Black tree will have maximum number of internal nodes if it has alternating red and black nodes and if it is a complete binary tree. So the black height will be half of the actual height. Let h = height of the R-B tree. So black height, $b = h/2$
 $\Rightarrow h = 2b$

So the largest possible number of internal nodes is
$$2^{h+1} - 1$$
$$= 2^{2b+1} - 1$$

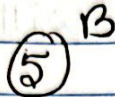
Ans: to the Ques: No: 2(b)

To have the minimum number of internal nodes, the R-B tree will have no red nodes as if there is a red node, there must be two black nodes. So the height of the tree will be black

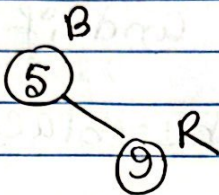
height, $h = b$. So the smallest possible number of internal nodes is $2^h - 1 = 2^b - 1$

Ans: to the Ques: No: 3

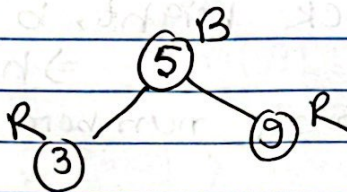
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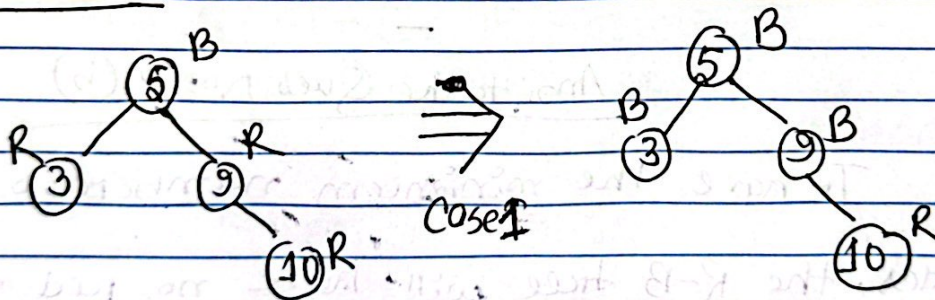
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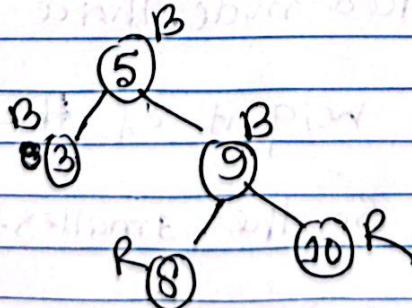
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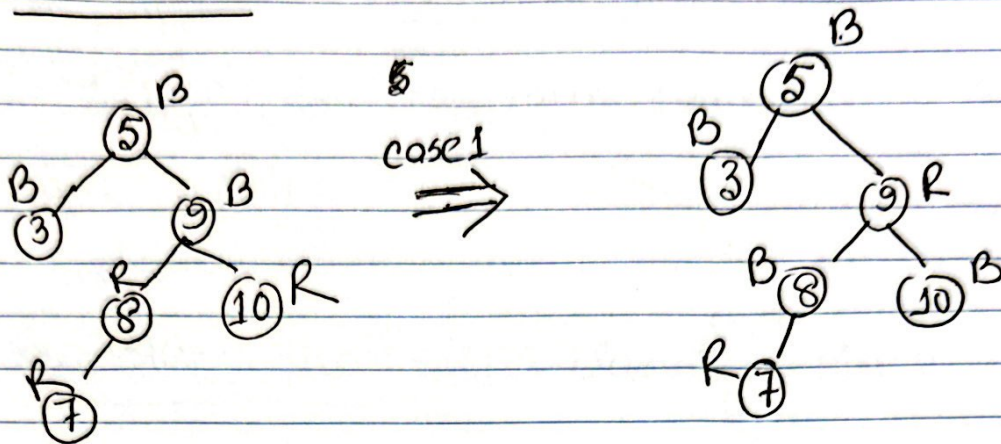
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Insert 8:



Insert 7:



Insert 6:

