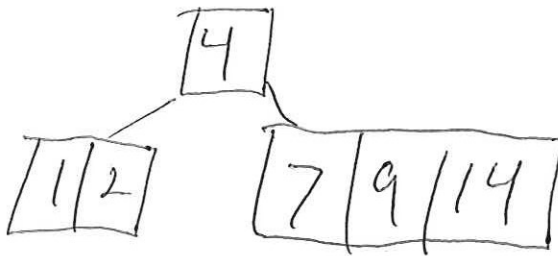


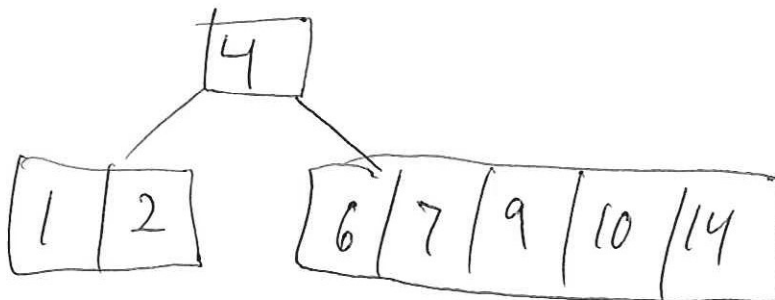
3. First 5 inserts:



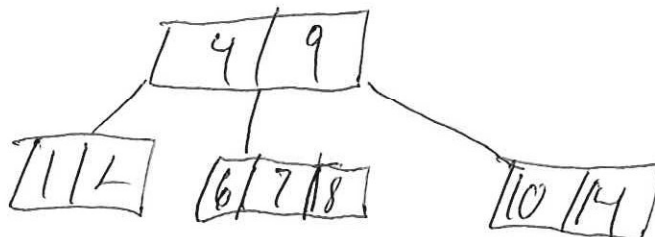
↓ Insert 7



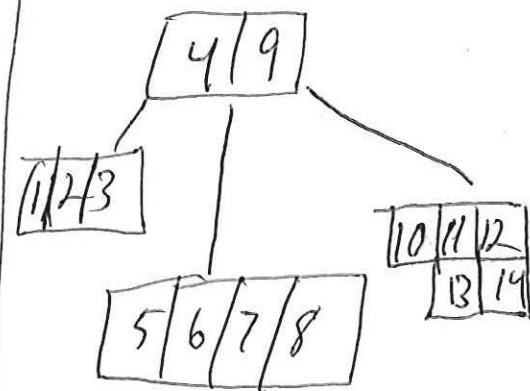
Before Inserting 8:



↓ Insert 8



Final Tree:



③ a) for a set of  $n$  possible locations number of possible placement of toll booths is  $2^n$

for each possible placement running time for calculating sum of money for all the tolls and checking for regulation (can be done alongside) is  $O(n)$

so running time for brute force algorithm is  $O(n 2^n)$ .

b) Recursive definition for  $a[j]$  :

Base case:  $a[0] = 0$

otherwise :  $a[j] = \max \{ a[\text{index of } l(j)] + l[j], a[j-1] \}$

c) DP  $(a, T, L)$  }

$a[0] = 0$ ;

for  $i = 1$  to  $n$

if  $(a[L[i]] + T[i] > a[i-1])$

$a[i] = a[L[i]] + T[i]$ ;

else

$a[i] = a[i-1]$ ;

end for

}

Here,  $L[i] = \text{index of } l(i)$  which can be computed as the following procedure.

```
computeL(a, L) {
```

```
    L[0] = 0;
```

```
    for i = 1 to n
```

```
        for j = i-1 to 1
```

```
            if (a[i] - a[j] ≥ 10) {
```

```
                L[i] = j;
```

```
                break;
```

```
            }
```

```
        end for
```

```
    end for
```

```
}
```

d) The algorithm make one call of the procedure computeL() and another call to DP() where they take  $O(n^2)$  and  $O(n)$  time respectively. Hence overall running time for the algorithm is  $O(n^2)$